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Manufacturing costs can be materially reduced by a production control system which assures that all demands for materials are promptly met with a minimum amount of inventory on hand. Here is a closed-loop control system, based on a computerized inventory model, which has been successfully applied in several industries. Last of a series on the Manufacturing Information and Control System (MIACS) developed by General Electric and now being integrated into Honeywell's FACTOR System (after the recent merger of the two companies' computer businesses).

A Computerized Inventory Model For Production Control

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The first article of this series (Sept. '70, p. 78) described MIACS as an integrated control system using three computer models to completely describe, monitor, and control the operations of a manufacturing or processing facility: a material inventory model, a process model, and a resource inventory model. The second article (Dec. '70, p. 60) described the Integrated Data Store (IDS) file management technique, which permits efficient maintenance of and access to the data required by MIACS modules.

This third and last part will develop the concepts of the materials inventory model, which would logically be the first of the three modules to be implemented in a projected complete system. Called GEIMS (Generalized Inventory Management System), it can by itself make enough contribution to economic control of the enterprise to pay back quite promptly the investment in it.

The GEIMS model calculates the demand for the final products (or replacement parts) on a manufacturing enterprise and converts that demand into detailed schedules of item requirements throughout the plant or factory. These requirements, the production schedules based on them, and the stored inventory data are continuously modified on the basis of reports of actual plant activity fed back to the computer.

The system is designed to reduce production costs and plant investment by:

- (1) Reducing the amount and value of inventory required to meet customer delivery criteria, while
- (2) Reducing the lead times required between customer order and delivery,
- (3) Eliminating equipment and personnel idle time

due to lack of raw materials or in-process inventory,

(4) Reducing shrinkage by close control of inventory and in-process materials.

(5) Timely notification to management of impending delays and out-of-stock conditions, and

(6) Maintenance of product documentation and orderly control of effectivity ("used when" data) when products are modified.

The central element of the system is a model based primarily on two descriptors: the indented parts list for each product, and the lead time required to produce the subassemblies called out at each level of this list. (Other, detailed parameters are such factors as cost data for economic order quantity calculations.)

The organization of a parts list, and the file structure used in MIACS to store such lists efficiently in a computer, were described in some detail in the second article of this series. Briefly, a parts list is a record of the structure of every item—subassembly and final product—inventoried within the enterprise (or the portion of it covered by the model). It has the general form of Figure 1, in which the parts list for item A shows that A consists of 2 of item B, 3 of item C, and 12 of item D. A separate callout chain shows that B, in turn, is made up of items E, F, and G in varying quantities. Links established in the Integrated Data Store create in effect an indented parts list showing the components of every item down to the lowest assembly and piece-part level, and listing all assemblies on which each part and subassembly is used.

Simplified model

The inventory model for a typical product is shown in

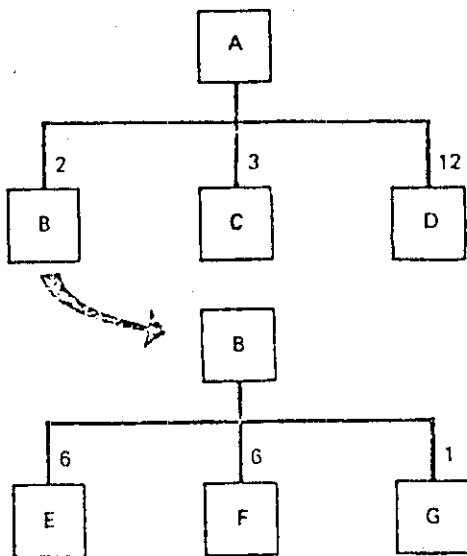


FIG. 1. Typical product structure. Finished product A consists of 2 of item B, 3 of item C, and 12 of item D. Subassembly B, in turn, is made up of items E, F, and G, in quantities of 6, 6, and 1, respectively.

Figure 2, where the simplifying assumption is that the item is not normally inventoried but is manufactured to order. (This assumption will be modified later.) The input to the model, a schedule of demand vs future date, is combined with the lead time required to assemble the final product using the first (highest) level of subassemblies. This generates a schedule of quantities vs time for the start of final assembly.

The latter schedule in turn is applied to the parts list, to generate a schedule of requirements of first-level subassemblies. This is applied to the lead times required to fabricate each first-level subassembly, which generates the starting schedules for those fabrications and the due dates for the next lower-level subassemblies.

This iteration is continued backward in time and downward in level of assembly until a schedule for procuring purchased materials and starting the lowest level of assembly or fabrication is generated. This process is referred to as the "explosion" of a demand into its many detailed intermediate demands.

As engineering changes are made, they are incorporated into the parts list by substituting new detail records at the appropriate levels. If required, both the old and the new parts can remain in the chain, with effectivity notes (either by part number or by effective date) for both. The computer will take effectivity into account when exploding the demand for a product.

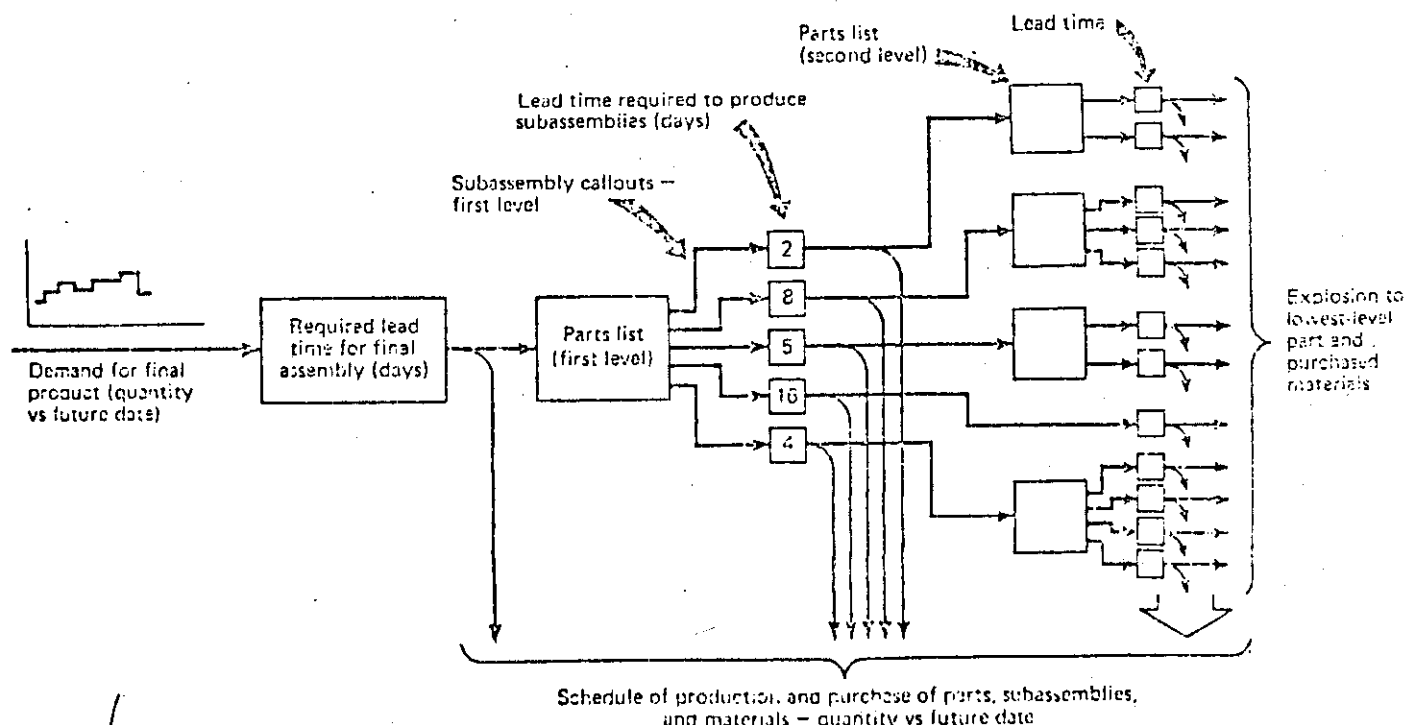
Supply orders

In the case of items which are inventoried, an additional step is introduced into this process: Demand is converted into "supply orders," as illustrated in Figure 3.

Here, "demand" is defined as the requirements (estimated or specifically ordered) from inventory for delivery or as input to production. "Supply" is defined as the amount actually produced by the plant in response to a supply order. Supply is a function of both demand and inventory on hand, as generated by an ordering rule.

This rule is calculated in terms of the economic goals of the enterprise, seeking the optimum tradeoff

FIG. 2. Inventory model generates exploded demand by (at each subassembly level) identifying the component parts from the stored parts list and offsetting the supply orders for each by the time required to make or procure the item.



between costs of manufacture in various lot sizes and the cost of ownership of inventory. Ordering rules may be stored in the inventory system in the simple form shown in Figure 3, in which case the only required inputs are demand and inventory status. Additionally, if the costs of manufacturing and of holding inventory are expected to vary significantly, the system can recalculate the ordering rules from the economic data as new costs are entered into the system.

The application of the ordering rule of Figure 3 for item B is indicated by Figure 4. Assume that two demands for product A are entered on January 10; a requirement for 50 on April 5 and for 30 on April 12. Assume that product A is not inventoried but is assembled in each order quantity for direct delivery to the customer. If final assembly of A requires 15 days, the ordering rule might call for the start of final assembly of each lot 15 days before the respective delivery dates. (The amount of lead time specified by the rule can be made to vary as a function of lot size, if required.)

All first-level parts and subassemblies must be available in inventory on the date final assembly is to start. This creates demands for parts B, C, and D on March 22 and 29. The demand for B is shown in Figure 4. As shown, the future inventory of B is calculated to be reduced below the reorder point by the demand of March 29. This finding triggers the calculation of a supply order for B, of the size and starting date called for in the ordering rule of Figure 3. Specifically, in the case illustrated, the calculation causes a lot of 560 of item B to be started on March 13, to meet a need for 60 on March 29 and to replenish inventory.

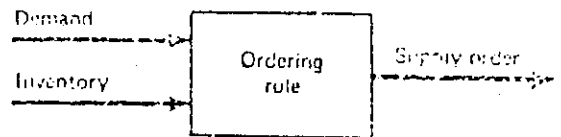
Note that supply orders will not automatically be generated for every item on the parts list of item A—no orders will be generated for items of which there is an adequate inventory to meet the projected demand.

And note also that the ordering rule calls for the consolidation of all demand over a 60-day period in calculating the supply order. The decision to consolidate, and the length of the consolidation period, would normally be based on the economics of lot size and cost of ownership.

Similarly, the ordering rule can incorporate provision for revising the order quantity up or down in response to new demands or cancellations after the order has been generated. Both the magnitude of the revision and the maximum allowable time between the initial order and the revision can be specified in the ordering rule.

Orders to the shop

The flow of information is illustrated in Figure 5. The plant or facility under control is divided into "administrative centers"—logical or functional groupings of resources (perhaps manufacturing departments such as a foundry or machine shop, or in some cases a single machine, such as an injection molding unit)—each assigned a number or code des-

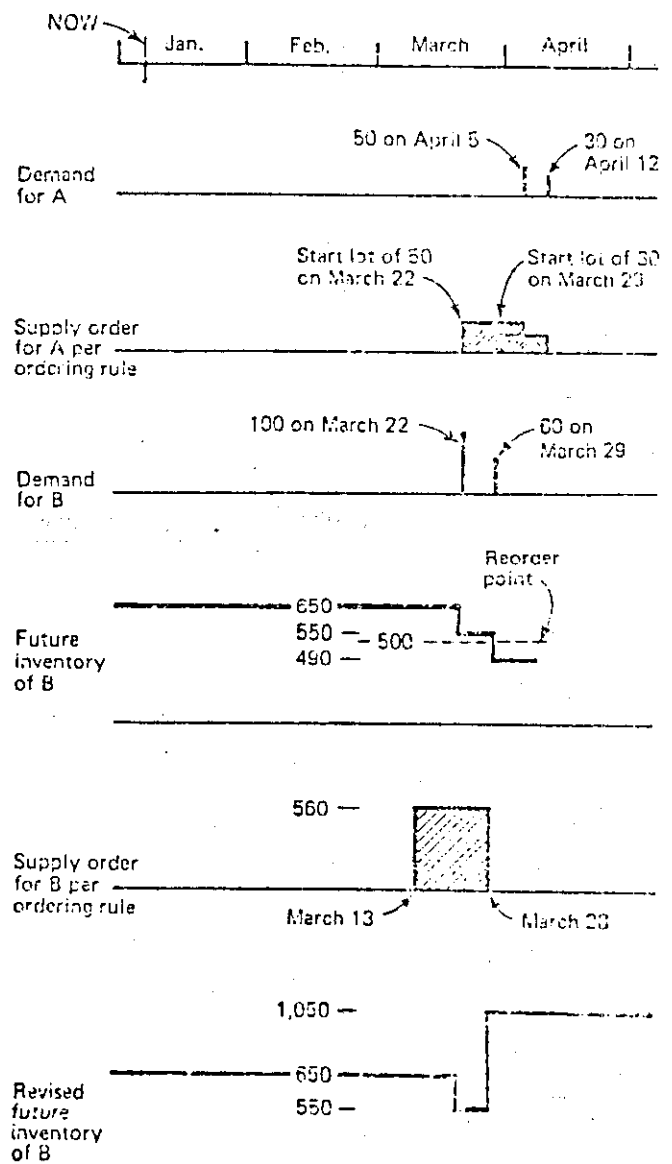


Ordering Rule for Item B

1. Reorder point: 500 items.
2. Due date of order: date inventory on hand is calculated to drop below reorder point, minus one.
3. Calculation of lot size: sum of demand for 60 days following due date, plus 500 items, and minus any items being made for inventory under existing orders in process. Minimum lot size: 500 items.
4. Starting date of work order: due date minus 15 days.

FIG. 3. Supply is a function of demand and inventory, as expressed by an ordering rule. Here is a hypothetical rule for item B of the example. If desired, the parameters determining lot size (step 3) can be recalculated within the model as costs undergo change.

FIG. 4. Demand for final product A generates corresponding demand for subassembly B, earlier in time. The model finds that the second demand will pull future inventory below the reorder point, so the ordering rule generates an order to start production of 560 of B on March 13, to be delivered to inventory on March 28.



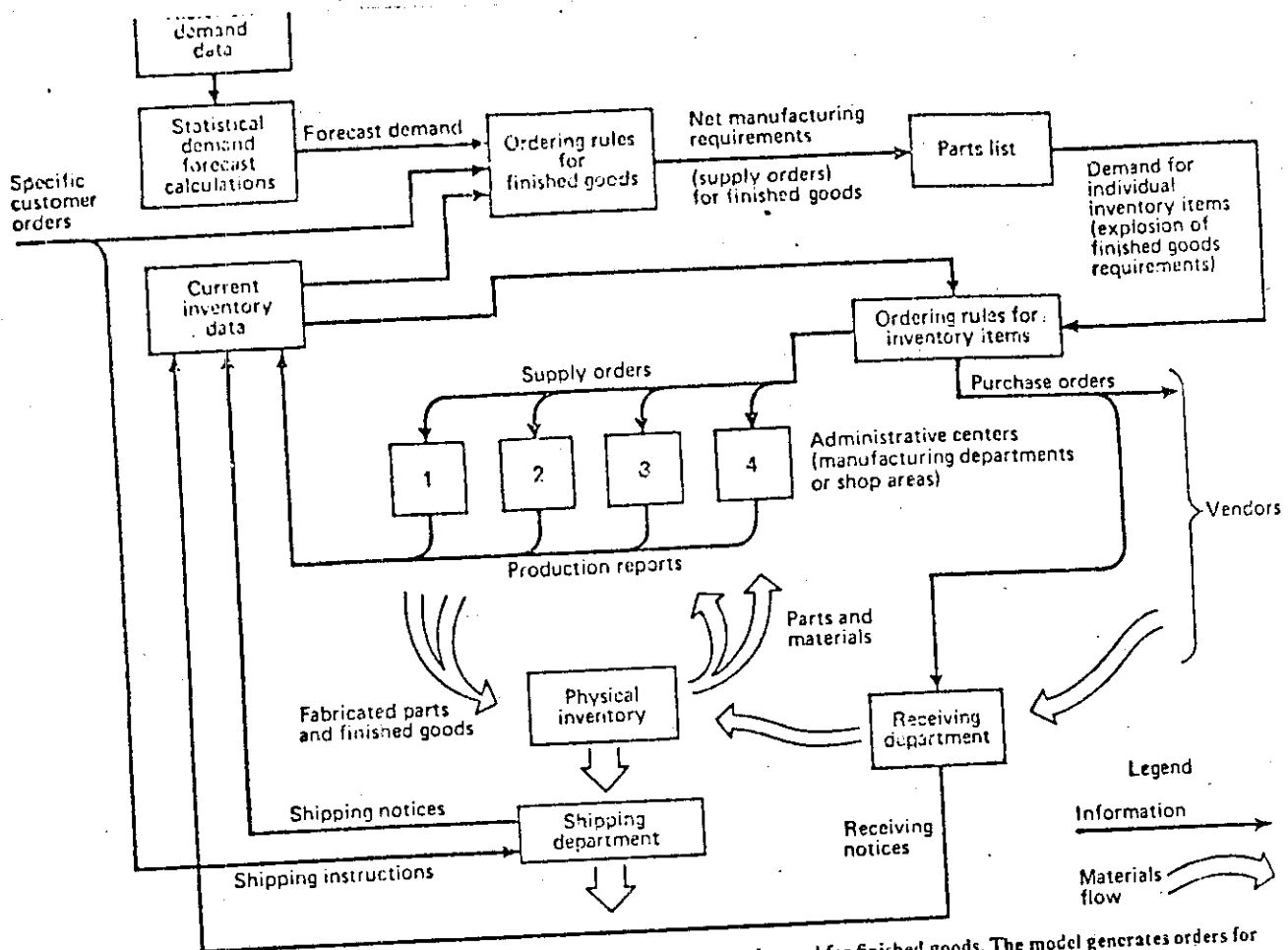


FIG. 5. Flow of information through the inventory control system starts with the demand for finished goods. The model generates orders for the continuous replenishment of inventory. Detail records within the inventory file enable the computer to direct the work orders to the administrative areas responsible for each task. Feedback from production, receiving, and shipping enables the computer to keep track of the inventory quantities used in the model.

ignation within the inventory system.

Using the IDS file structure, it is easy to link to the master for each inventory item the identification of the administrative center or centers responsible for making or procuring the item. Thus as the computer generates supply orders it can direct these orders to the appropriate administrative centers. This means that the entire plant is loaded in an economical and orderly way, with one center producing today the in-process inventory which another center will require tomorrow, without the development either of shortages or of excessive inventory.

As work progresses, actual plant performance is fed back to the computer: Production personnel report on supply orders completed, the receiving department enters arriving goods into the system, and the shipping department reports on goods dispatched out of the facility. This input may be collected manually and entered periodically as a batch, if the computer is a batch system or, if it is an on-line system with remote terminals in the production, receiving, and shipping areas, each transaction can be entered as it occurs.

This feedback makes it simple for the computer program to detect any emerging discrepancies be-

tween schedule and performance, and to generate timely expedite reports to management of impending delays or shortages.

The model can be made to accept inquiries by management or production supervisors and to provide the supply status of any inventory item or the shipping status of any customer order.

Users' experience

The inventory model has been implemented by several industrial users and is producing the expected economic gains. One interesting example is General Electric's Power Production and Conversion Div., Philadelphia, a manufacturer of power generation and distribution equipment. The division has implemented these inventory capabilities on a Honeywell 600 computer, with direct communication to a plant in Burlington, Iowa, via a satellite G-100 computer. Over a two-year period in which plant size doubled, customer delivery cycle has been brought down from 28-32 weeks to 10-12 weeks. The time required for a cycle of product design and manufacturing startup by the engineering department has been reduced by 75 percent—an increase, according to the user's calculations, of 60 percent in engineering productivity. □

[54] **METHOD AND ARRANGEMENT FOR OPTICALLY REPRESENTING INDUSTRIAL MANAGEMENT DATA**

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[51] Int. Cl.**G06f 15/20**

[58] Field of Search.....340/324, 166 EL; 307/311; 313/108; 250/217; 324/114; 235/151.3

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Primary Examiner—Malcolm A. Morrison

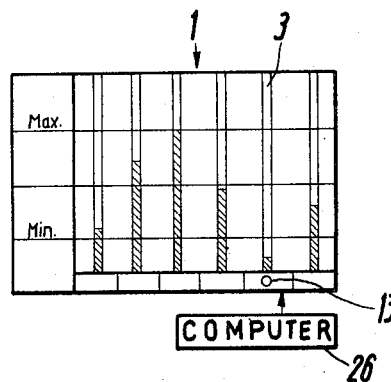
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[57] **ABSTRACT**

Data obtained in industrial production or office work, such as rate or quantity of output of a machine or a worker, are applied to computers for deriving electric signals characterizing the operation with regard to efficiency, profitableness, etc. Indicator boards are provided having electroluminescent strips or discs, or cathode ray tubes the extent of the luminous portions of which are controlled by the computers.

25 Claims, 8 Drawing Figures



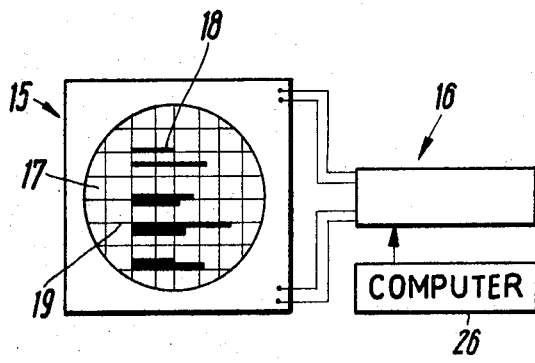


Fig. 6

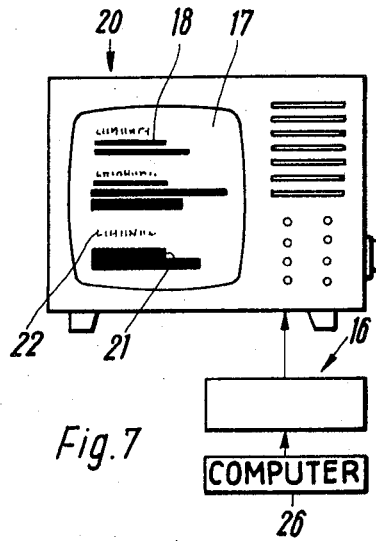


Fig. 7

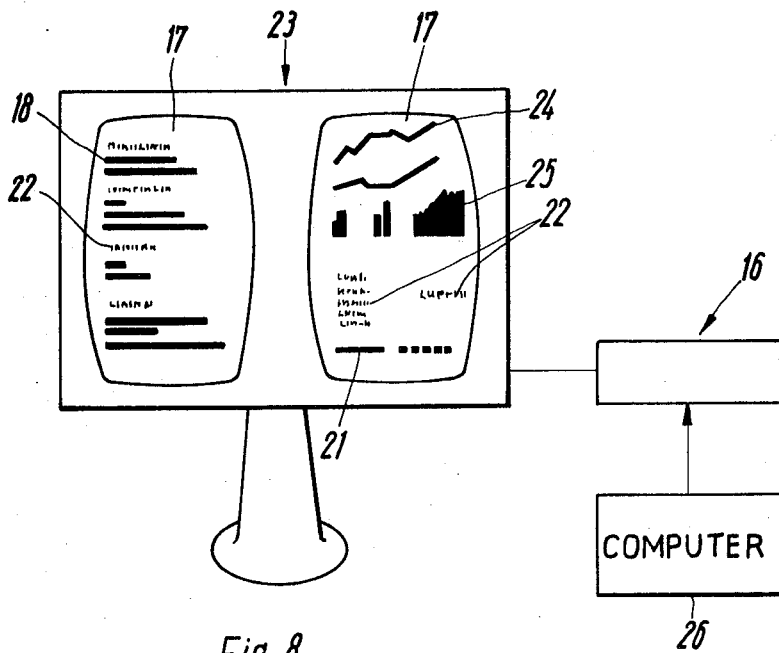


Fig. 8

METHOD AND ARRANGEMENT FOR OPTICALLY REPRESENTING INDUSTRIAL MANAGEMENT DATA

FIELD OF THE INVENTION

The invention relates to a method for optically representing data characterizing a production process and to an arrangement for performing the method.

BACKGROUND OF THE INVENTION

In offices as well as in industrial enterprises, factories, etc., it is imperative to have access to a number of production data, such as the quantities of planned, actually produced and sold products, prices, prime costs, number of work hours, and the like. These values are collected for storage and for further processing. In the interest of clarity, use is made for this purpose of graphic diagrams, tables, charts and the like.

It is previously known to utilize indicator boards, on which are placed strips of paper, which may be in different colors, or magnetic markers. Also known are boards provided with colored tapes which can be pulled off from reels. These methods of representing production data are complex and time-consuming and are not suitable for automatic processing, for instance, in computers.

OBJECTS AND SUMMARY OF THE INVENTION

It is an object of the invention to provide an arrangement and a method for handling industrial or other data, which makes it possible to represent them optically in a rapid and dependable manner with as little use of personnel and time as possible, and to facilitate storage and other processing of such data with the aid of a computer.

According to the invention, data characterizing a production process such as an industrial process or other operation or a quantity which has been derived from such data with the aid of calculators, preferably an electronically controlled calculator, such as a computer, is obtained and transformed into electrical signals. The signals are applied to a digitally controllable optical display unit such as an electroluminescent device, and cause the same to respond to and indicate the magnitudes of the quantities. This method can be carried out very simply and easily with the aid of a digitally controllable optical display unit comprising an electroluminescent device having at least one element which is responsive to produce an electroluminescent output signal of variable magnitude and which responds directly or indirectly to a digital signal-generating means, each element being provided with a scale for the corresponding quantity.

In a modified embodiment of the invention, the digitally controlled optical display unit is a luminescent screen device, such as a conventional TV image tube, the input circuit of which is connected, preferably via an encoder, to the calculator, the control circuits of the tube being arranged to supply electrical signals for displaying the magnitude of a quantity representing a characteristic of an industrial process or an optical signal corresponding to such magnitude.

An advantage of the use of electroluminescent strips consists in that they can be controlled digitally from a computer. Also, adjacent strips produce indications

that are easily compared, and the range of electroluminescent response is capable of relatively fine subdivision.

It is also possible to use as indicators counter-tubes or cathode ray oscillographs, as well as combinations thereof. Since it is generally preferred in office equipment to use analog representation, the invention in its preferred form uses devices including electroluminescent strips.

The representation may also be in the form of strips on a luminescent screen, i.e. on a TV image-reproducing tube. This provides, in addition, a possibility of recording and displaying diagrams, charts, symbols and characters and to illustrate the time variations of processes.

BRIEF DESCRIPTION OF THE DRAWING

The invention will be described below with reference to the drawing, which shows schematically by way of example various embodiments thereof.

FIG. 1 shows an arrangement for representing productivity,

FIG. 2 is a modification thereof,

FIG. 3 shows an arrangement adapted for production planning,

FIG. 4 shows an arrangement for representing production data in connection with cost calculation and marketing,

FIG. 5 is an arrangement adapted for capacity planning,

FIG. 6 is an arrangement comprising an oscillograph,

FIG. 7 is a modified form including a TV apparatus, and

FIG. 8 is an arrangement with two image screens.

DESCRIPTION OF EMBODIMENTS

FIG. 1 shows schematically an arrangement for representing the characteristic quantity "productivity" of an industrial production process. The productivity is a characteristic data or quantity which is derived from other data such as the types and quantities of products as well as the number of working hours used for their production, and it is usually indicated as a percentage. The arrangement comprises an indicator board which may have a number of electroluminescent devices, or cathode ray tubes, such as conventional TV tubes thereon. In the embodiment shown, there are provided a number of electroluminescent strips 3. The luminescent strips may be of the type referred to in the magazine "Flugwelt" 19 (1967), No. 9, page 638, published by Krausskopf-Flugwelt-Verlang GmbH., Mainz, Germany, or in the magazine "Luftfahrt-technik-Raumfahrttechnik", Vol. 13 (1967), No. 11, pp. 276-279, published by VDI Verlag, Dusseldorf, Germany. The strips are responsive to a calculator or computer 26, preferably of the electronic type, and the output signals of which are applied as input signals to the strips. Provided at the lower edge of the board are markings indicating fields or categories of production, such as departments of an enterprise, premium or bonus groups, accounting information, e.g. accounting sections, etc. for the various strips 3. There are supplied manually or automatically to the calculator or computer a number of data and the computer calculates therefrom the corresponding maximum,

minimum, or average productivity values over a predetermined period of time and indicates them on the board. A scale may be provided on the board for enabling reading of numerical values therefrom. Optical alarm signals 13 are provided, which are activated when a quantity decreases below a predetermined value, such as an average or a minimum. The instantaneous values may be correlated with values stored in the calculator and derived from indications of the indicating device during the foregoing comparison intervals. In addition to the optical representation of the production data, it is also possible to provide digital representation, for instance, by means of digital counting tubes.

FIG. 2 shows a modification of the arrangement, in which the indicating devices comprise disc-type elements. According to the magnitude of the quantity to be represented, a greater or lesser sector is rendered luminescent in response to the output signal of the computer.

FIG. 3 shows the combination of two arrangements adapted for production planning. An indicator board 1 is provided with a number of pairs of strips 3a, 3b and with a pair of marginal fields 5, 6. Field 5 contains a tabulation of the machines used in the company or plant and field 6 contains a time scale. Strip 3b may show the total running time of a machine, whereas strip 3a shows the time during which the machine was loaded or not idling. A second board 2 with electroluminescent strips 3 is located adjacent to the first one and is provided in a marginal field 7 with designations of the products produced by the machines and in a second marginal field 8 also with a time scale. The indicating devices of both boards are controlled by a computer 26 to which there may be applied input signals indicative of the quantities of different products via a digital signal generator 9 in combination with an encoder.

FIG. 4 shows an arrangement serving to facilitate cost calculation or analysis and to make it possible to perform such operations more rapidly than was heretofore possible, and to facilitate taking steps to modify marketing procedure. An indicator board 1 is divided into two regions 10 and 11 having corresponding scales for indicating production volume, turnover, sales or the like and for indicating marginal cost 11. The electroluminescent strips 3 are continuous but preferably separate for the two fields and are paired together, each lower pair 3e, 3f being connected to the output side of the corresponding upper pair 3c, 3d. The lower scale may be provided with symbols designating representatives' or agents' districts corresponding to the lower pairs. The strips serve to indicate desired (nominal or theoretical) values as well as actual values of production volume, turnover etc., 3c, 3d, as well as of the marginal cost 3e, 3f and are controlled by a computer. It is obvious that a greater number than two of coordinated strips could be provided, which might differ from each other in width and/or color or the like to facilitate distinguishing them from each other.

FIG. 5 shows another combination of two arrangements which facilitates the capacity planning of an enterprise in a very much simplified and rationalized manner. In this case, also, a pair of indicator boards 1, 2 are provided, one of which has a plurality of pairs of

strips 3g, 3h, as well as a time scale 14. The pairs of strips correspond to locations or operating units or the like, and the strips themselves indicate the desired (nominal or theoretical) and the actual value of the degree of utilization, such as the number of persons or machines assigned to a work, rate of output, or the like of the operating equipment or personnel of the locations or units. The strips are preferably controlled by a computer 26, as are also the strips 3i and 3k of the second board, there being applied to computer 26 the required input values, which are processed therein and transformed into output signals. The second board 2 is also provided with a time scale 12, which may indicate chronologically dates and/or weekdays and/or months. Provided in one field of board 2 are designations of various products corresponding to strips 3i, 3k, the strips showing the desired (nominal) and the actual amount of time spent in production. The embodiment shown possesses only pairs of strips 3i and 3k. However, it is just as possible, if several shifts are operating, to provide a corresponding number of strips for each product designation as well as an arrangement for indicating the desired (nominal) and the actual value for each shift. The strips may be of different construction, such as of different widths, colors or the like. For instance, all strips indicating actual values may be of the same type. Preferably, there is integrally combined with each indicator board a decoder for transforming the output signals of the computer, which may be in machine code, into a form suitable for application to the electroluminescent strips. Owing to the storage capability of a computer, it is also possible to represent on the boards values corresponding to arbitrary past periods in time. Furthermore, it is possible to monitor a continuously running machine, if the quantity to be measured is derived immediately from the machine itself. The arrangements described are therefore suitable for obtaining all sorts of values or quantities that may be of interest in connection with planning or monitoring the production of a plant or an office and for processing and indicating the same as rapidly as possible, there being no limits in this regard other than those depending on the size and storage capacity of the computer.

FIG. 6 shows an arrangement in which the display unit for optically displaying data or quantities representing characteristic features of a production process comprises an oscillograph 15 and a computer 26 connected thereto. The connection may be direct, but in the embodiment shown there is inserted between the computer and oscillograph 15 an encoder 16. The computer has stored therein all data that may be of interest, such as production and other quantities and amounts, work hours, machine or equipment capacities, prices, prime costs and the like together with corresponding programs and instructions which may represent the organization of a production process, time planning, pre-planning and so on. To represent or display the data of interest on the image screen 17 of the oscillograph, they are processed in the computer and transformed into electric signals, such as pulse sequences, which are supplied to the control circuits of the oscillograph. However, in the preferred embodiment shown, the computer supplies its output to an encoder 16, which performs a further processing and sup-

plies its output signal to the control circuit of the image screen device. Decoder 16 is preferably provided with a repeating device, such as an endless magnetic tape storage device, to make it possible to repeat indefinitely the input signals supplied to the control circuits. The electric signals which may be pulse sequences, are rendered visible on image screen 17 as optical signals, which may take the form of lines 18 of a length corresponding to the measure of the corresponding data or quantity to be represented. Depending on the optical inertia of the image screen, several lines can be displayed simultaneously to the human eye. Since most oscillograph screens are provided with a raster pattern 19, this can be used as a reference scale for measuring the displayed quantities.

FIG. 7 shows a modified arrangement, in which a T.V. apparatus 20 is used as an image screen device, on the screen 17 of which a measure of the data or quantity characterizing an industrial or other process is displayed in the form of optical signals such as lines 18, beams 21 or the like. It is also possible to represent a quantity as a number or numeral 22 or to bring a corresponding legend to luminescence. The computer or the encoder 16, which is preferably connected thereto may in this case also be provided with a repeating unit and gives off electrical signals, such as pulse sequences, forming a video signal which is applied to the input control circuit of the apparatus. The pulse sequence is put together by the computer in accordance with the desired type of T.V. representation. In most cases, it will be sufficient to represent the optical signals as corresponding points of the image screen, which are brought to luminescence, and to leave all the remaining points dark. The reverse process, where the optical signals are represented as dark points on an otherwise luminous screen, would also be possible.

If the T.V. apparatus is capable of reproducing colors, data or quantities of particular interest, such as monthly balances, total productivity figures, optimal prices or the like could be made to stand out by means of differently colored optical signals. If no composite colors are used, this, of course, simplifies the control process of the apparatus.

It is also possible to arrange for predetermined sections of the screen to emit optical signals of a corresponding color. The incoming electrical signal may be automatically supplied to the electron gun corresponding to a preselected color to be used and to activate the gun to produce the optical signal when the corresponding image point falls in the desired section of the screen, such as the last hundred lines thereof. In a simplified modification, the tube may be provided with a screen having at least one section thereof provided with different phosphors for emitting different colors.

If production data are to be displayed in correlation with months or years, the optical signals corresponding to months could be displayed on a section of the screen which has a phosphor of conventional TV type and makes the corresponding signals appear in white or slightly bluish color, whereas the data corresponding to years are represented by points, for instance, in the section comprising the last 50 lines, which may comprise a phosphor of the type used in oscillographs, so as to make these image points appear green.

In any case, there is obtained a clear and easily evaluated representation of a number of quantities or corresponding optical displays, which are distinguished by different colors and are therefore easily grouped together.

However, the representation of data is not limited to the use of lines, beams, numbers or letters, as will be explained in the following.

FIG. 8 shows an embodiment having two image screens 17 and which is particularly suitable for conference purposes. The computer 26 and/or encoder 16 applies to the control circuits of an image screen device 23, electric signals such as pulse sequences, in the manner already described. The representation of data may be in the form of lines 18, beams 21, numbers, numerals, letters or symbols 22. However, it is also possible to represent the time variation of data as a linear graph 24 or a shaded area 25. If the image device is adapted for reproducing different colors, displays of particular interest can be either repeated once more or be represented separately on the second screen. Both screens can be of the type used in color television receivers. The two screens can be adapted to reproduce different unicolored optical signals or the screens may be of the type described in connection with FIG. 7, or a combination of both is possible.

The arrangement, therefore, makes it possible to obtain a modern, easily interpreted display of data representing production or other processes to the exclusion of human error. Owing to the great storage capacity and calculating speed of the computer, it is possible to represent to the viewer a large number of data in a minimum of time. These possibilities will be illustrated by the following.

The production planning of an enterprise gives rise to a problem of the following type:

The enterprise produces n products $X_1, X_2 \dots X_n$ with the aid of m production facilities $V_1, V_2 \dots V_m$. The term "production facilities" is to be taken in its broadest sense and includes machines, personnel, raw materials, intermediate products, electric energy, storage space, transport vehicles, etc.

Normally, the most important facilities are machines and workers, which may form organized production units. The capacity of such a unit is the maximum amount produced in a certain interval.

Usually, it is desired to maximize the overall profit of the enterprise, which is a function of the quantities and types of products. For different products, the price includes different values of the components which together constitute the price quotation: variable cost, fixed cost and profit. The first two together form the prime cost and the last two the marginal cost.

Variable cost includes, for example, the cost of electric energy, raw materials, proportionate wages, etc. Fixed cost, which includes, for example, depreciation of equipment, rents, insurances, fixed wages, etc., remains unchanged regardless of whether equipment or labor is used productively or not.

As a simple example, assume that two products X_1 and X_2 are manufactured in quantities (so far unknown) of x and y and the corresponding profits per unit are a and b , respectively. The total profit is then expressed by the linear equation $Z = ax + by$; since the production capacity is limited, both x and y are limited

and if further factors are taken into account, such as the fixed and variable costs of the products, marketing considerations, etc., additional conditions are to be met and new variables are introduced, such as the degree of utilization of different facilities. These additional conditions may also be expressed by equations, e.g. of linear character.

The totality of these equations constitutes an optimum model, especially a maximum profit model or minimum cost model indicating how the production should be planned. It is possible to arrive at values for the several variables by solving the usually very large number of equations of the optimum model obtained by known linear programming methods, preferably by means of a computer, to which there is fed the mathematical program for solving the system of equations as well as the input data representing the known properties of the production, such as capacities, profits of products, fixed costs, marketing considerations, etc.

The methods of obtaining the optimum models and solving the same by means of computers may be, for instance, of the types described in one of the following well-known textbooks:

Dorfman-Samuelson-Solow: *Linear Programming and Economic Analysis*, New York, 1958

Ferguson-Sargent: *Linear Programming*, New York, 1958

Manne-Markowitz: *Studies in Process Analysis*, New York, 1963

The theoretical value of the quantity of a product which has been obtained by the computer may now be multiplied with the corresponding marginal cost. This operation can be performed by the computer, which has already been supplied with information about the marginal costs of different products. The resulting value may be represented on a luminescent strip.

The marginal cost value supplied to the computer is, on the one hand, the so-called planned marginal cost obtained from the planned cost accounting or from a roughly estimated production program and, on the other hand, the actual marginal cost which is obtained as the difference between the effective net product sales and the product-variable costs. This is done since the quotation prices of the product arrived at by the cost accounting (with the product costs calculated by the cost accounting) usually do not correspond to the established actual product market prices which may differ also because of competitive grounds. Since, as a general rule, the product line formed of a number of different products is not permanent but, because of numerous influencing factors may be, and rationally should be, varied, it is advantageous — as contemplated by the invention — to indicate the actual marginal cost per product as well as its planned marginal cost not only as a mathematical but also as an optical (i.e. visual) magnitude. If the actual marginal costs are shown visually to be larger than the planned marginal cost of the products, an immediate overall visual survey meaningfully indicates not only the particularly profitable products, but also those products which, as a general rule, may be sold cheaper by the difference between the actual marginal cost and the planned marginal cost or those which, with their already well-established prices, particularly contribute to the profit buildup of the enterprise and therefore their produc-

tion and sale should be encouraged. The same, but converse considerations apply for those products whose actual marginal cost is smaller than their planned marginal cost; a condition which is also visually recognizable according to the invention. These indications, which may thus also be represented visually, have a great significance for the determination of the product line and for market and price policy considerations. This significance of such indications is even more intensified by the fact that — based on the marginal cost values per product supplied to the computer and based on the work plan per product and further, based on data which are also fed to the computer and which relate to available production capacities pertaining to equipment and labor per work station — according to the invention, it may also be visually shown next to the visual indication of the production capacity, whether and to what measure does in a planned production program the production capacity fulfill the requirements in its entirety and also per operating work stations for both a single shift and a multiple shift operation (it is noted that a work plan discloses which work station is used to make the product and what average production times are needed for equipment and labor). In this manner, again, important indications are obtained for further investment policies since the operational bottlenecks are immediately visually recognizable in a positive manner. Since on the opposite side of the exemplarily described arrangement of the invention, not only the planned and actual marginal cost per product (and thus its meaningful difference) may be recognized, but there are also visually shown the absolute magnitudes of each planned product quantity or actual marginal costs derived from product quantities calculated and proposed by the computer based on a given optimal model. Such indications may furnish the management with further important keys for decision making. Since previously the computer also calculated from the given program the entire sum of the fixed costs relating to the operation (that is, the so-called actual total fixed cost), there is immediately shown not only digitally but also in a meaningful manner visually, with what type of product and product quantities (and the marginal costs resulting therefrom) will the actually present total fixed costs be fully covered. As a further result, the apparatus described by way of example, visually and digitally shows to what extent does the total profit of the enterprise increase if the available production capacities are fully balanced. Such result is obtained principally because, subsequent to covering the total actual fixed cost, the excess of the marginal cost values of the products gives the actual profit of the enterprise for the planned or computer-calculated optimal production program. These data are also optically and digitally visible. Thus, as a result, the invention assists in a particular manner the product line planning and production planning for a maximum profit.

The aforementioned considerations thus show the importance of a rapid indication of the actual and planned marginal cost as well as the magnitude obtained by multiplying the product quantity by its marginal cost. It is similarly important to show rapidly the utilization of operational capacity of the production program which is either planned in a conventional manner or is calculated by a computer based on an optimum model.

If the computer, based on the optimum model, calculates the quantities of the individual products, then, since the manufacturing steps of the individual products and the production factors necessary therefor are fixed, the utilization of the production capacity is determined. The speed of calculation and also a rapid visual indication of the results is of great importance since it is then possible to obtain various simulated programs which are close to reality. In practice, the determination of the production program is never left solely to the computer, not even if it works with mathematical optimum models. The reason is that there are constantly changes in the magnitudes of the influencing variables so that only a cooperation of the human spirit (usually taking effect as a result of teamwork of management), sales and operational leadership, cost accounting and further operational units, may result in a determination of the optimum product line and production program by utilizing the calculating speed and capacity of the computer and further utilizing suitable rapid visual indicator devices to illustrate the results automatically which is the purpose and the subject of the present invention.

In practice, there are thus preponderantly variable simulated situations which are "played through" mostly as a result of the teamwork of the departments in charge around the conference table utilizing a continuously changing quantity of operational data and accordingly, supplying different data quantities to the computer (which may be located on or off the premises). Thus, continuously changing products are often preplanned for production in determined quantities, because the sales department may have already fixed orders which have to be fulfilled in any case irrespective of whether these products are particularly profitable or not. Since such fixed planned product species and quantities (which may be determined for example in the course of a conference) are supplied as a so-called condition in an optimal model, the problem is thus always to optimize the remainder of the production program, that is, to optimize the still undecided product species and product quantities. In such cases, from the product line some preplanned individual products, based on definite sales considerations are included in the mathematical model as conditions and thus have an effect on the remainder of the optimized calculated results.

Thus, since such product line planning and production planning take place mostly as a teamwork and usually around a conference table, it is particularly important that a rapid calculating process and also a rapid optical indication of the results take place. The speed of obtaining results is not only dependent upon the calculating speed of the computer, but first of all, upon the type and extent of the program which, in turn, is dependent upon the number of the individual calculating steps which are necessarily tied with the basic mathematical optimum model. By weighing all circumstances, it may be said that a larger tolerance in the accuracy of the results is of lesser importance than a smaller tolerance which latter, however, gives rise to unproportionately longer calculating times of the computer.

For a solving process in case of a practically still representable, somewhat at greater tolerance of the

calculator results, the so-called "transporting method" and thereamong preferably the so-called "Hungarian Method" is very rapid; it requires only one-tenth of the calculating period compared with other solving processes. The so-called Hungarian Method, as a variant of the so-called "transporting method" is described in detail particularly in H. C. Joksch: *Lineares Programmieren*, Tübingen, 1965, pp. 159-167.

These considerations illustrate the importance of a rapid display of nominal and actual marginal costs.

If the computer calculates optimal values of the quantities of different products, it may be programmed to derive therefrom the degree of utilization of the corresponding production facilities. This makes it possible to ascertain if adjustment of the running program should be made, e.g. it may turn out that one production unit can do the job originally envisaged for two units.

A visual inspection of the board may indicate that a correction should be made in the data or the program supplied to the computer, and the result of such an alteration will then be immediately visible.

It is clear from the above that not only the maximum profit but other optimal conditions may be obtained in similar manner. In many enterprises, such as base material industries, oil shipping and other transport companies, the counterpart of maximum production would be optimum transportation of goods. Assume, for instance, that a company possesses m mines W_1, W_2, \dots, W_m and n blast furnaces H_1, H_2, \dots, H_n at different locations, and the maximum output of the mines and the capacities of the furnaces are known, as well as the cost of transport from any mine to any furnace. It is then of importance to find a transportation plan representing a minimum of cost.

The solution of a model of this type of problem can be found preferably by means of the aforementioned Hungarian Method, described in the cited reference. The result is a set of output values from the computer that may be represented on the indicator board of the present invention.

The term "production process", as used in the present specification and claims, is, therefore, to be taken in its broadest sense to include the production not only of material products but of any kind of effort or work, such as transportation work.

That which is claimed is:

1. A method of monitoring a production process by optically representing industrial management data characterizing the production process, comprising the steps of:

obtaining a first quantity derived from data characterizing quantity and cost figures related to individual products of an industrial process and a second quantity derived from data characterizing the manufacturing steps on the products during operation of said process

deriving electrical signals representing said quantities, and

applying said signals simultaneously to digitally controllable optical display units to cause said units to respond and provide a visual indication corresponding to the magnitudes of the quantities and the utilization of production capacity.

2. A method according to claim 1, in which said quantities are derived from production data by means of a computer.

3. A method according to claim 1, in which said signals are applied to electroluminescent devices for visually comparing the results corresponding to the first and the second quantities.

4. A method according to claim 1, in which said signals are applied to cathode ray tubes.

5. A method according to claim 1, in which electric pulses are derived from one display unit for controlling the combining of the optically displayed quantities and applying them to another optical display unit.

6. A method according to claim 1, in which at least one of said quantities is derived from the movement of an object.

7. A method according to claim 6, in which said object is a machine.

8. A method according to claim 6, in which said object is the product obtained from a process.

9. A system for monitoring a production process by optically representing industrial management data characterizing the production process, comprising:

a general purpose computer including means responsive to data characterizing quantity and cost figures related to individual products and manufacturing steps on the products during the operation of a process for obtaining a quantity derived from said data,

signal-deriving means responsive to said data responsive means for deriving an electrical signal representing said quantity.

a digital signal generator responsive to said electrical signal, and

a digitally controllable optical display unit responsive to said electrical signal to produce an optical signal and further including an additional digitally controllable optical display unit responsive to a switching unit for transforming said electrical signals into secondary signals representing a re-arrangement of said optical signals to provide a visual display of utilization of production capacity.

10. An arrangement according to claim 9, in which the display comprise units an electroluminescent ele-

ment having a luminescent area of controllable extent and a scale for measuring said extent.

11. An arrangement according to claim 10, in which said element is an electroluminescent strip.

12. An arrangement according to claim 10, in which said scale is exchangeable.

13. An arrangement according to claim 10, in which said element is a strip forming a closed configuration.

14. An arrangement according to claim 13, in which said strip is circular.

15. An arrangement according to claim 13, in which a plurality of parallel strips is provided.

16. An arrangement according to claim 10, in which said element is a circular disc and said luminescent area is a sector of said disc.

17. An arrangement according to claim 10, comprising a plurality of electroluminescent elements of different shape.

18. An arrangement according to claim 17, in which said elements are of different widths.

19. An arrangement according to claim 10, comprising a plurality of electroluminescent elements of different colors.

20. An arrangement according to claim 9, in which said computer is connected via an encoder to each said display unit.

21. An arrangement according to claim 9, in which said digitally controllable optical display units are luminescent screen devices having an input circuit connected to a computer means for producing an electric output signal representing the measure of a quantity characterizing a production process.

22. An arrangement according to claim 21, in which said luminescent screen device is a T.V. image tubes.

23. An arrangement according to claim 21, wherein said luminescent screen devices are connected via electrical switching means for deriving and displaying additional quantities characterizing a production process.

24. An arrangement according to claim 23, in which said luminescent screen devices are adapted for displaying a plurality of signals in different colors.

25. An arrangement according to claim 24, in which the screen have different sections displaying said signal in different colors thereon.

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[54] **APPARATUS AND METHOD FOR SCHEDULE MONITORING AND CONTROL**

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[51] Int. Cl.² **G09B 19/18; G09F 3/18; G06C 3/00**

[58] Field of Search **35/24 A, 24 B; 40/19.5; 58/149, 151; 116/135; 235/89**

[56] **References Cited**

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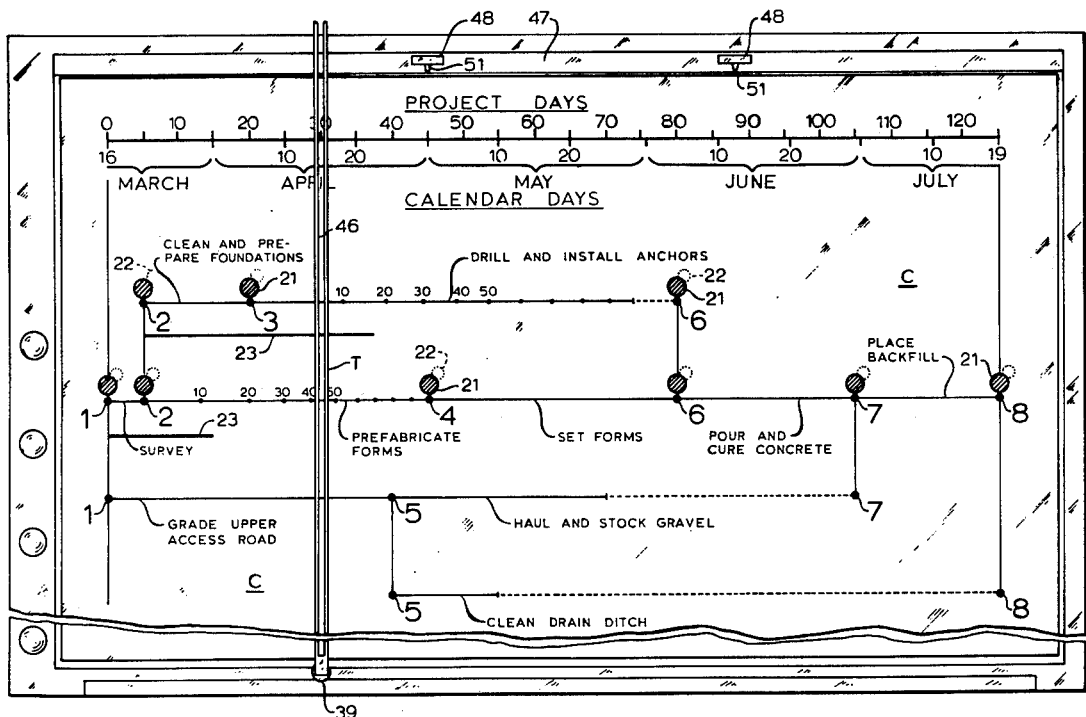
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Primary Examiner—E. S. Jackmon
Attorney, Agent, or Firm—Stanley Bialos; Alvin E. Hendricson

[57] **ABSTRACT**

For the scheduling of activities, particularly construction activities by the Critical Path Method, an apparatus is provided with a mechanical simulation of the network associated with the method, and is adapted to support a network chart of activities and events. It has a progress bar for monitoring such activities and events, which is power driven at a constant rate for indicating the scheduled progress of the entire project at any given time, whereby the scheduler can determine at a glance at the board what adjustments, if any, need be made in the timing and progress of the various activities of the project.

11 Claims, 14 Drawing Figures



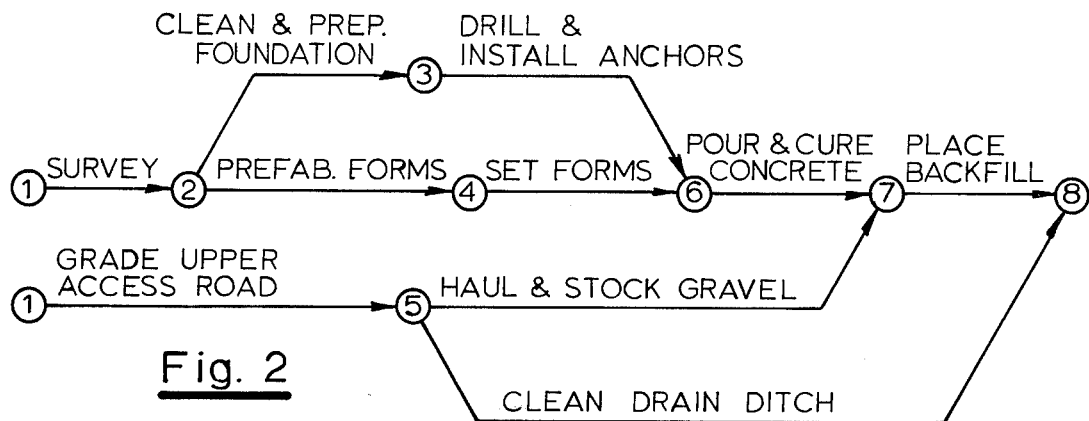


Fig. 2

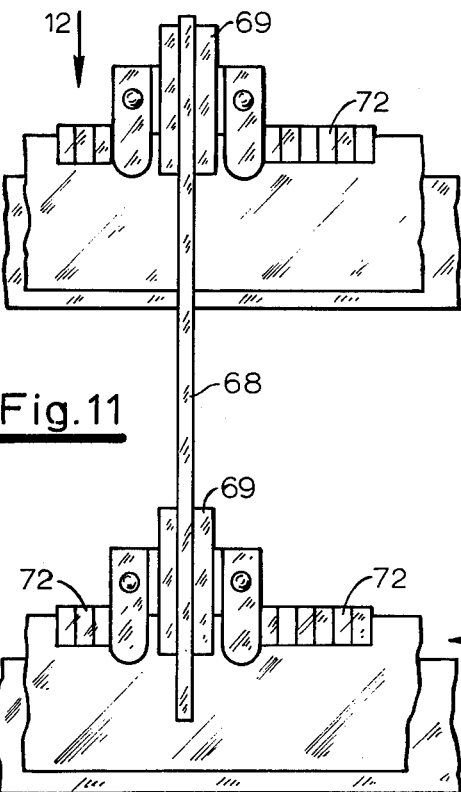
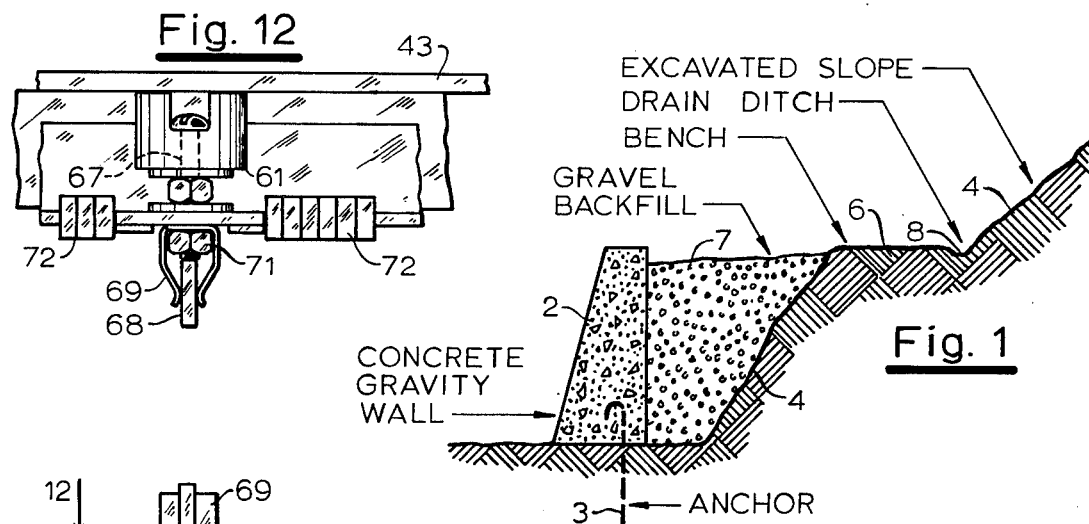


Fig. 14

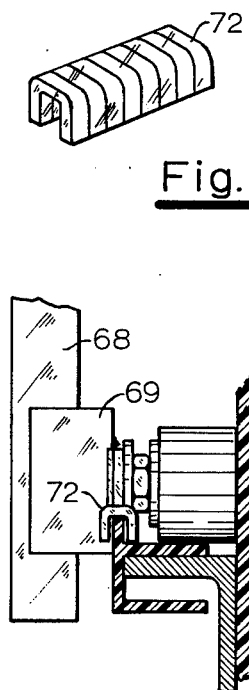


Fig. 13

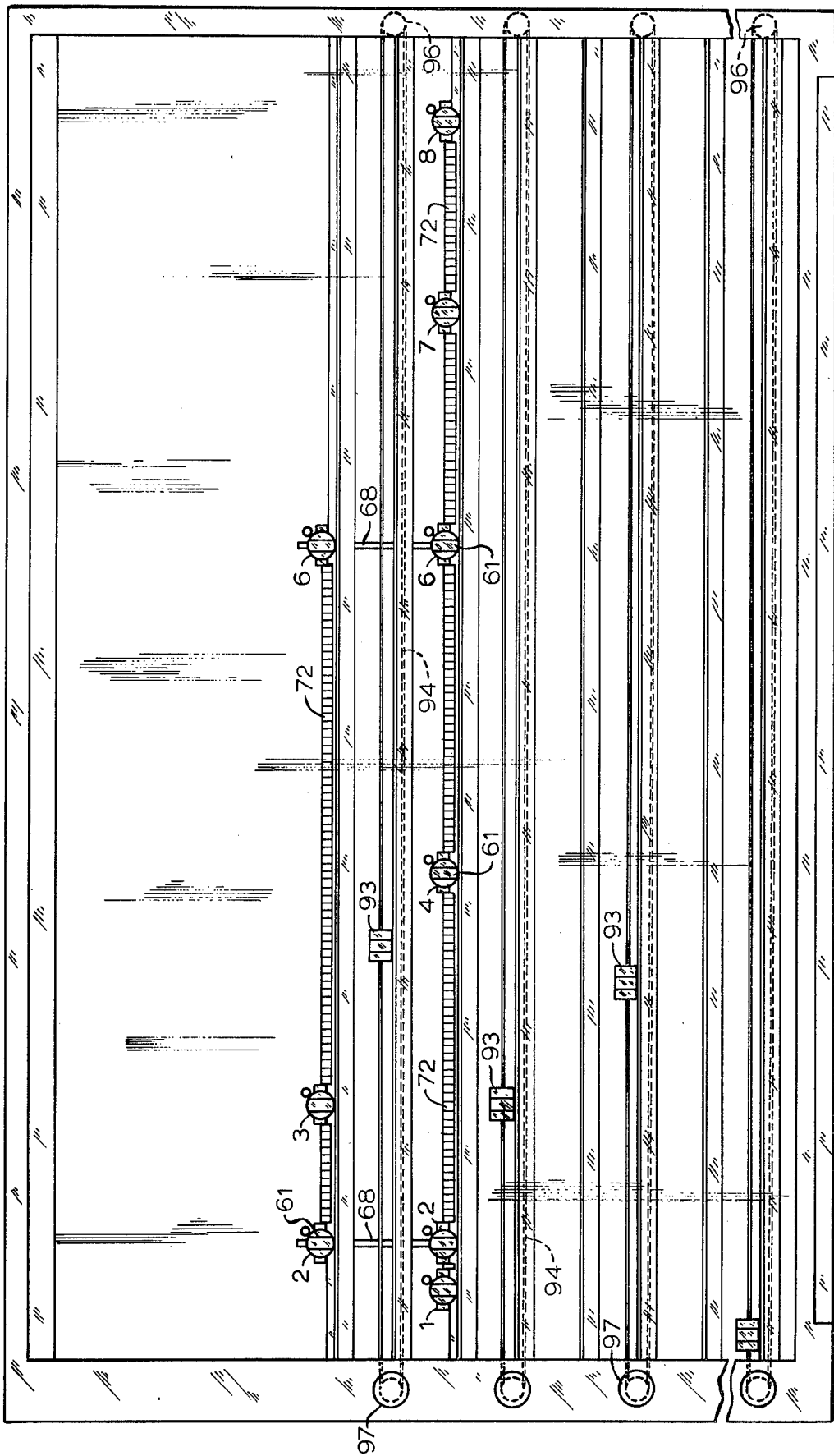


Fig. 4

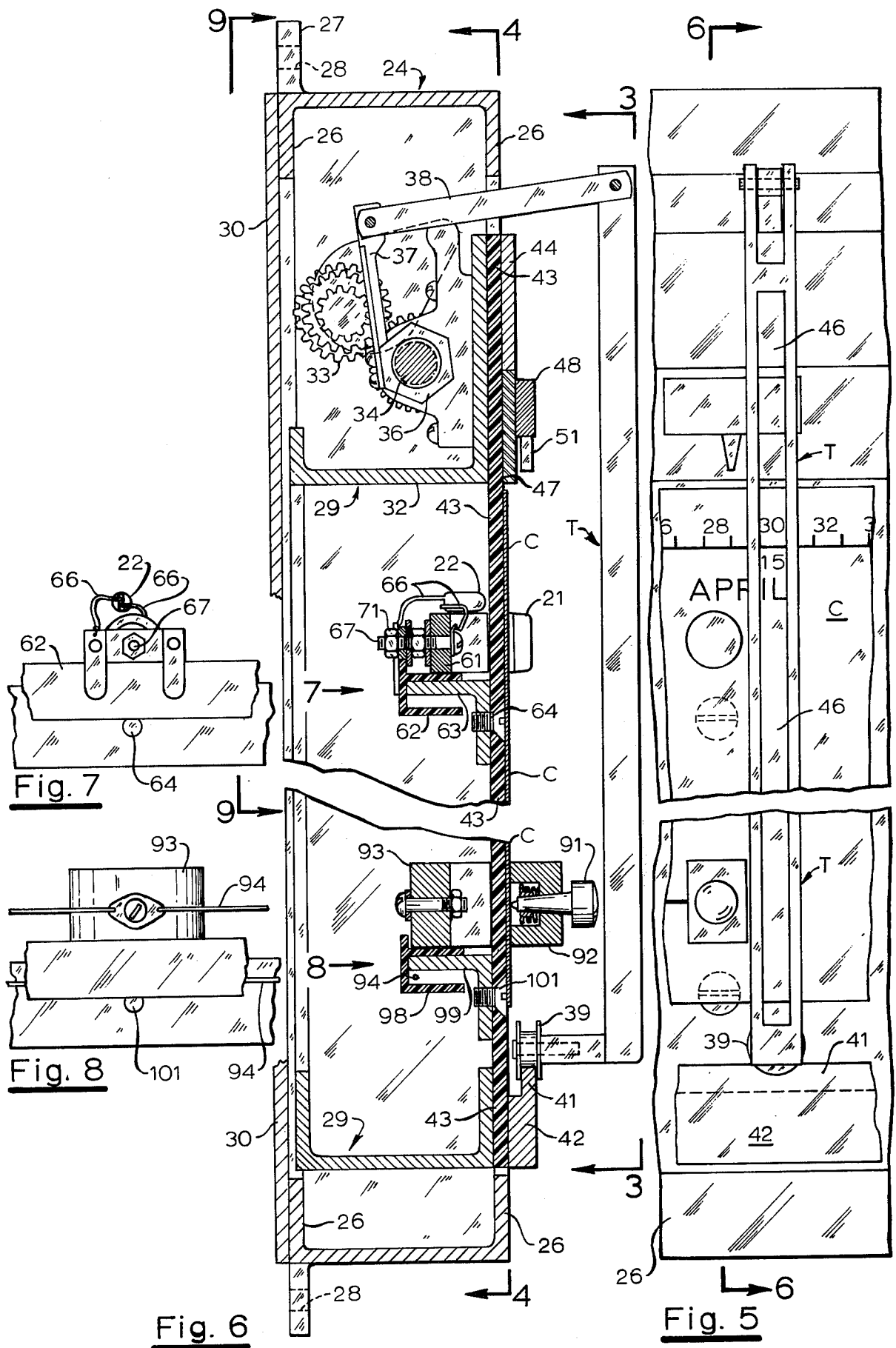


Fig. 7

Fig. 8

Fig. 6

Fig. 5

Fig. 10

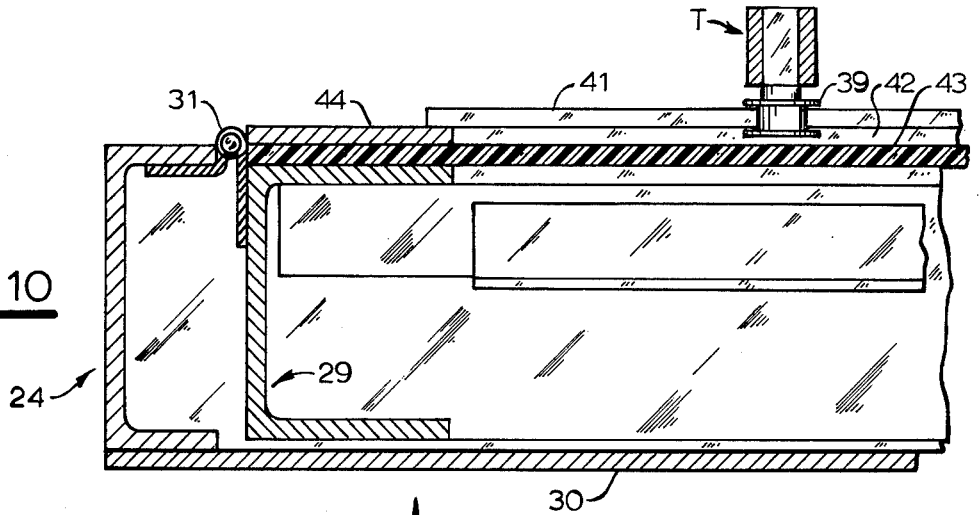
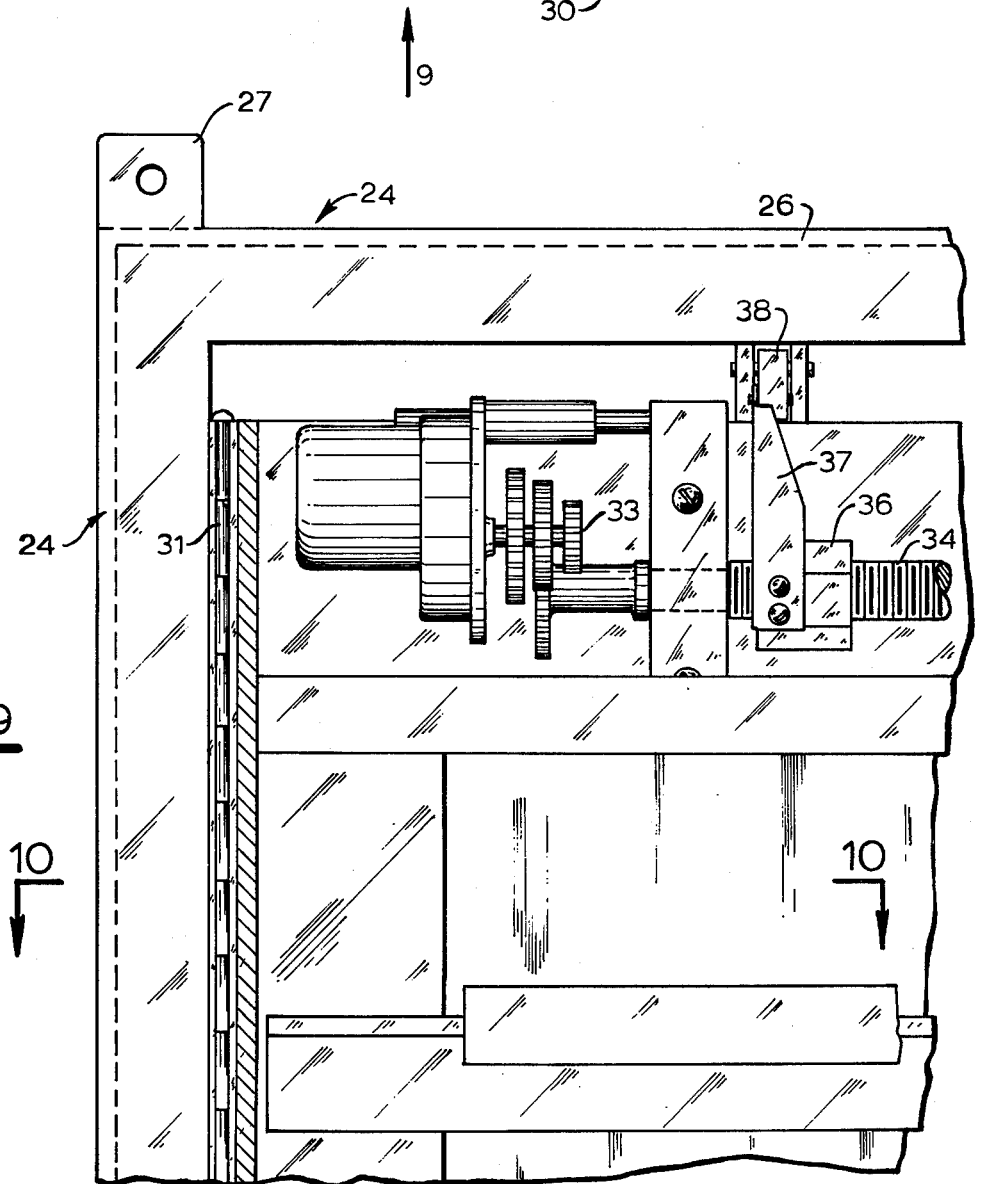


Fig. 9



APPARATUS AND METHOD FOR SCHEDULE MONITORING AND CONTROL

BACKGROUND OF THE INVENTION

The critical path method (CPM) is a comparatively recent engineering development particularly adapted for use in the construction industry for the planning and scheduling of construction activities. For a detailed analysis of CPM, reference is made to the book entitled, CPM IN CONSTRUCTION MANAGEMENT, Second Edition, by James J. O'Brien, published by McGraw-Hill Book Company, Copyright 1971.

In construction projects there are activities which may be scheduled to start at certain events (times) and end at other events later on. An activity is a work item leading from or to an event; and an event is a point of time for completion or starting of an activity. In the CPM method, the activities are usually drawn on a chart indicated by lines or arrows and which start or terminate where two or more activities meet.

The engineers responsible for the particular scheduling of a project usually have a network chart of events and activities made up beforehand which is based on estimates. When reports from the field are received after the project has started, this network chart is examined regularly (which may be daily or weekly) to determine the progress of the entire project and of its individual activities, which is noted on the chart. If a critical activity is behind schedule, the Project Manager may employ additional resources (e.g. personnel or equipment), and from the chart he can tell whether the entire schedule has to be updated in order to complete it on time. This is a tedious job; and quite a complex one in situations where the project is complex.

In the invention hereof, a monitoring apparatus which will be described in detail hereinafter is provided whereby the project engineers or management in charge of the project may at a glance determine the status thereof so that they may take appropriate action when required.

TYPICAL EXAMPLE OF CPM

As background for explanation of CPM, the following typical example is illustrative. A scale network chart therefor, for the apparatus hereof will be explained later.

With reference to such example, FIG. 1 is a schematic sectional-elevation of a hypothetical project; and FIG. 2 is a network and activity schedule chart drawn for such example. The project of the sample comprises construction of a concrete, gravity wall 2 anchored at 3 near the toe of an excavated slope 4 on which a bench 6 exists, and back filled with gravel 7 between the slope and wall 2 up to the elevation of bench 6 in order to provide a widened road at such elevation. Drain ditch 8 is provided adjacent to the slope and the bench.

Conditions require that the project must be completed within a maximum period of 140 calendar days, including a 15 day contingency allowance for foul weather. Excluding the 15 day allowance for such foul weather, the schedule is 125 calendar days. With these limitations in mind, the management develops a network with activity data for the project on a calendar day basis. The activity data are hand calculated but complex networks often justify computer assistance.

FIG. 2 illustrates the CPM activity and event network for the example of FIG. 1, which is drawn after the

activities are determined and their durations estimated by the personnel in charge. In the FIG. 2 network, the activities indicated by straight lines with arrows and the events indicated by numbered circles, are not drawn to a time scale. This is done later when a time bar chart or schedule used in this invention is made up. It will be described later.

The following is a table of activities and events of the FIG. 2 network, as estimated by the management:

Event Interval	Activity	D	ES	EF	LS	LF	F
1 - 2	Survey	5	0	5	0	5	0
1 - 5	Grade Upper Access Rd.	40	0	40	35	75	35
2 - 3	Clean & Prep. Foundation	14	5	19	11	25	6
2 - 4	Prefabricate Forms	40	5	45	5	45	0
3 - 6	Drill & Install Anchors	55	19	74	25	80	6
4 - 6	Set Forms	35	45	80	45	80	0
5 - 7	Haul & Stock Gravel	30	40	70	75	105	35
6 - 7	Pour & Cure Concrete	25	80	105	80	105	0
5 - 8	Clean Drain Ditch	15	40	55	110	125	70
7 - 8	Place Backfill	20	105	125	105	125	0

D - Duration
 ES - Early Start
 LS - Late Start
 F - Float
 EF - Early Finish
 LF - Late Finish

The activities are listed in a column and are indicated between event numbers. The duration of the entire project is estimated at 125 days as was indicated previously. The estimated duration of each activity in project days is indicated in column D. ES designates the earliest possible start which can be made for any particular activity determined from the time the project is commenced; LS the latest possible estimated start without delaying the entire program; EF the estimated earliest possible finish of an activity; LF the latest possible estimated finish without delaying the entire program; and F represents float time in project days.

With respect to float, it should be kept in mind that certain activities are not critical, and therefore they will not interfere with completion of the project if the activity is not started or finished within a reasonable time. However, certain activities are critical and have no spare time or float. For example, the "Survey" (1 - 2) must be completed before the two activities "Prefabricate Forms" (2 - 4) and "Clean and Prepare Foundation" (2 - 3), which lead from (1 - 2), can be started. Also, as a further example the activity "Prefabricate Forms" (2 - 4) must be completed before the activity "Set Forms" (4 - 6) can be started. If critical activities, such as 1 - 2, 2 - 4, or 4 - 6 are completed after latest finish (LF) times, the entire project will be delayed.

There are various ways in which "Float" can be calculated. It can be determined by the formula $F = LF - EF$; or $LS - ES$; or $LF - ES - D$.

Those activities which have no float time (zero float) are critical because they must be completed in the scheduled time if the project is not to be delayed. Those which have very little float in days are considered near critical. A near critical activity may be considered one in which the float is within about 10% of the calculated duration of the entire project. The more float there is in an activity, the less critical such activity

becomes; and as previously explained, the entire project may be delayed if a critical activity is not started or completed in time. Near critical activities are usually considered the same as critical to provide a factor of safety.

From the above, it will be noted that the critical and near critical activities of the illustrative Example are 1 - 2; 2 - 4; 3 - 6; 4 - 6; 6 - 7; and 7 - 8. Of these, activities 1 - 2; 2 - 4; 4 - 6; 6 - 7; and 7 - 8 comprises the critical path which is defined as the longest route through the CPM network.

SUMMARY AND OBJECTS OF THE INVENTION

Summarizing the invention hereof, it comprises performing the usual CPM estimates and calculations for a particular project as explained above, and making a scaled network chart or schedule in time-bar form, on which the events are plotted at their calculated early start (ES) positions. In cases in which two or more activities share the same total float, this float may be distributed among the affected activities. Thus in the example herein shared float is allocated equally to Activities 2 - 3 and 3 - 6. Therefore the scheduled start of Activity 3 - 6 is 3 days later than its early start, or Day 22.

An apparatus is provided which has a hinged front panel, advantageously transparent, on which the network schedule is mounted. Markers in the form of discs or buttons for various events are removably positioned at the front of the schedule at the various event positions. The markers are of magnetic material, and are held in position by magnets movably supported at the back of the panel so that the markers can be shifted by manually shifting the magnets.

A vertical indicator bar is mounted for movement along the panel, and is moved continuously by a clock driving mechanism which is timed for the entire calculated project. The chart has a horizontal scale drawn thereof which indicates in calendar days as read by the position of the indicator bar, the number of days the project has progressed as scheduled. Horizontal scales are drawn parallel to various activity lines and indicate the estimated percent completion of the respective activities as determined by the position of the indicator bar. Thus, by noting the position of the indicator bar with reference to the activity scales, the Project Manager in charge (who usually must depend on information conveyed to him by subordinates or have to wade through a mass of written reports or data) can by himself determine at a glance the overall progress of the project as compared to the scheduled progress. In this connection, the actual progress of any particular activity can be estimated and recorded on the schedule at any particular time by means of a marking pen. Desirably, such marking pen is provided for each critical or near critical activity on the critical path; and means is provided to hold the marking pen and move it manually along the chart.

Should it be found at any particular time that a critical or near critical activity is behind schedule, job management can readily recognize this from the position of the indicator bar of the apparatus with reference to the estimated actual progress, and take appropriate action such as by adding more personnel to work on the activity. For example, should it be noted that activity "Pre-fabricate Forms" (2 - 4) is substantially behind schedule even after increased personnel has been employed to work on this activity, the entire project can be re-

scheduled by rescheduling one or more of the succeeding critical activities for shorter duration or durations, or accept the lost time as probably unrecoverable. If the latter is opted, means is provided on the apparatus for prolonging and updating this schedule, for example, by increasing the overall duration time.

For the purpose of updating or adjusting the event markers, means is provided at the rear of the schedule control panel for manually simultaneously shifting all of the critical event markers at one time so that they need not be individually shifted. This is done by interconnecting such event markers so that when one of the critical markers is shifted, the ones interconnected therewith are also shifted. Shifting of the event markers to update the schedule results in altering the "Float" of the non-critical activities.

From the preceding, it is seen that the invention has as its objects, among others, the provision of an improved simple method and apparatus for determining at a glance the status of any particular project which is programmed by the critical path method, which is simple to operate and perform, and economical. Other projects will become apparent from the following more detailed description and accompanying drawings, in which:

DESCRIPTION OF DRAWINGS

FIGS. 1 and 2 have been previously described.

FIG. 3 is a more or less schematic front elevation of the apparatus with the CPM scale chart thereon, drawn from the information illustrated in FIGS. 1 and 2; the view looking in a direction of line 3 — 3 in FIG. 6;

FIG. 4 is a front elevation of a light transmitting panel for mounting the CPM chart, and which is mounted on a hinged frame of the apparatus, with the chart, the indicator bar and event marker discs of magnetic material omitted to illustrate the construction more clearly; the view looking in the direction of line 4 — 4 in FIG. 6, with parts of the apparatus in the same relationship as shown in FIG. 3;

FIG. 5 is a fragmentary front elevational view of the apparatus looking at an indicator bar portion thereof;

FIG. 6 is a transverse vertical section taken in a plane indicated by line 6 — 6 in FIG. 5;

FIG. 7 is a fragmentary rear elevation of the magnet arrangement for holding or retaining event marker buttons; the view looking for holding or retaining event marker buttons; the view looking in the direction of arrow 7 in FIG. 6;

FIG. 8 is a fragmentary rear elevation of a magnet structure for holding recorder mechanism, looking in the direction of arrow 8 in FIG. 6;

FIG. 9 is a fragmentary rear elevation of the apparatus looking in the direction of arrow 9 in FIG. 10, with a back cover panel removed to illustrate the support frame arrangement and drive mechanism for the indicator bar, and with a portion of the structure shown in section to illustrate the construction more clearly;

FIG. 10 is a fragmentary horizontal section taken in a plane indicated by line 10 — 10 in FIG. 9;

FIG. 11 is a fragmentary rear elevation view illustrating a form of vertical connection between magnet mechanism, looking at the marker 6 portion of the apparatus in FIG. 3;

FIG. 12 is a fragmentary top plan view of the portion of the mechanism shown in FIG. 11, looking in the direction of arrow 12 in FIG. 11;

FIG. 13 is a isometric view of a type of spacer bar between magnet mechanism; and

FIG. 14 is a fragmentary end elevation, looking in the direction of arrow 14 in FIG. 11.

DETAILED DESCRIPTION

As was noted previously, after the activities have been determined and outlined as indicated in FIGS. 1 and 2, and by the table for FIG. 2, the time bar chart or schedule of the entire project is drawn to scale for various duration times. A suitable scale may be:

Scale	Schedule Duration Range
0.1 in. = 1 day	1 to 2 yrs.
0.2 in. = 1 day	6 mos. to 1 yr.
0.4 in. = 1 day	6 mos. or less

For schedules longer than two years' duration or should it be desired to use a larger scale, two or more monitoring boards or panels may be employed, or the schedule may be divided into two or more successive networks.

A suitable paper size for the schedule is about 40 by 84 inches; and it has been found that lines of activities which extend horizontally should be spaced about 3 inches apart vertically. The board may be of any suitable vertical height to accommodate the activity lines spaced apart vertically; a suitable number being about 10 to 12. Shared float among the activities where this occurs is distributed by allocation.

FIG. 3 is a schedule of the project previously explained, drawn from the data depicted in FIG. 1 and 2. Events for critical and near critical activities are indicated by circles but non-critical activities are shown on the chart solely to complete the network and to enable overview thereof. It will be noted from FIG. 3 that the activities follow the order indicated in the aforementioned table. Where more than one activity terminates or originates from an event, such events are separated vertically on the schedule. Float times for non-critical activities are indicated by dotted lines. A scale in calendar days is drawn on the chart indicating the date of termination (duration). This same scale also indicates the estimated duration of the project (project days), commencing with 0 start and ending at the estimated end of the project (125 days).

The activity lines between events are scaled to indicate the estimated duration of the respective activities divided into estimated percentage of scheduled completion of the particular activity at any particular time during its estimated duration determined with reference to the position of the indicator bar. Thus, it will be noted for example that the activity "Drill and Install Anchors" (3 - 6) having an original estimated duration of 55 days is divided into increments indicating percent completion. These are estimated for all activities by the project management as previously related; and it will be observed that the spacing for the percentages decreases with time on the theory that as work on the activity progresses, the personnel become more familiar with it and can work more efficiently at the end than at the start of the project. A time bar T, to be described later, on the apparatus is provided which is moved by clock drive mechanism to indicate scheduled progress of the project in calendar days.

As will be described later in greater detail when the apparatus structure is described, the prepared schedule

or chart is mounted on a panel; and means is provided for removably positioning critical or near critical event markers 21, which are in the form of discs, at the points such events had been calculated to occur. For example, event 2 occurs after the "Survey" (1 - 2) has been completed, which is 5 days after commencement of the project, and the marker therefor is positioned at such point. The marker for event 3 is placed at a position 14 days after event 2, etc. For clarity, these event markers 21 are schematically illustrated in FIG. 3 in offset relationship with reference to the event points, and are desirably illuminated as indicated by the dotted circles 22 in FIG. 3. As will be described later, the markers are held by magnetic means at the back of the panel.

Also, means is provided for recording the overall progress of each of the critical or near critical activities by providing a recorder which is moved manually along the panel and draws a recorder line 23 adjacent the activity line. This record is made by the management based upon an estimate of how the work for such activity has progressed. Instead of providing such recorder on the schedule itself, the same effect may be obtained by the Project Manager, by drawing the line on the schedule by a hand held pencil or pen.

As illustrative of how the schedule is used in the apparatus, assume that on Apr. 15, 30 days after the project start, the panel and board control schedule appear after examination as shown in FIG. 3. It will be noted that the project is in difficulty because critical activity 2 - 4 is behind schedule 15 days. This is ascertained from the fact that the time bar T is at 30 project days at which the activity should be about 45% completed, but the actual completion, as indicated by recorder line 23 just below activity line 2 - 4, is only about 10% which should have occurred about 15 days after the project commenced instead of 30 days. Therefore, management has to make a decision whether to prolong the completion date of the project or to recover lost time.

Assuming that a decision is made to add a carpenter crew in order to increase the production rate of prefabricating the forms (activity 2 - 4), and after this is done the activity is still behind schedule by 10 days, then management must make another decision of whether to reschedule one or more of the succeeding critical activities for shorter duration or durations, or accept the lost time as probably unrecoverable. The desirable form of apparatus will now be described, including its means for rearranging the schedule or timing by prolonging of updating it.

Referring to FIGS. 3 through 14 which illustrate the apparatus structure, it comprises a support frame 24 of rectangular box shape having vertical flanges 26 which form sides channel-shaped in cross section. Outwardly extending lugs 27 having apertures 28 provide means for attaching the frame to a suitable support such as a wall or easel. Frame 24 houses an inner frame structure 29 which carries apparatus components and which is pivotally or hingedly connected along one side edge thereof by means of hinge 31 to a side flange of frame 24. Frame 29 provides a support for the chart and also for driving means for indicator bar T which is mounted on frame 29. The rear of frame 24 is covered by detachably connected backing panel 30.

As can be seen best from FIGS. 6 and 9, frame 29 has an upper horizontal channel 32 above which is mounted clock drive gearing 33 connected to drive a screw 34. A nut 36 on screw 34 is attached to an up-

standing link 37 to which is pivotally connected a link 38 which is turn is pivotally connected to the top end of indicator bar T. As the clock gearing 33 is driven, indicator bar T is moved longitudinally along the chart at a constant rate. In this connection, gearing 33 comprises

replaceable speed changed gearing so that the rate may be fixed in accordance with the scale of the chart desired. The lower end of indicator bar T carries a freely journalled rotatable roller 39 which rides on a rail 41 forming part of a bar 42 secured to the bottom of frame 29. Bar 42 also attaches a light transmitting support panel 43 (desirably of plastic material, such as "Plexiglas") to inner frame 29; suitable additional bars 44 also being provided to secure light transmitting panel 43 to frame 29. The aforementioned pivotally connected link 38 spaces indicator bar T away from panel 43 and away from the schedule chart C which is attached to panel 43 by any suitable means such as "Scotch Tape" along its edges. The chart is desirably of light transmitting material, such as tracing paper or the like.

For ease of reading, indicator bar T is formed with a longitudinally extending open narrow space 46 through which indicia on the scales may be more readily seen. A longitudinally extending bar 47 of magnetic material is provided above chart C upon which so-called milestone event markers are removably positioned, each comprising a magnet 48 and a pointer 51. Milestone events are those of special importance.

As was previously related, event markers 21 are of magnetic material. They are removably held over the desired critical or near critical events. Means for movably retaining markers 21 for freely slidable movement along the chart C when desired, comprises magnets 61 each of which is supported on an insulating U-shaped track member 62, desirably of plastic material, in turn mounted on an angle member 63 secured to light transmitting panel 43 by countersunk screws 64. Because each magnet 61 is supported for slidable movement along the track, it may be readily positioned at a desired event point on the chart at which an event marker 21 may be held by the magnet.

Each magnet 61 supports the aforementioned light source 22 which illuminates the associated marker 21 by light transmitted through transparent support panel 43. Light transmitting panel 43 also serves to allow the scheduler to see where to set up the magnets 61 in desired positions to hold markers 21. All the light sources, desirably electric light bulbs 22, are connected by any conventional circuitry including wires 66 suitably insulated from the magnets and held in position by bolts 67.

As was previously noted, where two or more activities start from the same event, a plurality of rows of vertically displaced event markers is provided such as 2 - 2 and 6 - 6 shown in FIGS. 3 and 4. In order that the vertically spaced magnets for such events may be shifted in unison together, they are rigidly connected together by a connecting bar 68, as can be seen from FIGS. 11, 12 and 14. For attaching the bar to a magnet, a U-shaped clip 69 is fastened to the magnet structure by a nut 71 which cooperates with bolt 67. The clip rigidly clamps the bar so that two of the magnet structures, one above the other, can be manually moved laterally in unison when desired.

In revising (updating) the schedule for reasons previously explained, a plurality of magnet structures is

manually shifted or moved to the right or left with reference to FIGS. 3 and 4. They are so interconnected that all of them that are to be shifted can be simultaneously moved in unison merely by pushing one of the magnet structures from left to right or vice versa, thus updating the schedule by prolonging or shortening it, respectively. Means for effecting such shifting includes readily convenient structure for initially properly spacing the magnets apart when the apparatus is first set up, so that the event markers 21 retained by magnets 61 will be properly spaced apart.

For such purpose, a premanufactured conventional bundle 72 of individual staples adhesively secured together edge to edge is most suitable because the bundle can be readily adjusted to the desired length by removing staples or by combining bundles. When so adjusted to space two event markers apart for example 4 - 6 as shown in FIG. 4, it will maintain the desired spacing whereby all of the markers in advance of a selected marker may be moved in unison merely by manually pushing a selected magnet structure which retains the marker. The magnet structures which are connected together by an upright bar 68 will also be shifted simultaneously. Not only do the bundles 72 act as spacers, but they are desirably of metal so that they can form part of the illumination circuitry which also includes aforementioned wires 66.

The structural mechanism has been described which relates to monitoring the actual progress of critical or near critical activities so as to relate them to the scheduled progress as forecast at the beginning when plans were initially made by the management. Accordingly actual progress as recorded on the chart from field reports can be compared readily at any time with scheduled progress shown by the indicator bar T.

Desirably, the aforementioned record line 23 should be drawn indicating the actual progress of the respective particular activities at the time. Line 23 may be drawn from time to time on the chart by hand held pencil or pen. However, recording means is provided forming part of the apparatus to enable drawing of such line conveniently.

Such recording means, as can be seen best from FIGS. 6 and 8, comprises a spring pressed writing instrument 91 slidably mounted on a button magnet 92 which is pressed against chart C by the spring. Magnet 92 is fixedly held against the chart by means of a cooperating horeshoe magnet 93 which is secured to an endless flexible strand 94 which runs about a pulley 96 at one end and about a pulley 97 at the opposite end having a manually controllable knob by which pen 91 may be moved across the chart. Magnet 93 can slide along insulating track 98 adhered to angle member 99 which supports the track; angle member 99 being secured to the transparent panel 43 by means of countersunk screws 101.

I claim:

1. Schedule control and monitoring apparatus for use in network type scheduling of a project having a plurality of interrelated activities with events at the start and termination thereof and wherein a succession of activities of the longest time path for the project are critical, said apparatus comprising a support frame, means for mounting on said frame a prepared chart on which a network schedule of such activities and events is written chronologically with the distance of consecutive critical activities drawn to scale representing scheduled times based on estimates of the durations to complete

the respective activities, a plurality of magnets movably disposed on said frame behind said chart whereby the positions thereof can be changed, a plurality of markers of magnetic material adapted for disposition on the front face of said chart at said event points and supported by said magnets, an indicator bar mounted on said frame over said chart, and means for moving said indicator bar at a constant rate past said markers from a position indicating the commencement of the project to provide a comparison between the actual progress of any given activity and the scheduled progress thereof.

2. The apparatus of claim 1 wherein means interconnects a plurality of said magnets for simultaneously shifting said magnets together along said chart when the schedule is to be changed.

3. The apparatus of claim 1 further comprising a plurality of recording instruments for various respective activities; each of said instruments being mounted on a body of magnetic material at the front of the chart, a magnet for each of said bodies disposed at the back of the chart for retaining said respective body and pressing the instrument thereof against chart, a track on which each magnet is slidably mounted, and manual means including a control knob for sliding said magnet along said track.

4. The apparatus of claim 3 wherein said manual sliding means includes an endless flexible member to which the magnet is attached.

5. The apparatus of claim 1 wherein the support frame includes a panel of light transmitting material to which the chart is attached at the front side of the panel, the magnets are slidably mounted on tracks at the rear side of the panel, and means carried by respective magnets is provided to illuminate and respective marks supported thereby of events at the start and termination of critical activities.

6. The apparatus of claim 2 wherein said support frame includes a panel to which the chart is attached at the front side of the panel, means is provided for spacing the indicator bar to clear said markers as the indicator bar moves, means is provided for mounting said panel for hinged movement along one edge about an upright axis to allow the panel to be swung for access to the interconnected magnets at the rear of said panel, and the means spacing said indicator bar to clear said markers includes linkage connected to the top of the bar for pivotal movement about a horizontal axis.

7. The apparatus of claim 6 wherein a track at the bottom of the panel is provided, and a roller is journaled at the bottom bar rides on said track.

8. Schedule control and monitoring apparatus for a critical path method project which has a succession of activities some of which are critical whereby they must be completed on time if the project is not to be delayed and also comprising events which are times at which activities start or terminate, said apparatus comprising a support frame including a transparent panel, means mounting on the front face of said panel a prepared chart on which a network schedule of such activities and events is written chronologically with the distance between consecutive critical activities drawn to scale

representing scheduled times based on estimates of the durations to complete the respective activities, magnets slidably mounted at the rear face of said panel for respectively supporting a plurality of markers of magnetic material on the front face of said chart at event points, means interconnecting a plurality of said magnets for simultaneously shifting said magnets along said chart when the schedule is to be changed, an indicator bar mounted on said frame over said chart, and means for moving said indicator bar at a constant rate past said markers from a position indicating the commencement of the project to provide a comparison between the actual progress of any given activity and the scheduled progress thereof.

9. A schedule control and monitoring method for use in network type scheduling wherein a plurality of interdependent activities culminating in events of a project are to be scheduled and monitoring during progress of the project comprising the steps of:

providing a chart on which a network schedule of activities and events of the project are presented chronologically with the distance between consecutive events drawn to a time scale representing scheduled times for performance of activities based on estimates of the durations thereof,

mounting said networks chart upon a front face of a panel,

forming a mechanical counterpart of the network of said chart on the back face of said panel by placing magnets at the locations of said events on the chart and mechanically linking the magnets between the events as related on the chart,

placing a magnetically responsive marker on said chart at each event and supporting said markers by said magnets,

drawing lines indicating the actual progress of said activities on the chart to the scale thereof, and moving an indicator bar at a constant rate related to the scale of the chart across the front face of said panel over said chart to provide a visual comparison of actual progress of the project to estimated times of completion of activities thereof.

10. The method of claim 9 further defined by providing illumination at each magnet, and electrically and mechanically linking said magnets in the manner events are connected on the chart and connecting said electrical linkage across a power supply whereby only markers located at events of a critical longest path of activities of the project are illuminated to visually depict the critical path of the project.

11. The method of claim 10 further defined by elongating the linkage between magnets corresponding to markers at the beginning and end of activities which are not to be completed in estimated times to thus reposition magnets and markers on said chart for illuminating the markers of events of an adjusted longest critical path of activities to provide a visual indication of adjusted network schedule and comparison to original network schedule and actual activities progress.

* * * * *

[54] WAREHOUSING MONITOR AND CONTROL SYSTEM

4,135,241 1/1979 Stanis et al. 235/385
 4,141,078 2/1979 Bridges et al. 364/403
 4,180,204 12/1979 Koenig et al. 235/385

[75] Inventors: Charles A. Smith; Robert T. Danevicz, both of Grand Rapids, Mich.

Primary Examiner—Errol A. Krass
 Attorney, Agent, or Firm—Price, Heneveld, Huizenga & Cooper

[73] Assignee: Rapistan Division, Lear Siegler, Inc., Grand Rapids, Mich.

[57] ABSTRACT

[21] Appl. No.: 138,250
 [22] Filed: Apr. 7, 1980

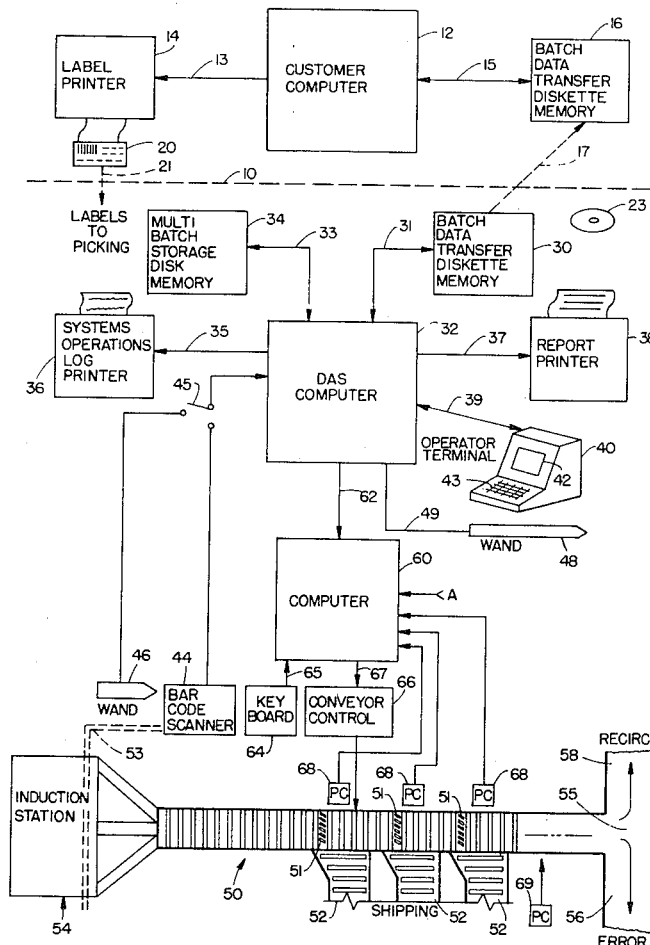
A warehouse product monitoring and control system includes a distribution audit system computer which receives batch picking information on a recording medium identifying each transaction or article to be picked with a unique number. Associated with each unique number is processing information for controlling the sorting of the article within the system. The same number, unique for each article, is contained on a label positioned on each article as it is picked. The computer control provides continuous monitoring of the article as it is scanned by one or more label reading devices such that the picking and sorting status of each article to be picked can be continuously monitored. The system includes operator interface terminals and printers for providing a variety of status reports to the operating personnel such that the operational status of the system can be continuously monitored.

[51] Int. Cl.³ G06F 15/24
 [52] U.S. Cl. 364/403; 198/347; 198/418; 364/478; 364/900
 [58] Field of Search 364/403, 478, 200, 900; 235/432, 385, 91 L; 340/150; 198/347, 418

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17 Claims, 9 Drawing Figures



PSC COMMUNICATIONS SUB-SYSTEMS

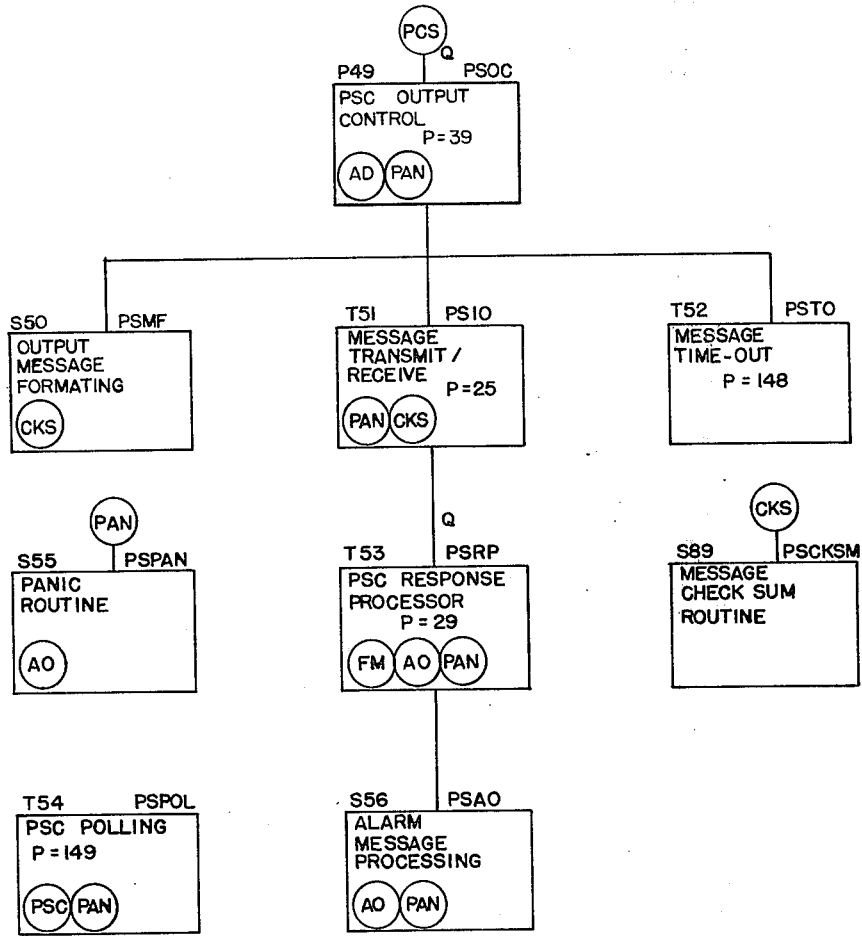


FIG. 7

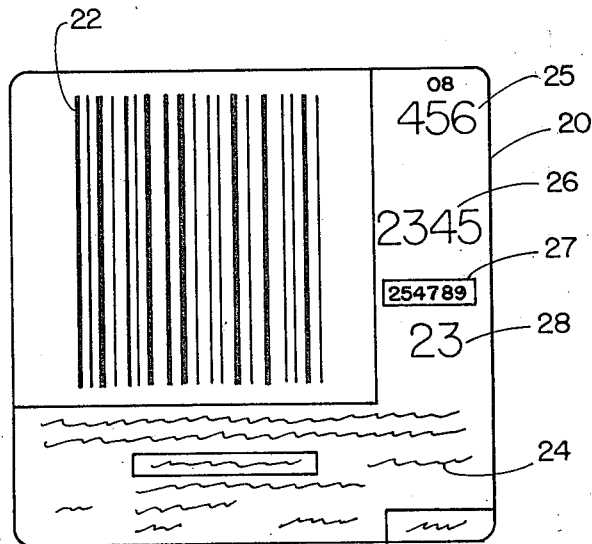


FIG. 2

OPERATOR INTERFACE SUBSYSTEM

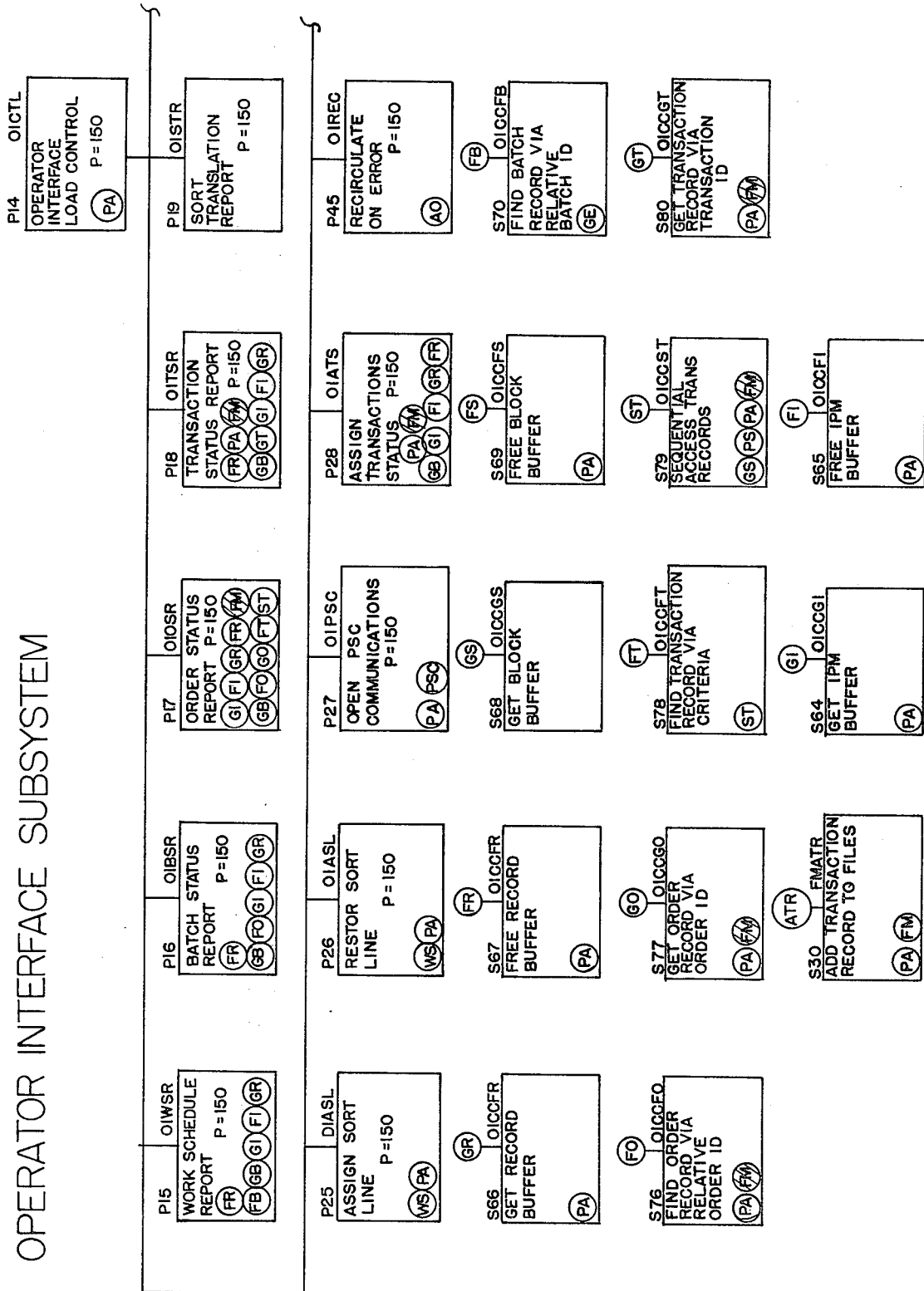


FIG. 3A

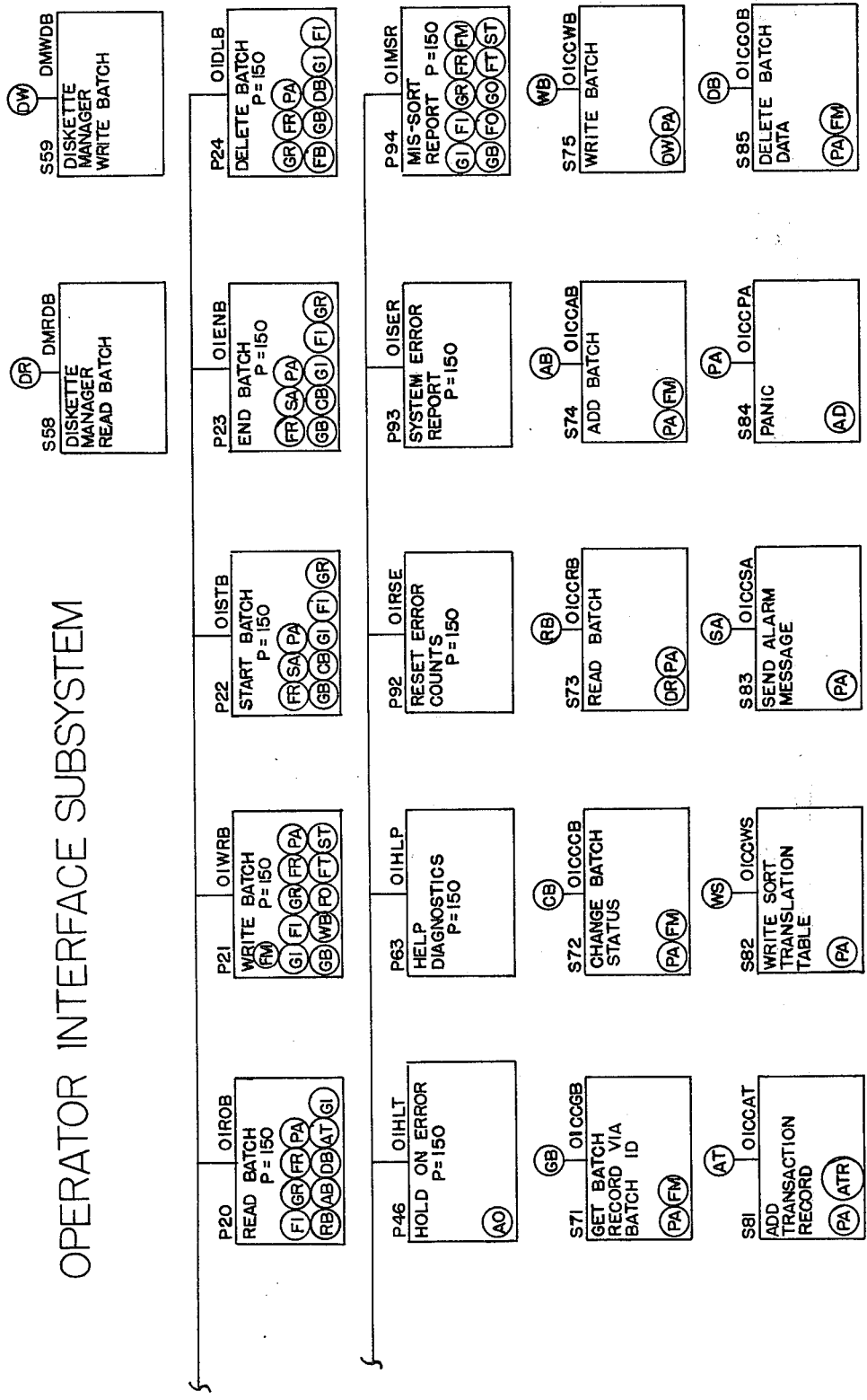


FIG. 3B

FILE MANAGER SUB-SYSTEM

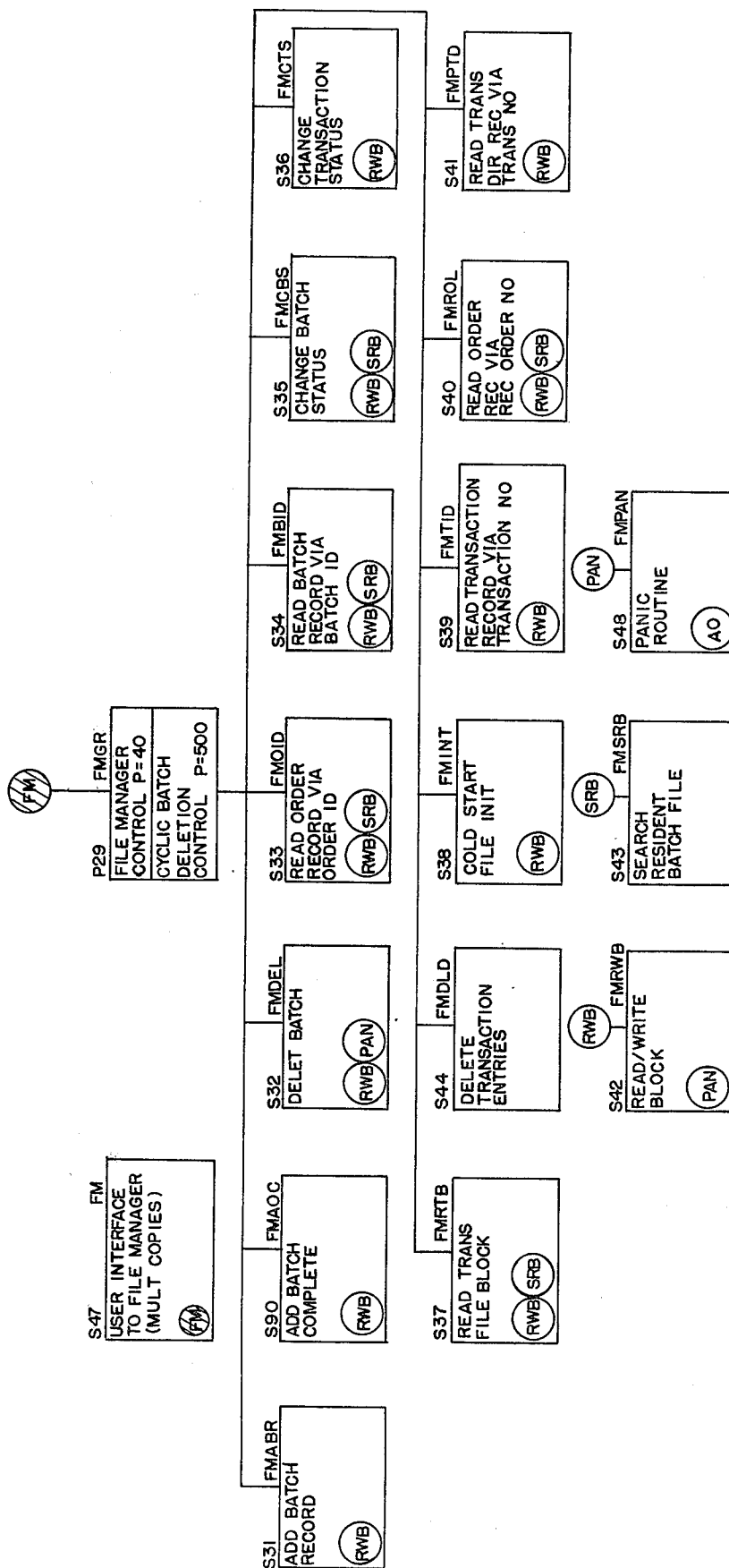


FIG. 4

SCANNER INPUT SUB-SYSTEM

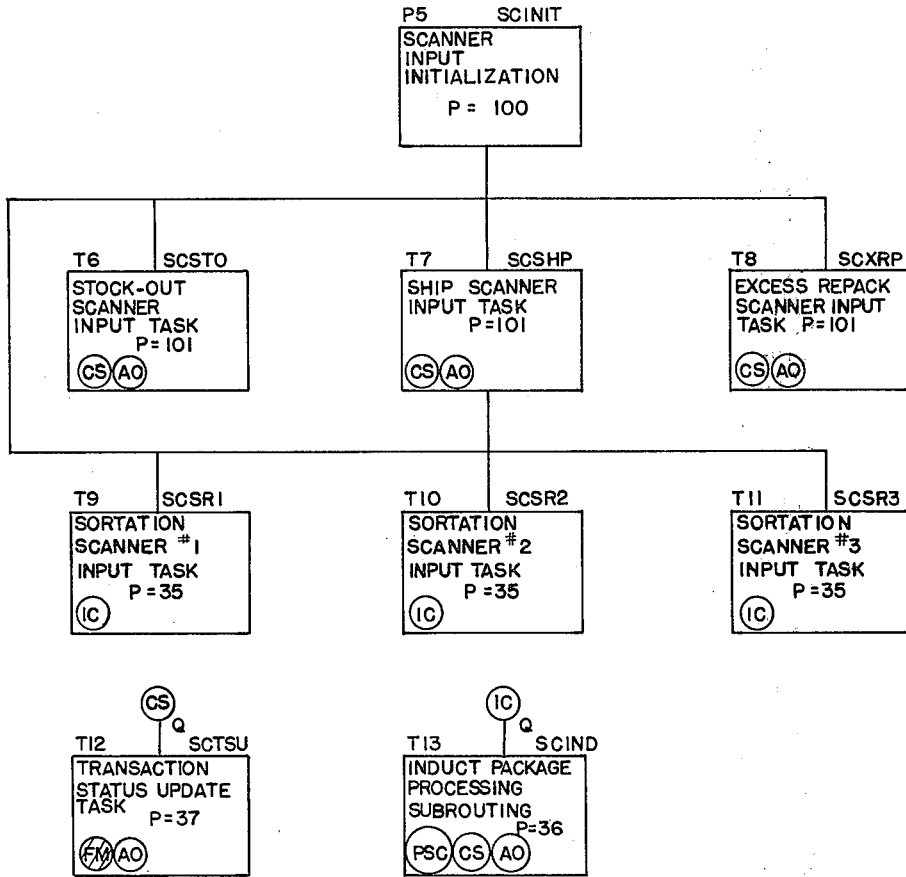


FIG. 8

AUXILIARY FUNCTIONS SUB-SYSTEM

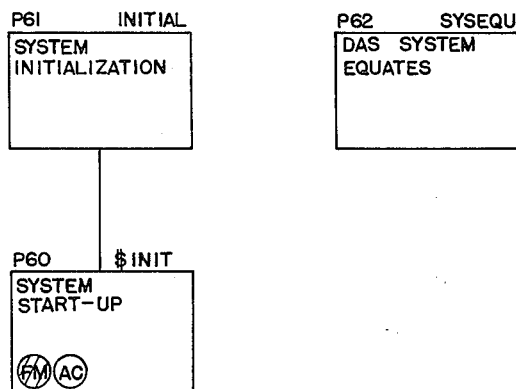
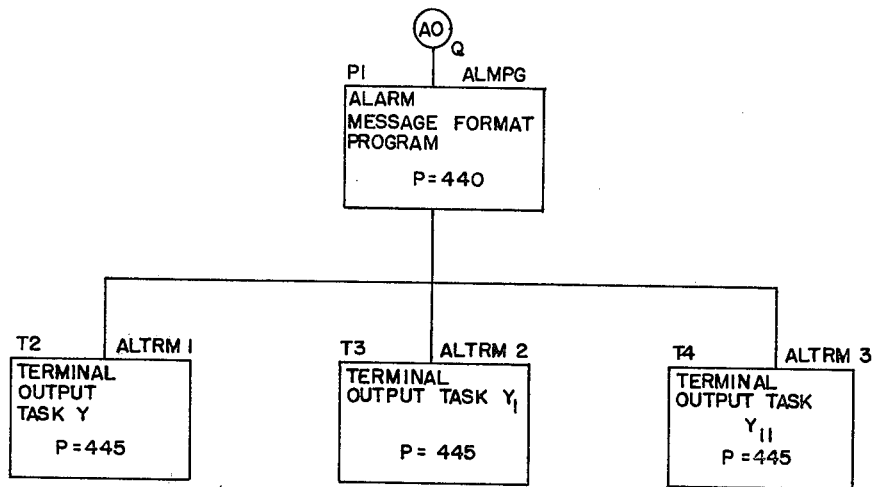


FIG. 5

ALARM MESSAGE SUB-SYSTEM



LEGEND

- P PROGRAM
- T TASK
- S SUBROUTINE
- D DATA
- P-n TASK/PROG. PRIORITY
- Q FIFO QUEUE
- ⊘ QUEUEABLE RESOURCE
- CONNECTOR

FIG. 6

WAREHOUSING MONITOR AND CONTROL SYSTEM

BACKGROUND OF THE INVENTION

The present invention relates to a method and apparatus for monitoring and controlling the flow of articles in a warehouse or other installation.

There exists a great variety of systems in which orders are filled by, for example, a central warehouse where items to be shipped to a given destination are manually, semiautomatically, or automatically picked, sorted, and routed to a shipping dock for shipment to, for example, retail stores. Typically, labels identifying the shipping destination are applied to articles as they are picked which labels are subsequently employed at a sorting location for the diverting of articles to a particular shipping area of the installation. U.S. Pat. No. 4,181,947 illustrates a sorting system which can be employed with such a system.

Where, however, a warehouse system is of relatively large size and it is capable of simultaneously handling a great number of orders, maintaining accurate information as to the status of individual orders being filled as well as the operational status of the entire system, while maintaining a maximum throughput efficiency, is virtually impossible with the systems of the prior art.

SUMMARY OF THE PRESENT INVENTION

It is an object of the present invention to provide a method and apparatus for maintaining an accurate account of each transaction occurring within a warehousing system such that at any given time during an operating day, the status of each transaction and order being filled can be monitored and controlled. In order to achieve this objective, each transaction, which in the preferred embodiment of the invention pertains to a given article to be transferred from storage to a dispatching area, is assigned a unique identification number. This number is subsequently applied to an article as it is picked in the form of a machine readable coded label. The number is also entered into a control system memory, and as the article travels through the system, the label is read such that at any given time, the status of each transaction and the operation of the entire system is accurately known and can be controlled.

In the preferred embodiment of the invention, the system incorporates a computer controlled sorting and control system in which each article to be removed from storage and shipped to a destination, is assigned a unique code which identifies a single transaction within the system. Orders to be filled by batch picking and transfer of articles from storage to a given dispatch location at the warehouse, are grouped together in the computer memory to form batches of orders with the customer identification and discharge location being associated with each transaction number. By reading only the transaction number from an encoded label on an article at one or more locations within the system, the computer can continuously update the order status information and provide the operators of the system with current status information as well as provide control information to the sorting system employed. By providing operator interface circuits, reports can be generated either in soft or hard copies to monitor the operational status of each transaction, order or batch of orders to facilitate the accurate and fast movement of articles within the system. Such information can be employed

to quickly identify and correct failures and breakdowns within the system or to rapidly reassign divert locations for the sorting system in the event of a breakdown of a sortation line.

These and other objects, features and advantages of the system will become apparent upon reading the following description thereof together with reference to the drawings in which:

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a pictorial representation of a warehousing system embodying the present invention including an electrical circuit diagram in block form of the control system;

FIG. 2 is typical label associated with each article and identifying a unique transaction within the system;

FIGS. 3A and 3B are block diagrams showing the module layout of the software employed for the operator interface portion of the control system;

FIG. 4 is a block diagram of the module layout of the software for the file manager portion of the control system;

FIG. 5 is a block diagram of the module layout of the software for the auxiliary functions of the control system;

FIG. 6 is a block diagram of the module layout of the software for the alarm message portion of the control system including the legend employed for FIGS. 3 through 8;

FIG. 7 is a block diagram of the module layout of the software for the interface between the sorting control system and the main control system; and

FIG. 8 is a block diagram of the module layout of the software for the code reading portion of the control system of the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

FIG. 1 represents a physical installation as well as the control system located within the installation. At the upper portion of the figure above dashed line 10 there is shown a central control area corresponding typically to an office area of a customer's facility, while the portion below line 10 represents the actual warehouse installation including the sortation conveyer equipment, and in which the actual articles to be transferred from storage to a dispatching dock or the like are stored in a conventional warehousing storage system. Such a system may include, for example, tiers of vertically and horizontally arranged storage bins with access aisles extending between adjacent tiers permitting either manual, semiautomatic, or automatic picking of articles from the storage bins to be achieved through the use of stacker cranes or other picking systems. Articles removed from storage are placed on conveying systems which ultimately connect with a sorting conveyer. Not shown in FIG. 1 are the storage bins, access aisles, or the conveyers leading to the induction station and sorting conveyer system. In some installations, the warehouse will be physically separate from the central control area while in other installations the two locations may share a common building at separate areas. Having briefly described the environment in which the present invention pertains, a description of the overall system is now presented.

In FIG. 1, the customer computer 12 typically will be a relatively large computer used by the customer for all

of its business operations. Such a computer may be, for example, a commercially available IBM System 3 type computer which is interfaced with a label printer 14 by means of a data link 13 for the printing of labels 20, such as shown in FIG. 2. The label printer can be any one of a number of dot matrix-type printers which are commercially available such as the Printronix Model 600. Also, interfaced with the customer computer 12, is a batch data transfer diskette memory 16 coupled to the computer by a data link 15 for the generation of diskettes containing information corresponding to one batch of orders. The unit 16 may be, for example, a commercially available IBM floppy diskette model 3540. The customer computer 12 is employed in connection with label printer 14 and disketter unit 16 through conventional programming techniques to arrange orders to any one of the several retail outlets of the customer in a logical picking sequence such that the articles can be batch picked. Thus, for example, if the customer is a large grocery chain, with several grocery stores located in the geographical area served by the warehouse, one batch of orders to be filled may include 50 different grocery items with the batch including orders for 15 different retail outlets. The orders are arranged such that each item which is common to the group of articles to be picked will be simultaneously picked by the warehousing picking system. Thus, the computer 12 will be programmed, for example, to print labels in order with all of the labels for example, pertaining to a case of a certain type of produce such as beans to be simultaneously picked for all of the orders. This organization provides for efficient batch picking of articles.

Assigned to each transaction, which corresponds to a single case of a given item, is a unique six digit number which is employed within the control system shown in FIG. 1, to uniquely identify that article such that its progress within the system can be monitored and controlled. This number is applied in machine readable form to the label 20 as seen in FIG. 2 in the form of a bar code 22 which occupies a large portion of the label. The label 20 also includes man readable information 24 at the bottom of the label indicating for example the product identification, pricing information, and shipment location to be transferred. Further, the label will include man readable indicia 25 indicating the storage location of the article within the warehousing system which is contained within the memory of computer 12. Also on label 20 is indicia 26, corresponding to the store number to which the article is to be shipped, indicia 27 showing the unique six digit number (which in the example shown in FIG. 2 is the number 254789); and finally, indicia 28 pertaining to the discharge chute of the sortation conveyor system to which the article is to be sent for shipment.

Each article to be picked corresponding to a single transaction thus receives a label 20 and each label printed by printer 14 in a given day has a unique number 22 as compared to all other labels printed. Information corresponding to the label information as well as additional information is simultaneously recorded on a floppy disc by unit 16 which includes the label information for a plurality of labels forming several orders forming, in turn, a single batch for processing within the control system. Each diskette therefore includes stored data corresponding to each label of a given batch including the six digit label number, the sort lane number,

the store number, the product code, which is assigned by the customer to any particular type of product, and information pertaining to where the product is stored in the warehouse. Further, each diskette includes a header record which provides an identification of the particular batch, as well as a description of the batch which can be any 50 character description desired by the customer. Thus, each group of labels provided by printer 14 will have associated with them a single diskette corresponding to a batch of items to be picked. Each batch will typically include several orders, each of which include several transactions with each transaction having a single label associated with it. The system of the present invention is capable of processing simultaneously up to 10 batches with each batch having a capability of 35 orders per batch and total of 9,000 transactions for each batch. Typically, however, only one or two batches will be picked simultaneously with the remaining batch information contained in memory in either a pending status or possibly a completed status, depending upon the operational status of the picking and sorting of the batches. The generation of the labels, as well as the diskette in a batch picking sequence, with the exception of the utilization of a unique number assigned to and associated with each label and transaction is conventional and achieved by computer 12. The labels are manually carried to the warehouse as indicated by dashed line 21 as are the diskettes 23 generated and associated with each batch as indicated by dashed line 17.

Each diskette from the customer computer 12 generated by the diskette unit 16 is read by a batch data transfer diskette memory unit 30 associated with and coupled to a distribution audit system (DAS) computer 32 by means of a data link 31. The diskette unit 30 may comprise for example a commercially available IBM diskette unit model number 4964, while computer 32 may for example be an IBM model 4955D computer. The function of unit 30 is to read the information from the diskette hand carried from unit 16 into a multi-batch storage disc memory unit 34 coupled to computer unit 32 through data link 33. Unit 34 is capable of storing not only the control program for computer 32, but also the data stored on each of the diskettes associated with each batch of articles to be picked. Unit 34 may, for example, comprise a commercially available IBM 4962 disc storage unit. Hard copies of alarm, status, and control information are provided by a systems operations log printer 36 coupled to computer 32 by means of a data link 35. Printer 36 may, for example, comprise a Digital Equipment Corporation matrix printer type LA 120. A second report printer 38 is also coupled to computer 32 through data link 37 and is employed for providing a variety of status reports as discussed below and may be for example an IBM model 4974 matrix printer. Also coupled to computer 32 by means of a data link 39 is an operator terminal 40. Terminal 40 includes a CRT display 42 and a digital keyboard 43 and may be a commercially available IBM model 4979 display station. Several such operator terminals and printers may be positioned at different locations within the system as conveniently desired. Other interface inputs to the computer 32 include one or more bar code scanners 44 selectively coupled to the computer 32 by means of a switch 45 which either couples the scanners 44 which are Accusort model number M scanners to an input of computer 32 or label wand scanners 46 to the computer 32. The wand scanners 46 can be Accusort model number

4600 scanner for detecting the bar code 22 on labels 20 of FIG. 2. Scanners 44 optically read labels on articles prior to transfer or induction onto the sortation conveyor 50 as shown by lines 53 in FIG. 1. Also coupled to computer 32 by means of a data line 49 is a second wand scanner 48 which is employed for scanning stock-out labels as described below and which can also be an Accusort model 4600 wand scanner. An additional wand scanner (not shown) is employed for hand scanning excess repack labels as also described below.

The DAS computer 32 interfaces with a programmable sort control (PSC) computer 60 by means of a data interface coupling 62. Associated with the PSC computer 60 is a backup keyboard 64 coupled thereto by means of a serial data link 65, and a conveyer control circuit 66 coupled thereto by means of a data link 67. A plurality of article detecting photo cells 68 are spaced along the sorting conveyer 50 downstream of each article diverter 51 for detecting whether an article has been properly diverted to its assigned discharge or shipping chute 52 associated with the sorting conveyer 50. Further, a line full photo cell detector 69 is also coupled to the PSC computer 60 and is associated with each of the discharge chutes 52 such that a control signal is applied to computer 60 in the event any of the discharge chutes are filled with articles and thereby requiring a different discharge chute to be assigned to a particular order being sorted and staged for shipment.

At the input end of the sorting conveyer 50, there is provided an induction station 54 for receiving articles from several feeder conveyers (not shown) extending throughout the warehousing system and converging the articles onto the sorting conveyer 50. The PSC computer and its associated interface and control circuits to provide the induction and sorting of articles is described in detail in the above identified U.S. Pat. No. 4,181,947, the disclosure of which is incorporated herein by reference.

Sorting conveyer 50 terminates in a T 55 leading to either an error chute 56 or a recirculation loop 58 such that articles not sorted can either be recirculated through the sortation control system as selectively controlled as described below or fed to a storage area associated with the error chute for manual attention to a nonreadable label.

The control of the hardware elements shown in FIG. 1 corresponding to the control system of the present invention is achieved by the programming of computer 32 which program is stored in memory 34 and which is organized according to the module layout chart shown in FIGS. 3 through 8. With reference to these figures, FIGS. 3A and 3B disclose the layout of the programming for the interface between computer 32 and printer 38 and computer 32 and the operator terminal 40. The programming represented by FIG. 4 pertains to the internal handling of data between computer 32 and memory 34 as does the programming diagrammed in FIG. 5. FIG. 6, relates to the alarm messages provided by printer 36 and accordingly the programming for providing such alarms. The programming represented by FIG. 7 pertains to the programming which interfaces computer 32 with computer 60. Finally, the programming represented by FIG. 8 pertains to the intercoupling of computer 32 with the scanning wands 46, 48, and the bar code scanners 44. The module layout of the program as represented by FIGS. 3 through 8 is directly correlated with the pseudo code and structural English printout of the software organization disclosed in ap-

pendix A which is attached hereto, made part of the specification and incorporated herein by reference. In FIGS. 3 through 8, each block represents a module which performs a specific operation within the system. Each block is identified by the letters P, S, T, or D (described in the legend accompanying FIG. 6) followed by a number identifying a particular portion of the program employed. In addition, each block includes a label name constituting an acronym which is correlated with the programming information disclosed in appendix A. As can be seen, for example in FIGS. 3A and 3B, many of the modules call up additional modules indicated by an encircled two letter connector. Thus, for example, the panic subroutine shown in FIG. 3B which serves to stop the program in the event of a major failure, is identified by the two letters PA which is called up by many of the modules in the event the module is unable to perform its particular programming function. The programming comprising the source code for the PSC computer is also incorporated herein by reference as appendix B which is attached hereto and made a part of this specification.

Having described the hardware and the broad correlation between the program organization and individual elements of the hardware, a description of the operation of the system to achieve its desired monitoring and control of each transaction, as well as the various status reports available employing the unique transaction number identification of each article to be picked, is now presented.

Before any of the articles to be picked can be recognized and sorted by the DAS control system, the batch and transaction information must first be transferred to the DAS system. This is achieved by the operator control terminal 40 entering a READ command which transfers the information from a diskette 23 placed in unit 30 into the computer memory 34. The program for this is shown in FIG. 3B indicated as READ BATCH. Once the data base has been constructed by the operator for each of the diskettes being processed during a day's operation, the actual picking, sorting, and monitoring functions available are commenced by a START command which changes the batch status from pending to active and enables the DAS computer to sort the product.

The picker applies labels to the correct products and places the unique six digit bar code containing label on each article as they are picked and transferred onto conveyers within the system which feed to the induction station 54. Each article is scanned by scanner 44 for its unique number and this information is supplied to the DAS computer which correlates it with sort location information for that number, and communicates with the PSC computer the sort destination assigned to the particular transaction. Depending upon the selected mode of operation, the article can be stopped and hand wanded by wand 46, recirculated through recirculation loop 58 or shunted into the error chute 56 upon failure of scanner 44 to read the article's number. The status of the article is changed from "not-picked" to "in-sortation" and the PSC performs the physical sorting of the article and informs the DAS computer of the actual status of the package on the sorting conveyer. When the label is transferred to its assigned destination chute 52, (FIG. 1) the status of the article is changed from "in-sortation" to "staged for shipment". If the article is unable to sort to its assigned sorting lane, the PSC informs the DAS of the actual disposition of the package.

Once all of the articles of a batch have been picked and sorted, including articles for several different orders typically, the operator who monitors the progress of the batch picking as described below, enters an END command in the operator terminal 40 which changes the batch status to "complete". At this time, the operator can enter a WRITE command in terminal 40 which generates a new diskette which is returned to the customer computer and which has information as to the completed status of each transaction in the batch which can be employed by the customer for inventory, billing, and other functions such as the printing of shipping manifests and other documents. Once the batch has been processed by the DAS control system, it can be deleted from the DAS data base, memory 34 by entering a DELETE command.

There are a variety of system batch and sortation control commands which are employed by the system of the present invention. The system control commands include a listing of system commands for status reports which is identified by the command HELP entered in keyboard 43 to generate a listing of available reports on CRT 42.

By entering the command ERROR a system error report as shown in appendix A is generated by printer 38 which permits the operator to visualize the number of missorts and transfer failures and take corrective action. Another system control command is RESET which is employed by the operator for resetting error counts which may correspond to missorts where an item is discharged to the wrong discharge chute 52 inadvertently, or transfer failures where the PSC computer fails to divert an article to any of the desired discharge chutes. The final system control command is the time setting function which is employed to provide a real time display or printout to all of the various reports and is entered by entering the command \$ T employed by the operator to set the system time.

In addition to the system control commands, there are several batch control commands which are entered by keyboard 43 some of which have been previously discussed. The first of these, is the READ command which effects the reading of the customer's diskette 23 by the DAS computer diskette unit 30 reading the information into the memory 34. The START command starts the batch processing monitor and control while an END command as previously discussed, ends the batch processing. The WRITE command causes the diskette to be reprogrammed with the updated status of the orders and transactions therein reported back to the main customer computer through a reprogrammed diskette while a DELETE command then is employed to delete the completed batch information from memory 34. The ASSIGN command is employed for example in the event a given product is out of stock and with the batch processing system, this could affect several transactions associated with several different orders. To provide a quick update of the status of the particular product and therefore the transactions in which it is involved, the operator can assign to the transaction numbers affected, a stockout status so that this condition can be instantaneously displayed at any of the report printers or operator terminals once the condition has been reported.

Finally, there are several sortation control commands entered into keyboard 43 to control the sorting of articles. One mode of operation is identified and controlled by the command HOLD. With the HOLD command,

the induction station 54 is stopped when a no-read occurs on scanner 44 and the operator must manually scan the label using wand 46 (FIG. 1) to read the label. He does that by actuating switch 45 such that the wand 46 output information is supplied to the DAS computer 32. The induction station 54 then restarts, and the article is sorted.

With the RECIRC command, articles with labels not read by scanner 44 are automatically directed to recirculation lane 58. Upon initial startup of the system, the command OPEN is employed to open the communications interface 62 between the DAS computer 32, the PSC computer 60. The command REROUTE is employed to assign an alternate sortation line for articles with a common destination contained within all batches being processed. This is employed for example, in the event a sort line 52 has a mechanical failure and it is temporarily shut down. The command RESTORE will automatically reassign rerouted articles not yet processed to the original sort line once it has been repaired.

In addition to the interface between the operator terminal 40 and the system employing the above identified commands, the system will automatically provide two different types of printouts available to management and control personnel for the monitoring and operational control of the system. The first type of printout is provided by the system's operational log printer 36, and provides indications of the status of the system as well as missorts and the like. The following is a typical printout during a short period of time representing a variety of operational conditions of the system.

```
04/01/80 14:30:34 System Started
04/01/80 14:30:46 PSC Communications Line Open
04/01/80 14:30:59 Start of Batch T02
04/01/80 14:31:04 Hold on Scanner Error
04/01/80 14:31:43 Scanner Error At Induction #1
Scanner
04/01/80 14:31:54 Transaction 984561 at Induction
#1 Scanner Not On File
04/01/80 14:32:17 Transaction 411589 at Induction
#1 Scanner Out of Batch
04/01/80 14:32:44 Sortation Line 1 Full
04/01/80 14:32:49 Transaction 598402 Missorted To
Sort Line 3
04/01/80 14:33:24 Unsuccessful Transfer At Sort
Line 1 for Transaction 632368
04/01/80 14:33:27 Scanner Error at Induction #2
Scanner
04/01/80 14:34:29 Transaction 581419 Has Been Lost
From Tracking
04/01/80 14:34:46 Recirculate On Scanner Error
04/01/80 14:35:37 End of Batch T02
04/01/80 14:36:08 Transaction 581419 on Sort Line 2
Not On File
04/01/80 14:36:27 PSC Communications Failure
00/00/00 00:00:10 Power Fail/Restart
```

As can be seen from the above report, the status of the system and its communications with the PSC is indicated as well as events such as scanning errors at the induction station, sortation lines being full, as detected by the line full photo detector 69 (FIG. 1), missorts and the like. This record is automatically printed without a specific operator command such that monitoring personnel can take corrective action as required. Naturally, many of the messages printed simply indicate the status of a particular batch, for example, at 14:30:59, batch T02

was beginning to be processed and at 14:35:37, batch T02 was completed.

The second type of printed information, is specifically requested results from the entry of the system control command HELP, which provides a variety of selectable status reports. These reports are called up using keyboard 44 selectively as indicated by the following table:

- Schedule—Work Schedule Batches
- Batch—Batch/Order Status
- Order—Order/Transaction Status
- Transact—Transaction Status
- Missort—Missort Exceptions
- Route—Sort Line Translations

The reports generated by the entry of the commands at keyboard 43 are provided by the report printer 38 (FIG. 1) and an example of each such report is shown in the attached appendix A which also includes a data dictionary which identifies the manner in which the modules shown in FIGS. 3 through 8 are called up in the program. Basically each report provides a complete current status of batches being processed by the system, orders included in each batch, and particular transactions included in each order. Further, these reports provide an indication as to the percentage of completion, the number of items which are out of stock, and an indication of items not picked, missorted, or in processing. This information is used by the supervisory personnel to monitor on a regular basis the operation of the picking and sorting system, as well as to provide particular commands through terminal 40 for effecting the efficient operation of the system.

In the event that a stockout condition exists where items to be picked are not available, this information can, as previously indicated, be entered through an operator terminal, but also wand 48 can be employed to

scan the labels which have not been applied to articles to be picked and the system automatically assigns a stockout status to the transactions and permits the batch to be completed noting the stockout status in the batch status report. An additional wand scanner (not shown) is provided for excess labels to cover the situation where incomplete case loads of materials are ordered and the efficiency of the picker in filling the cartons is unknown. To cover this situation, several additional labels are provided in the event that a greater number of cartons are required due to the packing efficiency. In the event these extra labels are not needed, they are scanned by the additional wand scanner to maintain accurate status information as to all labels and therefore all transactions within the system.

With the system of the present invention, therefore, high speed article processing can be achieved and in the environment of the preferred embodiment a picking and sorting operation for warehouses is provided with a control system for efficiently monitoring the current operational status of a plurality of batch picks each including several orders, in turn, including several transactions. By assigning each transaction a unique identifying number which is physically associated with the article to be picked in the form of a machine readable label, and which is entered into the control system memory in the form of stored data, the current status of each transaction, order, and batch can be continuously monitored and controlled. The resultant system results in a more efficiently operated system with a higher throughput than previously achievable through conventional picking systems.

It will become apparent to those skilled in the art that various modifications to the preferred embodiment of the invention described and disclosed herein can be made without departing from the spirit or scope of the invention as defined by the appended claims.

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*****00000010
* 00000020
* AUXILIARY FUNCTIONS DATA DICTIONARY 00000030
* 00000040
*****00000050
00000060
I/O # DESCRIPTION 00000070
00000080
01AX CONTROL DATA SET FORMAT 00000090
02AX START-UP CODE 00000100
03AX SYSTEM INITIALIZATION BATCH TABLE 00000110
04AX SYSTEM INITIALIZATION BUFFER IN USE FLAG 00000120
05AX SYSTEM INITIALIZATION DISK BUFFERS 00000130
06AX SYSTEM INITIALIZATION ORDER TABLE 00000140
07AX SYSTEM INITIALIZATION BATCH ERROR FLAG 00000150
08AX SYSTEM INITIALIZATION DISK ERROR FLAG 00000160
09AX SYSTEM CONFIGURATION PROGRAM DIALOGUE 00000170

```

```

*****00000180
* 00000190
* AUXILIARY FUNCTIONS DATA REFERENCES 00000200
* 00000210
*****00000220
00000230
01AX AXINT 00000240
02AX $INIT,$INITIAL 00000250
03AX $INITIAL 00000260
04AX $INITIAL 00000270
05AX $INITIAL 00000280
06AX $INITIAL 00000290
07AX $INITIAL 00000300
08AX $INITIAL 00000310
09AX AXCNFG 00000320

*****00000330
* 00000340
* AUXILIARY DATA DICTIONARY ENTRIES 00000350
* 00000360
*****00000370
00000380
01AX CONTROL DATA SET FORMAT FOR AXINT.P 00000390
00000400
PROGRAM AXINT IS USED TO INITIALIZE DATA SETS TO A PREDEFINED 00000410
HEX VALUE. IT'S OPERATION IS CONTROLLED BY INITIALIZ- 00000420
ATION COMMANDS IN A CONTROL DATA SET. THESE COMMANDS 00000430
CAN BE CREATED USING THE TEXT EDITOR WITH TAB SETS AT 00000440
10 AND 20. THE COMMAND FORMAT IS AS FOLLOWS: 00000450
00000460
COL. 00000470
1.....8 10.....15 20.....22 00000480
DSNAME VOLSER VALUE 00000490
00000500
DSNAME: DATA SET MEMBER NAME 00000510
00000520
VOLSER: VOLUME ID WHERE DATA SET IS LOCATED 00000530
00000540
VALUE: THE DECIMAL VALUE OF A HEX BYTE YOU WANT TO 00000550
INITIALIZE THE FILE TO. FOR INSTANCE 00000560
54 = X'40' 00000570
0 = X'00' 00000580
255 = X'FF' 00000590
00000600
NOTE: ALL PARAMETERS MUST BE SPECIFIED. 00000610
00000620
THE LAST COMMAND MUST BE AN 'END' STARTING IN COL. 1. 00000630
00000640
02AX START-UP CODE 00000650
00000660
THE EDX OPERATING SYSTEM PASSES A PARAMETER TO THE 00000670
INITIALIZATION PROGRAM NAMED "$INITIAL." IF $INITIAL HAS 00000680
A PARAMETER DEFINED IN ITS PROGRAM STATEMENT. (PARM=1) 00000690
THIS PARAMETER CONTAINS A START-UP CODE INDICATING THE 00000700
TYPE OF START-UP BEING PERFORMED. THE START-UP CODE 00000710
VALUES FOLLOW: 00000720
00000730
'0' = MANUAL IPL 00000740
'1' = POWER FAIL/RESTART 00000750
00000760
THIS CODE IS SAVED IN $SYSCOM BY PROGRAM $INITIAL AND 00000770
CHECKED BY THE SYSTEM START-UP PROGRAM. 00000780
00000790
03AX SYSTEM INITIALIZATION BATCH TABLE 00000800
00000810
THIS TABLE IS DEFINED WITHIN MODULE $INITIAL. 00000820
IT'S PURPOSE IS TO HOLD THE BATCH ID, RELATIVE BATCH, 00000830
AND NUMBER OF TRANSACTIONS WITHIN THE BATCH, FOR ALL 00000840
BATCHES WITH A STATUS OF PENDING, ACTIVE OR COMPLETE. 00000850
THIS TABLE IS BUILT AT SYSTEM START UP TIME AND WILL 00000860
BE USED TO UPDATE THE ORDER RECORDS FOUND IN THE ORDER 00000870
DISK FILE. 00000880
00000890
04AX SYSTEM INITIALIZATION BUFFER IN USE FLAG. 00000900
00000910

```

THIS FLAG IS DEFINED WITHIN MODULE \$INITIAL.
 THE FLAG INDICATES WHICH OF THE TWO DISK I/O BUFFERS
 WAS LAST USED. WHEN READING THE TRANSACTION FILE,
 MODULE \$INITIAL WILL ALTERNATE THE USE OF THESE BUFFERS.

05AX

SYSTEM INITIALIZATION DISK BUFFERS.

THERE ARE TWO DISK BUFFERS DEFINED WITHIN MODULE \$INITIAL.
 EACH BUFFER IS CAPABLE OF HOLDING 50 DISK SECTORS.
 THEIR USE IS DEFINED IN 04AX.

06AX

SYSTEM INITIALIZATION ORDER TABLE.

THIS TABLE IS DEFINED WITHIN MODULE \$INITIAL.
 IT CONTAINS COUNTS OF TRANSACTIONS BY THEIR STATUS
 FOR ALL ORDERS IN THE BATCH BEING PROCESSED. THIS
 TABLE IS ONLY USED AT SYSTEM START UP TIME TO UPDATE
 THE ORDER RECORDS ON DISK.

07AX

SYSTEM INITIALIZATION BATCH ERROR FLAG.

THIS FLAG IS DEFINED WITHIN MODULE \$INITIAL.
 WHEN SET, THIS FLAG INDICATES THAT AN UNUSED
 RECORD WAS FOUND IN THE TRANSACTION FILE BEFORE
 ALL TRANSACTIONS IN THE BATCH HAVE BEEN PROCESSED.

08AX

SYSTEM INITIALIZATION DISK ERROR FLAG.

THIS FLAG IS DEFINED WITHIN MODULE \$INITIAL.
 THIS FLAG IS SET WHENEVER A DISK I/O ERROR
 HAS OCCURRED DURING THE UPDATE OF THE ORDER
 FILE AT SYSTEM START UP TIME.

09AX

SYSTEM CONFIGURATION PROGRAM DIALOGUE

WHEN USING THE SYSTEM CONFIGURATION UTILITY, PARAMETERS
 THAT WILL REMAIN THE SAME AS DISPLAYED DO NOT NEED TO BE
 ENTERED. JUST PRESSING THE "ENTER" OR "RETURN" KEY WILL
 INDICATE NO CHANGE. WHEN SPECIFYING MESSAGE ROUTING, ALL
 LOG DEVICES SHOULD BE ENTERED ON THE SAME LINE, SEPARATED
 BY SPACES OR COMMAS. A ZERO (0) IN THE FIRST POSITION IN-
 DICATES THAT THE MESSAGE IS NOT TO BE PRINTED ON ANY LOG
 DEVICE.

00000920
 00000930
 00000940
 00000950
 00000960
 00000970
 00000980
 00000990
 00001000
 00001010
 00001020
 00001030
 00001040
 00001050
 00001060
 00001070
 00001080
 00001090
 00001100
 00001110
 00001120
 00001130
 00001140
 00001150
 00001160
 00001170
 00001180
 00001190
 00001200
 00001210
 00001220
 00001230
 00001240
 00001250
 00001260
 00001270
 00001280
 00001290
 00001300
 00001310
 00001320
 00001330
 00001340
 00001350

 *
 * SCANNER DATA DICTIONARY
 *

I/O #	DESCRIPTION		
01SC	SCSTFLG	SCANNER STOP FLAG	00000010
02SC		INITIALIZATION	00000020
04SC		SCANNER DATA TRANSMISSION	00000030
05SC	STJPBUFQ	STATUS UPDATE BUFFER QUEUE	00000040
06SC		SCANNER ID	00000050
08SC	TSTATUPQ	TRANSACTION STATUS UPDATE QUEUE	00000060
09SC		RECIRCULATION DESTINATION	00000070
10SC	STOPEVT	SCANNER PROGSTOP EVENT	00000080
11SC		ERROR ON READ	00000090
12SC		PASSING INFORMATION TO SCIND	00000100
13SC		ERROR DESTINATION	00000110
14SC	STPNR	STOP ON NO READ FLAG	00000120
15SC		SLOW DOWN EVENT	00000130
16SC	TSTATUPR	RESERVE STATUS UPDATE QUEUE	00000140
17SC	STJPBUFR	RESERVE STATUS UPDATE BUFFER QUEUE	00000150
18SC	SCINDEV	PACKAGE PROCESSING TASK EVENT	00000160
19SC	SCTSUEV	STATUS UPDATE TASK EVENT	00000170
20SC	INDQUEBF	INDUCTION BUFFER QUEUE	00000180
21SC	INDQUE	INDUCTION QUEUE	00000190
22SC		NOT USED	00000200
23SC		ROUTING OF ERROR CONDITIONS	00000210
			00000220
			00000230
			00000240
			00000250
			00000260
			00000270
			00000280
			00000290
			00000300

TO RETURN A BUFFER TO THE QUEUE USE:

NEXTQ STUPBUFQ,LOC,FULL=ROUTINE ,WHERE:

LOC = LABEL ON THE WORD THAT CONTAINS THE ADDRESS OF THE BUFFER.

ROUTINE = NAME OF THE INSTRUCTION TO BRANCH TO IF STUPBUFQ IS FULL.

06SC SCANNER ID'S

SCANNER ID'S USED IN THE SCANNER SJB-SYSTEM ARE:

ID	SCANNER	
		00001070
		00001080
		00001090
		00001100
		00001110
		00001120
		00001130
		00001140
		00001150
		00001160
		00001170
		00001180
		00001190
		00001200
		00001210
		00001220
1	INDUCTION #1	00001230
2	INDUCTION #2	00001240
3	INDUCTION #3	00001250
4	EXCESS REPACK	00001260
5	STOCK OUT	00001270
6	STAGED FOR SHIPMENT	00001280

08SC TRANSACTION STATUS UPDATE QUEUE

THIS TRANSACTION STATUS UPDATE QUEUE IS USED BY SCTSU TO GET THE TRANSACTION NUMBER AND NEW STATUS OF A TRANSACTION TO BE UPDATED. IT CONTAINS ADDRESSES OF BUFFERS THAT CONTAIN THE INFORMATION NEEDED.

TSTATUPQ DEFINEQ COUNT=50

TO PLACE AN ENTRY ON THE QUEUE USE:

NEXTQ TSTATJPQ,LOC,FULL=ROUTINE ,WHERE:

LOC = LABEL ON WORD CONTAINING ADDRESS OF BUFFER

ROUTINE = LABEL OF INSTRUCTION TO BRANCH TO IF QUEUE IS FULL.

TO GET AN ENTRY OFF THE QUEUE USE:

FIRSTQ TSTATJPQ,LOC,EMPTY=ROUTINE ,WHERE:

LJC=LABEL OF A WORD TO RECEIVE ADDRESS OF BUFFER

ROUTINE=LABEL OF INSTRUCTION TO BRANCH TO IF QUEUE IS EMPTY.

09SC RECIRCULATION DESTINATION

THE RECIRCULATION DESTINATION WILL BE DESTINATION ZERO (0).

10SC SCANNER PROGSTOP EVENT

THIS EVENT IS WAITED ON BY THE SCANNER INITIALIZATION ROUTINE AFTER THE SCANNER STOP FLAG HAS BEEN SET. THIS EVENT WILL NEVER BE POSTED, SO A PROGRAM STOP WILL NEVER BE EXECUTED.

11SC INDUCTION SCANNER READ ERROR FLAG

THIS FLAG WILL BE INTERNAL TO THE SCANNER SUB-SYSTEM. WHEN A READ ERROR OCCURS, THE STOP ON 'NO READ' FLAG (14SC) IN SYSCOM IS PASSED AS THE ERROR FLAG. IF NO ERROR IS DETECTED THE READ ERROR FLAG IS SET TO -1.

PASSING INFORMATION TO SCIND

PASSING INFORMATION TO SCIND IS DONE THROUGH THE USE OF A QUEUE. FIRST A BUFFER IS OBTAINED FROM THE INDUCTION BUFFER QUEUE (20SC). THE SCAN INFORMATION IS THEN PUT INTO THIS 8-BYTE BUFFER. THE ADDRESS OF THIS BUFFER IS THEN PUT ON THE INDUCTION QUEUE (21SC).

THE FORMAT OF THE DATA IN THE BUFFER IS:

WORD#	DATA	
		00001700
		00001710
		00001720
		00001730
		00001740
		00001750
		00001760
		00001770
		00001780
		00001790
		00001800
1	SCANNER ID	00001810

2-3	TRANSACTION NUMBER	00001820
4	ERROR FLAG	00001830
13SC	ERROR DESTINATION	00001840
	ONLY SPECIAL DEFAULT TRANSACTION #000000, GOES TO THIS SPECIAL DESTINATION. THIS DESTINATION MAY VARY FROM SYSTEM TO SYSTEM.	00001850
		00001860
		00001870
		00001880
		00001890
		00001900
14SC	STOP ON 'NO READ' FLAG: (STPNR)	00001910
	THIS FLAG WILL RESIDE IN SYSCOM AND BE INITIALIZED BY THE OPERATOR INTERFACE SUB-SYSTEM. IF THIS FLAG IS A "0", ALL 'NO READS' AT INDUCTION ARE SENT TO RECIRCULATION. IF IT IS A "1", THEN NO MESSAGE IS SENT TO THE PSC WHEN A 'NO READ' CONDITION OCCURS. THIS STOPS THE PACKAGE AT THE INDUCTION POINT.	00001920
		00001930
		00001940
		00001950
		00001960
		00001970
		00001980
		00001990
	TO EXAMINE THIS FLAG USE THE FOLLOWING PROCESS:	00002000
	MOVE #1,\$SYSCOM * GET ADDRESS OF SYSCOM	00002010
	MOVE #2,(+SCVR,#1) * GET ADDRESS OF STPNR	00002020
	MOVE FLAG,(0,#2) * GET STPNR	00002030
		00002040
	"FLAG" NOW CONTAINS THE STOP ON 'NO READ' FLAG.	00002050
		00002060
		00002070
15SC	SLOW DOWN FLAG	00002080
	THIS EVENT WILL BE DEFINED IN SYSCOM. IT IS A LINK TO THE PSC COMMUNICATIONS SUB SYSTEM. THIS EVENT WHEN RESET WILL SIGNAL A NEED TO SLOW DOWN THE RATE OF TRANSACTIONS BEING INDUCTED. WHEN POSTED, EVERYTHING RUNS NORMALLY. WHEN RESET, NO INDUCTION MESSAGE IS SENT TO THE PSC.	00002090
		00002100
		00002110
		00002120
		00002130
		00002140
		00002150
		00002160
16SC	RESERVE STATUS UPDATE QUEUE	00002170
	THIS QUEUE IS USED FOR STATUS UPDATES WHEN THE SCANNER SUB SYSTEM GOES INTO SLOW DOWN MODE. IT IS USED THE SAME WAY AS THE PRIMARY UPDATE QUEUE. SEE (08SC).	00002180
		00002190
		00002200
		00002210
		00002220
17SC	RESERVE STATUS UPDATE BUFFER QUEUE	00002230
	THIS QUEUE IS USED TO SUPPLY BUFFERS WHEN THE SCANNER SUB SYSTEM GOES INTO SLOW DOWN MODE. IT IS USED THE SAME WAY AS THE PRIMARY STATUS UPDATE BUFFER QUEUE. SEE 05SC.	00002240
		00002250
		00002260
		00002270
		00002280
		00002290
18SC	PACKAGE PROCESSING TASK EVENT (SCINDEV)	00002300
	THIS EVENT IS POSTED WHEN SOMETHING IS PUT IN THE INDUCTION QUEUE. IT IS RESET WHEN SCIND HAS EMPTIED THE QUEUE. IT IS ALSO POSTED BY SCINIT EVERY 5 SECONDS.	00002310
		00002320
		00002330
		00002340
		00002350
19SC	STATUS UPDATE TASK EVENT (SCTSUEV)	00002360
	THIS EVENT IS POSTED WHEN SOMETHING IS PUT IN THE STATUS UPDATE QUEUE. IT IS RESET WHEN SCTSU HAS EMPTIED THE QUEUE. IT IS ALSO POSTED BY SCINIT EVERY 5 SECONDS.	00002370
		00002380
		00002390
		00002400
		00002410
20SC	INDUCTION BUFFER QUEUE (INDQJEBF)	00002420
	THIS QUEUE IS USED TO SUPPLY BUFFERS FOR SENDING INFORMATION TO SCIND. (SEE 12SC). THERE ARE 50 BUFFERS, EACH A SIZE OF 8 BYTES.	00002430
		00002440
		00002450
		00002460
		00002470
21SC	INDUCTION QUEUE (INDQJE)	00002480
	THIS QUEUE IS USED TO SEND INFORMATION TO SCIND. THIS QUEUE CONTAINS 50 1-WORD ELEMENTS.	00002490
		00002500
		00002510
		00002520
22SC	NOT USED	00002530
		00002540
23SC	ROUTING OF ERROR CONDITIONS	00002550
		00002560

TRANSACTION SCANNED	RECIRC MODE	HOLD MODE	
1. NOT ON FILE	: ERROR CHUTE	: ERROR CHUTE	00002570
2. ON FILE, BATCH PENDING	: *RECIRC	: *RECIRC	00002580
3. ON FILE, BATCH COMPLETE	: *ERROR CHUTE	: *ERROR CHUTE	00002590
4. TRANS# FJJND, BUT BATCH-	:	:	00002600
-IS ADDING OR DELETING	: ERROR CHUTE	: ERROR CHUTE	00002610
5. NO READ (SCANNER ERROR)	: RECIRC	: HOLD (14SC)	00002620
6. TRANS# 00000 (13SC)	: ERROR CHUTE	: ERROR CHUTE	00002630
			00002640
			00002650
			00002660
			00002670
			00002680

* NOTE - CHANGE STATUS TO 'IN SORTATION'

```

*****00000010
* 00000020
* PSC DATA DICTIONARY INDEX 00000030
* 00000040
*****00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380

```

I/O #	DESCRIPTION	
01PS	PSC OUTPUT BUFFER FORMAT	00000090
02PS	PSC OUTPUT QUEUE AND EVENT	00000100
03PS	PSC CALL CODES	00000110
04PS	INDUCT PACKAGE	00000120
05PS	OPEN LINK TO PSC	00000130
06PS	SEND IDLE MESSAGE TO PSC	00000140
07PS	COMMUNICATIONS STATUS FLAG	00000150
08PS	IDLE MESSAGE COUNT	00000160
09PS	RESPONSE COUNT TABLE	00000170
10PS	RESPONSE DUMP FLAG	00000180
11PS	PSC COMMUNICATIONS PANIC CALL	00000190
12PS	COMMUNICATIONS SEQUENCE BYTES	00000200
13PS	NUMBER OF TIME-OUTS	00000210
14PS	INVALID SEQUENCE COUNTER	00000220
15PS	INVALID CHECKSUM COUNTER	00000230
16PS	PSC MESSAGE FORMATTING	00000240
17PS	TRANSMISSION COMPLETION EVENT	00000250
18PS	PSC IDLE MESSAGE FORMAT	00000260
19PS	PSC OPEN LINK MESSAGE FORMAT	00000270
20PS	PSC INDUCT PACKAGE MESSAGE FORMAT	00000280
21PS	TRANSMISSION COMPLETION CODE WORD	00000290
22PS	PSC RESPONSE MESSAGE FORMAT	00000300
23PS	PSC RESPONSE BUFFERS	00000310
24PS	PSC RESPONSE QUEUE	00000320
25PS	PSC RESPONSES	00000330
26PS	ALARM PROCESSING ROUTINE	00000340
27PS	PSC MESSAGE CHECKSUM CALCULATION	00000350
28PS	PSC RESPONSE OVERFLOW EVENT	00000360
29PS	PSC INPUT/OUTPUT EVENT	00000370
30PS	PSC RESPONSE PROCESSING EVENT	00000380

```

*****00000390
* 00000400
* PSC DATA REFERENCES 00000410
* 00000420
*****00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580

```

I/O #	MODULES REFERENCING	
01PS	SCIND, PSPDL, OIPSC	00000470
02PS	SCIND, PSPDL, OIPSC	00000480
03PS	SCIND	00000490
04PS	SCIND	00000500
05PS	OIPSC	00000510
06PS	PSPDL	00000520
07PS	PSOC, PSPAN, OIPSC, SCIND	00000530
08PS	PSOC, PSPDL	00000540
09PS	PSOC, PSAO	00000550
10PS	PSOC, PSAO, PSID	00000560
11PS	PSOC, PSID, PSRP, PSPOL, PSPAN, PSAO, PSCKSM	00000570
12PS	PSOC, PSMF	00000580

13PS	PSDC,PSID	00000590
14PS	PSDC,PSID	00000500
15PS	PSDC,PSID	00000510
16PS	PSDC,PSMF	00000620
17PS	PSDC,PSTD	00000630
18PS	PSMF	00000640
19PS	PSMF	00000650
20PS	PSMF	00000660
21PS	PSID,PSTD	00000670
22PS	PSID,PSRP	00000680
23PS	PSID,PSRP,PSAD	00000690
24PS	PSID,PSRP	00000700
25PS	PSRP,PSAD	00000710
26PS	PSRP,PSAD	00000720
27PS	PSMF,PSCKSM	00000730
28PS	PSID,PSRP,PSJC	00000740
29PS	PSDC,PSID	00000750
30PS	PSJI,PSRP	00000760

```

*****00000770
* 00000780
* PSC DATA DICTIONARY ENTRIES 00000790
* 00000800
*****00000810
00000820
01PS PSC OUTPUT BUFFER FFORMAT 00000830
      * CALL CODE * 00000840
      * PARAMETER 1 * 00000850
      * . * 00000860
      * . * SIZE = 5 WORDS 00000870
      * . * 00000880
      * . * 00000890
      * PARAMETER N * 00000900
      00000910
      FOR VALID CALL CODES SEE: 03PS 00000920
      00000930
      NOTE: ALL BUFFERS ARE TO BE OBTAINED FROM A BUFFER QUEUE 00000940
      DEFINED IN $SYSCOM. 00000950
      00000960
02PS PSC OUTPUT QUEUE 00000970
      00000980
      THE PSC OUTPUT QUEUE IS DEFINED IN $SYSCOM AND IS USED TO PASS 00000990
      THE ADDRESS OF THE CALL BUFFER TO THE PSC COMMUNICATIONS 00010000
      SUB-SYSTEM. AFTER QUEUEING THE CALL BUFFER, THE PSC 00010100
      RUN EVENT ALSO DEFINED IF $SYSCOM MUST BE POSTED. 00010200
      00010300
03PS PSC CALL CODES 00010400
      00010500
      POPN OPEN LINK TO PSC 00010600
      PIDL SEND IDLE MESSAGE TO PSC 00010700
      PIND SEND INDJCT PACKAGE MESSAGE TO PSC 00010800
      00010900
04PS INDJCT PACKAGE 00011000
      00011100
      CODE: PIND 00011200
      00011300
      INPUT BUFFER LAYOUT: 00011400
      00011500
      WORD # CONTENTS 00011600
      00011700
      1 CALL CODE 00011800
      2-3 TRANSACTION NUMBER 00011900
      4 INDJCT SCANNER I.D. 00012000
      5 SORT DESTINATION 00012100
      00012200
05PS OPEN LINK TO PSC 00012300
      00012400
      CODE: POPN 00012500
      00012600
      INPUT BUFFER LAYOUT: 00012700
      00012800
      WORD # CONTENTS 00012900
      00013000
      1 CALL CODE 00013100
  
```


06PS	SEND IDLE MESSAGE TO PSC	00001320
		00001330
	CODE: PIDL	00001340
		00001350
		00001360
		00001370
		00001380
		00001390
		00001400
		00001410
		00001420
07PS	COMMUNICATIONS STATUS FLAG	00001430
		00001440
	THE COMMUNICATIONS STATUS FLAG IS A GLOBAL VARIABLE DEFINED	00001450
	IN \$SYSCOM. IT CONTAINS THE CURRENT STATUS OF THE	00001460
	COMMUNICATIONS LINK BETWEEN DAS AND PSC.	00001470
	THE STATUS VALUES ARE:	00001480
		00001490
	ZERO = "OPEN"	00001500
	OTHER = "CLOSED"	00001510
		00001520
08PS	IDLE MESSAGE COUNT	00001530
		00001540
	THE IDLE MESSAGE COUNT IS A LOCAL VARIABLE DEFINED WITHIN	00001550
	THE PSC COMMUNICATIONS SUB-SYSTEM. IT CONTAINS THE NUMBER	00001560
	OF COMMUNICATION INACTIVITY CYCLES TO BE PERFORMED BEFORE	00001570
	THE PSC REQUIRES AN IDLE MESSAGE TO BE SENT.	00001580
		00001590
09PS	RESPONSE COUNT TABLE	00001600
		00001610
	THE PSC RESPONSE COUNT TABLE CONTAINS ALL POSSIBLE PSC	00001620
	RESPONSE CODES (EXCEPT TRANS. COMPLETE) AND A COUNT THAT	00001630
	INDICATES THE NUMBER OF TIMES DAS HAS RECEIVED THIS	00001640
	RESPONSE FROM THE PSC. IN ADDITION TO THESE COUNTS, ONE	00001650
	ADDITIONAL COUNTER IS DEFINED FOR ALL UNKNOWN RESPONSES	00001660
	RECEIVED FROM THE PSC.	00001670
		00001680
10PS	RESPONSE DUMP FLAG	00001690
		00001700
	THE PSC RESPONSE DUMP FLAG IS A LOCAL VARIABLE DEFINED	00001710
	WITHIN THE PSC COMMUNICATIONS SUB-SYSTEM. WHEN SET,	00001720
	ALL ABNORMAL RESPONSES RECEIVED BY DAS FROM PSC WILL BE	00001730
	OUTPUT ON THE LOG DEVICES.	00001740
		00001750
11PS	PSC COMMUNICATIONS PANIC CALL	00001760
		00001770
	TO PROCESS A PANIC WITHIN THE PSC COMMUNICATIONS SUB-SYSTEM	00001780
	THE USER SHOULD CALL THE SUBROUTINE "PSPAN" WITH THE PANIC CODE	00001790
	AS A PARAMETER. NO RETURN WILL BE MADE FROM THIS CALL.	00001800
		00001810
12PS	COMMUNICATION SEQUENCE BYTES	00001820
		00001830
	THERE ARE TWO COMMUNICATION SEQUENCE BYTES, ONE FOR INPUT	00001840
	FROM THE PSC AND THE OTHER FOR NEXT OUTPUT TO THE PSC. AFTER	00001850
	EACH SUCCESSFUL INTERCHANGE BETWEEN DAS AND PSC, THESE	00001860
	SEQUENCE BYTES WILL BE UPDATED TO REFLECT THE LAST SEQUENCE	00001870
	VALUE RECEIVED AND LAST SEQUENCE VALUE TRANSMITTED. THE	00001880
	INITIAL VALUES OF THESE SEQUENCE BYTES FOLLOW:	00001890
		00001900
	INPUT = EBCDIC "5"	00001910
	OUTPUT = EBCDIC "5"	00001920
		00001930
	THE ALTERNATE SEQUENCE BYTES VALUES = EBCDIC "J".	00001940
		00001950

13PS NUMBER OF TIME-OUTS
 THIS COUNTER IS A LOCAL VARIABLE DEFINED WITHIN THE PSC COMMUNICATIONS SUB-SYSTEM. IT IS INCREMENTED EACH TIME A PSC RESPONSE HAS TIMED OUT.

14PS INVALID SEQUENCE COUNTER
 THIS COUNTER IS A LOCAL VARIABLE DEFINED WITHIN THE PSC COMMUNICATIONS SUB-SYSTEM. IT IS INCREMENTED EACH TIME THE PSC RESPONDS WITH AN INVALID SEQUENCE BYTE.

15PS INVALID CHECKSUM COUNTER
 THIS COUNTER IS A LOCAL VARIABLE DEFINED WITHIN THE PSC COMMUNICATIONS SUB-SYSTEM. IT IS INCREMENTED EACH TIME THE PSC RESPONDS WITH AN INVALID CHECKSUM.

16PS PSC MESSAGE FORMATTING
 TO FORMAT A PSC OUTPUT MESSAGE, THE SUBROUTINE "PSMF" IS CALLED. PARAMETERS PASSED TO THIS SUBROUTINE ARE:
 DATA INPUT FROM PSC OUTPUT QUEUE.
 PARAMETERS RETURNED FROM THIS SUBROUTINE ARE:
 COMPLETION CODE (SUCCESSFUL, NOT SUCCESSFUL).

17PS TRANSMISSION COMPLETION EVENT
 THIS EVENT IS DEFINED WITHIN THE PSC COMMUNICATIONS SUB-SYSTEM AND IS USED TO PASS ONE OF THE FOLLOWING CODES BACK TO THE PSC OUTPUT CONTROL TASK "PSOC."
 RETRY TRANSMISSION
 DATA RECEIVED
 IDLE RECEIVED

18PS PSC IDLE MESSAGE FORMAT

BYTE #	CONTENTS
1	CURRENT OUTPUT SEQUENCE VALUE (EBCDIC)
2	IDLE CONTROL CODE "I" (EBCDIC)
3-13	SPACES (EBCDIC)
14	CHECKSUM (EBCDIC)

NOTE: CHECKSUM = ADDITION OF BYTES 1-13 THEN FORCED TO BE A PRINTABLE CHARACTER.

19PS PSC OPEN LINK MESSAGE FORMAT

BYTE #	CONTENTS
1	CURRENT OUTPUT SEQUENCE VALUE (EBCDIC)
2	OPEN LINK MESSAGE CODE "L" (EBCDIC)
3-13	SPACES (EBCDIC)
14	CHECKSUM

NOTE: CHECKSUM = ADDITION OF BYTES 1-13 THEN FORCED TO BE A PRINTABLE CHARACTER.

20PS PSC INDUCT PACKAGE MESSAGE FORMAT

BYTE #	CONTENTS
1	CURRENT OUTPUT SEQUENCE VALUE (EBCDIC)
2	INDUCT PACKAGE CODE "D" (EBCDIC)
3-8	TRANSACTION NUMBER (EBCDIC)
9-10	INDUCT NUMBER (EBCDIC)
11-13	SORT DESTINATION (EBCDIC)
14	CHECKSUM (EBCDIC)

NOTE: CHECKSUM = ADDITION OF BYTES 1-13 THEN FORCED TO BE A PRINTABLE CHARACTER.

21PS TRANSMISSION COMPLETION CODE EVENT 00002700
 00002710
 00002720
 00002730
 00002740
 00002750
 00002760
 00002770
 00002780
 00002790
 00002800
 00002810
 00002820
 00002830
 00002840
 00002850
 00002860
 00002870
 00002880
 00002890
 00002900
 00002910
 00002920
 00002930
 00002940
 00002950
 00002960
 00002970
 00002980
 00002990
 00003000
 00003010
 00003020
 00003030
 00003040
 00003050
 00003060
 00003070
 00003080
 00003090
 00003100
 00003110
 00003120
 00003130
 00003140
 00003150
 00003160
 00003170
 00003180
 00003190
 00003200
 00003210
 00003220
 00003230
 00003240
 00003250
 00003260
 00003270
 00003280
 00003290
 00003300
 00003310
 00003320
 00003330
 00003340
 00003350
 00003360
 00003370
 00003380
 00003390
 00003400
 00003410

21PS TRANSMISSION COMPLETION CODE EVENT
 THE COMPLETION CODE ECB IS A LOCAL EVENT DEFINED WITHIN
 THE PSC COMMUNICATIONS SUB-SYSTEM. IT IS USED TO HOLD THE
 TRANSMISSION COMPLETION EVENT CODE THAT IS POSTED BACK TO
 THE PSC OUTPUT CONTROL TASK "PSOC."

22PS PSC RESPONSE MESSAGE FORMAT
 BYTE # CONTENTS
 1 SEQUENCE BYTE VALUE (EBCDIC)
 2 RESPONSE MESSAGE TYPE (EBCDIC)
 3-13 RESPONSE PARAMETERS (EBCDIC)
 14 CHECKSUM (EBCDIC)
 15 CARRIAGE RETURN (EBCDIC)
 16 LINE FEED (EBCDIC)
 NOTE: CHECKSUM = ADDITION OF BYTES 1-13 THEN FORCED TO
 BE A PRINTABLE CHARACTER.
 FOR VALID RESPONSE MESSAGE TYPES SEE: 25PS

23PS PSC RESPONSE BUFFERS
 BYTE # CONTENTS
 1 SEQUENCE VALUE RECEIVED (EBCDIC)
 2 RESPONSE MESSAGE TYPE (EBCDIC)
 3-13 RESPONSE PARAMETERS (EBCDIC)
 14 CHECKSUM (EBCDIC)
 NOTE: ALL BUFFERS ARE TO BE OBTAINED FROM A BUFFER QUEUE
 DEFINED WITHIN THE PSC COMMUNICATIONS SUB-SYSTEM.
 FOR VALID RESPONSE MESSAGE TYPES SEE: 25PS

24PS PSC RESPONSE QUEUE
 THIS QUEUE IS DEFINED WITHIN THE PSC COMMUNICATIONS SUB-
 SYSTEM AND IS USED TO PASS THE PSC RESPONSE MESSAGE TO
 THE PSC RESPONSE PROCESSOR ROUTINE.

25PS PSC RESPONSES
 FOR PSC RESPONSES SEE: SOFTWARE MANUAL FOR THE PROGRAMMABLE
 SORT CONTROLLER, SECTION 9, DOCUMENT NUMBER
 7G-01730-600-0094-??.

26PS ALARM PROCESSING ROUTINE
 THIS ROUTINE IS USED TO PRINT MESSAGES ON THE LOG DEVICES
 FOR PSC RESPONSES OTHER THAN A TRANSFER COMPLETE. PARAMETERS
 PASSED TO THIS ROUTINE ARE: DATA RECEIVED FROM THE PSC.

27PS PSC MESSAGE CHECKSUM CALCULATION
 TO CALCULATE A MESSAGE CHECKSUM, CALL ROUTINE 'PSCKSM'
 WITH THE ADDRESS OF THE MESSAGE AND THE ADDRESS WHERE
 THE CHECKSUM IS TO BE PLACED.

28PS PSC RESPONSE OVERFLOW EVENT
 THIS EVENT INDICATES WHEN THERE HAS BEEN A PSC RESPONSE
 OVERFLOW CONDITION DO TO THE RESPONSE QUEUE BEING FULL.
 WHEN THE EVENT IS RESET, ALL NEW TRANSMISSIONS TO THE PSC
 ARE DELAYED UNTIL THE PSC RESPONSE QUEUE IS CLEANED OUT
 AND THE EVENT IS POSTED. THIS EVENT IS DEFINED WITHIN
 THE PSC COMMUNICATIONS SUB-SYSTEM.

29PS PSC INPUT/OUTPUT EVENT

THIS EVENT IF DEFINED WITHIN THE PSC COMMUNICATIONS SUB-SYSTEM. WHEN POSTED, IT INITIATES THE PSC INPUT/OUTPUT SEQUENCE. THIS EVENT IS RESET AT THE END OF A PSC INPUT/OUTPUT SEQUENCE.

30PS PSC RESPONSE PROCESSING EVENT

THIS EVENT ID DEFINED WITHIN THE PDC COMMUNICATIONS SUB-SYSTEM. WHEN POSTED, IT INITIATES THE PSC RESPONSE PROCESSING SEQUENCE. THIS EVENT IS RESET AFTER ALL CURRENT PSC RESPONSES HAVE BEEN PROCESSED.

*****			00000010
*			00000020
*	* ALARM MESSAGE DATA DICTIONARY INDEX		00000030
*			00000040
*****			00000050
I/O #	DESCRIPTION		00000060
01AL	ALMQUE	ALARM MESSAGE QUEUE	00000070
02AL	BUFFQUE	BUFFER POOL QUEUE	00000080
03AL		ALARM MESSAGES	00000090
04AL	TCW	TERMINAL CONTROL WORD	00000100
05AL	OTADTBL	OUTPUT TASK START EVENT TABLE	00000110
06AL	OTEETBL	OUTPUT TASKS END-EVENT TABLE	00000120
07AL		PANIC	00000130
08AL		FATAL ERROR	00000140
09AL	TIMRTSK	TIME OUT TASK	00000150
10AL	DATES	DATE AND TIME SUBROUTINE	00000160
11AL	CRFLG	MESSAGE COMPRESS FLAG	00000170
12AL		NON-COMPRESSIBLE MESSAGES	00000180
13AL	MOVTXT	MESSAGE COMPRESSING SUBROUTINE	00000190
			00000200
			00000210
			00000220

*****			00000230
*			00000240
*	* ALARM MESSAGES DATA REFERENCES		00000250
*			00000260
*****			00000270
I/O #	MODULES REFERENCING		00000280
01AL	ALMPG, SCTSU, SCIND, SCSTO, SCSHP, SCXRP, PSOC, PSRP, PSAO, DICCPA		00000290
	\$INIT, PSIO, DIHLT, DIREC, DICCSA		00000300
02AL	ALMPG, SCTSU, SCIND, SCSTO, SCSHP, SCXRP, PSOC, PSRP, PSAO, DICCSA		00000310
	\$INIT, PSIO, DIHLT, DIREC		00000320
03AL.01	SCIND		00000330
.02	SCIND, SCTSU		00000340
.03	SCIND, SCSHP, SCSTO, SCXRP		00000350
.04	PSAO		00000360
.05	PSOC		00000370
.06	PSAD		00000380
.07	DISTB		00000390
.08	DIENB		00000400
.09	PSOC		00000410
.10	PSAO		00000420
.11	PSAO		00000430
.12	\$INIT		00000440
.13	\$INIT		00000450
.14	PSIO		00000460
.15	PSIO		00000470
.16	PSIO		00000480
.17	PSRP		00000490
.18	PSAD		00000500
.19	DIHLT		00000510
.20	DIREC		00000520
.21	PSAO		00000530
			00000540
			00000550

```

04AL ALMPG,AXCNFG 00000560
05AL ALMPG,AXCNFG 00000570
06AL ALMPG,ALTRM1,ALTRM2,ALTRM3 00000580
07AL ALMPG,SCIND,SCTSU,SCSTO,SCSHP,SCXRP,PSPAN,FMPAN,$INIT,OICCPA, 00000590
    DIREC,OIHLT 00000600
08AL ALMPG,SCIND,SCTSU,SCSTO,SCSHP,SCXRP,PSPAN,FMPAN,$INIT,OICCPA, 00000610
    SCSR1,SCSR2,SCSR3 00000620
09AL ALMPG 00000630
10AL ALMPG 00000640
11AL ALMPG 00000650
12AL ALMPG 00000660
13AL ALMPG 00000670
        00000680
    
```

```

*****00000690
* 00000700
* ALARM MESSAGE DATA DICTIONARY ENRIFS 00000710
* 00000720
*****00000730
    
```

```

01AL ALARM MESSAGE QJEUE 00000740
    ALMQUE 00000750
        00000760
        00000770
        00000780
    
```

```

    THE ALARM MESSAGE QUEUE RESIDES IN SYSCOM. EACH ENTRY 00000790
    IN THIS QUEUE IS 1-WORD LONG AND IS JUST A POINTER 00000800
    TO A DATA BUFFER. THIS 8-WORD BUFFER IS WHERE THE 00000810
    DATA FOR THE MESSAGE IS STORED. THIS BUFFER IS ALSO 00000820
    IN SYSCOM. SEE Doc. 12. 00000830
    
```

```

    THE MAXIMUM NUMBER OF ENTRIES THAT CAN BE PUT IN 00000840
    ALMQUE IS 110. 00000850
    
```

TO PUT A MESSAGE ON THE ALARM QUEUE USE:

```

LABEL NEXTQ *,LOC,FULL=ROUTINE,P1=ALMQUE WHERE: 00000860
    
```

```

    LOC=LABEL ON A WORD CONTAINING THE ADDRESS OF THE BUFFER 00000870
    THAT CONTAINS THE MESSAGE DATA. 00000880
    
```

```

    ROUTINE=LABEL OF STATEMENT TO EXECUTE IF THE QUEUE IS FULL. 00000890
    ALMQUE=THE ADDRESS OF THE ALARM BUFFER 00000900
    
```

TO GET THE ADDRESS OF ALMQUE, USE THE FOLLOWING PROCESS:

```

    MOVE #1,$SYSCOM * GET ADDRESS OF SYSCOM 00000910
    MOVE ALMQUE,(+ALMS,#1) * GET ADDRESS OF ALMQUE 00000920
    
```

```

    00000930
    00000940
    00000950
    00000960
    00000970
    00000980
    00000990
    00010000
    00010010
    00010020
    00010030
    00010040
    00010050
    00010060
    00010070
    00010080
    00010090
    00010100
    00010110
    00010120
    00010130
    00010140
    00010150
    00010160
    00010170
    00010180
    00010190
    00010200
    00010210
    00010220
    00010230
    00010240
    00010250
    00010260
    00010270
    
```

```

02AL BUFFQUE 00010280
    
```

```

    BUFFQUE IS A QUEUE THAT CONTAINS ALARM MESSAGE OUTPUT 00010290
    BUFFERS. THESE BUFFERS ARE 3-WORDS LONG. TO OBTAIN A 00010300
    BUFFER USE: 00010310
    
```

```

    LABEL FIRSTQ *,LOC,EMPTY=ROUTINE,P1=BUFFQUE 00010320
    
```

WHERE:

```

    LOC=NAME OF A VARIABLE TO RECEIVE THE ADDRESS OF THE BUFFER 00010330
    ROUTINE=NAME OF STATEMENT TO BRANCH TO IF BUFFQUE IS EMPTY. 00010340
    BUFFQUE=THE ADDRESS OF THE BUFFER QUEUE 00010350
    
```

```

    00010360
    00010370
    00010380
    00010390
    00010400
    00010410
    00010420
    00010430
    00010440
    00010450
    00010460
    00010470
    00010480
    00010490
    00010500
    00010510
    00010520
    00010530
    00010540
    00010550
    00010560
    00010570
    00010580
    00010590
    00010600
    00010610
    00010620
    00010630
    00010640
    00010650
    00010660
    00010670
    00010680
    00010690
    00010700
    00010710
    00010720
    00010730
    00010740
    00010750
    00010760
    00010770
    00010780
    00010790
    00010800
    00010810
    00010820
    00010830
    00010840
    00010850
    00010860
    00010870
    00010880
    00010890
    00010900
    00010910
    00010920
    00010930
    00010940
    00010950
    00010960
    00010970
    00010980
    00010990
    00011000
    00011010
    00011020
    00011030
    00011040
    00011050
    00011060
    00011070
    00011080
    00011090
    00011100
    00011110
    00011120
    00011130
    00011140
    00011150
    00011160
    00011170
    00011180
    00011190
    00011200
    00011210
    00011220
    00011230
    00011240
    00011250
    00011260
    00011270
    
```

WHEN FILLING THIS BUFFER WITH ALARM MESSAGE DATA USE THE FOLLOWING FORMAT.

```

WORD # CONTENTS 00011280
    
```

```

    1 MSG # SEE DATA DICTIONARY 13 00011290
    2-8 MESSAGE PARAMETERS 00011300
    
```

TO RESTORE A BUFFER USE:

```

    NEXTQ *,LOC,FULL=ROUTINE,P1=BUFFQUE 00011310
    
```

WHERE:

LDC=LABEL ON THE ADDRESS OF THE BUFFER
ROUTINE=INSTRUCTION TO BRANCH TO IF BUFFQUE IS FULL.
BUFFQUE=THE ADDRESS OF THE BUFFER QUEUE

TO OBTAIN THE ADDRESS OF BUFFQUE USE:

```

MOVE #1,$SYSCOM * GET ADDRESS OF SYSCOM
MOVE BUFFQUE,(+ALBF,#1) * GET ADDRESS OF BUFFQUE
    
```

00001280
00001290
00001300
00001310
00001320
00001330
00001340
00001350
00001360
00001370
00001380
00001390
00001400
00001410
00001420
00001430
00001440
00001450
00001460
00001470
00001480
00001490
00001500
00001510
00001520
00001530
00001540
00001550
00001560
00001570
00001580
00001590
00001600
00001610
00001620
00001630
00001640
00001650
00001660
00001670
00001680
00001690
00001700
00001710
00001720
00001730
00001740
00001750
00001760
00001770
00001780
00001790
00001800
00001810
00001820
00001830
00001840
00001850
00001860
00001870
00001880
00001890
00001900
00001910
00001920
00001930
00001940
00001950
00001960
00001970
00001980
00001990
00002000
00002010
00002020

03AL

ALARM MESSAGES

MSG #	MESSAGE
1	TRANSACTION (TRANS #) AT (SCANNER NAME) SCANNER OUT OF BATCH.
2	TRANSACTION (TRANS #) AT (SCANNER NAME) SCANNER NOT ON FILE.
3	SCANNER ERROR AT (SCANNER NAME) SCANNER.
4	DATA AS RECEIVED FROM PSC = (DATA)
5	PSC COMMUNICATIONS LINE OPEN
6	TRANSACTION (TRANS #) MIS-SORTED TO SORT LINE (LINE #)
7	START OF BATCH (BATCH ID)
8	END OF BATCH (BATCH ID)
9	PSC COMMUNICATIONS FAILURE
10	SORTATION LINE (DESTINATION) FULL.
11	UNSUCCESSFUL TRANSFER AT SORT LINE (DESTINATION) FOR TRANSACTION (TRANS #)
12	SYSTEM STARTED
13	POWER FAIL/RESTART
14	INVALID CHECKSUM FROM PSC, DATA = (DATA)
15	INVALID SEQUENCE FROM PSC, DATA = (DATA)
16	PSC TIME-OUT
17	TRANSACTION (TRANS #) ON SORT LINE (LINE #) NOT ON FILE
18	DATA RECEIVED FROM PSC OUT-OF-RANGE, DATA = (DATA)
19	HOLD ON SCANNER ERROR
20	RECIRCULATE ON SCANNER ERROR
21	TRANSACTION (TRANS #) HAS BEEN LOST FROM TRACKING

ALARM MESSAGE INPUT BUFFER

*MESSAGE #1 TRANSACTION OUT OF BATCH

INPUT BUFFER LAYOUT

WORD #	CONTENTS
1	MESSAGE # (1)
2-3	TRANSACTION #
4-5	SCANNER ID
6-8	UNUSED

*MESSAGE #2	TRANSACTION NOT ON FILE	00002030
	INPUT BUFFER LAYOUT	00002040
	WORD # CONTENTS	00002050
		00002060
		00002070
		00002080
	1 MESSAGE # (2)	00002090
	2-3 TRANSACTION #	00002100
	4-5 SCANNER ID	00002110
	6-8 UNUSED	00002120
		00002130
*MESSAGE #3	NO READ AT (SCANNER ID) SCANNER	00002140
	INPUT BUFFER LAYOUT	00002150
	WORD # CONTENTS	00002160
		00002170
		00002180
		00002190
	1 MESSAGE # (3)	00002200
	2-3 UNUSED	00002210
	4-5 SCANNER ID	00002220
	6-8 UNUSED	00002230
		00002240
*MESSAGE #4	PSC DATA MESSAGE	00002250
	INPUT BUFFER LAYOUT	00002260
	WORD # CONTENTS	00002270
		00002280
		00002290
		00002300
	1 MESSAGE # (4)	00002310
	2-8 DATA RECEIVED FROM PSC (EBCDIC)	00002320
		00002330
*MESSAGE #5	PSC COMMUNICATIONS MESSAGE	00002340
	INPUT BUFFER LAYOUT	00002350
	WORD # CONTENTS	00002360
		00002370
		00002380
		00002390
	1 MESSAGE # (5)	00002400
	2-8 UNUSED	00002410
		00002420
*MESSAGE #6	MIS-SORTED CARTON MESSAGE	00002430
	INPUT BUFFER LAYOUT	00002440
	WORD # CONTENTS	00002450
		00002460
		00002470
		00002480
	1 MESSAGE # (6)	00002490
	2-4 TRANSACTION # (EBCDIC)	00002500
	5-6 MIS-SORTED SHIPPING LINE # (EBCDIC)	00002510
	7-8 UNUSED	00002520
		00002530
*MESSAGE #7	START OF BATCH	00002540
	INPUT BUFFER LAYOUT	00002550
	WORD # CONTENTS	00002560
		00002570
		00002580
		00002590
	1 MESSAGE # (7)	00002600
	2-3 BATCH ID (RIGHT JUSTIFIED)	00002610
	4-8 UNUSED	00002620
		00002630
*MESSAGE #8	END OF BATCH	00002640
	INPUT BUFFER LAYOUT	00002650
	WORD # CONTENTS	00002660
		00002670
		00002680
		00002690
	1 MESSAGE # (8)	00002700
	2-3 BATCH ID (RIGHT JUSTIFIED)	00002710
	4-8 UNUSED	00002720
		00002730
MESSAGE #9	PSC COMMUNICATION FAILURE	00002740
	INPUT BUFFER LAYOUT	00002750
		00002760
		00002770

WORD #	CONTENTS	00002780
1	MESSAGE # (9)	00002790
2-8	UNUSED	00002800
		00002810
		00002820
*MESSAGE #10	SHORTATION LINE FULL	00002830
		00002840
	INPUT BUFFER LAYOUT	00002850
		00002860
WORD #	CONTENTS	00002870
1	MESSAGE # (10)	00002880
2-3	DESTINATION (EBCDIC)	00002890
4-8	UNUSED	00002900
		00002910
		00002920
*MESSAGE #11	TRANSFER FAILURE	00002930
		00002940
	INPUT BUFFER LAYOUT	00002950
		00002960
WORD #	CONTENTS	00002970
1	MESSAGE # (11)	00002980
2-4	TRANSACTION # (EBCDIC)	00002990
5-6	DESTINATION (EBCDIC)	00003000
7-8	UNUSED	00003010
		00003020
		00003030
*MESSAGE #12	SYSTEM STARTED	00003040
		00003050
	INPUT BUFFER LAYOUT	00003060
		00003070
		00003080
WORD #	CONTENTS	00003090
1	MESSAGE # (12)	00003100
2-8	UNUSED	00003110
		00003120
*MESSAGE #13	POWER FAIL/RESTART	00003130
		00003140
	INPUT BUFFER LAYOUT	00003150
		00003160
WORD #	CONTENTS	00003170
1	MESSAGE # (13)	00003180
2-8	UNUSED	00003190
		00003200
		00003210
*MESSAGE #14	INVALID CHECKSUM FROM PSC	00003220
		00003230
	INPUT BUFFER LAYOUT	00003240
		00003250
WORD #	CONTENTS	00003260
1	MESSAGE # (14)	00003270
2-8	DATA RECEIVED FROM PSC (EBCDIC)	00003280
		00003290
		00003300
*MESSAGE #15	INVALID SEQUENCE FROM PSC	00003310
		00003320
	INPUT BUFFER LAYOUT	00003330
		00003340
WORD #	CONTENTS	00003350
1	MESSAGE # (15)	00003360
2-8	DATA RECEIVED FROM PSC (EBCDIC)	00003370
		00003380
		00003390
*MESSAGE #16	PSC TIME-OUT	00003400
		00003410
	INPUT BUFFER LAYOUT	00003420
		00003430
WORD #	CONTENTS	00003440
1	MESSAGE # (16)	00003450
2-8	UNUSED	00003460
		00003470
		00003480
*MESSAGE #17	TRANSACTION ON SORT LINE NOT ON FILE	00003490
		00003500
	INPUT BUFFER LAYOUT	00003510

	WORD #	CONTENTS	00003520
			00003530
			00003540
	1	MESSAGE # (17)	00003550
	2-4	TRANSACTION # (EBCDIC)	00003560
	5-6	DESTINATION (EBCDIC)	00003570
	7-8	UNUSED	00003580
			00003590
	*MESSAGE #18	DATA RECEIVED FROM PSC OUT-OF-RANGE	00003600
			00003610
		INPUT BUFFER LAYOUT	00003620
			00003630
	WORD #	CONTENTS	00003640
			00003650
	1	MESSAGE # (18)	00003660
	2-8	DATA RECEIVED FROM PSC (EBCDIC)	00003670
			00003680
	*MESSAGE #19	HOLD ON SCANNER ERROR	00003690
			00003700
		INPUT BUFFER LAYOUT	00003710
			00003720
	WORD #	CONTENTS	00003730
			00003740
	1	MESSAGE # (19)	00003750
	2-8	UNUSED	00003760
			00003770
	*MESSAGE #20	RECIRCULATE ON SCANNER ERROR	00003780
			00003790
		INPUT BUFFER LAYOUT	00003800
			00003810
	WORD #	CONTENTS	00003820
			00003830
	1	MESSAGE # (20)	00003840
	2-8	UNUSED	00003850
			00003860
	*MESSAGE #21	TRANSACTION HAS BEEN LOST FROM TRACKING	00003870
			00003880
		INPUT BUFFER LAYOUT	00003890
			00003900
	WORD #	CONTENTS	00003910
			00003920
	1	MESSAGE # (21)	00003930
	2-4	TRANSACTION # (EBCDIC)	00003940
	5-8	UNUSED	00003950
			00003960

04AL TERMINAL CONTROL WORD (TCW)

THE TERMINAL CONTROL WORDS ARE LOCATED IN \$SYSCOM (SEE 39FM). ASSOCIATED WITH EACH MESSAGE THAT ALMPG WILL OUTPUT IS A TCW WHICH SPECIFIES WHICH TERMINAL(S) ANY SPECIFIC MESSAGE WILL BE SENT TO.

EACH BIT IN THE TCW CORRESPONDS TO AN OUTPUT DEVICE. IF THE BIT IS SET 'ON', TO A ONE, THEN THE ASSOCIATED MESSAGE IS SENT TO THAT DEVICE.

TCW

0/1/2/3/4/5/6/7/8/9/10/11/12/13/14/15 BITS

BIT	OUTPUT DEVICE	00004100
		00004110
		00004120
		00004130
0	LOG DEVICE #1	00004140
1	LOG DEVICE #2	00004150
2	LOG DEVICE #3	00004160
3		00004170
4		00004180
5		00004190
6		00004200
7		00004210
8		00004220
9		00004230
10		00004240
11		00004250
12		00004260
13		00004270

14
15

TERMINAL CONTROL WORD TABLE (MESSAGE ROUTING TABLE)

THE TERMINAL CONTROL WORDS FOR ALL OF THE ALARM MESSAGES ARE IN THE TCW TABLE IN \$SYSCOM. TO REFERENCE A TCW FOR A MESSAGE GET THE ADDRESS OF THE TABLE (SEE 39FM). SUBTRACT ONE FROM THE MESSAGE NUMBER, DOUBLE IT, AND ADD IT TO THIS ADDRESS.

TCW1	TCW FOR MESSAGE 1
TCW1	TCW FOR MESSAGE 2
.	.
.	.
.	.
TCW N	TCW FOR MESSAGE N

05AL OUTPUT TASK START EVENT TABLE (OTADTBL)

THIS TABLE RESIDES IN THE ALMPG PROGRAM. IT IS A TABLE OF ECB ADDRESSES FOR THE OUTPUT TASKS. THESE TASKS ARE USED FOR PRINTING ALARM MESSAGES.

OUTPUT TASK START EVENT TABLE:

OTADTBL	EQU	*
	DATA	A(ALTRM1S)
	DATA	A(ALTRM2S)
	.	.
	.	.
	DATA	A(ALTRMNS)

THERE WILL BE ONE OF THESE EVENTS FOR EACH DEVICE THAT AN ALARM MESSAGE IS TO BE SENT TO.

05AL OUTPUT TASKS END-EVENT TABLE (OTEETBL)

THIS TABLE IS A TABLE OF ADDRESSES. THESE ADDRESSES ARE THE ADDRESSES OF THE ECB'S FOR THE END EVENTS OF THE OUTPUT TASKS.

OUTPUT TASK END EVENT TABLE:

OTEETBL	EQU	*
	DATA	A(ALTRM1E)
	DATA	A(ALTRM2E)
	.	.
	.	.
	DATA	A(ALTRMNE)

07AL PANIC

IF A FATAL SYSTEM ERROR IS ENCOUNTERED, A PROGRAM CAN PANIC BY PUTTING A NUMBER ON THE ALARM MESSAGE QUEUE (ALMQUE) THAT IS LESS THAN THE ADDRESS OF \$SYSCOM. THIS WILL CAUSE A FATAL ERROR MESSAGE TO BE PRINTED ON ALL LOG PRINTERS. THE NUMBER PUT ON THE QUEUE WILL ALSO BE PRINTED.

WHEN USING THIS TECHNIQUE THE PROGRAM PANICING MUST MAKE MAKE SURE THAT IT BRINGS ITSELF DOWN.

FOR A LIST OF PANIC CODES SEE SECTION 08AL.

08AL FATAL ERROR

IF A PROGRAM HAS A FATAL ERROR IT WILL PUT A UNIQUE ERROR CODE ON THE ALARM MESSAGE QUEUE (ALMQUE). WHEN THE ALARM MESSAGE PROGRAM DEQUEUES THE ERROR CODE, IT WILL FORMAT THE ERROR CODE INTO A MESSAGE OF THE FORM:

00004280
00004290
00004300
00004310
00004320
00004330
00004340
00004350
00004360
00004370
00004380
00004390
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00004690
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00004970
00004980
00004990
00005000

FATAL ERROR DD

WHERE:

DD = THE PANIC CODE

THIS MESSAGE WILL BE PRINTED ON ALL LCG DEVICES.

FILE MANAGER FATAL ERRORS

ERROR # MODULE ERROR

ERROR #	MODULE	ERROR
01	FMRWB	DISK ERROR WHILE ADDING BATCH RECORD
02	FMRWB	DISK ERROR WHILE ADDING TRANS. RECORD
03	FMRWB	DISK ERROR WHILE DELETING BATCH
04	FMRWB	DISK ERROR WHILE CHANGING BATCH STATUS
05	FMRWB	DISK ERROR WHILE CHANGING TRANS. STATUS (TERMINALS)
06	FMRWB	DISK ERROR WHILE CHANGING TRANS. STATUS (SCANNERS)
07	FMRWB	DISK ERROR WHILE READING BATCH REC. VIA BATCH I.D.
08	FMRWB	DISK ERROR WHILE READING TRANS. REC. VIA TRANS. NO.
09	FMRWB	DISK ERROR WHILE READING TRANS. DIR. REC.
10	FMRWB	DISK ERROR WHILE READING TRANS. FILE BLOCK
11	FMRWB	DISK ERROR WHILE READING ORDER REC. VIA ORDER I.D.
12	FMRWB	DISK ERROR WHILE READING ORDER REC. VIA REL. ORDER NO.
13	FMRWB	DISK ERROR WHILE CODE START FILE INIT.
14	FMRWB	DISK ERROR WHILE DELETING TRANS. FILE REC.
15	FMATR	INVALID COMPLETION CODE RETURNED
16	FMDL	BATCH DELETION QUEUE FULL
17	FMRWB	INVALID INPUT PARAMETERS

OPERATOR INTERFACE FATAL ERRORS

ERROR # MODULE ERROR

ERROR #	MODULE	ERROR
101	OIWSR	IPM BUFFER QUEUE EMPTY (OICCGI)
102	OIWSR	IPM BUFFER QUEUE FULL (OICCFI)
103	OIWSR	SRB BUFFER QUEUE EMPTY (OICCGI)
104	OIWSR	SRB BUFFER QUEUE FULL (OICCFI)
105	OIWSR	INVALID CALL CODE (OICCGB)
106	OIWSR	INVALID COMPLETION CODE (OICCGB)
107	OIBSR	IPM BUFFER QUEUE EMPTY (OICCGI)
108	OIBSR	IPM BUFFER QUEUE FULL (OICCFI)
109	OIBSR	SRB BUFFER QUEUE EMPTY (OICCGI)
110	OIBSR	SRB BUFFER QUEUE FULL (OICCFI)
111	OIBSR	INVALID CALL CODE (OICCGB)
112	OIBSR	INVALID COMPLETION CODE (OICCGB)
113	OIBSR	INVALID CALL CODE (OICCFI)
114	OIBSR	INVALID RELATIVE ORDER NUMBER (OICCFI)
115	OIBSR	INVALID COMPLETION CODE (OICCFI)
116	OIOSR	IPM BUFFER QUEUE EMPTY (OICCGI)
117	OIOSR	IPM BUFFER QUEUE FULL (OICCFI)
118	OIOSR	SRB BUFFER QUEUE EMPTY (OICCGI)
119	OIOSR	SRB BUFFER QUEUE FULL (OICCFI)
120	OIOSR	INVALID CALL CODE (TRANS. DIR)
121	OIOSR	SBB BUFFER QUEUE FULL (OICCFI)
122	OIOSR	INVALID CALL CODE (OICCGB)
123	OIOSR	INVALID COMPLETION CODE (OICCGB)
124	OIOSR	INVALID CALL CODE (OICCFI)
125	OIOSR	INVALID RELATIVE ORDER NUMBER (OICCFI)
126	OIOSR	INVALID COMPLETION CODE (OICCFI)
127	OIOSR	INVALID CALL CODE (OICCGI)
128	OIOSR	INVALID COMPLETION CODE (OICCGI)
129	OIOSR	INVALID CALL CODE (OICCFI)
130	OIOSR	INVALID RELATIVE TRANSACTION NUMBER (OICCFI)
131	OIOSR	INVALID COMPLETION CODE (OICCFI)
132	OITSR	IPM BUFFER QUEUE EMPTY (OICCGI)
133	OITSR	IPM BUFFER QUEUE FULL (OICCFI)
134	OITSR	SRB BUFFER QUEUE EMPTY (OICCGI)

00005010
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00005680
00005690
00005700
00005710
00005720
00005730

135	OITSR	SRB BUFFER QUEJE FULL (OICCFR)	00005740
136	OITSR	INVALID CALL CODE (TRANS DIR)	00005750
137	OITSR	INVALID COMPLETION CODE (TRANS DIR)	00005760
138	OITSR	INVALID CALL CODE (OICCSB)	00005770
139	OITSR	INVALID COMPLETION CODE (OICCSB)	00005780
140	OITSR	INVALID CALL CODE (OICCGT)	00005790
141	OITSR	INVALID COMPLETION CODE (OICCGT)	00005800
142	OISTB	IPM BUFFER QUEJE EMPTY (OICCGI)	00005810
143	OISTB	IPM BUFFER QUEUE FULL (OICCFI)	00005820
144	OISTB	SRB BUFFER QUEJE EMPTY (OICCGR)	00005830
145	OISTB	SRB BUFFER QUEJE FULL (OICCFR)	00005840
146	OISTB	INVALID CALL CODE (OICCSB)	00005850
147	OISTB	INVALID COMPLETION CODE (OICCSB)	00005860
148	OISTB	INVALID CALL CODE (OICCCB)	00005870
149	OISTB	INVALID STATUS CODE (OICCCB)	00005880
150	OISTB	INVALID COMPLETION CODE (OICCCB)	00005890
151	OISTB	ALARM QUEUE FULL (OICCSA)	00005900
152	OISTB	ALARM BUFFER EMPTY (OICCSA)	00005910
153	OIEVB	IPM BUFFER QUEJE EMPTY (OICCGI)	00005920
154	OIEVB	IPM BUFFER QUEUE FULL (OICCFI)	00005930
155	OIEVB	SRB BUFFER QUEJE EMPTY (OICCGR)	00005940
156	OIEVB	SRB BUFFER QUEJE FULL (OICCFR)	00005950
157	OIEVB	INVALID CALL CODE (OICCSB)	00005960
158	OIEVB	INVALID COMPLETION CODE (OICCSB)	00005970
159	OIEVB	INVALID CALL CODE (OICCCB)	00005980
160	OIEVB	INVALID STATUS CODE (OICCCB)	00005990
161	OIEVB	INVALID COMPLETION CODE (OICCCB)	00006000
162	OIEVB	ALARM QUEUE FULL (OICCSA)	00006010
163	OIEVB	ALARM BUFFER QUEUE EMPTY (OICCSA)	00006020
164	OIASL	DISK WRITE ERROR SORT TABLE (OICCWS)	00006030
165	OIRSL	DISK WRITE ERROR SORT TABLE (OICCWS)	00006040
166	OIRDB	IPM BUFFER QUEJE EMPTY (OICCGI)	00006050
167	OIRDB	IPM BUFFER QUEUE FULL (OICCFI)	00006060
168	OIRDB	SRB BUFFER QUEJE EMPTY (OICCGR)	00006070
169	OIRDB	SRB BUFFER QUEJE FULL (OICCFR)	00006080
170	OIRDB	BATCH ALREADY BEING ADDED (OICCAB)	00006090
171	OIRDB	INVALID CALL CODE (OICCAB)	00006100
172	OIRDB	INVALID COMPLETION CODE (OICCAB)	00006110
173	OIRDB	INVALID CALL CODE (OICCAT)	00006120
174	OIRDB	INVALID COMPLETION CODE (OICCAT)	00006130
175	OIRDB	SRB BUFFER QUEJE EMPTY (OICCGR)	00006140
176	OIRDB	INVALID CALL CODE (OICCRB)	00006150
177	OIRDB	INVALID COMPLETION CODE (OICCRB)	00006160
178	OIRDB	INVALID RELATIVE BATCH NUMBER (OICCCB)	00006170
179	OIRDB	INVALID CALL CODE (OICCCB)	00006180
180	OIRDB	INVALID COMPLETION CODE (OICCCB)	00006190
181	OIPSC	PSC QUEUE FULL	00006200
182	OIPSC	PSC BUFFER QUEJE EMPTY	00006210
183	OIATS	INVALID CALL CODE (OICCSB)	00006220
184	OIATS	INVALID COMPLETION CODE (OICCSB)	00006230
185	OIATS	IPM BUFFER QUEJE EMPTY (OICCGI)	00006240
186	OIATS	IPM BUFFER QUEJE FULL (OICCFI)	00006250
187	OIATS	SRB BUFFER QUEJE EMPTY (OICCGR)	00006260
188	OIATS	SRB BUFFER QUEJE FULL (OICCFR)	00006270
189	OIATS	INVALID CALL CODE (OICCGT)	00006280
190	OIATS	INVALID COMPLETION CODE (OICCGT)	00006290
191	OIATS	INVALID CALL CODE (TRANS STATUS)	00006300
192	OIATS	INVALID STATUS CODE (TRANS STATUS)	00006310
193	OIATS	INVALID COMPLETION CODE (TRANS STATUS)	00006320
194	OIWRB	IPM BUFFER QUEJE EMPTY (OICCGI)	00006330
195	OIWRB	IPM BUFFER QUEUE FULL (OICCFI)	00006340
196	OIMSR	IPM BUFFER QUEJE EMPTY (OICCGI)	00006350
197	OIWRB	SBB BUFFER QUEJE FULL (OICCFI)	00006360
198	OIWRB	SRB BUFFER QUEJE FULL (OICCFR)	00006370
199	OIWRB	INVALID CALL CODE (OICCW)	00006380
200	OIWRB	INVALID COMPLETION CODE (OICCW)	00006390
601	OIWRB	INVALID CALL CODE (OICCFI)	00006400
602	OIWRB	INVALID RELATIVE ORDER NUMBER (OICCFI)	00006410
603	OIWRB	INVALID COMPLETION CODE (OICCFI)	00006420
604	OIWRB	INVALID CALL CODE (OICCGT)	00006430
605	OIWRB	INVALID RELATIVE TRANSACTION NUMBER (OICCGT)	00006440
606	OIWRB	INVALID COMPLETION CODE (OICCGT)	00006450
607	OIWRB	INVALID CALL CODE (OICCSB)	00006460
608	OIWRB	INVALID COMPLETION CODE (OICCSB)	00006470

609	OIDL3	IPM BUFFER QUEJE EMPTY (OICCGI)	00006480
610	OIDL3	IPM BUFFER QUEJE FULL (OICCFI)	00006490
611	OIDL3	SRB BUFFER QUEJE EMPTY (OICCGR)	00006500
612	OIDL3	SRB BUFFER QUEJE FULL (OICCFR)	00006510
613	OIDL3	INVALID CALL CODE (OICCSB)	00006520
614	OIDL3	INVALID COMPLETION CODE (OICCSB)	00006530
615	OIDL3	INVALID RELATIVE BATCH NUMBER (OICCSB)	00006540
616	OIDL3	INVALID CALL CODE (OICCSB)	00006550
617	OIDL3	INVALID COMPLETION CODE (OICCSB)	00006560
618	OICTL	INVALID LOAD CODE (LJADIT)	00006570
619	OIHLT	ALARM BUFFER QUEJE EMPTY	00006580
620	OIHLT	ALARM MESSAGE QUEJE FULL	00006590
621	OIREC	ALARM BUFFER QUEJE EMPTY	00006600
622	OIREC	ALARM MESSAGE QUEJE FULL	00006610
623	FMATR	IPM BUFFER QUEJE EMPTY	00006620
624	FMATR	ERROR RETURNED FROM F.M. WHILE CHANGING BATCH STATUS TO PENDING.	00006630
625	FMATR	IPM BUFFER QUEJE FULL	00006640
626	FMATR	ERROR RETURNED FROM F.M. WHILE DELETING A BATCH.	00006650
627	OITSR	INVALID RELATIVE ORDER	00006660
628	OITSR	INVALID FM CALL CODE	00006670
629	OIMSR	IPM BUFFER QUEJE FULL (OICCFI)	00006700
630	OIMSR	SRB BUFFER QUEJE EMPTY (OICCGR)	00006710
631	OIMSR	SRB BUFFER QUEJE FULL (OICCFR)	00006720
632	OIMSR	SBB BUFFER QUEJE FULL (OICCSF)	00006730
633	OIMSR	INVALID CALL CODE (OICCSB)	00006740
634	OIMSR	INVALID COMPLETION CODE (OICCSB)	00006750
635	OIMSR	INVALID CALL CODE (OICCFI)	00006760
636	OIMSR	INVALID RELATIVE ORDER NUMBER (OICCFI)	00006770
637	OIMSR	INVALID COMPLETION CODE (OICCFI)	00006780
638	OIMSR	INVALID CALL CODE (OICCSF)	00006790
639	OIMSR	INVALID RELATIVE TRANSACTION NUMBER (OICCSF)	00006800
640	OIMSR	INVALID COMPLETION CODE (OICCSF)	00006810
641	OIMSR	INVALID CALL CODE (OICCSF)	00006820
642	OIMSR	INVALID COMPLETION CODE (OICCSF)	00006830
643	OIMSR	INVALID CALL CODE (TRANS. DIR)	00006840

ALARM MESSAGE FATAL ERRORS

ERROR #	MODULE	ERROR	
			00006850
			00006860
			00006870
			00006880
			00006890
201	ALMPG	BAD MESSAGE NUMBER	00006900
202	ALMPG	BUFFER QUEUE FULL	00006910
			00006920

SCANNER INPUT FATAL ERRORS

ERROR #	MODULE	ERROR	
			00006930
			00006940
			00006950
			00006960
301	SCIND	DIRECTORY READ ERROR	00006970
302	SCIND	BUFFER QUEJE EMPTY (ALARM)	00006980
303	SCIND	PSC OUTPUT QUEJE FULL	00006990
304	SCIND	PSC BUFFER QUEJE EMPTY	00007000
305	SCIND	STATUS UPDATE QUEUE FULL	00007010
306	SCIND	STATUS UPDATE BUFFER QUEJE EMPTY	00007020
307	SCTSU	ERROR RETURNED FROM FILE MANAGER	00007030
308	SCTSU	STATUS UPDATE BUFFER QUEJE FULL	00007040
309	SCTSU	ALARM QUEJE FULL	00007050
310	SCTSU	ALARM BUFFER QUEJE EMPTY	00007060
311	SCTSU	FILE MANAGER BUFFER QUEJE FULL	00007070
312	SCTSU	FILE MANAGER BUFFER QUEJE EMPTY	00007080
313	SCSTO	STATUS UPDATE QUEUE FULL	00007090
	SCSHP		00007100
	SCXRP		00007110
314	SCSTO	STATUS UPDATE BUFFER QUEJE EMPTY	00007120
	SCSHP		00007130
	SCXRP		00007140
315	SCSTO	ALARM QUEJE FULL	00007150
	SCSHP		00007160
	SCXRP		00007170
316	SCSTO	ALARM BUFFER QUEJE EMPTY	00007180
	SCSHP		00007190
	SCXRP		00007200

317	SCIND	ALARM QJEU E FJLL	00007210
318	SCSR1	INDUCTION QJEU E FULL	00007220
319	SCSR1	INDUCTION BUFFER QJEU E EMPTY	00007230
320	SCSR2	INDUCTION QJEU E FULL	00007240
321	SCSR2	INDUCTION BUFFER QJEU E EMPTY	00007250
322	SCSR3	INDUCTION QJEU E FULL	00007260
323	SCSR3	INDUCTION BUFFER QJEU E EMPTY	00007270
324	SCIND	INDUCTION BUFFER QJEU E FULL	00007280

PSC COMMUNICATIONS FATAL ERRORS

ERROR #	MODULE	ERROR	
			00007300
			00007310
			00007320
			00007330
401	PSOC	ALARM MESSAGE QJEU E FULL	00007340
402	PSOC	ALARM MSG. BUFFER QJEU E EMPTY	00007350
403	PSIO	PSC RESPONSE QJEU E FULL	00007360
404	PSIO	PSC RESPONSE BUFFER QJEU E EMPTY	00007370
405	PSRP	F.M. BUFFER QJEU E FJLL	00007380
406	PSRP	ALARM MESSAGE QJEU E FULL	00007390
407	PSRP	ALARM MESSAGE BUFFER QJEU E EMPTY	00007400
408	PSRP	UNEXPECTED F.M. COMPLETION CODE RETURNED	00007410
409	PSRP	F.M. BUFFER QJEU E EMPTY	00007420
410	PSRP	PSC RESPONSE BUFFER QJEU E FULL	00007430
411	PSPJL	PSC OUTPUT QJEU E FJLL	00007440
412	PSPJL	PSC OUTPUT BUFFER QJEU E EMPTY	00007450
413	PSOC	PSC OUTPUT BUFFER QJEU E FULL	00007460
414	PSOC	PSC MESSAGE FORMATTING ERROR FOUND	00007470
415	PSAJ	ALARM MESSAGE BUFFER QJEU E FULL	00007480
416	PSAJ	ALARM MESSAGE QJEU E FULL	00007490
417	PSAJ	ALARM MESSAGE BUFFER QJEU E EMPTY	00007500
418	PSIO	ALARM MESSAGE BUFFER QJEU E EMPTY	00007510
419	PSIO	ALARM MESSAGE QJEU E FULL	00007520
420		NOT USED	00007530
421	PSRP	TRANS. NO. CONVERSION ERROR	00007540
			00007550

AUXILIARY FUNCTIONS FATAL ERRORS

ERROR #	MODULE	ERROR	
			00007560
			00007570
			00007580
			00007590
501	\$INIT	FILE MANAGER SUB-SYSTEM LOAD ERROR	00007600
502	\$INIT	FILE MANAGER BUFFER QJEU E EMPTY	00007610
503	\$INIT	UNEXPECTED F.M. COMPLETION CODE RECEIVED	00007620
504	\$INIT	FILE MANAGER BUFFER QJEU E FULL	00007630
505	\$INIT	PSC COMM. SUB-SYSTEM LOAD ERROR	00007640
506	\$INIT	SCANNER INPUT SUB-SYSTEM LOAD ERROR	00007650
507	\$INIT	OPERATOR INTERFACE SUB-SYSTEM LOAD ERROR	00007660
508	\$INIT	ALARM MESSAGE BUFFER QJEU E EMPTY	00007670
509	\$INIT	ALARM MESSAGE QJEU E FULL	00007680
			00007690

09AL TIME OUT TASK (TIMRTSK)

WHEN THE OUTPUT TASKS ARE STARTED A TIME OUT TASK IS ALSO STARTED. IF THE TIME OUT TASK TIMES OUT THE OUTPUT TASKS, IT IS ASSUMED THAT THERE IS A HARDWARE PROBLEM. TO AVOID HANGING UP THE OTHER ALARM MESSAGES THE PRINTTEXTS ARE TERMINATED THROUGH POSTS OF THE PRINT ECB'S OF EACH OUTPUT TASK.

10AL DATE AND TIME SUBROUTINE (DATES)

THIS SUBROUTINE IS CALLED TO PUT THE DATE AND TIME INTO THE ALARM OUTPUT BUFFER. THE SUBROUTINE USES THE FIRST 18 BYTES OF THE BUFFER FOR THE DATE AND TIME. THE FORMAT IS:

```

      . TIME   DATE
      XX:XX:XX XX/XX/XX

```

11AL MESSAGE COMPRESS FLAG

THIS FLAG (CRFLS) IS USED TO INDICATE IF A MESSAGE SHOULD BE COMPRESSED OR NOT. IF THE FLAG IS A '1', THE MESSAGE IS COMPRESSED. IF NOT, IT'S NOT.

12AL NON-COMPRESSABLE MESSAGES

00007700
00007710
00007720
00007730
00007740
00007750
00007760
00007770
00007780
00007790
00007800
00007810
00007820
00007830
00007840
00007850
00007860
00007870
00007880
00007890
00007900
00007910
00007920
00007930
00007940

NON-COMPRESSABLE MESSAGES ARE THOSE DISPLAYING BLANKS THAT ARE SIGNIFICANT. THESE MESSAGES ARE THOSE CONTAINING DATA SENT FROM THE PSC.

13AL MESSAGE COMPRESSING SUBROUTINE (M3VXTX) 00007950
 00007960
 00007970
 00007980
 00007990
 00008000
 THIS SUBROUTINE WILL COMPRESS A STRING OF BLANKS IN AN ALARM MESSAGE DOWN TO A SINGLE BLANK. ALL OCCURRENCES OF BLANKS IN THE MESSAGES ARE COMPRESSED. 00008010
 00008020
 00008030
 00008040

*****00000010
 * 00000020
 * FILE MANAGER DATA DICTIONARY INDEX 00000030
 * 00000040
 *****00000050

I/O #	DESCRIPTION	
01FM	BATCH FILE DEFINITION	00000060
02FM	RESIDENT BATCH FILE DEFINITION	00000070
03FM	TRANSACTION DIRECTORY DEFINITION	00000080
04FM	TRANSACTION FILE DEFINITION	00000090
05FM	ORDER FILE DEFINITION	00000100
06FM	SHORT TRANSLATION FILE DEFINITION	00000110
07FM	DISKETTE FILE DEFINITION	00000120
08FM	FILE MANAGER USER CALLING SEQUENCE	00000130
09FM	FILE MANAGER INPUT BUFFER FORMAT	00000140
10FM	FILE MANAGER CALL CODES	00000150
11FM	FILE MANAGER COMPLETION CODES	00000160
12FM	ADD BATCH RECORD CALL FORMAT	00000170
13FM	ADD TRANSACTION RECORD CALL FORMAT	00000180
14FM	DELETE BATCH CALL FORMAT	00000190
15FM	CHANGE BATCH STATUS CALL FORMAT	00000200
16FM	CHANGE TRANSACTION STATUS FROM TERMINAL CALL FORMAT	00000210
17FM	CHANGE TRANSACTION STATUS FROM SCANNERS CALL FORMAT	00000220
18FM	READ BATCH RECORD VIA BATCH I.O. CALL FORMAT	00000230
19FM	READ TRANSACTION RECORD VIA TRANSACTION NO. CALL FORMAT	00000240
20FM	READ TRANSACTION DIRECTORY RECORD VIA TRANS. NO. CALL FORMAT	00000250
21FM	READ TRANSACTION FILE BLOCK CALL FORMAT	00000260
22FM	READ ORDER RECORD VIA ORDER I.O. CALL FORMAT	00000270
23FM	READ ORDER RECORD VIA RELATIVE ORDER NO. CALL FORMAT	00000280
24FM	COLD START FILE INITIALIZATION CALL FORMAT	00000290
25FM	FILE MANAGER ENTRY EVENT	00000300
26FM	CALL CODE DISPATCH TABLE	00000310
27FM	FILE MANAGER COMPLETION EVENT	00000320
28FM	READ/WRITE FILE BLOCK	00000330
29FM	ORDER I.O. TABLE	00000340
30FM	NOT USED	00000350
31FM	BATCH ADDITION COUNTS	00000360
32FM	BATCH DELETION QUEUE	00000370
33FM	RESIDENT BATCH FILE SEARCH	00000380
34FM	FILE MANAGER RESOURCE DEFINITION	00000390
35FM	FILE MANAGER PANIC CALL	00000400
36FM	RECORD BUFFER QUEUE	00000410
37FM	SECTOR BUFFER QUEUE	00000420
38FM	ADD BATCH COMPLETE	00000430
39FM	SYSTEM PARAMETER FILE DEFINITION	00000440
40FM	SYSTEM ERROR COUNTS	00000450

*****00000490
 * 00000500
 * FILE MANAGER DATA REFERENCES 00000510
 * 00000520
 *****00000530

I/O #	MODULES REFERENCING	
01FM	FMATR,FMABR,FMDEL,FMOID,FMBID,FMCBS,FMRTB,FMINT,FMROL,DICCGB	00000540
	FMOLD,FMADC,JIATS,DIBSR,JIATS,\$INITIAL	00000550
		00000560
		00000570
		00000580

02FM	FMATR,FMABR,FMDEL,FMCBS,FMINT,FMOLD,FMCTS,FMDRB,SCIND,FMADC, DIRDB,OITSR,JIWSR,OICCFB	0000590 0000600
03FM	FMATR,FMCTS,FMID,FMRTD,SCIND,FMTR	0000610
04FM	FMATR,FMCTS,FMRTB,FMID,SCTSU,SCIND,FMOLD,FMRTD,OICCFB, \$INITIAL	0000620 0000630
05FM	FMATR,FMCTS,FMRTB,FMID,SCTSU,OICCFB,OICCGD,FMADC,\$INITIAL	0000640
06FM	FMINT,SCIND,DIASL,OICCWS,DIRSL,DIRSTR,AXCNFG	0000650
07FM	AXBLDD,DIRDB,JIWSR	0000660
08FM	FM,PSRP,SCTSU,\$INIT,FMATR,DIATS,DIRSTR,OICCAB,OICCAT,OICCCB, OICCCB,OICCFB,OICCCB,OICCGD,OICCGT,OICCST,OIOSR	0000670 0000680
09FM	FMGR,FMRTB,PSRP,\$INIT,FMATR,OICCFI,OICCSI,SCTSU	0000690
10FM	SCTSU	0000700
11FM	SCTSU	0000710
12FM	FMABR,OICCAB	0000720
13FM	FMATR,OICCAT	0000730
14FM	FMATR,FMDEL,OICCCB	0000740
15FM	FMCBS,OICCCB	0000750
16FM	FMCTS,DIATS	0000760
17FM	FMCTS,PSRP,SCTSU	0000770
18FM	FMBID,OICCCB	0000780
19FM	FMID,OICCGT	0000790
20FM	FMID,OITSR,OIOSR	0000800
21FM	FMRTB,OICCST	0000810
22FM	FMID,OICCGD	0000820
23FM	FMPOL,OICCFB	0000830
24FM	FMINT,\$INIT	0000840
25FM	FMGR,FM	0000850
26FM	FMGR	0000860
27FM	FMGR,FM	0000870
28FM	FMABR,FMDEL,FMID,FMBID,FMCBS,FMCTS,FMRTB,FMINT FMID,FMRLD,FMRTD,FMRTB,FMADC	0000880 0000890
29FM	FMATR	0000900
30FM		0000910
31FM	FMABR	0000920
32FM	FMDEL,FMINT,FMOLD	0000930
33FM	FMID,FMBID,FMCBS,FMRTB,FMRLD,FMRTB	0000940
34FM	FMOLD,FM	0000950
35FM	FMPAN,FMATR,FMDEL,FMRTB	0000960
36FM	OICCGR,OICCFB	0000970
37FM	OICCGS,OICCFB	0000980
38FM	FMADC,FMATR	0000990
39FM	\$INITIAL,PSOC,DIASL,DIRDB,DIRSL,DIRSTR,AXCNFG,SCINIT,ALMPG,OISER	0001000
40FM	OISER,DIRSE,SCIND,PSRP	0001010

```
*****0001020
* 0001030
* FILE MANAGER DATA DICTIONARY ENTRIES 0001040
* 0001050
*****0001060
```

I/O #	DESCRIPTION	
01FM	BATCH FILE DEFINITION	0001070
	THE BATCH FILE CONTAINS 64 BYTE RECORDS WITH	0001080
	4 RECORDS PER SECTOR. TOTAL FILE FOR 11 BATCHES	0001090
	OCCUPIES 3 SECTORS. THIS FILE RESIDES ON DISK AS	0001100
	A CONTIGUOUS FILE CALLED "BATCH". THE LAYOUT OF	0001110
	EACH 64 BYTE RECORD FOLLOWS.	0001120
		0001130
		0001140
		0001150
		0001160
		0001170
		0001180
		0001190
		0001200
		0001210
		0001220
		0001230
		0001240
		0001250
		0001260
		0001270
		0001280

NOTE: VALID BATCH STATUS VALUES FOLLOW:

EQJ MEANING

BADD ADDING
 BDEL DELETION
 BFRE UNUSED RECORD
 SPEN PENDING
 BACT ACTIVE
 BCMP COMPLETE

00001290
 00001300
 00001310
 00001320
 00001330
 00001340
 00001350
 00001360
 00001370
 00001380
 00001390
 00001400
 00001410
 00001420
 00001430
 00001440
 00001450
 00001460
 00001470
 00001480
 00001490
 00001500
 00001510
 00001520
 00001530
 00001540
 00001550
 00001560
 00001570
 00001580
 00001590
 00001600

02FM RESIDENT BATCH FILE DEFINITION

THE RESIDENT BATCH FILE CONTAINS 4 BYTE RECORDS.
 TOTAL SPACE ALLOCATED FOR 11 BATCHES IS 44 BYTES.
 THIS FILE RESIDES IN \$SYSCOM AS A CONTIGUOUS
 FILE. THE LAYOUT FOR EACH 4 BYTE RECORD FOLLOWS:

INDEX	FIELD	SIZE (BYTES)	TYPE
IBSTA	BATCH STATUS	1	BINARY
IBBID	BATCH I.D.	3	EBCDIC

NOTE: VALID BATCH STATUS VALUES FOLLOW:

-EQU- -MEANING-

BADD ADDING
 BDEL DELETING
 BFRE UNUSED RECORD
 BPEN PENDING
 BACT ACTIVE
 BCMP COMPLETE

03FM TRANSACTION DIRECTORY DEFINITION

TRANSACTION DIRECTORY CONTAINS 4 BYTE RECORDS WITH
 64 RECORDS PER SECTOR. TOTAL DIRECTORY FOR
 999,999 ENTRIES OCCUPIES 15,625 SECTORS. THIS
 DIRECTORY RESIDES ON DISK AS A CONTIGUOUS FILE
 CALLED "XDIR". THE LAYOUT OF EACH 4 BYTE RECORD
 FOLLOWS:

INDEX	FIELD	SIZE (BYTES)	TYPE
IDRLX	REL. TRANSACTION NO.	2	BINARY
IDRLB	REL. BATCH NO.	1	BINARY
IDJST	DESTINATION	1	BINARY

NOTE: IF RECORD IS UNUSED, ALL BITS ARE SET.

04FM TRANSACTION FILE DEFINITION

THE TRANSACTION FILE CONTAINS 18 BYTE RECORDS WITH
 14 RECORDS PER SECTOR. EACH BATCH MAY CONTAIN 9,002
 RECORDS (643 SECTORS). THIS FILE RESIDES ON DISK
 AS A CONTIGUOUS FILE CALLED "XACT" CONTAINING SPACE
 FOR 11 BATCHES FOR A TOTAL OF 7,073 SECTORS.
 THE LAYOUT OF EACH 18 BYTE RECORD FOLLOWS:

INDEX	FIELD	SIZE (BYTES)	TYPE
IXXNO	TRANSACTION NO.	4	BINARY
IXSTA	STATUS	1	BINARY
IXRLO	RELATIVE ORDER NO.	1	BINARY
IXPRD	PRODUCT	5	EBCDIC
IXJST	DESTINATION	1	BINARY
IXSRC	SOURCE	5	EBCDIC

NOTE: VALID TRANSACTION STATUS VALUES FOLLOW:

-EQU- -MEANING-

XNPK NOT PICKED
 XINS IN-SORTATION

00001610
 00001620
 00001630
 00001640
 00001650
 00001660
 00001670
 00001680
 00001690
 00001700
 00001710
 00001720
 00001730
 00001740
 00001750
 00001760
 00001770
 00001780
 00001790
 00001800
 00001810
 00001820
 00001830
 00001840
 00001850
 00001860
 00001870
 00001880
 00001890
 00001900
 00001910
 00001920
 00001930
 00001940
 00001950
 00001960
 00001970
 00001980
 00001990
 00002000
 00002010

XSTG STAGED FOR SHIPPING
 XSTO STOCKED-OUT
 XXSR EXCESS REPACK

00002020
 00002030
 00002040
 00002050
 00002060
 00002070
 00002080
 00002090
 00002100
 00002110
 00002120
 00002130

05FM ORDER FILE DEFINITION

ORDER FILE CONTAINS 20 BYTE RECORDS WITH 12 RECORDS PER SECTOR. TOTAL FILE FOR 36 RECORDS PER BATCH WITH SPACE ALLOCATED FOR 11 BATCHES OCCUPIES 33 SECTORS. THIS FILE RESIDES ON DISK AS A CONTIGUOUS FILE CALLED "ORDER". THE LAYOUT OF EACH 20 BYTE RECORD FOLLOWS.

INDEX	FIELD	SIZE (BYTES)	TYPE
IDRLD	RELATIVE ORDER NO.	2	BINARY
IDDID	ORDER I.D.	6	EBCDIC
IDQNP	QTY. NOT PICKED	2	BINARY
IDQIS	QTY. IN SORTATION	2	BINARY
IDQSS	QTY. STAGED FOR SHIP.	2	BINARY
IDQXS	QTY. EXCESS REPACK	2	BINARY
IDQSD	QTY. STOCK-OUT	2	BINARY
IDQMS	QTY. MIS-SORT	2	BINARY

00002140
 00002150
 00002160
 00002170
 00002180
 00002190
 00002200
 00002210
 00002220
 00002230

06FM SORT TRANSLATION FILE DEFINITION

THE SORT TRANSLATION FILE CONSISTS OF 1 SECTOR OF 256 BYTES ON DISK CALLED "XLATE" OF WHICH THE FIRST 100 BYTES ARE USED. THESE ENTRIES CORRESPOND DIRECTLY TO THE SORT DESTINATION VALUES 0 THRU 99. (DESTINATION 0 IS NOT USED) NORMALLY THESE ENTRIES ARE EACH ZERO. IF A TRANSLATION IS IN EFFECT, THE NEW (NON-ZERO) DESTINATION (1-99) IS PLACED IN THE RELATIVE BYTE CORRESPONDING TO THE OLD DESTINATION.

ON IPL, THE DISK-RESIDENT SORT TRANSLATION FILE DATA IS COPIED INTO \$SYSCOM AS A 50 WORD (100 BYTE) MEMORY RESIDENT TABLE. SUBSEQUENT UPDATES TO THIS DATA ARE MADE TO THE MEMORY-RESIDENT AND DISK-RESIDENT FILES.

00002240
 00002250
 00002260
 00002270
 00002280
 00002290
 00002300
 00002310
 00002320
 00002330
 00002340

07FM DISKETTE FILE DEFINITION

EACH DISKETTE CONTAINING DATA TO BE READ BY THE DAS OR DATA WRITTEN BY THE DAS FOLLOWS THE IBM STANDARD FOR INFORMATION INTERCHANGE (ONE-SIDED, 128 BYTE SECTORS). EACH DISKETTE FILE CONTAINS 24 BYTE RECORDS WITH 5 RECORDS PER SECTOR. RECORDS 1, 2, AND 3 HAVE UNIQUE LAYOUTS WITH RECORDS 4 THRU N HAVING AN IDENTICAL FORMAT. THESE LAYOUTS FOLLOW:

REC.#	FIELD	SIZE (BYTES)	TYPE
1	NO. OF RECS. IN FILE	4	BINARY
	BATCH I.D.	3	EBCDIC
	HEADER (PART 1)	17	EBCDIC
2	HEADER (PART 2)	24	EBCDIC
3	HEADER (PART 3)	9	EBCDIC
	UNUSED	15	EBCDIC
4-N	TRANSACTION NO.	4	BINARY
	ORDER I.D.	6	EBCDIC
	PRODUCT	5	EBCDIC
	DESTINATION	2	EBCDIC
	SOURCE	5	EBCDIC
	STATUS	1	EBCDIC

00002350
 00002360
 00002370
 00002380
 00002390
 00002400
 00002410
 00002420
 00002430
 00002440
 00002450
 00002460
 00002470
 00002480
 00002490
 00002500
 00002510
 00002520
 00002530
 00002540
 00002550
 00002560
 00002570
 00002580
 00002590
 00002600
 00002610
 00002620
 00002630
 00002640
 00002650
 00002660
 00002670
 00002680
 00002690
 00002700
 00002710
 00002720
 00002730
 00002740

NOTE: NO. OF RECS IN FILE IN RECORD 1 INCLUDES RECORDS 1, 2, AND 3.

08FM FILE MANAGER JSER CALLING SEQUENCE

FOR ANY TASK TO USE THE FILE MANAGEMENT PROGRAM THE FOLLOWING CALLING SEQUENCE MUST BE USED:

```
MOVE #1, BUFADR
CALL FM
```

WHERE BUFADR IS THE ADDRESS OF A 16 WORD BUFFER CONTAINING INPUT DATA AND FM IS A SUBROUTINE INCLUDED IN THE CALLING PROGRAM WHICH PERFORMS ALL INTERFACE FUNCTIONS. NOTE: #1 WILL STILL CONTAIN THE INPUT BUFFER ADDRESS UPON RETURN FROM SUBROUTINE FM.

09FM FILE MANAGER INPUT BUFFER FORMAT

```
*      CALL CODE      *
*  COMPLETION CODE  *
*   PARAMETER 1     *
*   PARAMETER 2     *
*   PARAMETER 3     *   SIZE = 16 WORDS
*   PARAMETER 4     *
*   PARAMETER N     *
*                   *
```

FOR VALID CALL CODES SEE: 10FM

NOTE: ALL BUFFERS ARE TO BE OBTAINED FROM A BUFFER QUEUE DEFINED IN \$SYSCOM FOR THIS PURPOSE

10FM FILE MANAGER CALL CODES

```
FABR  ADD BATCH RECORD
FATR  ADD TRANSACTION RECORD
FOEL  DELETE BATCH
FCBS  CHANGE BATCH STATUS
FCTST CHANGE TRANSACTION STATUS (FROM TERMINALS)
FCTSS CHANGE TRANSACTION STATUS (FROM SCANNERS)
FBID  READ BATCH RECORD VIA BATCH I.D.
FTID  READ TRANSACTION RECORD VIA TRANSACTION NUMBER
FRTD  READ TRANSACTION DIRECTORY RECORD VIA TRANSACTION NO.
FRTB  READ TRANSACTION FILE BLOCK
FOID  READ ORDER RECORD VIA ORDER I.D.
FRDL  READ ORDER RECORD VIA RELATIVE ORDER NO.
FINT  COLD START FILE INITIALIZATION
```

11FM FILE MANAGER COMPLETION CODES

```
FOK   SUCCESSFUL
FBX   BATCH FILE SPACE EXHAUSTED
FTX   TRANSACTION FILE SPACE EXHAUSTED
FBA   BATCH ALREADY ON FILE
FTA   TRANSACTION ALREADY ON FILE
FIS   INVALID STATUS
FBN   BATCH NOT ON FILE
FIR   INVALID RELATIVE RECORD/BLOCK NUMBER
FNB   TRANSACTION NOT IN BATCH
FTN   TRANSACTION NOT ON FILE
FOX   ORDER FILE SPACE EXHAUSTED
FON   ORDER NOT ON FILE
FCC   INVALID CALL CODE
FAD   BATCH ALREADY BEING ADDED
```

12FM CODE: FABR ADD BATCH RECORD

INPUT BUFFER LAYOUT:

00002750
00002760
00002770
00002780
00002790
00002800
00002810
00002820
00002830
00002840
00002850
00002860
00002870
00002880
00002890
00002900
00002910
00002920
00002930
00002940
00002950
00002960
00002970
00002980
00002990
00003000
00003010
00003020
00003030
00003040
00003050
00003060
00003070
00003080
00003090
00003100
00003110
00003120
00003130
00003140
00003150
00003160
00003170
00003180
00003190
00003200
00003210
00003220
00003230
00003240
00003250
00003260
00003270
00003280
00003290
00003300
00003310
00003320
00003330
00003340
00003350
00003360
00003370
00003380
00003390
00003400
00003410
00003420
00003430
00003440

63

64

WORD #	CONTENTS	
		00003450
		00003460
1	CALL CODE (FABR)	00003470
2	IGNORE	00003480
3	\$SYSCOM ADR. OF BATCH I.D. DATA (SEE BELOW FOR DATA DEFINITION)	00003490
		00003500
		00003510

RETURN BUFFER LAYOUT:

WORD #	CONTENTS	
		00003520
		00003530
		00003540
		00003550
1	UNCHANGED	00003560
2	COMPLETION CODE	00003570
3	UNCHANGED	00003580
		00003590

COMPLETION CODES:

CODE	MEANING	
		00003600
		00003610
		00003620
		00003630
FDK	SUCCESSFUL	00003640
FBX	BATCH FILE SPACE EXHAUSTED	00003650
FBA	BATCH ALREADY ON FILE	00003660
FCC	INVALID CALL CODE	00003670
FAD	BATCH ALREADY BEING ADDED	00003680
		00003690

BATCH IDENTIFICATION DATA DEFINITION

THE BATCH IDENTIFICATION DATA IS USED TO ADD A BATCH RECORD TO THE BATCH FILE. THE MEMORY ADDRESS OF THIS DATA IS PASSED TO THE FILE MANAGER IN THE ADD BATCH CALL. THE SIZE OF THIS DATA AREA IS 54 BYTES OF CONTIGUOUS MEMORY. THE LAYOUT OF THIS DATA AREA FOLLOWS:

FIELD	SIZE (BYTES)	TYPE	
SPARE	1		00003700
BATCH I.D.	3	EBCDIC	00003710
BATCH HEADER	50	EBCDIC	00003720

13FM

CODE: FATR ADD TRANSACTION RECORD

INPUT BUFFER LAYOUT:

WORD #	CONTENTS	
		00003730
		00003740
1	CALL CODE (FATR)	00003750
2	IGNORE	00003760
3	\$SYSCOM ADR. OF TRANSACTION I.D. DATA (SEE BELOW FOR DATA DEFINITION)	00003770
4	ZERO = FIRST/NEXT REC. NON-ZERO = LAST REC.	00003780
		00003790
		00003800
		00003810
		00003820
		00003830
		00003840
		00003850

RETURN BUFFER LAYOUT:

WORD #	CONTENTS	
		00003860
		00003870
		00003880
		00003890
1	CALL CODE (FATR)	00003900
2	IGNORE	00003910
3	\$SYSCOM ADR. OF TRANSACTION I.D. DATA (SEE BELOW FOR DATA DEFINITION)	00003920
4	ZERO = FIRST/NEXT REC. NON-ZERO = LAST REC.	00003930
		00003940
		00003950
		00003960
		00003970
		00003980
		00003990
		0004000
1	UNCHANGED	0004010
2	COMPLETION CODE	0004020
3-4	UNCHANGED	0004030
		0004040

COMPLETION CODES:

CODE	MEANING	
		0004050
		0004060
		0004070
		0004080
FDK	SUCCESSFUL	0004090
FTA	TRANSACTION ALREADY ON FILE	0004100
FTX	TRANSACTION FILE SPACE EXHAUSTED	0004110
FDX	ORDER FILE SPACE EXHAUSTED	0004120
FCC	INVALID CALL CODE	0004130
		0004140

NOTE: FTA, FTX, FDX CAUSE AN AJTD. DELETE OF THE BATCH BEING ADDED.

0004150
0004160
0004170

TRANSACTION IDENTIFICATION DATA DEFINITION

THE TRANSACTION IDENTIFICATION DATA IS USED TO ADD A TRANSACTION RECORD TO THE TRANSACTION FILE. THE MEMORY ADDRESS OF THIS DATA AREA IS PASSED TO THE FILE MANAGER IN THE ADD TRANSACTION CALL. THE SIZE OF THIS DATA AREA IS 24 BYTES OF CONTIGUOUS MEMORY. THE LAYOUT OF THIS DATA AREA FOLLOWS:

FIELD	SIZE (BYTES)	TYPE
TRANSACTION NO.	4	BINARY
ORDER I.D.	6	EBCDIC
PRODUCT I.D.	6	EBCDIC
SORT DESTINATION	1	BINARY
SOURCE	5	EBCDIC
STATUS	2	BINARY

14FM

CODE: FDEL

DELETE BATCH

INPUT BUFFER LAYOUT:

WORD #	CONTENTS
1	CALL CODE (FDEL)
2	IGNORE
3-N	RELATIVE BATCH NO(S) OF BATCH (ES) TO BE DELETED
N+1	ZERO

RETURN BUFFER LAYOUT:

WORD #	CONTENTS
1	UNCHANGED
2	COMPLETION CODE
3-N+1	UNCHANGED

COMPLETION CODES:

CODE	MEANING
FOK	SUCCESSFUL
FIR	INVALID RELATIVE RECORD/ BLOCK NUMBER
FCC	INVALID CALL CODE

15FM

CODE: FCBS

CHANGE BATCH STATUS

INPUT BUFFER LAYOUT:

WORD #	CONTENTS
1	CALL CODE (FCBS)
2	IGNORE
3-4	BATCH I.D.
5	NEW STATUS (BPEN=PENDING, BACT=ACTIVE, BCMP=COMPLETE)

RETURN BUFFER LAYOUT:

WORD #	CONTENTS
1	UNCHANGED
2	COMPLETION CODE
3-5	UNCHANGED

COMPLETION CODES:

CODE	MEANING
FOK	SUCCESSFUL
FIS	INVALID STATUS

00004180
00004190
00004200
00004210
00004220
00004230
00004240
00004250
00004260
00004270
00004280
00004290
00004300
00004310
00004320
00004330
00004340
00004350
00004360
00004370
00004380
00004390
00004400
00004410
00004420
00004430
00004440
00004450
00004460
00004470
00004480
00004490
00004500
00004510
00004520
00004530
00004540
00004550
00004560
00004570
00004580
00004590
00004600
00004610
00004620
00004630
00004640
00004650
00004660
00004670
00004680
00004690
00004700
00004710
00004720
00004730
00004740
00004750
00004760
00004770
00004780
00004790
00004800
00004810
00004820
00004830
00004840
00004850
00004860
00004870
00004880
00004890
00004900

FBN
FCC

16FM

CODE: FCTST

INPUT BUFFER LAYOUT:

WORD #
1
2
3-4
5-6
7

BATCH NOT ON FILE
INVALID CALL CODE

CHANGE TRANSACTION STATUS
FROM TERMINAL

CONTENTS

CALL CODE
IGNORE
TRANSACTION NUMBER
BATCH I.D.
STATUS (XNPK=NOT PICKED,XINS
=IN SORTATION,XSTO=STOCK-
OUT,XXS=EXCESS REPACK,XSTG
=STAGED FOR SHIPMENT)

00004910
00004920
00004930
00004940
00004950
00004960
00004970
00004980
00004990
00005000
00005010
00005020
00005030
00005040
00005050
00005060
00005070
00005080
00005090
00005100
00005110
00005120
00005130
00005140
00005150
00005160
00005170
00005180
00005190
00005200
00005210
00005220
00005230
00005240
00005250
00005260
00005270
00005280
00005290
00005300
00005310
00005320
00005330
00005340
00005350
00005360
00005370
00005380
00005390
00005400
00005410
00005420
00005430
00005440
00005450
00005460
00005470
00005480
00005490
00005500
00005510
00005520
00005530
00005540
00005550
00005560
00005570
00005580
00005590
00005600
00005610
00005620
00005630

RETURN BUFFER LAYOUT:

WORD #
1
2
3-7

CONTENTS

UNCHANGED
COMPLETION CODE
UNCHANGED

COMPLETION CODES:

CODE
FJK
FIS
FTN
FCC

MEANING

SUCCESSFUL
INVALID STATUS
TRANSACTION NOT IN BATCH
TRANSACTION NOT ON FILE
INVALID CALL CODE

17FM

CODE: FCTSS

INPUT BUFFER LAYOUT:

WORD #
1
2
3-4
5

6-7
8

CONTENTS

CALL CODE (FCTSS)
IGNORE
TRANSACTION NUMBER
STATUS (XINS=IN SORTATION,
XSTO=STOCK-OUT,XXS=EXCESS
REPACK,XSTG=STAGED FOR
SHIPMENT,XMIS=MIS-SORT)
TRANSACTION DIRECTORY ENTRY
FOR STATUS OF IN-SORTATION.
MIS-SORT DEST. FOR STATUS
OF MIS-SORT

RETURN BUFFER LAYOUT:

WORD #
1
2
3-5

CONTENTS

UNCHANGED
COMPLETION CODE
UNCHANGED

COMPLETION CODES:

CODE
FJK
FIS
FTN
FCC

MEANING

SUCCESSFUL
INVALID STATUS
TRANSACTION NOT ON FILE
INVALID CALL CODE

18FM

CODE: FBID

READ BATCH RECORD VIA
BATCH I.D.

00005640
00005650
00005660
00005670
00005680
00005690
00005700
00005710
00005720
00005730
00005740
00005750
00005760
00005770

INPUT BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-4
5

CALL CODE (FBID)
IGNORE
BATCH I.D.
RECORD BUFFER ADDRESS (IN
\$SYSCOM)(BUFFER SIZE > OR =
64 BYTES)

00005780
00005790
00005800
00005810
00005820
00005830
00005840

RETURN BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-5

UNCHANGED
COMPLETION CODE
UNCHANGED

00005850
00005860
00005870
00005880
00005890
00005900
00005910
00005920

COMPLETION CODES:

CODE

MEANING

FBK
FBN
FCC

SUCCESSFUL
BATCH NOT ON FILE
INVALID CALL CODE

00005930
00005940
00005950
00005960

19FM

CODE: FTID

READ TRANSACTION RECORD
VIA TRANSACTION NO.

00005970
00005980
00005990
00006000
00006010
00006020
00006030
00006040
00006050
00006060
00006070

INPUT BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-4
5

CALL CODE (FTID)
IGNORE
TRANSACTION NUMBER
RECORD BUFFER ADDRESS (IN
\$SYSCOM)(BUFFER > OR = 18
BYTES)

RETURN BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-5

UNCHANGED
COMPLETION CODE
UNCHANGED

00006150
00006160
00006170
00006180
00006190
00006200
00006210
00006220

COMPLETION CODES:

CODE

MEANING

FBK
FTN
FCC

SUCCESSFUL
TRANSACTION NOT ON FILE
INVALID CALL CODE

00006230
00006240
00006250

20FM

CODE: FRTD

READ TRANSACTION DIRECTORY
RECORD VIA TRANSACTION NO.

00006260
00006270

INPUT BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-4
5

CALL CODE (FRTD)
IGNORE
TRANSACTION NUMBER
RECORD BUFFER ADDRESS (IN
\$SYSCOM)(BUFFER > OR = 4 BYTES)

00006280
00006290
00006300
00006310
00006320
00006330
00006340
00006350
00006360

RETURN BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-5UNCHANGED
COMPLETION CODE
UNCHANGED

COMPLETION CODES:

CODE

MEANING

FOK
FTN
FCCSUCCESSFUL
TRANSACTION NOT ON FILE
INVALID CALL CODE

21FM

CODE: FRTB

READ TRANSACTION FILE BLOCK

INPUT BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-4
5
6CALL CODE (FRTB)
IGNORE
BATCH I.D.
RELATIVE BLOCK NUMBER (1-NO.
OF TRANS. BLOCKS USED)
BLOCK BUFFER ADDRESS (IN
\$SYSCOM)(BUFFER > OR = 256 BYTE

RETURN BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-6UNCHANGED
COMPLETION CODE
UNCHANGED

COMPLETION CODES:

CODE

MEANING

FOK
FBN
FIR
FCCSUCCESSFUL
BATCH NOT ON FILE
INVALID RELATIVE RECORD/
BLOCK NUMBER
INVALID CALL CODE

22FM

CODE: FOID

READ ORDER RECORD VIA
ORDER I.D.

INPUT BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-4
5-7
8CALL CODE (FOID)
IGNORE
BATCH I.D.
ORDER I.D.
RECORD BUFFER ADDRESS (IN
\$SYSCOM)(BUFFER SIZE > OR =
18 BYTES)

RETURN BUFFER LAYOUT:

WORD #

CONTENTS

1
2
3-8UNCHANGED
COMPLETION CODE
UNCHANGED

COMPLETION CODES:

00006370
00006380
00006390
00006400
00006410
00006420
00006430
00006440
00006450
00006460
00006470
00006480
00006490
00006500
00006510
00006520
00006530
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00006550
00006560
00006570
00006580
00006590
00006600
00006610
00006620
00006630
00006640
00006650
00006660
00006670
00006680
00006690
00006700
00006710
00006720
00006730
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00006790
00006800
00006810
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00006880
00006890
00006900
00006910
00006920
00006930
00006940
00006950
00006960
00006970
00006980
00006990
00007000
00007010
00007020
00007030
00007040
00007050
00007060
00007070
00007080
00007090

26FM FILE MANAGER CALL CODE DISPATCH TABLE FORMAT

* ADRS. OF WDRK ROUTINE FOR CALL CODE 1 *
 * ADRS. OF WDRK ROUTINE FOR CALL CODE 2 *
 * ADRS. OF WDRK ROUTINE FOR CALL CODE 3 *
 * ADRS. OF WDRK ROUTINE FOR CALL CODE 4 *
 * ADRS. OF WDRK ROUTINE FOR CALL CODE 5 *
 * ADRS. OF WDRK ROUTINE FOR CALL CODE N *

27FM FILE MANAGER COMPLETION EVENT

THIS EVENT IS DEFINED IN \$SYSCOM AND IS USED TO PASS THE ADDRESS OF A BUFFER CONTAINING THE FILE MANAGER OUTPUT DATA BACK TO THE CALLER. THIS BUFFER IS THE SAME ONE USED BY THE CALLER FOR THE FILE MANAGER ENTRY EVENT.

28FM READ/WRITE FILE BLOCK

TO READ OR WRITE A SECTOR FROM/TO ONE OF THE DISK FILES SPECIFIED BELOW, SJBRoutine FMRWB IS CALLED WITH THE FOLLOWING INPUT PARAMETERS:

CALL PARAMETERS

READ/WRITE INDICATOR
 BUFFER ADDRESS (256 BYTE BUFFER) FOR READ/WRITE
 RELATIVE BLOCK NUMBER
 FILE CODE (SEE BELOW)
 BUFFER ADDRESS SUPPLIED BY FM USER

FILE CODES:

CODE	FILE	
XDIRF	TRANSACTION DIRECTORY	00008100
XACTF	TRANSACTION FILE	00008100
BATCHF	BATCH FILE	00008100
ORDERF	ORDER FILE	00008110
XLATEF	TRANSLATION FILE	00008120
ORDERI	BUILD RESIDENT ORDER FILE	00008130
ORDERU	UPDATE RESIDENT ORDER FILE	00008140
		00008150
		00008160
		00008170
		00008180
		00008190
		00008200
		00008210
		00008220
		00008230
		00008240
		00008250
		00008260
		00008270
		00008280
		00008290
		00008300
		00008310
		00008320
		00008330
		00008340
		00008350
		00008360
		00008370
		00008380
		00008390
		00008400
		00008410
		00008420
		00008430
		00008440
		00008450
		00008460
		00008470
		00008480
		00008490
		00008500
		00008510
		00008520
		00008530
		00008540

29FM ORDER I.D. TABLE

THE ORDER I.D. TABLE IS A LOCAL TABLE CONTAINING ORDER I.D.'S AND COUNTS OF TRANSACTIONS PER ORDER FOR ALL ORDERS FOUND IN BATCH BEING ADDED TO SYSTEM.

31FM BATCH ADDITION COUNTS

TWO COUNTERS ARE USED BY ADD TRANSACTION (FMATR) TO INDICATE THE NO. OF TRANSACTIONS AND NO. OF ORDERS ADDED THUS FAR IN THE BATCH ADDITION PROCESS.

32FM BATCH DELETION QUEUE

A FIFO QUEUE OF ONE-WORD ENTRIES CONTAINING RELATIVE BATCH NOS. OF BATCHES TO BE DELETED BY FMOLD.

33FM RESIDENT BATCH FILE SEARCH

TO OBTAIN THE RELATIVE BATCH NO. GIVEN A BATCH I.D. CALL FMSRB WITH THE BATCH I.D. AS A PARAMETER. FMSRB SEARCHES THE RESIDENT BATCH FILE AND RETURNS EITHER THE RELATIVE BATCH NO. OR A ZERO VALUE IF THE BATCH IS NOT ON FILE.

34FM FILE MANAGER RESOURCE DEFINITION 00008550
 00008560
 A QUEUE CONTROL BLOCK DEFINED IN \$SYSCOM USED TO
 FORCE SERIAL USE OF THE FILE MANAGER PROGRAM. 00008570
 00008580
 00008590

35FM FILE MANAGER PANIC CALL 00008600
 00008610
 TO PROCESS A PANIC WITHIN THE FILE MANAGER SUB-SYSTEM
 THE USER SHOULD CALL THE SUBROUTINE FMPAN WITH THE
 PANIC CODE IN REGISTER 1. (#1) NO RETURN WILL BE
 MADE FROM THIS CALL. 00008620
 00008630
 00008640
 00008650
 00008660

36FM RECORD BUFFER QUEUE 00008670
 00008680
 THE RECORD BUFFER QUEUE IS DEFINED IN \$SYSCOM
 AND CONTAINS A POOL OF 64 BYTE BUFFERS. 00008690
 00008700
 THESE BUFFERS ARE USED TO PASS COPIES OF RECORDS
 TO AND FROM THE FILE MANAGER. 00008710
 00008720
 00008730

37FM SECTOR BUFFER QUEUE 00008740
 00008750
 THE SECTOR BUFFER QUEUE IS DEFINED IN \$SYSCOM
 AND CONTAINS A POOL OF 256 BYTE BUFFERS. 00008760
 00008770
 THESE BUFFERS ARE USED TO PASS COPIES OF
 DISK FILE SECTORS TO AND FROM THE FILE MANAGER. 00008780
 00008790
 00008800

38FM CODE: FADC ADD BATCH COMPLETE 00008810
 00008820
 INPUT BUFFER LAYOUT: 00008830
 00008840
 WORD # CONTENTS 00008850
 00008860
 1 CALL CODE (FADC) 00008870
 2 IGNORE 00008880
 00008890
 RETURN BUFFER LAYOUT: 00008900
 00008910
 WORD # CONTENTS 00008920
 00008930
 1 UNCHANGED 00008940
 2 COMPLETION CODE 00008950
 00008960
 00008970
 00008980
 00008990
 00009000
 00009010
 00009020
 00009030
 00009040

COMPLETION CODES: 00009050
 00009060
 00009070
 00009080
 00009090
 00009100
 00009110
 00009120

CODE	MEANING
FJK	SUCCESSFUL
FBN	BATCH NOT ON FILE
FCC	INVALID CALL CODE

00009130
 00009140
 00009150
 00009160
 00009170
 00009180
 00009190
 00009200
 00009210
 00009220
 00009230
 00009240
 00009250

39FM SYSTEM PARAMETER FILE DEFINITION
 THE SYSTEM PARAMETER FILE CONSISTS OF 1 SECTOR OF 256
 BYTES ON DISK CALLED "PARMFILE" OF WHICH THE FIRST 58
 BYTES ARE USED. THE LAYOUT OF THIS FILE FOLLOWS.
 00009100
 00009110
 00009120
 00009130
 00009140
 00009150
 00009160
 00009170
 00009180
 00009190
 00009200
 00009210
 00009220
 00009230
 00009240
 00009250

INDEX	FIELD	SIZE(BYTES)	TYPE
SNOS	# OF SHIPPING LINES	2	BINARY
SREC	RECIRC DESTINATION	2	BINARY
SERD	ERROR DESTINATION	2	BINARY
SNOI	# OF INDUCTION LINES	2	BINARY
SOUT	# OF STOCK-JIT SCANNERS	2	BINARY
XRPK	# OF EXCESS REPACK SCNRS	2	BINARY
SHIP	# OF SHIPPING SCANNERS	2	BINARY
LOGD	# OF LOG DEVICES	2	BINARY
LOGMR	LOG MESSAGE ROUTING TABLE	42	BINARY

FOR FURTHER DEFINITION OF THE LOG MESSAGE ROUTING TABLE
 ENTRIES, SEE 04AL.

ON IPL, THE DISK-RESIDENT SYSTEM PARAMETER FILE DATA IS COPIED INTO \$SYSCJM AS A 29 WORD (58 BYTE) MEMORY RESIDENT TABLE. THE INDEX INTO THE MEMEORY RESIDENT TABLE IS THE SAME AS THE DISK FILE INDEX.

00009260
00009270
00009280
00009290
00009300
00009310
00009320
00009330
00009340
00009350
00009360
00009370
00009380
00009390
00009400
00009410
00009420
00009430
00009440
00009450
00009460

40FM

SYSTEM ERROR COUNTS

THE SYSTEM ERROR COUNTS RESIDE IN \$SYSCJM AND CONTAIN THE NUMBER OF ERRORS FOR EACH CONDITION. THE MIS-SORT COUNT TABLE CONTAINS 1 WORD FOR EACH SORT LINE IN THE SYSTEM. THESE COUNTS CAN BE CLEARED BY OPERATOR COMMAND OR BY AN IPL. THE LAYOUT OF THESE COUNTS FOLLOJW:

Table with 4 columns: INDEX, FIELD, SIZE(BYTES), TYPE. Rows include MSRT, NRED, RJCT, PREC, SCRC with their respective field names, sizes, and binary types.

*****00000010
* 00000020
* OPERATOR INTERFACE DATA DICTIONARY INDEX 00000030
* 00000040
*****00000050

Table with 2 columns: I/O #, DESCRIPTION. Lists various system components and their I/O addresses from 010I to 420I.

430I	BATCH VERIFICATION DISPLAY LAYOUT	00000510
440I	TRANSACTION VERIFICATION DISPLAY LAYOUT	00000520
450I	CHANGE BATCH STATUS CALL	00000530
460I	WRITE SORT TRANSLATION TABLE CALL	00000540
470I	HELP COMMAND LAYOUT	00000550
480I	GET IPM CALL	00000560
490I	GET RECORD BUFFER CALL	00000570
500I	FREE IPM CALL	00000580
510I	FREE RECORD BUFFER CALL	00000590
520I	SEND ALARM MESSAGE CALL	00000600
530I	DELETE BATCH CALL	00000610
540I	PANIC CALL	00000620
550I	GET BLOCK BUFFER CALL	00000630
560I	FREE BLOCK BUFFER CALL	00000640
570I	DISKETTE COMMAND WORD FLAG DEFINITIONS	00000650
580I	SYSTEM ERROR REPORT LAYOUT	00000660
590I	SYSTEM ERROR REPORT LOCAL DATA	00000670
600I	MIS-SORT REPORT LOCAL DATA	00000680
610I	MIS-SORT REPORT LAYOUT	00000690
		00000700

```
*****00000710
* 00000720
* OPERATOR INTERFACE MODULE REFERENCES 00000730
* 00000740
*****00000750
00000760
I/O # DESCRIPTION 00000770
010I 0ICTL 00000780
020I:1 0ICTL 00000790
:2 0ICCB(0IRDB) 00000800
:3 0ICCAT(0IRDB) 00000810
:4 0ICCB(0IRDB) 00000820
:5 0ICCAT(0IRDB) 00000830
:6 0IATS 00000840
:7 0IBSR,0IOSR,0ITSR,0IWRB,0ISTB,0IENB,0IDLB,0IATS,0IMSR 00000850
:8 0ITSR,0IATS 00000860
:9 0ICCAT(0IRDB) 00000870
:10 0ICCGO(0IOSR),0ICCGO(0IMSR) 00000880
:11 0IASL,0IRSL 00000890
:12 0IBSR,0IOSR,0ICCGO(0IOSR),0ICCGO(0IMSR),0IMSR 00000900
:13 0IATS,0ITSR 00000910
:14 0ICCRB(0IRDB),0ICCW(0IWRB) 00000920
:15 0IRDB 00000930
:16 0IRDB 00000940
:17 0IDLB 00000950
:18 0IDLB 00000960
:19 0IATS 00000970
:20 INITIAL 00000980
:21 INITIAL 00000990
:22 INITIAL 0001000
:23 0IASL 0001010
:24 0IRSL 0001020
030I:1 0IWSR,0IBSR,0IOSR,0ITSR,0ISTR,0ISER,0IMSR 0001030
:2 0IBSR,0IOSR,0IWRB,0ISTB,0IENB,0IDLB,0IATS,0IMSR 0001040
:3 0IOSR,0IMSR 0001050
:4 0IOSR,0IMSR 0001060
:5 0IOSR 0001070
:6 0IRDB,0IWRB 0001080
:7 0IRDB 0001090
:8 0IWRB 0001100
:9 0ISTB 0001110
:10 0IENB 0001120
:11 0IDLB 0001130
:12 0IDLB 0001140
:13 0ITSR,0IATS 0001150
:14 0IATS 0001160
:15 0IASL,0IRSL 0001170
:16 0IASL 0001180
:17 0IATS 0001190
:18 0ISER 0001200
0001210
```

:19	OIRSE	00001220
:20	OIRSE	00001230
:21	OIRSE	00001240
:22	OIRSE	00001250
:23	OIRSE	00001260
:24	OIRSE	00001270
040I:	OICTL	00001280
:2	OICTL	00001290
:3	OIWSR	00001300
:4	OIWSR,OIDLB	00001310
:5	OIBSR	00001320
:6	OIOSR	00001330
:7	OITSR	00001340
:8	OISTR	00001350
:9	OIRDB	00001360
:10	OIWRB	00001370
:11	OISTB	00001380
:12	OIENB	00001390
:13	OIDLB	00001400
:14	OIATS	00001410
:15	OIASL	00001420
:16	OIRSL	00001430
:17	OIPSC	00001440
:18	OIPSC	00001450
:19	OIREC	00001460
:20	OIHLT	00001470
:21	OIHLP	00001480
:22	OIASL,OIRSL	00001490
:23	OIRDB	00001500
:24	OIWRB	00001510
:25	OIRDB,OIWRB	00001520
:25	OIATS	00001530
:27	INIT	00001540
:28	INIT	00001550
:29	INITIAL	00001560
:30	INITIAL	00001570
:31	OISER	00001580
:32	OISER	00001590
:33	OIRSE	00001600
:34	OIRSE	00001610
:35	OIMSR	00001620
:36	OIMSR	00001630
:37	OIOSR	00001640
050I	OIWSR	00001650
060I	OIBSR	00001660
070I	OIOSR	00001670
080I	OITSR	00001680
090I	OISTR	00001690
100I	OICTL,OIWSR,OIBSR,OIOSR,OITSR,OISTR,OIRDB,OIWRB,OISTB,OIENB, OIDLB,OIATS,OIASL,OIRSL,OIPSC,OIMSR	00001700
110I	OICTL,OIRDB,OIWRB	00001710
120I	OIWSR,OIBSR,OIOSR,OITSR,OIRDB,OIWRB,OISTB,OIENB,OIDLB,OIATS, OIMSR	00001720
130I	OIBSR,OIOSR,OIWRB,OIMSR	00001730
140I	OIOSR,OITSR,OIRDB,OIWRB,OIATS,OIMSR	00001740
150I	OIWSR	00001750
160I	OICCFB,OIBSR,OIWRB,OISTB,OIENB,OIDLB,OIATS,OIMSR	00001760
170I	OIBSR,OIOSR,OIWRB,OIDLB,OIMSR	00001770
180I	OIOSR,OIMSR	00001780
190I	OIOSR,OIWRB,OIMSR	00001790
200I	OICCFB	00001800
210I	OITSR,OIATS	00001810
220I	OIRDB	00001820
230I	OIWRB	00001830
240I	OIRDB	00001840
250I	OIRDB	00001850
260I	OICTL	00001860
270I	OIWSR	00001870
280I	OIBSR	00001880
290I	OIOSR	00001890
300I	OITSR	00001900
310I	OISTR	00001910
320I	OIRDB	00001920
		00001930
		00001940

330I	OIWRB	00001950
340I	OISTB	00001960
350I	OIENB	00001970
360I	OIDLB	00001980
370I	OIASL	00001990
380I	OIRSL	00002000
390I	OIPSC	00002010
400I	OIATS	00002020
410I	OIREC	00002030
420I	OILT	00002040
430I	OIRDB,OIWRB,OISTB,OIENB,OIDLB,OIATS	00002050
440I	OIATS	00002060
450I	OISTB,OIENB	00002070
460I	OIASL,OIRSL	00002080
470I	OIHLP	00002090
480I	OIWSR,OIBSR,OIOSR,OITSR,OIRDB,OISTB,OIENB,OIDLB,OIATS,OIWRB, OIMSR	00002100
490I	OIWSR,OIBSR,OIOSR,OITSR,OIRDB,OISTB,OIENB,OIDLB,OIATS,OIMSR	00002120
500I	OIWSR,OIBSR,OIOSR,OITSR,OIRDB,OISTB,OIENB,OIDLB,OIATS,OIWRB, OIMSR	00002130
510I	OIWSR,OIBSR,OIOSR,OITSR,OIRDB,OISTB,OIENB,OIDLB,OIATS,OIMSR	00002150
520I	OISTB,OIENB	00002160
530I	OIDLB,OIRDB	00002170
540I	OITSR,OIPSC,OIATS,OICCGB,OICCCB,OICCRB,OICCAB,OICCWB,OICCFD, OICCGO,OICCST,OICCGT,OICCAT,OICCWS,OICCSA,OICCDB,FMATR,OIMSR, OIOSR	00002180
550I	OICCST	00002190
560I	OICCST	00002200
570I	OIRDB,OIWRB	00002210
580I	OISER	00002220
590I	OISER	00002230
600I	OIMSR	00002240
610I	OIMSR	00002250
		00002260
		00002270
		00002280

*****00002290
 * 00002300
 * OPERATOR INTERFACE DATA DICTIONARY ENTRIES 00002310
 * 00002320
 *****00002330

I/O #	DESCRIPTION		
010I	OPERATOR INTERFACE COMMANDS:		00002340
	CATEGORY:	COMMAND: DESCRIPTION:	00002350
			00002360
	STATUS REPORTS	BATCH BATCH STATUS REPORT	00002370
		ORDER ORDER STATUS REPORT	00002380
		TRANSACT TRANSACTION STATUS REPORT	00002390
		SCHEDULE WORK SCHEDULE REPORT	00002400
		ROUTE SORT LINE ROUTING REPORT	00002410
		MISSORT MISSORT EXCEPTIONS	00002420
			00002430
	BATCH CONTROL	READ READ BATCH FROM DISKETTE	00002440
		START START BATCH PROCESSING	00002450
		END END BATCH PROCESSING	00002460
		WRITE WRITE BATCH TO DISKETTE	00002470
		DELETE DELETE BATCH INFORMATION	00002480
			00002490
			00002500
			00002510
			00002520
			00002530
	TRANSACTION	ASSIGN ASSIGN TRANSACTION STATUS	00002540
			00002550
	SCANNER INPUT	RECIRC RECIRCULATE ON SCANNER ERROR	00002560
		HOLD HOLD ON SCANNER ERROR	00002570
			00002580
	SORT LINE	REROJTE ASSIGN ALTERNATE SORT LINE	00002590
		RESTORE RESTORE SORT LINE	00002600
			00002610
	PSC CONTROL	OPEN OPEN PSC COMMUNICATIONS	00002620
			00002630
	DIAGNOSTIC	HELP OPERATOR COMMAND LIST	00002640
		RESET RESET SYSTEM ERROR COUNTS	00002650
		ERROR DISPLAY SYSTEM ERROR COUNTS	00002660
			00002670
020I	ERROR MESSAGES:		00002680

1: PROGRAM LOAD ERROR XX	00002690
	00002700
	00002710
UNABLE TO LOAD AN OPERATOR COMMAND PROGRAM FOR EXECU-	00002720
TION. EDX COMPLETION CODE (XX) FOR LOAD STATEMENT	00002730
DESCRIBES CAUSE OF ERROR.	00002740
2: BATCH FILE SPACE EXHAUSTED	00002750
	00002760
	00002770
ATTEMPT TO ADD TOO MANY BATCH RECORDS. BATCH NOT	00002780
ADDED.	00002790
	00002800
3: TRANSACTION FILE SPACE EXHAUSTED	00002810
	00002820
ATTEMPT TO ADD TOO MANY TRANSACTIONS RECORDS. BATCH	00002830
AUTOMATICALLY DELETED.	00002840
	00002850
4: BATCH ALREADY ON FILE XXX	00002860
	00002870
ATTEMPT TO ADD BATCH ID (XXX) FAILED BECAUSE A DUP-	00002880
LICATE BATCH ID WAS ENCOUNTERED.	00002890
	00002900
5: TRANSACTION ALREADY ON FILE XXXXXX	00002910
	00002920
ATTEMPT TO ADD TRANSACTION ID (XXXXXX) FAILED BECAUSE	00002930
A DUPLICATE TRANSACTION ID WAS ENCOUNTERED.	00002940
	00002950
6: INVALID STATUS X	00002960
	00002970
ATTEMPT TO CHANGE EITHER BATCH OR TRANSACTION STATUS	00002980
FAILED BECAUSE OF INVALID STATUS CODE. OLD STATUS NOT	00002990
MODIFIED.	00003000
	00003010
7: BATCH NOT ON FILE XXX	00003020
	00003030
ATTEMPT TO ACCESS BATCH ID (XXX) INFORMATION NOT CUR-	00003040
RENTLY AVAILABLE.	00003050
	00003060
8: TRANSACTION NOT ON FILE XXXXXX	00003070
	00003080
ATTEMPT TO ACCESS TRANSACTION ID (XXXXXX) INFORMATION	00003090
NOT CURRENTLY AVAILABLE.	00003100
	00003110
9: ORDER FILE SPACE EXHAUSTED	00003120
	00003130
ATTEMPT TO ADD TOO MANY ORDER RECORDS. BATCH AUTOMA-	00003140
TICALLY DELETED.	00003150
	00003160
10: ORDER NOT ON FILE XXXXXX	00003170
	00003180
ATTEMPT TO ACCESS ORDER ID (XXXXXX) INFORMATION NOT	00003190
CURRENTLY AVAILABLE.	00003200
	00003210
11: INVALID SORT LINE NUMBER XX	00003220
	00003230
ATTEMPT TO ASSIGN OR RESTORE A SORT LINE OUTSIDE	00003240
KNOWN RANGE.	00003250
	00003260
12: BATCH NO LONGER ON FILE - REPORT TERMINATED	00003270
	00003280
BATCH INFORMATION WAS DELETED DURING REPORT GENER-	00003290
ATION BUT AFTER BATCH ID WAS VERIFIED.	00003300
	00003310
13: INVALID TRANSACTION NUMBER	00003320
	00003330
TRANSACTION NUMBER OUT OF RANGE.	00003340
	00003350
14: DISKETTE I/O ERROR OCCURRED	00003360
	00003370
ATTEMPT TO READ OR WRITE A DISKETTE DATA BLOCK	00003380
FAILED.	00003390
	00003400
15: TRANSACTION VALIDATION ERROR	00003410
	00003420

	ATTEMPT TO ADD A TRANSACTION RECORD FAILED BECAUSE OF A DATA VALIDATION CHECK REJECTED EITHER; TRANSACTION ID, DESTINATION, OR STATUS CODE.	00003430
		00003440
		00003450
		00003460
16:	DISKETTE FORMAT ERROR	00003470
	ATTEMPT TO ACCESS A NON-STANDARD IBM SYSTEM 3 DISKETTE.	00003480
		00003490
		00003500
17:	BATCH MUST BE PENDING OR COMPLETE TO BE DELETED - COMMAND ABORTED	00003510
		00003520
		00003530
	ATTEMPTED TO DELETE AN ACTIVE BATCH. THE COMMAND IS ABORTED	00003540
		00003550
18:	ALL CURRENT BATCHES ARE ACTIVE - COMMAND ABORTED	00003560
		00003570
	ATTEMPTED TO DELETE ALL BATCHES WHEN ALL BATCHES ON FILE WERE ACTIVE. THE COMMAND IS ABORTED.	00003580
		00003590
		00003600
19:	TRANSACTION XXXXXX NOT IN BATCH XXX	00003610
		00003620
	ATTEMPT TO CHANGE A TRANSACTION STATUS THAT IS NOT IN THE REQUESTED BATCH.	00003630
		00003640
		00003650
20:	ERROR IN BATCH XXX	00003660
		00003670
	ERROR WAS FOUND IN THE TRANSACTION FILE AT SYSTEM START UP TIME. STATUS OF BATCH IS QUESTIONABLE.	00003680
		00003690
		00003700
21:	DISK ERROR WHILE PROCESSING BATCH XXX	00003710
		00003720
	ATTEMPT TO READ A PORTION OF THE TRANSACTION FILE AT SYSTEM START UP TIME FAILED. STATUS OF BATCH IS QUESTIONABLE.	00003730
		00003740
		00003750
		00003760
22:	ERROR WHILE READING BATCH FILE	00003770
		00003780
	ATTEMPT TO READ THE BATCH FILE AT SYSTEM START UP TIME FAILED. STATUS OF SYSTEM IS QUESTIONABLE.	00003790
		00003800
		00003810
23:	ATTEMPT TO ASSIGN A SORT LINE TO ITSELF	00003820
		00003830
	AN ATTEMPT WAS MADE TO REROUTE A SORT LINE TO ITSELF USING THE REROUTE COMMAND. NO ACTION IS TAKEN AND THE COMMAND IS ABORTED.	00003840
		00003850
		00003860
		00003870
24:	NO PRIOR ASSIGNMENT OF SORT LINE	00003880
		00003890
	AN ATTEMPT WAS MADE TO RESTORE A SORT LINE THAT HAD NOT BEEN REROUTED USING THE RESTORE COMMAND. NO ACTION IS TAKEN AND THE COMMAND IS ABORTED.	00003900
		00003910
		00003920
		00003930
0301	REQUEST MESSAGES:	00003940
		00003950
1:	DO YOU WANT HARD COPY?	00003960
		00003970
	Y USER TERMINAL DEQUEUED AND SYSTEM PRINTER ENQUEUED. ALL FUTURE MESSAGES LISTED ON SYSTEM PRINTER.	00003980
		00003990
		00004000
	N NO ACTION.	00004010
		00004020
		00004030
2:	ENTER BATCH ID:	00004040
		00004050
	XXX 3-CHARACTER ALPHANUMERIC BATCH ID.	00004060
		00004070
3:	ALL ORDERS ?	00004080
		00004090
	Y ALL ORDER INFORMATION FOR A PREVIOUSLY REQUESTED BATCH ID, WILL BE ACCESSED.	00004100
		00004110
		00004120
	N NO ACTION.	00004130
		00004140
4:	ENTER ORDER ID:	00004150
		00004160

XXXXXX	6-CHARACTER ALPHANUMERIC ORDER ID.	00004170
5: LIST EXCEPTIONS ?		00004180
		00004190
		00004200
Y	ALL TRANSACTION INFORMATION NOT "STAGED FOR SHIPMENT" FOR A PREVIOUSLY REQUESTED BATCH ID AND ORDER ID, WILL BE LISTED.	00004210
		00004220
		00004230
		00004240
N	NO ACTION.	00004250
		00004260
6: DISKETTE MOUNTED ?		00004270
		00004280
Y	ACCESS OF DISKETTE ATTEMPTED.	00004290
		00004300
N	NO ACTION.	00004310
		00004320
7: OK TO READ ?		00004330
		00004340
Y	DISKETTE BATCH INFORMATION READ.	00004350
		00004360
N	READ BATCH (FROM DISKETTE) COMMAND TERMINATED.	00004370
		00004380
		00004390
8: CORRECT BATCH ?		00004400
		00004410
Y	BATCH INFORMATION FOR A PREVIOUSLY REQUESTED BATCH ID IS WRITTEN TO DISKETTE.	00004420
		00004430
		00004440
N	WRITE BATCH (TO DISKETTE) COMMAND TERMINATED.	00004450
		00004460
		00004470
9: START BATCH ?		00004480
		00004490
Y	BATCH STATUS FOR PREVIOUSLY REQUESTED BATCH ID IS CHANGED.	00004500
		00004510
		00004520
N	BATCH STATUS NOT CHANGED.	00004530
		00004540
10: END BATCH ?		00004550
		00004560
Y	BATCH STATUS FOR PREVIOUSLY REQUESTED BATCH ID IS CHANGED.	00004570
		00004580
		00004590
N	BATCH STATUS NOT CHANGED.	00004600
		00004610
11: ALL BATCHES ?		00004620
		00004630
Y	ALL BATCH INFORMATION CURRENTLY AVAILABLE MAY BE DELETED.	00004640
		00004650
		00004660
N	NO ACTION.	00004670
		00004680
12: DELETE BATCH ?		00004690
		00004700
Y	ALL BATCH INFORMATION ASSOCIATED WITH THE PREVIOUSLY DISPLAYED BATCH ID WILL BE DESTROYED.	00004710
		00004720
		00004730
		00004740
N	NO ACTION.	00004750
		00004760
13: ENTER TRANSACTION ID:		00004770
		00004780
XXXXXX	6-DIGIT NUMERIC TRANSACTION ID, IN THE RANGE 1 THROUGH 999999.	00004790
		00004800
14: ENTER NEW STATUS NUMBER:		00004810
		00004820
		00004830
X	SINGLE DIGIT NUMERIC TRANSACTION STATUS, IN THE RANGE 1 THROUGH 5.	00004840
		00004850
		00004860
15: ENTER OLD SORT LINE:		00004870
		00004880
XX	2-DIGIT NUMERIC SORT LINE NUMBER, IN THE RANGE 1 THROUGH 99.	00004890
		00004900
		00004910

16: ENTER NEW SORT LINE:		00004920
		00004930
XX	2-DIGIT NUMERIC SORT LINE NUMBER, IN THE RANGE 1 THROUGH 99.	00004940
		00004950
		00004960
17: TYPE "CANCEL" TO TERMINATE:		00004970
		00004980
CANCEL	COMMAND ABORTED	00004990
		00005000
<ENTER>	CONTINUE PROCESSING	00005010
		00005020
18: LIST MISSORT COUNTS ?		00005030
		00005040
Y	MISSORTS COUNTS LISTED	00005050
		00005060
N	NO ACTION	00005070
		00005080
19: CLEAR ALL COUNTS ?		00005090
		00005100
Y	ALL ERROR COUNTS CLEARED	00005110
		00005120
N	NO ACTION	00005130
		00005140
20: CLEAR MISSORT COUNTS ?		00005150
		00005160
Y	MISSORT COUNTS CLEARED	00005170
		00005180
N	NO ACTION	00005190
		00005200
21: CLEAR SCANNER NO-READ COUNTS ?		00005210
		00005220
Y	NO-READ COUNTS CLEARED	00005230
		00005240
N	NO ACTION	00005250
		00005260
22: CLEAR REJECT COUNT ?		00005270
		00005280
Y	REJECT COUNT CLEARED	00005290
		00005300
N	NO ACTION	00005310
		00005320
23: CLEAR RECIRCULATION COUNTS ?		00005330
		00005340
Y	RECIRCULATION COUNTS CLEARED	00005350
		00005360
N	NO ACTION	00005370
		00005380
24: CLEAR LOST FROM TRACKING COUNT ?		00005390
		00005400
Y	LFT COUNT CLEARED	00005410
		00005420
N	NO ACTION	00005430
		00005440
		00005450
		00005460
		00005470
		00005480
		00005490
		00005500
		00005510
		00005520
		00005530
		00005540
		00005550
		00005560
		00005570
		00005580
		00005590
		00005600
		00005610
		00005620
		00005630
		00005640
		00005650

04DI DESCRIPTIVE MESSAGES:

- 1: DISKETTE ALREADY IN USE
- 2: RETRY COMMAND LATER
- 3: WORK SCHEDULE REPORT
- 4: NO CURRENT BATCH INFORMATION
- 5: BATCH STATUS REPORT
- 6: ORDER STATUS REPORT
- 7: TRANSACTION STATUS REPORT
- 8: SORT TRANSLATION REPORT
- 9: READ BATCH COMMAND
- 10: WRITE BATCH COMMAND
- 11: START BATCH COMMAND
- 12: END BATCH COMMAND
- 13: DELETE BATCH COMMAND
- 14: ASSIGN TRANSACTION STATUS COMMAND
- 15: ASSIGN SHIPPING LINE COMMAND
- 16: RESTORE SHIPPING LINE COMMAND
- 17: OPEN PSC COMMUNICATIONS COMMAND
- 18: PSC COMMUNICATIONS INITIATED
- 19: RECIRC ON SCANNER ERROR

20: HOLD ON SCANNER ERROR
 21: HELP COMMAND
 22: SHIPPING LINE ASSIGNED
 23: READ BATCH COMMAND ABORTED
 24: WRITE BATCH COMMAND ABORTED
 25: REMOVE DISKETTE
 26: STATUS CHANGE COMPLETE
 27: SYSTEM STARTED
 28: SET DATE AND TIME USING COMMAND \$T
 29: SYSTEM INITIALIZATION IN PROGRESS
 30: STATUS OF SYSTEM IS QUESTIONABLE
 31: SYSTEM ERROR REPORT
 32: M I S - S O R T S
 33: RESET ERROR COUNTS COMMAND
 34: COUNTS CLEARED
 35: MIS-SORT REPORT
 36: NO MIS-SORTS IN THIS ORDER
 37: NO EXCEPTIONS IN THIS ORDER

00005660
 00005670
 00005680
 00005690
 00005700
 00005710
 00005720
 00005730
 00005740
 00005750
 00005760
 00005770
 00005780
 00005790
 00005800
 00005810
 00005820
 00005830
 00005840

0501 WORK SCHEDULE REPORT LAYOUT:

REPORT FIELD:	FROM:
1: DATE	2701 TIME
2: TIME	2701 TIME
3: PAGE	2701 PAGECNT
4: BATCH ID	1201 BID
5: BATCH STATUS	1201 BSTATUS, BSTBL
6: ORDER COUNT	1201 BOCNT
7: TRANSACTION COUNT	1201 BTCNT
8: HEADER INFORMATION	1201 BHDR

00005850
 00005860
 00005870
 00005880
 00005890
 00005900
 00005910
 00005920
 00005930
 00005940
 00005950
 00005960
 00005970

* SEE REPORT LAYOUT SHEET

00005980
 00005990

0601 BATCH STATUS REPORT LAYOUT:

REPORT FIELD:	FROM:
1: DATE	2801 TIME
2: TIME	2801 TIME
3: PAGE	2801 PAGECNT
4: BATCH ID	1201 BID
5: BATCH STATUS	1201 BSTATUS, BSTBL
6: ORDER COUNT	1201 BOCNT
7: TRANSACTION COUNT	1201 BTCNT
8: HEADER INFORMATION	1201 BHDR
9: BATCH % COMPLETE	1201 BPCOM

00006000
 00006010
 00006020
 00006030
 00006040
 00006050
 00006060
 00006070
 00006080
 00006090
 00006100
 00006110
 00006120

* SEE REPORT LAYOUT SHEET

00006130
 00006140

0701 ORDER STATUS REPORT LAYOUT:

REPORT FIELD:	FROM:
1: DATE	2901 TIME
2: TIME	2901 TIME
3: PAGE	2901 PAGECNT
4: BATCH ID	1201 BID
5: BATCH STATUS	1201 BSTATUS, BSTBL
6: ORDER COUNT	1201 BOCNT
7: TRANSACTION COUNT	2901 RTCNT
8: HEADER INFORMATION	1201 BHDR
9: ORDER NUMBER	1301 OCNT
10: ORDER ID	1301 OID
11: QUANTITY PLANNED	1301 OTCNT
12: QUANTITY STAGED	1301 OQSTG
13: QUANTITY NOT SHIPPED	1301 OQNSTG
14: QUANTITY STOCK-OUT	1301 OQSTO
15: QUANTITY EXCESS REPACK	1301 OQXRP
16: QUANTITY IN SORTATION	1301 OQISORT
17: QUANTITY NOT PICKED	1301 OQNPCKD
18: QUANTITY MIS-SORT	1301 OQMIS
19: ORDER % COMPLETE	1301 OPCOM
20: TRANSACTION NUMBER	2901 RTCNT

00006150
 00006160
 00006170
 00006180
 00006190
 00006200
 00006210
 00006220
 00006230
 00006240
 00006250
 00006260
 00006270
 00006280
 00006290
 00006300
 00006310
 00006320
 00006330
 00006340
 00006350
 00006360
 00006370
 00006380
 00006390

21: TRANSACTION ID	14DI TID	00006400
22: TRANSACTION STATUS	14DI TSTATUS	00006410
23: PRODUCT ID	14DI TPROD	00006420
24: SOURCE LOCATION	14DI TSOURCE	00006430
25: DESTINATION	14DI TDEST	00006440
26: MIS-SORT DESTINATION		00006450

* SEE REPORT LAYOUT SHEET

08DI TRANSACTION STATUS REPORT LAYOUT:

REPORT FIELD:	FROM:	
1: DATE	30DI TIME	00006500
2: TIME	30DI TIME	00006510
3: PAGE	30DI PAGECNT	00006520
4: BATCH ID	12DI BID	00006530
5: BATCH STATUS	12DI BSTATUS, BSTBL	00006540
6: ORDER COUNT	12DI BOCNT	00006550
7: HEADER INFORMATION	12DI BHDR	00006560
8: TRANSACTION ID	14DI TID	00006570
9: TRANSACTION STATUS	14DI TSTATUS	00006580
10: PRODUCT ID	14DI TPROD	00006590
11: SOURCE LOCATION	14DI TSOURCE	00006600
12: DESTINATION	14DI TDEST	00006610
13: MIS-SORT DESTINATION	TMDEST	00006620

* SEE REPORT LAYOUT SHEET

09DI SORT LINE ROUTING REPORT LAYOUT:

REPORT FIELD:	FROM:	
1: DATE	31DI TIME	00006660
2: TIME	31DI TIME	00006670
3: PAGE	31DI PAGECNT	00006680
4: SORT LINE	05FM	00006690
5: ALTERNATE	05FM	00006700

* SEE REPORT LAYOUT SHEET

10DI LOCAL FLAG DEFINITIONS:

LABEL:	DESCRIPTION:	VALUE:	BYTES:	TYPE:	
1: FLGDEF	LOCAL FLAG DEFINITION		2	BINARY	00006710
2: FLGOK	(GENERAL PURPOSE "JK")	0001	0	EQUATE	00006720
3: FLGHCY	HARDCOPY DEVICE	0002	0	EQUATE	00006730
4: FLGBFD	BATCH FOUND	0004	0	EQUATE	00006740
5: FLBSPQB	SPECIFIC ORDER OR BATCH	0008	0	EQUATE	00006750
6: FLGOSQ	ORDER SEQUENCE PROCESSING	0010	0	EQUATE	00006760
7: FLGEXC	EXCEPTION PROCESSING	0020	0	EQUATE	00006770
8: FLGDEL	DELETION PROCESSING	0040	0	EQUATE	00006780
9: FLGLOR	LIST ORDER RECORD	0080	0	EQUATE	00006790
10: FLGLTR	LIST TRANSACTION RECORD	0100	0	EQUATE	00006800
11: FLGABO	BATCH READ/WRITE ABORTED	0200	0	EQUATE	00006810
12: FLGLR	LAST RECORD PROCESSING	0400	0	EQUATE	00006820
13: FLGRAB	REPORT ABORTED	0800	0	EQUATE	00006830

11DI GLOBAL FLAG DEFINITIONS:

LABEL:	DESCRIPTION:	VALUE:	BYTES:	TYPE:	
1: OIFLG	GLOBAL FLAG DEFINITION		2	BINARY	00006840
2: OIDIU	DISKETTE IN USE	0001	0	EQUATE	00006850

MODULE REFERENCE: DICTL, OIRDB, OIWRB.

12DI BATCH LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:	
1: BRDA	RESIDENT DIRECTORY ADDRESS	2	BINARY	00007000
2: BMAX	MAXIMUM BATCHES ALLOWED	2	BINARY	00007010

3:	BCNT	BATCH COUNTER	2	BINARY	00007140
4:	BID	BATCH ID	6	EBCDIC	00007150
5:	BSTATUS	BATCH STATUS	2	BINARY	00007160
6:	BHDR	BATCH HEADER	52	EBCDIC	00007170
7:	BCCNT	BATCH ORDER COUNT	2	BINARY	00007180
8:	BCNT	BATCH TRANSACTION COUNT	2	BINARY	00007190
9:	BRDES	RESIDENT DIRECTORY ENTRY SIZE	2	BINARY	00007200
10:	BIPM	IPM BUFFER ADDRESS	2	BINARY	00007210
11:	BSRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00007220
12:	BPANIC1	PANIC CODE 1	2	BINARY	00007230
13:	BPANIC2	PANIC CODE 2	2	BINARY	00007240
14:	BPANIC3	PANIC CODE 3	2	BINARY	00007250
15:	BPCOM	BATCH PERCENT COMPLETE	8	BINARY	00007260
16:	BESTBL	BATCH ENGLISH STATUS TABLE	48	EBCDIC	00007270
17:	BHDR1	BATCH REPORT HEADER	82	EBCDIC	00007280
18:	BHDR2	BATCH REPORT HEADER	30	EBCDIC	00007290
19:	BHDR3	BATCH REPORT HEADER	50	EBCDIC	00007300
20:	BOTL	BATCH REPORT DETAIL INFORMATION	32	EBCDIC	00007310
21:	BFTN	BATCH REPORT FOOTNOTE	34	EBCDIC	00007320

1301 ORDER LOCAL DATA:

	LABEL:	DESCRIPTION	BYTES:	TYPE:	
					00007330
					00007340
					00007350
					00007360
					00007370
1:	OID	ORDER ID	8	EBCDIC	00007380
2:	OCNT	ORDER COUNTER	2	BINARY	00007390
3:	ORNO	ORDER RELATIVE NUMBER	2	BINARY	00007400
4:	OTCNT	ORDER TRANSACTION COUNT	2	BINARY	00007410
5:	OQSTG	ORDER QUANTITY STAGED	2	BINARY	00007420
6:	OQNSTG	ORDER QUANTITY NOT STAGED	2	BINARY	00007430
7:	OQNPKD	ORDER QUANTITY NOT PICKED	2	BINARY	00007440
8:	OQISORT	ORDER QUANTITY IN SORTATION	2	BINARY	00007450
9:	OQXRP	ORDER QUANTITY EXCESS REPACK	2	BINARY	00007460
10:	OQSTO	ORDER QUANTITY STOCK OUT	2	BINARY	00007470
11:	OQMIS	ORDER QUANTITY MIS-SORT	2	BINARY	00007480
12:	OPCOM	ORDER QUANTITY PERCENT COMPLETE	8	BINARY	00007490
13:	OIPM	IPM BUFFER ADDRESS	2	BINARY	00007500
14:	OSRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00007510
15:	OPANIC1	PANIC CODE 1	2	BINARY	00007520
16:	OPANIC2	PANIC CODE 2	2	BINARY	00007530
17:	OPANIC3	PANIC CODE 3	2	BINARY	00007540
18:	OHDR1	ORDER REPORT HEADER	82	EBCDIC	00007550
19:	OHDR2	ORDER REPORT HEADER	82	EBCDIC	00007560
20:	ODTL	ORDER REPORT DETAIL INFORMATION	80	EBCDIC	00007570

1401 TRANSACTION LOCAL DATA:

	LABEL:	DESCRIPTION:	BYTES:	TYPE:	
					00007580
					00007590
					00007600
					00007610
					00007620
1:	TCNT	TRANSACTION COUNTER	2	BINARY	00007630
2:	TID	TRANSACTION ID	8	BINARY	00007640
3:	TSTATUS	TRANSACTION STATUS	2	BINARY	00007650
4:	TRON	TRANSACTION RELATIVE ORDER NUMBER	2	BINARY	00007660
5:	TPROD	TRANSACTION PRODUCT ID	6	EBCDIC	00007670
6:	TSOURCE	TRANSACTION SOURCE	6	EBCDIC	00007680
7:	TDEST	TRANSACTION DESTINATION	2	BINARY	00007690
8:	TFBLK	TRANSACTION FILE BLOCK NUMBER	2	BINARY	00007700
9:	TFOFF	TRANSACTION FILE ENTRY OFFSET	2	BINARY	00007710
10:	TFRS	FILE RECORD SIZE	2	BINARY	00007720
11:	TFBF	TRANSACTION FILE BLOCKING FACTOR	2	BINARY	00007730
12:	TIPM	IPM BUFFER ADDRESS	2	BINARY	00007740
13:	TSRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00007750
14:	TPANIC1	PANIC CODE 1	2	BINARY	00007760
15:	TPANIC2	PANIC CODE 2	2	BINARY	00007770
16:	TPANIC3	PANIC CODE 3	2	BINARY	00007780
17:	TESTBL	TRANSACTION ENGLISH STATUS TABLE	65	EBCDIC	00007790
18:	THDR1	TRANSACTION REPORT HEADER	64	EBCDIC	00007800
19:	THDR2	TRANSACTION REPORT HEADER	78	EBCDIC	00007810
20:	TOTL	TRANSACTION DETAIL INFORMATION	74	EBCDIC	00007820

1501 FIND BATCH CALL: DICCFB

	LABEL:	DESCRIPTION:	BYTES:	TYPE:	
					00007830
					00007840
					00007850
					00007860
					00007870

INPUT:

1:	BRDA	RESIDENT DIRECTORY ADDRESS	2	BINARY	00007880
2:	BRDES	RESIDENT DIRECTORY ENTRY SIZE	2	BINARY	00007890
3:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00007900
4:	BMAX	MAXIMUM BATCHES ALLOWED	2	BINARY	00007910
5:	BCNT	BATCH COUNTER TO START TABLE SEARCH	2	BINARY	00007920

LOCAL:

1:	SAVE.1	TEMPORARY SAVE #1	2	BINARY	00007930
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OUTPUT:

1:	TFLG	TEMPORARY FLAG	2	BINARY	00007940
2:	SAVE.2	TEMPORARY SAVE #2	2	BINARY	00007950

OUTPUT:

1:	BID	BATCH ID	3	EBCDIC	00007960
2:	FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00007970
3:	BCNT	UPDATED BATCH COUNTER	2	BINARY	00007980

16JI GET BATCH CALL: DICCGS

LABEL:	DESCRIPTION:	BYTES:	TYPE:	
				00007990
				00008000
				00008010
				00008020
				00008030
				00008040
				00008050
				00008060
				00008070
				00008080
				00008090
				00008100
				00008110
				00008120
				00008130
				00008140
				00008150
				00008160
				00008170
				00008180
				00008190
				00008200
				00008210
				00008220
				00008230
				00008240
				00008250
				00008260
				00008270
				00008280
				00008290
				00008300
				00008310
				00008320
				00008330
				00008340
				00008350
				00008360
				00008370
				00008380
				00008390
				00008400
				00008410
				00008420
				00008430
				00008440
				00008450
				00008460
				00008470
				00008480
				00008490
				00008500
				00008510
				00008520
				00008530
				00008540
				00008550
				00008560
				00008570
				00008580
				00008590
				00008600
				00008610

INPUT:

1:	BID	BATCH ID	3	EBCDIC	00008180
2:	BIPM	IPM BUFFER ADDRESS	2	BINARY	00008190
3:	BSRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00008200
4:	BPANIC1	BAD CALL CODE PANIC CODE	2	BINARY	00008210
5:	BPANIC2	BAD RETURN CODE PANIC CODE	2	BINARY	00008220
6:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00008230

LOCAL:

1:	SAVE.1	TEMPORARY SAVE #1	2	BINARY	00008240
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OUTPUT:

1:	BSTATUS	BATCH STATUS	2	BINARY	00008250
2:	BHDR	BATCH HEADER	50	EBCDIC	00008260
3:	BDCNT	BATCH ORDER COUNTER	2	BINARY	00008270
4:	BTCNT	BATCH TRANSACTION COUNTER	2	BINARY	00008280
5:	PANIC	UPDATED PANIC WORK AREA	2	BINARY	00008290
6:	FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00008300

17JI FIND ORDER CALL: DICCFJ

LABEL:	DESCRIPTION:	BYTES:	TYPE:	
				00008310
				00008320
				00008330
				00008340
				00008350
				00008360
				00008370
				00008380
				00008390
				00008400
				00008410
				00008420
				00008430
				00008440
				00008450
				00008460
				00008470
				00008480
				00008490
				00008500
				00008510
				00008520
				00008530
				00008540
				00008550
				00008560
				00008570
				00008580
				00008590
				00008600
				00008610

LOCAL:

1:	SAVE.1	TEMPORARY SAVE #1	2	BINARY	00008370
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OUTPUT:

1:	OJD	ORDER ID	6	EBCDIC	00008380
2:	ORNO	ORDER RELATIVE NUMBER	2	BINARY	00008390
3:	OQSTB	ORDER QUANTITY STAGED	2	BINARY	00008400
4:	OQNPKE	ORDER QUANTITY NOT PICKED	2	BINARY	00008410
5:	OQISRT	ORDER QUANTITY IN SORTATION	2	BINARY	00008420
6:	OQXRP	ORDER QUANTITY EXCESS REPACK	2	BINARY	00008430

103

104

Label	Description	Bytes	Type	Address
7: QJSTG	ORDER QUANTITY STOCK OUT	2	BINARY	00008520
8: QJMIS	ORDER QUANTITY MIS-SORT	2	BINARY	00008530
9: FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00008540
10: QCNT	UPDATED ORDER NUMBER	2	BINARY	00008550
11: PANIC	UPDATED PANIC WORK AREA	2	BINARY	00008560

180I GET ORDER CALL: DICCGO

Label	Description	Bytes	Type	Address
INPUT:				
1: OIPM	IPM BUFFER ADDRESS	2	BINARY	00008700
2: BID	BATCH ID	3	EBCDIC	00008710
3: OI	ORDER ID	6	EBCDIC	00008720
4: JSRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00008730
5: FLGDEF	LOCAL FLAG WORD	2	BINARY	00008740
6: OPANIC1	BAD CALL CODE PANIC CODE	2	BINARY	00008750
7: OPANIC2	BAD RETURN CODE PANIC CODE	2	BINARY	00008760

LOCAL:

1: SAVE.1	TEMPORARY SAVE #1	2	BINARY	00008770
2: BNRTXT	BATCH NO LONGER THERE MESSAGE	50	EBCDIC	00008780
3: DNTXT	ORDER NOT ON FILE MESSAGE	32	EBCDIC	00008790

OUTPUT:

1: OI	ORDER ID	6	EBCDIC	00008800
2: ORNO	ORDER RELATIVE NUMBER	2	BINARY	00008810
3: QJSTG	ORDER QUANTITY STAGED	2	BINARY	00008820
4: QJNPCK	ORDER QUANTITY NOT PICKED	2	BINARY	00008830
5: QJISORT	ORDER QUANTITY IN SORTATION	2	BINARY	00008840
6: QJXRP	ORDER QUANTITY EXCESS REPACK	2	BINARY	00008850
7: QJSTJ	ORDER QUANTITY STOCK OUT	2	BINARY	00008860
8: QJMIS	ORDER QUANTITY MIS-SORT	2	BINARY	00008870
9: FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00008880
10: PANIC	UPDATED PANIC WORK AREA	2	BINARY	00008890

190I FIND TRANSACTION CALL: DICGFT

Label	Description	Bytes	Type	Address
INPUT:				
1: TCNT	TRANSACTION COUNTER TO START ACCESS	2	BINARY	00009000
2: BT CNT	BATCH TRANSACTION COUNTER	2	BINARY	00009010
3: TSBB	SYSCOM BLOCK BUFFER ADDRESS	2	BINARY	00009020
4: TFOFF	TRANSACTION FILE ENTRY OFFSET	2	BINARY	00009030
5: ORNO	ORDER RELATIVE NUMBER	2	BINARY	00009040
6: FLGDEF	LOCAL FLAG WORD	2	BINARY	00009050

LOCAL:

1: TIFLG	TEMPORARY OK FLAG	2	BINARY	00009060
2: T2FLG	TEMPORARY ORDER SEQUENCE FLAG	2	BINARY	00009070
3: SAVE.2	TEMPORARY SAVE #2	2	BINARY	00009080

OUTPUT:

1: TID	TRANSACTION ID	4	BINARY	00009090
2: TSTATUS	TRANSACTION STATUS	2	BINARY	00009100
3: TRON	TRANSACTION RELATIVE ORDER NUMBER	2	BINARY	00009110
4: TPROD	TRANSACTION PRODUCT ID	6	EBCDIC	00009120
5: TSDURCE	TRANSACTION SOURCE	6	EBCDIC	00009130
6: TDEST	TRANSACTION DESTINATION	2	BINARY	00009140
7: FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00009150
8: TCNT	UPDATED TRANSACTION COUNTER	2	BINARY	00009160

200I SEQUENTIAL TRANSACTION CALL: DICGST

Label	Description	Bytes	Type	Address
				00009170
				00009180
				00009190
				00009200
				00009210
				00009220
				00009230
				00009240
				00009250
				00009260
				00009270
				00009280
				00009290
				00009300
				00009310
				00009320
				00009330
				00009340

INPUT:

1:	TCNT	TRANSACTION COUNTER TO ACCESS FILE	2	BINARY
2:	TFBF	TRANSACTION FILE BLOCKING FACTOR	2	BINARY
3:	TFOFF	TRANSACTION FILE ENTRY OFFSET	2	BINARY
4:	TFRS	TRANSACTION FILE RECORD SIZE	2	BINARY
5:	TFBLK	TRANSACTION FILE BLOCK NUMBER	2	BINARY
6:	TIPM	IPM BUFFER ADDRESS	2	BINARY
7:	BID	BATCH ID	3	EBCDIC
8:	SBB	SYSCOM BLOCK BUFFER ADDRESS	2	BINARY
9:	TPANIC1	BAD CALL CODE PANIC CODE	2	BINARY
10:	TPANIC2	BAD BLOCK NUMBER PANIC CODE	2	BINARY
11:	TPANIC3	BAD RETURN CODE PANIC CODE	2	BINARY
12:	BDCNT	BATCH ORDER COUNT	2	BINARY
13:	BTCNT	BATCH TRANSACTION COUNT	2	BINARY

LOCAL:

1:	TMPI	TEMPORARY WORK AREA	2	BINARY
2:	SAVE.1	TEMPORARY SAVE #1	2	BINARY
3:	SAVE.2	TEMPORARY SAVE #2	2	BINARY
4:	DATBUF	TRANSACTION SECTOR BUFFER	256	BINARY

OUTPUT:

1:	TCNT	UPDATED TRANSACTION COUNTER	2	BINARY
2:	TFBLK	UPDATED FILE BLOCK NUMBER	2	BINARY
3:	SBB	UPDATED SYSCOM BLOCK BUFFER ADDR.	2	BINARY
4:	TSBB	UPDATED TRANSACTION BLOCK BUFR	2	BINARY
5:	PANIC	UPDATED PANIC WORK AREA	2	BINARY
6:	FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY
7:	OCNT	UPDATED ORDER COUNTER	2	BINARY

2101

SET TRANSACTION CALL: DICCGT

LABEL:	DESCRIPTION:	BYTES:	TYPE:
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INPUT:

1:	TIPM	IPM BUFFER ADDRESS	2	BINARY
2:	TID	TRANSACTION ID	4	BINARY
3:	TSRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY
4:	TPANIC1	BAD CALL CODE PANIC CODE	2	BINARY
6:	TPANIC2	BAD RETURN CODE PANIC CODE	2	BINARY
7:	FLGDEF	LOCAL FLAG WORD	2	BINARY

LOCAL:

1:	SAVE.1	TEMPORARY SAVE #1	2	BINARY
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OUTPUT:

1:	TSTATUS	TRANSACTION STATUS	2	BINARY
2:	TRON	TRANSACTION RELATIVE ORDER NUMBER	2	BINARY
3:	TPROD	TRANSACTION PRODUCT ID	6	EBCDIC
4:	TSOURCE	TRANSACTION SOURCE	6	EBCDIC
5:	TDEST	TRANSACTION DESTINATION	2	BINARY
6:	PANIC	UPDATED PANIC WORK AREA	2	BINARY
7:	FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY

2201

READ BATCH CALL: DICCRB

LABEL:	DESCRIPTION:	BYTES:	TYPE:
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INPUT:

1:	DCMD	DISKETTE COMMAND WORD	2	BINARY
2:	TEMP	TEMPORARY WORK AREA	2-8	BINARY
3:	DPANIC1	BAD CALL CODE PANIC CODE	2	BINARY
4:	DPANIC2	BAD RETURN CODE PANIC CODE	2	BINARY
5:	BSRB	SYSCOM RECORD BUFFER	2	BINARY
6:	FLGDEF	LOCAL FLAG WORD	2	BINARY

00009350
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00009990
00010000
00010010
00010020
00010030
00010040
00010050
00010060
00010070

LOCAL:					00010080
					00010090
1:	RBTXT1	FORMAT ERROR MESSAGE	28	EBCDIC	00010100
2:	RBTXT1	I/O ERROR MESSAGE	34	EBCDIC	00010110
					00010120
OUTPUT:					00010130
					00010140
1:	BSRB	UPDATED RECORD BUFFER	2	BINARY	00010150
2:	PANIC	UPDATED PANIC WORK AREA	2	BINARY	00010160
3:	FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00010170
					00010180
2301	WRITE BATCH CALL:	OICCW8			00010190
					00010200
	LABEL:	DESCRIPTION:	BYTES:	TYPE:	00010210
					00010220
INPUT:					00010230
					00010240
1:	DCMD	DISKETTE COMMAND WORD	2	BINARY	00010250
2:	OIWRB	PROGRAM TASK WORD ADDRESS	2	BINARY	00010260
3:	DPANIC1	BAD CALL CODE PANIC CODE	2	BINARY	00010270
4:	DPANIC2	BAD RETURN CODE PANIC CODE	2	BINARY	00010280
5:	BID	BATCH ID	3	EBCDIC	00010290
6:	BHDR	BATCH HEADER	50	EBCDIC	00010300
7:	TID	TRANSACTION ID	4	BINARY	00010310
8:	OID	ORDER ID	6	EBCDIC	00010320
9:	TPROD	TRANSACTION PRODUCT ID	6	EBCDIC	00010330
10:	TDEST	TRANSACTION DESTINATION	2	BINARY	00010340
11:	TSOURCE	TRANSACTION SOURCE	6	EBCDIC	00010350
12:	TSTATUS	TRANSACTION STATUS	2	BINARY	00010360
13:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00010370
					00010380
LOCAL:					00010390
					00010400
1:	TEMP	TEMPORARY WORK AREA	2	BINARY	00010410
2:	SAVE.1	TEMPORARY SAVE #1	2	BINARY	00010420
3:	WBTXT1	I/O ERROR MESSAGE	34	EBCDIC	00010430
					00010440
OUTPUT:					00010450
					00010460
1:	PANIC	UPDATED PANIC WORK AREA	2	BINARY	00010470
2:	TCNT	UPDATED TRANSACTION COUNTER	2	BINARY	00010480
3:	FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00010490
					00010500
2401	ADD BATCH CALL:	OICCA8			00010510
					00010520
	LABEL:	DESCRIPTION:	BYTES:	TYPE:	00010530
					00010540
INPUT:					00010550
					00010560
1:	BIPM	IPM BUFFER ADDRESS	2	BINARY	00010570
2:	BID	BATCH ID	3	EBCDIC	00010580
3:	BHDR	BATCH HEADER	50	EBCDIC	00010590
4:	BSRB	SYS COM RECORD BUFFER ADDRESS	2	BINARY	00010600
5:	BPANIC1	BATCH ALREADY ADDING PANIC CODE	2	BINARY	00010610
6:	BPANIC2	BAD CALL CODE PANIC CODE	2	BINARY	00010620
7:	BPANIC3	BAD RETURN CODE PANIC CODE	2	BINARY	00010630
8:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00010640
					00010650
LOCAL:					00010660
					00010670
1:	SAVE.1	TEMPORARY SAVE #1	2	BINARY	00010680
2:	ABTXT1	SPACE EXHAUSTED MESSAGE	34	EBCDIC	00010690
3:	ABTXT2	ALREADY ON FILE MESSAGE	32	EBCDIC	00010700
					00010710
OUTPUT:					00010720
					00010730
1:	PANIC	UPDATED PANIC WORK AREA	2	BINARY	00010740
2:	FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00010750
					00010760
2501	ADD TRANSACTION CALL:	OICCAT			00010770
					00010780
	LABEL:	DESCRIPTION:	BYTES:	TYPE:	00010790
					00010800

INPUT:

1:	TIPM	IPM BUFFER ADDRESS	2	BINARY
2:	TSR3	SYSCOM RECORD BUJFFER ADDRESS	2	BINARY
3:	TPANIC1	BAD CALL CODE PANIC CODE	2	BINARY
4:	TPANIC2	BAD RETURN CODE PANIC CODE	2	BINARY
5:	TID	TRANSACTION ID	4	BINARY
6:	OID	ORDER ID	6	EBCDIC
7:	TPROD	TRANSACTION PRODUCT ID	6	EBCDIC
8:	TDEST	TRANSACTION DESTINATION	2	BINARY
9:	TSOURCE	TRANSACTION SOURCE	6	EBCDIC
10:	TSTATUS	TRANSACTION STATUS	2	BINARY
11:	FLGDEF	LOCAL FLAG WORD	2	BINARY

LOCAL:

1:	TFLG	TEMPORARY FLAG WORD	2	BINARY
2:	SAVE.1	TEMPORARY SAVE #1	2	BINARY
3:	ATTXT1	ALREADY ON FILE MESSAGE	36	EBCDIC
4:	ATTXT2	TRANSACTION SPACE EXHAUSTED MSG	40	EBCDIC
5:	ATTXT3	ORDER SPACE EXHAUSTED MESSAGE	34	EBCDIC

OUTPUT:

1:	PANIC	UPDATED PANIC WORK AREA	2	BINARY
2:	FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY

26JI OPERATOR INTERFACE CONTROL LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:
1:	OIEND	END ECB OF OPERATOR INTERFACE	3 BINARY
2:	MSGLODR	LOAD ERROR MESSAGE	26 EBCDIC
3:	MSGDIU	DISKETTE IN USE MESSAGE	30 EBCDIC
4:	MSGRTY	RETRY COMMAND MESSAGE	22 EBCDIC
5:	PANIC	PANIC WORK AREA	2 BINARY
6:	OITSK1CD	LOAD TASK #1 CODE	2 BINARY
7:	OITSK2CD	LOAD TASK #2 CODE	2 BINARY
8:	OITSK3CD	LOAD TASK #3 CODE	2 BINARY
9:	OITSK4CD	LOAD TASK #4 CODE	2 BINARY
10:	LDRES	LOAD PROGRAM QUEUE CONTROL BLOCK	10 BINARY

27JI WORK SCHEDULE REPORT LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:
1:	FLGDEF	LOCAL FLAG WORD	2 BINARY
2:	IPM	IPM BUFFER ADDRESS	2 BINARY
3:	IPMPC1	IPM PANIC CODE 1	2 BINARY
4:	IPMPC2	IPM PANIC CODE 2	2 BINARY
5:	SRB	SYSCOM RECORD BUJFFER ADDRESS	2 BINARY
6:	SRBPC1	SRB PANIC CODE 1	2 BINARY
7:	SRBPC2	SRB PANIC CODE 2	2 BINARY
8:	TIME	TABLE:HH,MM,SS,MO,DA,YR	12 BINARY
9:	PAGECNT	REPORT PAGE COUNT	2 BINARY
10:	LINECNT	REPORT LINE COUNT	2 BINARY
11:	TEMP	TEMPORARY WORK AREA	4 BINARY
12:	PANIC	PANIC CODE WORK AREA	2 BINARY
13:	MSGWSR	BANNER	22 EBCDIC
14:	FLG1	AUXILIARY FLAG WORD	2 BINARY
15:	CRTCNT	CRT MAXIMUM SCREEN SIZE	2 BINARY

28JI BATCH STATUS REPORT LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:
1:	FLGDEF	LOCAL FLAG WORD	2 BINARY
2:	IPM	IPM BUFFER ADDRESS	2 BINARY
3:	IPMPC1	IPM PANIC CODE 1	2 BINARY
4:	IPMPC2	IPM PANIC CODE 2	2 BINARY
5:	SRB	SYSCOM RECORD BUJFFER ADDRESS	2 BINARY
6:	SRBPC1	SRB PANIC CODE 1	2 BINARY
7:	SRBPC2	SRB PANIC CODE 2	2 BINARY
8:	TIME	TABLE:HH,MM,SS,MO,DA,YR	12 BINARY

00010810
 00010820
 00010830
 00010840
 00010850
 00010860
 00010870
 00010880
 00010890
 00010900
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 00010980
 00010990
 00011000
 00011010
 00011020
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 00011080
 00011090
 00011100
 00011110
 00011120
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 00011150
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 00011170
 00011180
 00011190
 00011200
 00011210
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 00011470
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 00011490
 00011500
 00011510
 00011520
 00011530
 00011540

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LINE	DESCRIPTION	BYTES	TYPE	ADDRESS
9:	PAGECNT	2	BINARY	00011550
10:	LINECNT	2	BINARY	00011560
11:	CRTCNT	2	BINARY	00011570
12:	PRINTCNT	2	BINARY	00011580
13:	TEMP	10	BINARY	00011590
14:	PANIC	2	BINARY	00011600
15:	MSGBSR	22	EBCDIC	00011610

290I ORDER STATUS REPORT LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:	ADDRESS:
1:	FLGDEF	2	BINARY	00011650
2:	IPM	2	BINARY	00011660
3:	IPMPC1	2	BINARY	00011670
4:	IPMPC2	2	BINARY	00011680
5:	SRB	2	BINARY	00011690
6:	SRBPC1	2	BINARY	00011700
7:	SRBPC2	2	BINARY	00011710
8:	SBB	2	BINARY	00011720
9:	SBBPC1	2	BINARY	00011730
10:	RTCNT	2	BINARY	00011740
11:	TIME	12	BINARY	00011750
12:	PAGECNT	2	BINARY	00011760
13:	LINECNT	2	BINARY	00011770
14:	CRTCNT	2	BINARY	00011780
15:	PRINTCNT	2	BINARY	00011790
16:	TEMP	10	BINARY	00011800
17:	PANIC	2	BINARY	00011810
18:	MSGOSR	22	EBCDIC	00011820
19:	MSGABORT	48	EBCDIC	00011830

300I TRANSACTION STATUS REPORT LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:	ADDRESS:
1:	FLGDEF	2	BINARY	00011850
2:	IPM	2	BINARY	00011870
3:	IPMPC1	2	BINARY	00011880
4:	IPMPC2	2	BINARY	00011890
5:	SRB	2	BINARY	00011900
6:	SRBPC1	2	BINARY	00011910
7:	SRBPC2	2	BINARY	00011920
8:	TIME	12	BINARY	00011930
9:	PAGECNT	2	BINARY	00011940
10:	LINECNT	2	BINARY	00011950
11:	TEMP	12	BINARY	00011960
12:	PANIC	2	BINARY	00011970
13:	MSGTSR	28	EBCDIC	00011980
14:	MSGFTN	32	EBCDIC	00011990
15:	CRTCNT	2	BINARY	00012000
16:	DZERO	4	BINARY	00012010
17:	TRANSMAX	4	BINARY	00012020

310I SORT TRANSLATION REPORT LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:	ADDRESS:
1:	TIME	12	BINARY	00012070
2:	PAGECNT	2	BINARY	00012080
3:	TEMP	10	BINARY	00012090
4:	MSGSTR	26	EBCDIC	00012100
5:	SLCNT	2	BINARY	00012110
6:	MAXSL	2	BINARY	00012120
7:	SLRBASE	2	BINARY	00012130
8:	ROWBASE	2	BINARY	00012140
9:	SLCBASE	2	BINARY	00012150
10:	COLBASE	2	BINARY	00012160
11:	SHDR1	82	EBCDIC	00012170
12:	SHDR2	16	EBCDIC	00012180
13:	SHDR3	16	EBCDIC	00012190
14:	SDTL	16	EBCDIC	00012200
15:	FLGDEF	2	BINARY	00012210
16:	CRTCNT	2	BINARY	00012220

3201

READ BATCH LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00012300
2:	IPM	IPM BUFFER ADDRES	2	BINARY	00012310
3:	IPMPC1	IPM PANIC CODE 1	2	BINARY	00012320
4:	IPMPC2	IPM PANIC CODE 2	2	BINARY	00012330
5:	SRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00012340
6:	SRBPC1	SRB PANIC CODE 1	2	BINARY	00012350
7:	SRBPC2	SRB PANIC CODE 2	2	BINARY	00012360
8:	TEMP	TEMPORARY WORK AREA	8	BINARY	00012370
9:	PANIC	PANIC CODE WORK AREA	2	BINARY	00012380
10:	DCMD	DISKETTE COMMAND WORD	2	BINARY	00012390
11:	DPANIC1	DISKETTE PANIC CODE 1	2	BINARY	00012400
12:	DPANIC2	DISKETTE PANIC CODE 2	2	BINARY	00012410
13:	DZERO	WORD CONSTANT 0	4	BINARY	00012420
14:	D999999	WORD CONSTANT 999999	4	BINARY	00012430
15:	TVERCNT	TRANS. VALIDATION ERROR ADDR. CNT.	2	BINARY	00012440
16:	FLAG	DISKETTE MOUNTED FLAG	2	BINARY	00012450
17:	COMMA	CONSTANT ","	2	EBCDIC	00012460
18:	MSGTID	TRANSACTION NO.	18	EBCDIC	00012470
19:	MSGDEST	DESTINATION	14	EBCDIC	00012480
20:	MSGSTAT	STATJS	8	EBCDIC	00012490
21:	MSGTVER	VALIDATION ERROR	38	EBCDIC	00012500

3301

WRITE BATCH LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00012550
2:	IPM	IPM BUFFER ADDRES	2	BINARY	00012560
3:	IPMPC1	IPM PANIC CODE 1	2	BINARY	00012570
4:	IPMPC2	IPM PANIC CODE 2	2	BINARY	00012580
5:	SBB	SYSCOM BLOCK BUFFER ADDRESS	2	BINARY	00012590
6:	SBBPC1	SBB PANIC CODE 1	2	BINARY	00012600
7:	SBBPC2	SBB PANIC CODE 2	2	BINARY	00012610
8:	TEMP	TEMPORARY WORK AREA	8	BINARY	00012620
9:	PANIC	PANIC CODE WORK AREA	2	BINARY	00012630
10:	DCMD	DISKETTE COMMAND WORD	2	BINARY	00012640
11:	DPANIC1	DISKETTE PANIC CODE 1	2	BINARY	00012650
12:	DPANIC2	DISKETTE PANIC CODE 2	2	BINARY	00012660
13:	ORTABLE	ORDER REFERENCE TABLE	210	EBCDIC	00012670
14:	FLAG	EXTRA FLAG WORD	2	BINARY	00012680
15:	SRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00012690
16:	SRBPC1	SRB PANIC CODE 1	2	BINARY	00012700
17:	SRBPC2	SRB PANIC CODE 2	2	BINARY	00012710

3401

START BATCH LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00012720
2:	IPM	IPM BUFFER ADDRES	2	BINARY	00012730
3:	IPMPC1	IPM PANIC CODE 1	2	BINARY	00012740
4:	IPMPC2	IPM PANIC CODE 2	2	BINARY	00012750
5:	SRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00012760
6:	SRBPC1	SRB PANIC CODE 1	2	BINARY	00012770
7:	SRBPC2	SRB PANIC CODE 2	2	BINARY	00012780
8:	TEMP	TEMPORARY WORK AREA	8	BINARY	00012790
9:	PANIC	PANIC CODE WORK AREA	2	BINARY	00012800
10:	KA0BFQ	ALARM BUFFER QUEUE ADDRESS	2	BINARY	00012810
11:	KA0ALQ	ALARM QUEUE ADDRESS	2	BINARY	00012820
12:	ALARM	ALARM CODE	2	BINARY	00012830
13:	ALARMPC1	ALARM PANIC CODE 1	2	BINARY	00012840
14:	ALARMPC2	ALARM PANIC CODE 2	2	BINARY	00012850

3501

END BATCH LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00012900
2:	IPM	IPM BUFFER ADDRES	2	BINARY	00012910

3:	IPMPC1	IPM PANIC CODE 1	2	BINARY	00013030
4:	IPMPC2	IPM PANIC CODE 2	2	BINARY	00013040
5:	SRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00013050
6:	SRBPC1	SRB PANIC CODE 1	2	BINARY	00013060
7:	SRBPC2	SRB PANIC CODE 2	2	BINARY	00013070
8:	TEMP	TEMPORARY WORK AREA	8	BINARY	00013080
9:	PANIC	PANIC CODE WORK AREA	2	BINARY	00013090
10:	KA0BFQ	ALARM BUFFER QUEUE ADDRESS	2	BINARY	00013100
11:	KAJALQ	ALARM QUEUE ADDRESS	2	BINARY	00013110
12:	ALARM	ALARM CODE	2	BINARY	00013120
13:	ALARMPC1	ALARM PANIC CODE 1	2	BINARY	00013130
14:	ALARMPC2	ALARM PANIC CODE 2	2	BINARY	00013140

360I DELETE BATCH LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00013150
2:	IPM	IPM BUFFER ADDRESS #1	2	BINARY	00013160
3:	IPMPC1	IPM PANIC CODE 1	2	BINARY	00013170
4:	IPMPC2	IPM PANIC CODE 2	2	BINARY	00013180
5:	SRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00013190
6:	SRBPC1	SRB PANIC CODE 1	2	BINARY	00013200
7:	SRBPC2	SRB PANIC CODE 2	2	BINARY	00013210
8:	TEMP	TEMPORARY WORK AREA	8	BINARY	00013220
9:	PANIC	PANIC CODE WORK AREA	2	BINARY	00013230
10:	IPM2	IPM BUFFER ADDRESS #2	2	BINARY	00013240
11:	DBTABLE	DELETION BATCH TABLE ADDRESS	2	BINARY	00013250
11:	BATCH	RELATIVE BATCH NUMBER	2	BINARY	00013260
12:	ACTBAT	ACTIVE BATCH FLAG WORD	2	BINARY	00013270
13:	QSTFLG	QUESTIONED FLAG WORD	2	BINARY	00013280
14:	PRTFLG	PRINT FLAG WORD	2	BINARY	00013290
15:	DMSG1	BATCH NOT IN FILE	26	EBCDIC	00013300
16:	DMSG2	BATCH MUST BE PENDING OR COMPLETE	72	EBCDIC	00013310
17:	DMSG3	NO CURRENT BATCH INFORMATION	36	EBCDIC	00013320
18:	DMSG4	ALL CURRENT BATCHES ARE ACTIVE	56	EBCDIC	00013330

370I REROUTE SHIPPING LINE LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	PANIC	PANIC CODE WORK AREA	2	BINARY	00013340
2:	STB	SORT TRANSLATION BUFFER	260	BINARY	00013350
3:	TEMP	TEMPORARY WORK AREA	2	BINARY	00013360
4:	SAVE	SAVE OLD SORT LINE	2	BINARY	00013370
5:	NEWSL	NEW SORT LINE ASSIGNMENT	2	BINARY	00013380
6:	SLTXT1	INVALID SORT LINE	34	EBCDIC	00013390

380I RESTORE SHIPPING LINE LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	PANIC	PANIC CODE WORK AREA	2	BINARY	00013400
2:	STB	SORT TRANSLATION BUFFER	260	BINARY	00013410

390I OPEN PSC COMMUNICATIONS LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	OIPSC1		38	EBCDIC	00013420
2:	OIPSC2		36	EBCDIC	00013430
3:	PANIC	PANIC CODE WORK AREA	2	BINARY	00013440
4:	PPANIC1	PSC PANIC CODE 1	2	BINARY	00013450
5:	PPANIC2	PSC PANIC CODE 2	2	BINARY	00013460

400I ASSIGN TRANSACTION STATUS LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:		
1:	FLGDEF	LOCAL FLAG WORD	2	BINARY	00013470
2:	IPM	IPM BUFFER ADDRESS	2	BINARY	00013480
3:	IPMPC1	IPM PANIC CODE 1	2	BINARY	00013490
4:	IPMPC2	IPM PANIC CODE 2	2	BINARY	00013500

5:	SRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00013760
6:	SRBPC1	SRB PANIC CODE 1	2	BINARY	00013770
7:	SRBPC2	SRB PANIC CODE 2	2	BINARY	00013780
8:	CRTCNT	CRT MAXIMUM SCREEN SIZE	2	BINARY	00013790
9:	LINECNT	REPORT LINE COUNT	2	BINARY	00013800
10:	TEMP	TEMPORARY WORK AREA	10	BINARY	00013810
11:	PANIC	PANIC CODE WORK AREA	2	BINARY	00013820
12:	TMX	END TRANSACTION RANGE ID	4	BINARY	00013830
13:	TNSTATUS	NEW TRANSACTION STATUS	2	BINARY	00013840
14:	MSGTNF	TRANSACTION NOT FOUND	32	EBCDIC	00013850
15:	MSGTRE	TRANSACTION RANGE ERROR	34	EBCDIC	00013860
16:	DZERO	DWORD CONSTANT 0	4	BINARY	00013870
17:	DWONE	DWORD CONSTANT 1	4	BINARY	00013880
18:	D999999	DWORD CONSTANT 999999	4	BINARY	00013890

410I RECIRC LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:	
1:	DIRECTXT BANNER	30	EBCDIC	00013900

420I HOLD LOCAL DATA:

LABEL:	DESCRIPTION:	BYTES:	TYPE:	
1:	DI-LLTXT BANNER	28	EBCDIC	00013970

430I BATCH VERIFICATION DISPLAY LAYOUT:

REPORT FIELD:	FROM:	
1: DATE	LOCAL TIME	00014000
2: TIME	LOCAL TIME	00014010
3: PAGE	LOCAL PAGECNT	00014020
4: BATCH ID	120I BID	00014030
5: BATCH STATUS	120I BSTATUS, BSTBL	00014040
6: ORDER COUNT	120I BOCNT	00014050
7: TRANSACTION COUNT	120I BTCNT	00014060
8: HEADER INFORMATION	120I BHDR	00014070

* SEE REPORT LAYOUT SHEET

440I TRANSACTION VERIFICATION DISPLAY LAYOUT:

REPORT FIELD:	FROM:	
1: DATE	LOCAL TIME	00014100
2: TIME	LOCAL TIME	00014110
3: PAGE	LOCAL PAGECNT	00014120
4: TRANSACTION ID	140I TID	00014130
5: TRANSACTION STATUS	140I TSTATUS	00014140
6: PRODUCT ID	140I TPROD	00014150
7: SOURCE LOCATION	140I TSOURCE	00014160
8: DESTINATION	140I TDEST	00014170

* SEE REPORT LAYOUT SHEET

450I CHANGE BATCH STATUS CALL: DICCB

LABEL:	DESCRIPTION:	BYTES:	TYPE:	
INPUT:				00014180
1: BIPM	IPM BUFFER ADDRESS	2	BINARY	00014190
2: BID	BATCH ID	3	EBCDIC	00014200
3: BSTATUS	BATCH STATUS	2	BINARY	00014210
4: BPANIC1	BAD CALL CODE PANIC CODE	2	BINARY	00014220
5: BPANIC2	BAD STATUS PANIC CODE	2	BINARY	00014230
6: BPANIC3	BAD RETURN CODE PANIC CODE	2	BINARY	00014240
7: FLGDEF	LOCAL FLAG WORD	2	BINARY	00014250

LOCAL:

00014260
00014270
00014280
00014290
00014300
00014310
00014320
00014330
00014340
00014350
00014360
00014370
00014380
00014390
00014400
00014410
00014420
00014430
00014440
00014450
00014460
00014470
00014480

	1: SAVE.1	TEMPORARY SAVE #1	2	BINARY	00014490
	2: BNTXT	NOT ON FILE MESSAGE	28	EBCDIC	00014500
					00014510
	OUTPUT:				00014520
					00014530
	1: PANIC	UPDATED PANIC WORK AREA	2	BINARY	00014540
	2: FLGDEF	UPDATED LOCAL FLAG WORD	2	BINARY	00014550
					00014560
460I	WRITE SORT TRANSLATION TABLE CALL: DICCSW				00014570
					00014580
	LABEL:	DESCRIPTION	BYTES:	TYPE:	00014590
					00014600
	INPUT:				00014610
					00014620
	1: STB	SORT TRANSLATION BUFFER	260	BINARY	00014630
	2: DSI	DISK DATA SET	2	BINARY	00014640
					00014650
	LOCAL:				00014660
					00014670
	1: SAVE.2	TEMPORARY SAVE #2	2	BINARY	00014680
					00014690
470I	HELP COMMAND LAYOUT:				00014700
	* SEE REPORT LAYOUT SHEET				00014710
					00014720
					00014730
480I	GET IPM CALL: DICCGI				00014740
					00014750
	USER MUST DEFINE IN PROGRAM LOCAL DATA:				00014760
					00014770
	LABEL:	DESCRIPTION:	BYTES:	TYPE:	00014780
					00014790
	INPUT:				00014800
					00014810
	1: IPMPC1	IPM BUFFER QUEUE EMPTY PANIC CODE	2	BINARY	00014820
					00014830
	OUTPUT:				00014840
					00014850
	1: IPM	IPM BUFFER ADDRESS	2	BINARY	00014860
	2: PANIC	UPDATED PANIC WORK AREA	2	BINARY	00014870
					00014880
490I	GET RECDRD BUFFER CALL: DICCSR				00014890
					00014900
	LABEL:	DESCRIPTION:	BYTES:	TYPE:	00014910
					00014920
	INPUT:				00014930
					00014940
	1: SRBPC1	SRB BUFFER QUEUE EMPTY PANIC CODE	2	BINARY	00014950
					00014960
	OUTPUT:				00014970
					00014980
	1: SRB	SYSCOM RECORD BUFFER ADDRESS	2	BINARY	00014990
	2: PANIC	UPDATED PANIC WORK AREA	2	BINARY	00015000
					00015010
500I	FREE IPM CALL: DICLFI				00015020
					00015030
	LABEL:	DESCRIPTION:	BYTES:	TYPE:	00015040
					00015050
	INPUT:				00015060
					00015070
	1: IPM	IPM BUFFER ADDRESS	2	BINARY	00015080
	2: IPMPC2	IPM BUFFER QUEUE FULL PANIC CODE	2	BINARY	00015090
					00015100
	LOCAL:				00015110
					00015120
	1: SAVE.2	TEMPORARY SAVE #2	2	BINARY	00015130
					00015140
	OUTPUT:				00015150
					00015160
	1: PANIC	UPDATED PANIC WORK AREA	2	BINARY	00015170
					00015180
510I	FREE RECORD BUFFER CALL: DICCFR				00015190
					00015200
	LABEL:	DESCRIPTION:	BYTES:	TYPE:	00015210
					00015220


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INPUT:
    1: SRB          SYSCOM RECORD BUFFER ADDRESS      2    BINARY
    2: SRBPC2       SRB BUFFER QUEUE FULL PANIC CODE  2    BINARY
LOCAL:
    1: SAVE.2       TEMPORARY SAVE #2                2    BINARY
OUTPUT:
    1: PANIC        UPDATED PANIC WORK AREA          2    BINARY
52DI SEND ALARM MESSAGE CALL: OICCSA
    LABEL:          DESCRIPTION                      BYTES: TYPE:
INPUT:
    1: ALARM        ALARM CODE                        2    BINARY
    2: ALARMPCC1   ALARM QUEUE EMPTY PANIC CODE                    2    BINARY
    3: ALARMPCC2   ALARM QUEUE FULL PANIC CODE                      2    BINARY
    4: BID         BATCH ID                          3    EBCDIC
LOCAL:
    1: SAVE.2       TEMPORARY SAVE #2                2    BINARY
OUTPUT:
    1: PANIC        UPDATED PANIC WORK AREA          2    BINARY
53DI DELETE BATCH CALL: OICCSB
    LABEL:          DESCRIPTION                      BYTES: TYPE:
INPUT:
    1: BIPM        IPM BUFFER ADDRESS                2    BINARY
    2: BPANIC1     BAD RELATIVE BATCH ID PANIC CODE                2    BINARY
    3: BPANIC2     BAD CALL CODE PANIC CODE                2    BINARY
    4: BPANIC3     BAD RETURN CODE PANIC CODE              2    BINARY
LOCAL:
    1: SAVE.1       TEMPORARY SAVE #1                2    BINARY
OUTPUT:
    1: PANIC        UPDATED PANIC WORK AREA          2    BINARY
54DI PANIC CALL: OICCPA
    LABEL:          DESCRIPTION                      BYTES: TYPE:
INPUT:
    1: PANIC        PAIC CODE                          2    BINARY
LOCAL:
    1: SAVE.2       TEMPORARY SAVE #2                2    BINARY
    2: NEVEVT      NEVER EVENT CONTROL BLOCK            6    BINARY
55DI GET BLOCK BUFFER CALL: OICCS
    LABEL:          DESCRIPTION                      BYTES: TYPE:
LOCAL:
    1: SAVE.2       TEMPORARY SAVE #2                2    BINARY
OUTPUT:

```

00015230
00015240
00015250
00015260
00015270
00015280
00015290
00015300
00015310
00015320
00015330
00015340
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00015360
00015370
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00015870
00015880
00015890
00015900
00015910
00015920
00015930
00015940
00015950
00015960

11: TIME	TABLE:HH,MM,SS,MO,DA,YR	12	BINARY	00016710
12: PAGECNT	REPORT PAGE COUNT	2	BINARY	00016720
13: LINECNT	REPORT LINE COUNT	2	BINARY	00016730
14: CRTCNT	CRT MAXIMUM SCREEN SIZE	2	BINARY	00016740
15: PRINTCNT	PRINTER MAXIMUM LINE SIZE	2	BINARY	00016750
16: TEMP	TEMPORARY WORK AREA	10	BINARY	00016760
17: PANIC	PANIC CODE WORK AREA	2	BINARY	00016770
18: MSGOSR	BANNER	22	EBCDIC	00016780
19: MSGABORT	REPORT ABORTED	48	EBCDIC	00016790

6101 MIS-SORT REPORT LAYOUT:

REPORT FIELD:	FROM:	
1: DATE	2901 TIME	00016800
2: TIME	2901 TIME	00016810
3: PAGE	2901 PAGECNT	00016820
4: BATCH ID	1201 BID	00016830
5: BATCH STATJS	1201 BSTATUS, BSTBL	00016840
6: ORDER COUNT	1201 BOCNT	00016850
7: TRANSACTION COUNT	2901 RTCNT	00016860
8: HEADER INFORMATION	1201 BHDR	00016870
9: ORDER NUMBER	1301 OCNT	00016880
10: ORDER ID	1301 OID	00016890
11: QUANTITY PLANNED	1301 OTCNT	00016900
12: QUANTITY STAGED	1301 OQSTG	00016910
13: QUANTITY NOT SHIPPED	1301 OQNSTG	00016920
14: QUANTITY STOCK-OUT	1301 OQSTO	00016930
15: QUANTITY EXCESS REPACK	1301 OQXRP	00016940
16: QUANTITY IN SORTATION	1301 OQISORT	00016950
17: QUANTITY NOT PICKED	1301 OQNPKD	00016960
18: QUANTITY MIS-SORT	1301 OQMIS	00016970
19: ORDER % COMPLETE	1301 OPCOM	00016980
20: TRANSACTION NUMBER	2901 RTCNT	00016990
21: TRANSACTION ID	1401 TID	00017000
22: TRANSACTION STATUS	1401 TSTATUS	00017010
23: PRODUCT ID	1401 TPROD	00017020
24: SOURCE LOCATION	1401 TSOURCE	00017030
25: DESTINATION	1401 TDEST	00017040
26: MIS-SORT DESTINATION		00017050

* SEE REPORT LAYOUT SHEET

HELP COMMAND

STATUS REPORT COMMANDS	BATCH CONTROL COMMANDS
SCHEDULE - WORK SCHEDULE BATCHES	: READ - READ BATCH FROM DISKETTE
BATCH * BATCH/ORDER STATUS	: START - START BATCH PROCESSING
ORDER * ORDER/TRANSACTION STATJS	: END - END BATCH PROCESSING
TRANSACT - TRANSACTION STATUS	: WRITE - WRITE BATCH TO DISKETTE
MISSORT * MISSORT EXCEPTIONS	: DELETE - DELETE BATCH INFORMATION
ROJTE - SORT LINE TRANSLATIONS	: ASSIGN - ASSIGN TRANSACTION STATUS
> CANCEL * (TO TERMINATE REPORTS)	:

SYSTEM CONTROL COMMANDS	SORTATION CONTROL COMMANDS
HELP - LIST SYSTEM COMMANDS	: HOLD - HOLD ON ERROR
RESET - RESET ERROR COUNTS	: RECIRC - RECIRCULATE ON ERROR
ERROR - DISPLAY SYSTEM ERROR COUNTS	: OPEN - OPEN PSC COMMUNICATIONS
\$T - SET SYSTEM TIME	: REROUTE - ASSIGN ALTERNATE SORT LINE
	: RESTORE - RESTORE ORIGINAL SORT LINE

-- BATCH -- NO. OF NO. OF
ID. STATUS ORDERS TRANS. HEADER INFORMATION

S01	ACTIVE	3	150	SAMPLE BATCH OF 150 TRANSACTIONS
T02	ACTIVE	1	9	TEST BATCH T02
T03	PENDING	1	8	TEST BATCH T03
T04	PENDING	1	9	TEST BATCH T04
T05	COMPLETE	9	9	TEST BATCH T05
T06	PENDING	1	9	TEST BATCH T06
T07	PENDING	9	9	TEST BATCH T07

T08 COMPLETE 1 9 TEST BATCH T08
 T09 COMPLETE 9 9 TEST BATCH T09
 T10 PENDING 1 9 TEST BATCH T10

-- BATCH -- NO. OF NO. OF
 ID. STATUS ORDERS TRANS. HEADER INFORMATION

S01 ACTIVE 3 150 SAMPLE BATCH OF 150 TRANSACTIONS

#	ORDER ID.	QUANTITIES: PLAN	STAGED	NOT SHIP	% COMP.	STOCK -OJT	EXCESS REPACK	IN SORT.	NOT PICKED	MIS-SORT
1	000001	50	26	24	74.0	2	9	3	7	3
2	000002	50	35	15	74.0	2	0	4	8	1
3	000003	50	50	0	100.0	0	0	0	0	0

BATCH PERCENT COMPLETE: 82.7

1 000001 50 26 24 74.0 2 9 3 7 3

#	- TRANSACTION - NUMBER	STATUS	PRODUCT IDENTIFICATION	SOURCE LOCATION	DESTINATION
1	16	MIS-SORTED	000007	00007	1 (2)
2	91	EXCESS REPACK	000082	00082	1
3	94	EXCESS REPACK	000085	00085	1
4	97	MIS-SORTED	000088	00088	1 (3)
5	100	EXCESS REPACK	000091	00091	1
6	103	EXCESS REPACK	000094	00094	1
7	106	EXCESS REPACK	000097	00097	1
8	109	EXCESS REPACK	000100	00100	1
9	112	EXCESS REPACK	000103	00103	1
10	115	EXCESS REPACK	000106	00106	1
11	118	EXCESS REPACK	000109	00109	1
12	121	STOCK-OUT	000112	00112	1
13	124	IN SORTATION	000115	00115	1
14	127	NOT PICKED	000118	00118	1
15	130	IN SORTATION	000121	00121	1
16	133	IN SORTATION	000124	00124	1
17	136	NOT PICKED	000127	00127	1
18	139	NOT PICKED	000130	00130	1
19	142	NOT PICKED	000133	00133	1
20	145	STOCK-OUT	000136	00136	1
21	148	NOT PICKED	000139	00139	1
22	151	MIS-SORTED	000142	00142	1 (2)
23	154	NOT PICKED	000145	00145	1
24	157	NOT PICKED	000148	00148	1

2 000002 50 35 15 74.0 2 0 4 8 1

#	- TRANSACTION - NUMBER	STATUS	PRODUCT IDENTIFICATION	SOURCE LOCATION	DESTINATION
1	113	IN SORTATION	000104	00104	2
2	116	IN SORTATION	000107	00107	2
3	119	IN SORTATION	000110	00110	2
4	122	STOCK-OJT	000113	00113	2
5	125	STOCK-OJT	000116	00116	2
6	128	NOT PICKED	000119	00119	2
7	131	NOT PICKED	000122	00122	2
8	134	NOT PICKED	000125	00125	2
9	137	NOT PICKED	000128	00128	2
10	140	NOT PICKED	000131	00131	2
11	143	IN SORTATION	000134	00134	2
12	149	NOT PICKED	000140	00140	2
13	152	NOT PICKED	000143	00143	2
14	155	NOT PICKED	000146	00146	2
15	158	MIS-SORTED	000149	00149	2 (1)

3 000003 50 50 0 100.0 0 0 0 0 0

**** NO EXCEPTIONS IN THIS ORDER

**** REPORT COMPLETE

4,336,589

129

130

TRANSACTION NUMBER	STATUS	ORDER ID	PRODUCT ID	SOURCE LOCATION	DESTINATION
109	EXCESS REPACK	000001	000100	00100	1

#	ORDER ID.	QUANTITIES: PLAN	STAGED	NOT SHIP	% COMP.	STOCK -DJT	EXCESS REPACK	IN SORT.	NOT PICKED	MIS-SORT
1	000001	50	26	24	74.0	2	9	3	7	3

#	TRANSACTION NUMBER	STATUS	PRODUCT IDENTIFICATION	SOURCE LOCATION	DESTINATION					
1	16	MIS-SORTED	000007	00007	1 (2)					
2	97	MIS-SORTED	000088	00088	1 (3)					
3	151	MIS-SORTED	000142	00142	1 (2)					
2	000002	50	35	15	74.0	2	0	4	8	1

#	TRANSACTION NUMBER	STATUS	PRODUCT IDENTIFICATION	SOURCE LOCATION	DESTINATION
1	158	MIS-SORTED	000149	00149	2 (1)

#	ORDER ID.	QUANTITIES: PLAN	STAGED	NOT SHIP	% COMP.	STOCK -DJT	EXCESS REPACK	IN SORT.	NOT PICKED	MIS-SORT
3	000003	50	50	0	100.0	0	0	0	0	0

**** NO MIS-SORTS IN THIS ORDER

**** REPORT COMPLETE

SORT LINE	CHANGE TO	SORT LINE	CHANGE TO	SORT LINE	CHANGE TO
1		11		21	19
2	3	12			
3		13			
4		14			
5		15	11		
6	8	16			
7		17			
8		18			
9		19			
10		20			

TRANSACTIONS LOST	REJECTED	BY PSC	SENT TO RECIRCULATION FROM INDUCTION	TOTAL
0	3	2	5	7

SCANNER NO-READS

SCANNER #1	SCANNER #2
0	1

M I S - S O R T S

LINE	COUNT	LINE	COJNT
1	00001	18	00000
2	00002	19	00000
3	00001	20	00000
4	00000	21	00000
5	00000		
6	00000		
7	00000		
8	00000		
9	00000		
10	00000		
11	00000		

12 0
 13 0
 14 0
 15 0
 16 0
 17 0

00000010
 00000020
 00000030

SYSTEM START-UP PROGRAM

00000040
 00000050
 00000060

00000070
 00000080
 00000090

- MODULE HISTORY -

00000100
 00000110
 00000120

PROJECT: DAS 75-01709

00000130

SUB-SYSTEM: AUXILIARY FUNCTIONS

00000140

00000150

MODULE: \$INIT 60

00000160

00000170

00000180

- MODULE ABSTRACT -

00000190

00000200

00000210

THIS MODULE IS THE SYSTEM START-JP PROGRAM. IT LOADS ALL
 OTHER PROGRAMS IN THE SYSTEM AND CALLS THE FILE MANAGER
 TO INITIALIZE THE FILES. AFTER IT HAS DONE ALL THIS, IT
 PERFORMS A PROGSTOP.

00000220

00000230

00000240

WAIT FOR 1 SECOND

00000250

LOAD THE ALARM MESSAGE PRINTING SUB-SYSTEM

00000260

IF NO LOAD ERRORS ENCOUNTERED THEN

00000270

00000280

LOAD THE FILE MANAGER SUB-SYSTEM

00000290

IF LOAD ERROR ENCOUNTERED THEN

00000300

07AL

08AL

PANIC!

00000310

(ELSE)

00000320

ENDIF

00000330

00000340

09FM

GET A FILE MANAGER INPUT BUFFER

00000350

IF NONE AVAILABLE THEN

00000360

07AL

08AL

PANIC!

00000370

(ELSE)

00000380

ENDIF

00000390

00000400

24FM

SET UP BUFFER TO INITIALIZE THE FILES

00000410

08FM

CALL THE FILE MANAGER

00000420

24FM

IF FILE MANAGER COMPLETION CODE NOT SUCCESSFUL THEN

00000430

07AL

08AL

PANIC!

00000440

(ELSE)

00000450

ENDIF

00000460

00000470

09FM

RETURN THE FILE MANAGER BUFFER TO FREE POOL

00000480

IF FREE POOL FULL THEN

00000490

07AL

08AL

PANIC!

00000500

(ELSE)

00000510

ENDIF

00000520

LOAD THE PSC COMMUNICATIONS SUB-SYSTEM

00000530

IF LOAD ERROR ENCOUNTERED THEN

00000540

07AL

08AL

PANIC!

00000550

(ELSE)

00000560

ENDIF

00000570

LOAD THE SCANNER INPUT SUB-SYSTEM

00000580

IF LOAD ERROR ENCOUNTERED THEN

00000590

07AL

08AL

PANIC!

00000600

(ELSE)

00000610

ENDIF

00000620

LOAD THE OPERATOR INTERFACE SUB-SYSTEM

00000630

IF LOAD ERROR ENCOUNTERED THEN

00000640

07AL

PANIC!

00000650

00000660

00000670

00000680

08AL	(ELSE)	00000690
	ENDIF	00000700
02AL	GET AN ALARM MESSAGE BUFFER	00000710
	IF NONE AVAILABLE THEN	00000720
07AL	PANIC!	00000730
08AL	(ELSE)	00000740
	ENDIF	00000750
02AX	IF START-UP CODE = NEW IPL THEN	00000760
03AL	SET BUFFER TO PRINT "SYSTEM STARTED" MSG.	00000770
	ELSE	00000780
03AL	SET BUFFER TO PRINT "POWER FAIL/RESTART" MSG.	00000790
	ENDIF	00000800
01AL	PLACE BUFFER ON ALARM MESSAGE INPUT QUEUE	00000810
	IF QUEUE FULL THEN	00000820
07AL	PANIC!	00000830
08AL	(ELSE)	00000840
	ENDIF	00000850
	(ELSE)	00000860
	ENDIF	00000870
	PROGSTOP	00000880
		00000890
		00000900
		00000910
		00000010
		00000020
		00000030
		00000040
		00000050
	SYSTEM INSTALLATION PROGRAM	00000060
		00000070
		00000080
		00000090
	- MODULE HISTORY -	00000100
		00000110
		00000120
	PROJECT: DAS 73-01709	00000130
		00000140
	SUB-SYSTEM: AUXILIARY FUNCTIONS	00000150
		00000160
	MODULE: \$INITIAL 61	00000170
		00000180
		00000190
	- MODULE ABSTRACT -	00000200
		00000210
	THIS MODULE REBUILDS THE ORDER FILE FROM THE TRANSACTION	00000220
	RECORDS IN THE TRANSACTION FILE FOR EVERY BATCH ON FILE. IT	00000230
	THEN LOADS THE SYSTEM START-UP PROGRAM INTO PARTITION 3 AND	00000240
	EXITS.	00000250
		00000260
	INITIALIZATION	00000270
02AX	SAVE START-UP CODE	00000280
	DISPLAY SYSTEM INITIALIZATION MESSAGE	00000290
39FM	READ SYSTEM PARAMETER FILE	00000300
	IF THERE WERE NO READ ERRORS THEN	00000310
39FM	GET ADDRESS OF \$SYSCOM PARAMETER TABLE	00000320
39FM	MOVE SYSTEM PARAMETERS TO \$SYSCOM	00000330
01FM	READ THE BATCH FILE	00000340
	IF THERE WERE NO READ ERRORS, THEN	00000350
01FM	DOUNTIL ALL RECORDS IN BATCH FILE HAVE BEEN TESTED	00000360
01FM	INDEX TO FIRST/NEXT BATCH RECORD	00000370
01FM	IF BATCH STATUS = ACTIVE, PENDING, OR COMPLETE, THEN	00000380
01FM	SAVE RELATIVE BATCH, TRANSACTION COUNT, AND BATCH ID	00000390
03AX		00000400
	(ELSE)	00000410
	ENDIF	00000420
	ENDDD	00000430
03AX	DOUNTIL END OF BATCH TABLE REACHED	00000440
03AX	GET FIRST/NEXT RELATIVE BATCH AND TRANS COUNT FROM TBL	00000450
03AX	IF THIS IS NOT THE END OF THE TABLE, THEN	00000460
04FM	CALCULATE FIRST SECTOR IN TRANS FILE FOR THIS BATCH	00000470
04FM	CALCULATE NUMBER OF SECTORS IN THIS BATCH	00000480
01FM	MOVE NUMBER OF SECTORS IN BATCH TO SECTORS LEFT	00000490

	IF 'SECTORS LEFT' <= MAX SECTORS, THEN	00000500
	MOVE 'SECTORS LEFT' TO 'SECTOR COUNT'	00000510
	ZERO 'SECTORS LEFT'	00000520
	ELSE	00000530
	MOVE MAX SECTORS TO 'SECTOR COUNT'	00000540
	SUBTRACT MAX SECTORS FROM 'SECTORS LEFT'	00000550
	UPDATE TRANSACTION FILE SECTOR POINTER	00000560
	ENDIF	00000570
04AX	SET FLAG TO FIRST BUFFER	00000580
04FM	READ TO FIRST BUFFER 'SECTOR COUNT' SECTORS FROM ---	00000590
05AX	--- THE TRANSACTION FILE	00000600
05AX	SET UP BUFFER POINTER FOR THE FIRST BUFFER	00000610
06AX	ZERO THE ORDER TABLE	00000620
07AX	DO UNTIL THE TRANS COUNT = 0, OR BATCH ERROR FLAG ---	00000630
08AX		00000640
	--- SET, OR DISK ERROR FLAG SET	00000650
	WAIT FOR THE READ TO COMPLETE	00000660
	IF NO READ ERRORS, THEN	00000670
	IF 'SECTORS LEFT' NOT = 0, THEN	00000680
	IF 'SECTORS LEFT' <= MAX SECTORS, THEN	00000690
	MOVE 'SECTORS LEFT' TO 'SECTOR COUNT'	00000700
	ZERO 'SECTORS LEFT'	00000710
	ELSE	00000720
	MOVE MAX SECTORS TO 'SECTOR COUNT'	00000730
	SUBTRACT MAX SECTORS FROM 'SECTORS LEFT'	00000740
	UPDATE TRANSACTION FILE SECTOR POINTER	00000750
	ENDIF	00000760
	IF FLAG SET TO FIRST BUFFER, THEN	00000770
05AX	READ TO SECOND BUFFER 'SECTOR COUNT' ---	00000780
	--- SECTORS FROM THE TRANSACTION FILE	00000790
04AX	SET FLAG TO SECOND BUFFER	00000800
	ELSE	00000810
05AX	READ TO FIRST BUFFER 'SECTOR COUNT' ---	00000820
	--- SECTORS FROM THE TRANSACTION FILE	00000830
04AX	SET FLAG TO FIRST BUFFER	00000840
	ENDIF	00000850
	(ELSE)	00000860
	ENDIF	00000870
05AX	GET BUFFER POINTER	00000880
05AX	DO UNTIL END OF BUFFER OR TRANS COUNT = 0	00000890
04FM	INDEX TO FIRST/NEXT TRANSACTION RECORD	00000900
04FM	IF THE RECORD IS USED, THEN	00000910
04FM	GET REL. ORDER FROM RECORD	00000920
04FM	GET TRANSACTION STATUS FROM RECORD	00000930
06AX	INDEX INTO ORDER TABLE BY REL. ORDER	00000940
06AX	INDEX INTO STATUS COUNT FOR THIS ORDER	00000950
06AX	INCREMENT STATUS COUNT	00000960
	DECREMENT TRANSACTION COUNT	00000970
	ELSE	00000980
	SET TRANSACTION COUNT TO ZERO	00000990
07AX	SET BATCH ERROR FLAG	00010000
	DISPLAY BATCH ERROR MESSAGE	00010100
	ENDIF	00010200
	ENDDO	00010300
	IF TRANSACTION COUNT NOT = 0, THEN	00010400
	IF FLAG SET TO SECOND BUFFER, THEN	00010500
05AX	SET POINTER TO SECOND BUFFER	00010600
	ELSE	00010700
05AX	SET POINTER TO FIRST BUFFER	00010800
	ENDIF	00010900
	(ELSE)	00011000
	ENDIF	00011100
	ELSE	00011200
08AX	SET DISK ERROR FLAG	00011300
	ENDIF	00011400
	ENDDO	00011500
07AX	IF BATCH ERROR FLAG AND DISK ERROR FLAG NOT SET, THEN	00011600
08AX		00011700
05FM	CALCULATE FIRST ORDER SECTOR FOR CURRENT BATCH	00011800
05FM	READ ORDER SECTORS FOR THIS BATCH	00011900
	IF NO READ ERRORS, THEN	00012000
05FM	DO UNTIL END OF ORDER RECORDS FOR THIS BATCH	00012100
06AX		00012200

05FM	INDEX TO FIRST/NEXT ORDER RECORD	00001230
06AX		00001240
05FM	UPDATE ORDER STATUS COUNTS IN RECORD	00001250
06AX		00001260
	ENDDJ	00001270
05FM	WRITE ORDER SECTORS FOR THIS BATCH	00001280
	IF WRITE ERRORS, THEN	00001290
08AX	SET DISK ERROR FLAG	00001300
	(ELSE)	00001310
	ENDIF	00001320
	ELSE	00001330
08AX	SET DISK ERROR FLAG	00001340
	ENDIF	00001350
	ELSE	00001360
07AX	RESET BATCH ERROR FLAG	00001370
	ENDIF	00001380
	(ELSE)	00001390
	ENDIF	00001400
08AX	IF DISK ERROR FLAG SET, THEN	00001410
	DISPLAY DISK ERROR MESSAGE FOR THIS BATCH	00001420
	DISPLAY WARNING	00001430
08AX	RESET DISK ERROR FLAG	00001440
	(ELSE)	00001450
	ENDIF	00001460
	ENDDJ	00001470
	ELSE	00001480
	DISPLAY DISK ERROR WHILE READING BATCH FILE MESSAGE	00001490
	DISPLAY WARNING MESSAGE	00001500
	ENDIF	00001510
	ELSE	00001520
	DISPLAY UNABLE TO READ SYSTEM PARAMETERS	00001530
	DISPLAY SYSTEM INITIALIZATION ABORTED	00001540
	ENDIF	00001550
	IF SYSTEM PARAMETERS READ J.K. THEN	00001560
	LOAD SYSTEM START-UP PROGRAM	00001570
	(ELSE)	00001580
	ENDIF	00001590
	PROGSTOP	00001600
		00000010
		00000020
		00000030
		00000040
		00000050
	PSC OUTPUT CONTROL TASK	00000060
		00000070
		00000080
		00000090
	- MODULE HISTORY -	00000100
		00000110
		00000120
	PROJECT: DAS 75-01709	00000130
		00000140
	SUB-SYSTEM: PSC COMMUNICATIONS	00000150
		00000160
	MODULE: PSOC 49	00000170
		00000180
		00000190
	- MODULE ABSTRACT -	00000200
		00000210
	THIS TASK PROCESSES ALL PSC OUTPUT MESSAGES. IT REMOVES THE	00000220
	INPUT BUFFER FROM THE PSC OUTPUT QUEUE AND IF THE COMMUNICATIONS	00000230
	LINE IS OPEN, FORMATES THE MESSAGE, AND SENDS IT TO THE PSC.	00000240
	IT THEN WAITS ON THE TRANSMISSION COMPLETION CODE TO BE POSTED,	00000250
	AND WILL RETRY THE MESSAGE IF AN ERROR WAS ENCOUNTERED OR WILL	00000260
	CONTINUE COMMUNICATIONS WITH THE NEXT MESSAGE FOUND IN THE	00000270
	PSC OUTPUT QUEUE.	00000280
07PS	SET COMMUNICATIONS LINK = CLOSED	00000290
	SET LAST COMPLETION CODE = IDLE RECEIVED	00000300
08PS	RESET IDLE MESSAGE COUNT	00000310
09PS	CLEAR PSC RESPONSE COUNTS	00000320
10PS	RESET RESPONSE DUMP FLAG	00000330
13PS	RESET TOTAL NO. OF TIME-OUTS	00000340
14PS	RESET INVALID SEQUENCE COUNTER	00000350
15PS	RESET INVALID CHECKSUM COUNTER	00000360

39FM	GET SORT PARAMETERS FROM SYSCOM	00000370
	ATTACH THE PSC POLL TASK	00000380
	ATTACH PSC RESPONSE PROCESSING TASK	00000390
	ATTACH PSC INPUT/OUTPUT TASK	00000400
	DDUNTIL CPU STOPPED	00000410
15SC	WAIT ON SCANNER SLOW DOWN EVENT	00000420
28PS	WAIT ON PSC RESPONSE OVERFLOW EVENT	00000430
	IF LAST COMPLETION CODE NOT = DATA RECEIVED OR --	00000440
07PS	-- COMMUNICATION LINK CLOSED THEN	00000450
02PS	DDUNTIL DATA RECEIVED FROM PSC OUTPUT QUEUE	00000460
02PS	WAIT FOR PSC RJN EVENT TO BE POSTED	00000470
02PS	RETRIEVE BUFFER FROM PSC OUTPUT QUEUE	00000480
02PS	IF QUEUE EMPTY THEN	00000490
02PS	RESET PSC RJN EVENT	00000500
	(ELSE)	00000510
	ENDIF	00000520
	ENDDO	00000530
	ELSE	00000540
02PS	RETRIEVE BUFFER FROM PSC OUTPUT QUEUE	00000550
01PS	IF QUEUE EMPTY THEN	00000560
	SET MESSAGE CODE = IDLE	00000570
	ELSE	00000590
01PS	SET MESSAGE CODE = INPUT MESSAGE CODE	00000600
	SAVE DATA FROM INPUT BUFFER	00000610
01PS	RETURN BUFFER TO FREE POOL	00000620
	IF FREE POOL FULL THEN	00000630
11PS	PAVIC!	00000640
	(ELSE)	00000650
	ENDIF	00000660
	ENDIF	00000670
	ENDIF	00000680
07PS	IF COMMUNICATIONS LINK = CLOSED THEN	00000690
	IF MESSAGE CODE = OPEN LINK THEN	00000700
12PS	INITIALIZE THE SEQUENCE BYTES	00000710
07PS	SET COMMUNICATIONS LINK = OPEN	00000720
	ELSE	00000730
	SET MESSAGE CODE = NO MESSAGE	00000740
	ENDIF	00000750
	ELSE	00000760
	IF MESSAGE CODE = OPEN LINK THEN	00000770
	SET MESSAGE CODE = NO MESSAGE	00000780
	(ELSE)	00000790
	ENDIF	00000800
	ENDIF	00000810
	IF MESSAGE CODE NOT = NO MESSAGE THEN	00000820
12PS	TOGGLE LAST TRANSMITTED SEQUENCE BYTE	00000830
16PS	CALL MESSAGE FORMATTING ROUTINE	00000840
	IF NO ERROR RETURNED THEN	00000850
	SET RETRY COUNT = MAX. NO. OF RETRIES	00000860
08PS	SET IDLE MESSAGE COUNT = MAX. IDLE COUNT	00000870
	DDWHILE RETRY COUNT > ZERO	00000880
17PS	RESET TRANSMISSION COMPLETION EVENT	00000890
	ATTACH MESSAGE TIME-OUT TASK	00000900
29PS	POST PSC INPUT/OUTPUT EVENT	00000910
17PS	WAIT FOR TRANSMISSION COMPLETION EVENT	00000920
	IF COMPLETION EVENT CODE = RETRY THEN	00000930
	SUBTRACT 1 FROM RETRY COUNT	00000940
	ELSE	00000950
	SET RETRY COUNT LESS THAN ZERO	00000960
	ENDIF	00000970
	ENDDO	00000980
	IF RETRY COUNT = 0 THEN	00000990
07PS	SET COMMUNICATIONS LINK = CLOSED	00001000
02AL	GET ALARM MESSAGE BUFFER	00001010
	IF ONE RECEIVED THEN	00001020
03AL	SET UP BUFFER TO PRINT "COMM. FAILURE" MSG.	00001030
01AL	PLACE BUFFER ADDR. ON ALARM MSG. QUEUE	00001040
	IF ALARM MSG. QUEUE FULL THEN	00001050
11PS	PAVIC!	00001060
	(ELSE)	00001070
	ENDIF	00001080
	ELSE	00001090

11PS	PANIC!	03001100
	ENDIF	03001110
	ELSE	03001120
12PS	TOGGLE LAST RECEIVED SEQUENCE BYTE	03001130
	ENDIF	03001140
	ELSE	03001150
11PS	PANIC!	03001160
	ENDIF	03001170
	(ELSE)	03001180
	ENDIF	03001190
	IF LINK WAS JUST OPENED THEN	03001200
02AL	GET ALARM MESSAGE BUFFER	03001210
	IF ONE RECEIVED THEN	03001220
03AL	SET UP BUFFER TO PRINT "COMM. LINK OPEN" MSG.	03001230
01AL	PLACE BUFFER ADDRESS IN ALARM MSG. QUEUE	03001240
	IF ALARM MSG. QUEUE FULL THEN	03001250
11PS	PANIC!	03001260
	(ELSE)	03001270
	ENDIF	03001280
	ELSE	03001290
11PS	PANIC!	03001300
	ENDIF	03001310
	(ELSE)	03001320
	ENDIF	03001330
	ENDDO	03001340
	PROGSTOP	03001350
		03000010
		03000020
		03000030
		03000040
		03000050
	PSC RESPONSE PROCESSOR ROUTINE	03000060
		03000070
		03000080
		03000090
	- MODULE HISTORY -	03000100
		03000110
		03000120
	PROJECT: DAS 75-01739	03000130
		03000140
	SUB-SYSTEM: PSC COMMUNICATIONS	03000150
		03000160
	MODULE: PSRP 53	03000170
		03000180
		03000190
	- MODULE ABSTRACT -	03000200
		03000210
	THIS TASK PROCESSES ALL VALID RESPONSES FROM THE PSC.	03000220
	IT CALLS THE FILE MANAGER TO UPDATE THE TRANSACTION	03000230
	STATUS IF NEEDED. IT ALSO CALLS THE MESSAGE PROCESSING	03000240
	ROUTINE IF A MESSAGE IS TO BE PRINTED.	03000250
	UNTIL CPU STOPPED	03000260
24PS	UNTIL DATA RECEIVED FROM INPUT QUEUE	03000270
30PS	WAIT FOR PSC RESPONSE PROCESSING EVENT	03000280
	REMOVE AN ENTRY FROM THE INPUT QUEUE	03000290
	IF QUEUE EMPTY THEN	03000300
28PS	IF PSC OVERFLOW FLAG SET THEN	03000310
23PS	GET A PSC RESPONSE BUFFER	03000320
	IF NONE AVAILABLE THEN	03000330
11PS	PANIC !	03000340
	ELSE	03000350
22PS	MOVE LAST DATA RECEIVED INTO BUFFER	03000360
28PS	POST PSC RESPONSE OVERFLOW EVENT	03000370
	ENDIF	03000380
	ELSE	03000390
30PS	RESET PSC RESPONSE PROCESSING EVENT	03000400
	ENDIF	03000410
	(ELSE)	03000420
	ENDIF	03000430
	ENDDO	03000440
25PS	IF MSG. TYPE RECEIVED = TRANSFER COMPLETE OR MIS-SORT THEN	03000450
	CLEAR STATUS FLAG	03000460
25PS	IF INPUT DATA WITHIN RANGE THEN	03000470
25PS	IF MSG. TYPE REC. = TRANSFER COMPLETE THEN	03000480

	SET NEW STATUS = STAGED FOR SHIPPMENT	00000490
25PS	IF DEST. RECEIVED = REJECT THEN	00000500
40FM	ADD 1 TO TOTAL TRANS. SENT TO REJECT	00000510
	ELSE	00000520
25PS	IF DEST. RECEIVED = RECIRCULATION THEN	00000530
40FM	ADD 1 TO TOTAL TRANS. SENT TO RECIRCULATION	00000540
	(ELSE)	00000550
	ENDIF	00000560
	ENDIF	00000570
	ELSE	00000580
	SET NEW STATUS = MIS-SORT	00000590
40FM	ADD 1 TO NUMBER OF MIS-SORTS AT THIS LINE	00000600
26PS	CALL ALARM PROCESSING ROUTINE	00000610
25PS	IF MIS-SORT DESTINATION = 0 THEN	00000620
	SET STATUS CHANGE FLAG	00000630
	(ELSE)	00000640
	ENDIF	00000650
	ENDIF	00000660
	IF STATUS CHANGE FLAG CLEAR THEN	00000670
09FM	SET A FILE MANAGER CALL BUFFER	00000680
	IF ONE RECEIVED THEN	00000690
17FM	SET UP BUFF. FOR A CHANGE STATUS CALL	00000700
17FM	SET STATUS = NEW STATUS	00000710
17FM	SET DESTINATION = DESTINATION RECEIVED FROM PSC	00000720
08FM	CALL THE FILE MANAGER	00000730
17FM	SAVE THE COMPLETION CODE	00000740
09FM	RETURN THE FILE MANAGER BUFFER TO THE FREE POOL	00000750
	IF FREE POOL FULL THEN	00000760
11PS	PANIC!	00000770
	ELSE	00000780
	IF COMPLETION CODE NOT SUCCESSFUL THEN	00000790
	IF COMP. CODE = TRANS. NOT ON FILE THEN	00000800
02AL	GET ALARM MESSAGE BUFFER	00000810
	IF ONE RECEIVED THEN	00000820
03AL	SET JP BUFF. FOR "NOT ON FILE MSG"	00000830
01AL	PLACE BUFF. ON ALARM MSG. QUEUE	00000840
	IF QUEUE FULL THEN	00000850
11PS	PANIC!	00000860
	(ELSE)	00000870
	ENDIF	00000880
	ELSE	00000890
11PS	PANIC!	00000900
	ENDIF	00000910
	ELSE	00000920
11PS	PANIC!	00000930
	ENDIF	00000940
	(ELSE)	00000950
	ENDIF	00000960
	ENDIF	00000970
	ELSE	00000980
11PS	PANIC!	00000990
	ENDIF	00010000
	(ELSE)	00010100
	ENDIF	00010200
	ELSE	00010300
	IF TRANS. NO. = ZERO AND DESTINATION NOT EQUAL TO --	00010400
	-- RECIRC OR ERROR LINE THEN	00010500
25PS	CALL ALARM PROCESSING ROUTINE	00010600
	(ELSE)	00010700
	ENDIF	00010800
	ENDIF	00010900
	ELSE	00011000
26PS	CALL ALARM PROCESSING ROUTINE	00011100
	ENDIF	00011200
24PS	RETURN BUFFER TO FREE POOL	00011300
	IF BUFFER FULL THEN	00011400
11PS	PANIC!	00011500
	(ELSE)	00011600
	ENDIF	00011700
	ENDDD	00011800
	EXIT	00011900
		00000010
		00000020

PSC CHECKSUM CALCULATION SUBROUTINE

- MODULE HISTORY -

PROJECT: DAS 75-01709
 SUB-SYSTEM: PSC COMMUNICATIONS
 MODULE: PCKSM 39

- MODULE ABSTRACT -

THIS SUBROUTINE TAKES EBCDIC CHARACTER INPUT, CONVERTS IT TO ASCII AND CALCULATES AN ASCII CHECKSUM CHARACTER. THIS CHARACTER IS THEN MADE PRINTABLE AND CONVERTED BACK EBCDIC AND PASSED BACK TO THE CALLER.

```

SET INPUTS
CLEAR CHECKSUM RESULT LOCATION
CLEAR ERROR FLAG
DO 13 TIMES
  SET NEXT EBCDIC CHARACTER FROM BUFFER
  DO WHILE EBCDIC CHAR NOT EQ TO TABLE VALUE AND NOT END OF TABLE
    BUMP TO NEXT TABLE ENTRY
  ENDDO
  IF NOT AT END OF TABLE, THEN
    ADD ASCII VALUE TO RESULT
  ELSE
    SET ERROR FLAG
  ENDIF
  BUMP TO NEXT BUFFER POSITION
ENDDO
SET BIT 0 OFF IN RESULT
COMPLEMENT BIT 2 IN RESULT
SET BIT 1 TO RESULT OF COMPLEMENT
DO WHILE ASCII CHARACTER NOT EQ TO TABLE AND NOT END OF TABLE
  BUMP TO NEXT TABLE ENTRY
ENDDO
IF NOT END OF TABLE AND ERROR FLAG NOT SET THEN
  MOVE EBCDIC VALUE TO CHECK SUM CHARACTER
ELSE
  MOVE A -1 TO CHECK SUM CHARACTER
ENDIF
RETURN

```

PSC MESSAGE TRANSMIT/RECEIVE TASK

- MODULE HISTORY -

PROJECT: DAS 75-01709
 SUB-SYSTEM: PSC COMMUNICATIONS
 MODULE: PSID 51

- MODULE ABSTRACT -

```

00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210

```

	THIS TASK TAKES THE FORMATTED MESSAGE AND TRANSMITS IT TO THE PSC THEN STARTS THE PSC RESPONSE TIME-OUT TASK AND WAITS FOR THE RESPONSE FROM THE PSC. IF THE RESPONSE HAS TIMED OUT, THE COMPLETION CODE IS SET TO "RETRY". IF THE RESPONSE SEQUENCE BYTE OR MESSAGE CHECKSUM IS INVALID, THE COMPLETION CODE IS SET TO "RETRY". IF THE RESPONSE IS VALID AND DATA WAS RECEIVED FROM THE PSC, THE COMPLETION CODE IS SET TO "DATA RECEIVED" AND THE DATA IS PASSED TO THE PSC RESPONSE PROCESSING TASK. IF THE RESPONSE IS VALID AND AN IDLE WAS RECEIVED FROM THE PSC, THE COMPLETION CODE IS SET TO "IDLE RECEIVED".	00000220
	ENQUEUE THE PSC TERMINAL	00000230
	UNTIL CPU STOPPED	00000240
	RESET ERROR CODE	00000250
29PS	WAIT FOR PSC INPUT/OUTPUT EVENT	00000260
21PS	SET COMPLETION CODE = UNUSED	00000270
	GET FORMATTED MESSAGE	00000280
	OUTPUT THE MESSAGE TO THE PSC	00000290
	READ THE RESPONSE FROM THE PSC	00000300
22PS	IF CORRECT NUMBER OF BYTES RECEIVED THEN	00000310
12PS	IF REC. SEQ. BYTE = TOGGLE OF LAST RECEIVED SEQ. BYTE THEN	00000320
22PS	COMPUTE MESSAGE CHECKSUM	00000330
22PS	IF INPUT CHECKSUM = COMPUTED CHECKSUM THEN	00000340
22PS	IF MESSAGE TYPE = IDLE THEN	00000350
21PS	SET COMPLETION CODE = IDLE RECEIVED	00000360
30PS	POST RESPONSE PROCESSING EVENT	00000370
	ELSE	00000380
21PS	SET COMPLETION CODE = DATA RECEIVED	00000390
	ENDIF	00000400
	ELSE	00000410
15PS	ADD 1 TO INVALID CHECKSUM COUNTER	00000420
21PS	SET COMPLETION CODE = RETRY	00000430
	SET ERROR CODE = INVALID CHECKSUM	00000440
	ENDIF	00000450
	ELSE	00000460
14PS	ADD 1 TO INVALID SEQUENCE COUNTER	00000470
21PS	SET COMPLETION CODE = RETRY	00000480
	SET ERROR CODE = INVALID SEQUENCE	00000490
	ENDIF	00000500
	ELSE	00000510
PS13	ADD 1 TO TOTAL NUMBER OF TIME-OUTS	00000520
21PS	SET COMPLETION CODE = RETRY	00000530
	SET ERROR CODE = TIME-OUT	00000540
	ENDIF	00000550
21PS	IF COMPLETION CODE = DATA RECEIVED THEN	00000560
23PS	GET PSC RESPONSE PROCESSING BUFFER	00000570
	IF ONE RECEIVED THEN	00000580
22PS	PLACE DATA INPUT INTO RESPONSE PROC. BUFFER	00000590
23PS	PLACE BUFFER ADRS. ON PSC RESPONSE QUEUE	00000600
24PS	IF QUEUE FULL THEN	00000610
11PS	PANIC!	00000620
	ELSE	00000630
30PS	POST RESPONSE PROCESSING EVENT	00000640
	ENDIF	00000650
	ELSE	00000660
28PS	RESET PSC RESPONSE OVERFLOW EVENT	00000670
	ENDIF	00000680
	ELSE	00000690
	IF ERROR CODE SET THEN	00000700
10PS	IF DUMP FLAG SET THEN	00000710
02AL	SET ALARM MESSAGE BUFFER	00000720
	IF ONE RECEIVED THEN	00000730
	PLACE ERROR CODE INTO BUFFER	00000740
	COPY LAST PSC MSG. RECEIVED INTO BUFFER	00000750
01AL	PLACE BUFFER ADRS. ON ALARM MSG. QUEUE	00000760
	IF ALARM MSG. QUEUE FULL THEN	00000770
11PS	PANIC!	00000780
	(ELSE)	00000790
	ENDIF	00000800
	ELSE	00000810
11PS	PANIC!	00000820
		00000830
		00000840
		00000850
		00000860
		00000870
		00000880
		00000890
		00000900
		00000910
		00000920
		00000930
		00000940

```

                ENDIF
                (ELSE)
                ENDIF
                (ELSE)
                ENDIF
            ENDIF
29PS      RESET PSC INPUT/OJTPUT EVENT
21PS      POST COMPLETION CODE TO TIME-OUT TASK
          ENDDO
          EXIT

```

PSC POLL TASK

- MODULE HISTORY -

```

PROJECT:      DAS                75-01709
SJB-SYSTEM:   PSC COMMUNICATIONS
MODULE:        PSPOL             54

```

- MODULE ABSTRACT -

THIS TASK DELAYS FOR A PERIOD OF TIME THEN PLACES A PSC IDLE MESSAGE INTO THE PSC OUTPUT QUEUE FOR TRANSMISSION TO THE PSC.

```

00UNTIL CPU STOPPED
08PS      DOWHILE IDLE MESSAGE COUNT LT. OR EQ. ZERO
          DELAY FOR ?? MS.
          ENDDO
08PS      DOWHILE IDLE MESSAGE COUNT > ZERO
          DELAY FOR ?? MS.
08PS      SUBTRACT 1 FROM IDLE MESSAGE COUNT
          ENDDO
01PS      GET PSC OUTPUT BUFFER
          IF ONE RECEIVED THEN
06PS          SET BUFFER FOR IDLE MESSAGE
02PS          PLACE BUFFER ADRS. ON PSC OJTPUT QUEUE
02PS          POST PSC R/JN EVENT
          IF QUEUE FULL THEN
11PS              PANIC!
                (ELSE)
                ENDIF
          ELSE
11PS              PANIC!
          ENDIF
          ENDDO
          EXIT

```

PSC PANIC PROCESSOR

- MODULE HISTORY -

```

PROJECT:      DAS                75-01709
SJB-SYSTEM:   PSC COMMUNICATIONS
MODULE:        PSPAN             55

```

```

00000950
00000960
00000970
00000980
00000990
00001000
00001010
00001020
00001030
00001040
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170

```

- MODULE ABSTRACT -

THIS SUBROUTINE IS CALLED FROM SEVERAL PSC
MODULES TO SEND A PANIC CODE TO THE ALARM MESSAGE SUB-SYSTEM
TO PRINT A FATAL ERROR MESSAGE AND TO SUSPEND EXECUTION OF
THE PSC SUB-SYSTEM.

```

11PS  GET INPJTS
07PS  SET COMMUNICATIONS LINK = CLOSED
01AL  PUT PANIC CODE INPUT INTO ALARM MSG. QUEJE
07AL
08AL
DDUNTIL NEW IPL
      WAIT FOR IMPOSSIBLE EVENT
ENDDD
RETURN TO CALLER (NOT REALLY)

```

PSC ALARM MESSAGE PROCESSING

- MODULE HISTORY -

```

PROJECT:      DAS              75-01709
SUB-SYSTEM:   PSC COMMUNICATIONS
MODULE:       PSA0             56

```

- MODULE ABSTRACT -

THIS SUBROUTINE PROCESSES ALL PSC RESPONSES OTHER THAN
A TRANSFER COMPLETE. IT UPDATES THE APPROPRIATE MESSAGE
COUNTER AND PRINTS THE APPROPRIATE MESSAGE AT THE
LOG DEVICES.

```

26PS  GET INPUTS
23PS  SET PSC RESPONSE CODE
09PS  GET RESPONSE COUNT TABLE
DDUNTIL RESPONSE CODE FOUND IN TABLE OR END OF TABLE
09PS  INDEX TO FIRST/NEXT TABLE ENTRY
23PS  IF RESPONSE CODE = TABLE ENTRY THEN
09PS
09PS      ADD 1 TO RESPONSE CODE COUNT
      (ELSE)
      ENDIF
ENDDD
IF END OF TABLE FOUND THEN
09PS  ADD 1 TO UNKNOWN RESPONSE COUNT
      (ELSE)
      ENDIF
02AL  GET AN ALARM MESSAGE BUFFER
IF ONE RECEIVED THEN
26SC  IF DATA RECEIVED WITHIN RANGE THEN
25SC
23PS  CASEENTRY (RESPONSE CODE)
      CASE RESPONSE CODE = SORTATION LINE FULL
03AL  BUILD BUFFER FOR "SORTATION LINE FULL" MSG.
      CASE RESPONSE CODE = SORTATION TRANSFER FAILURE
03AL  BUILD BUFFER FOR "SORTATION TRANSFER FAILURE" MSG.
04PM  ADD 1 TO TRANSFER FAILURES AT THIS LINE
      CASE RESPONSE CODE = MISSORT
26SC  IF SORT DESTINATION = ZERO THEN
03AL  BUILD BUFFER FOR "TRANSACTION LOST" MSG.
      ELSE
03AL  BUILD BUFFER FOR "MISSORT" MSG.
      ENDIF

```



```

-----CASE RESPONSE CODE = SUCCESSFUL TRANSFER-----00000570
03AL      BUILD BUFFER FOR "OUT-OF-RANGE" MSG.          00000580
          CASE RESPONSE CODE = OTHER                  00000590
10PS     IF JUMP INPJT FLAG SET THEN                  00000600
03AL      BUILD BUFFER FOR PRINTING THE DATA RECEIVED 00000610
          ELSE                                         00000620
02AL     RETURN ALARM MESSAGE BUFFER TO FREE POOL    00000630
          IF FREE POOL FULL THEN                      00000640
11PS     PANIC!                                       00000650
          (ELSE)                                       00000660
          ENDIF                                        00000670
          ENDIF                                        00000680
          ENDCASE--                                    00000690
          ELSE                                         00000700
03AL      BUILD BUFFER FOR "OUT-OF-RANGE" MSG.          00000710
          ENDIF                                        00000720
          IF ALARM MESSAGE BUFFER NOT RETURNED TO FREE POOL THEN 00000730
01AL      PLACE BUFFER ADDR. IN ALARM MESSAGE QUEUE    00000740
          IF ALARM MESSAGE QUEUE FULL THEN            00000750
11PS     PANIC!                                       00000760
          (ELSE)                                       00000770
          ENDIF                                        00000780
          (ELSE)                                       00000790
          ENDIF                                        00000800
          ELSE                                         00000810
11PS     PANIC!                                       00000820
          ENDIF                                        00000830
26PS     RETURN TO CALLER                              00000840
-----00000010
-----00000020
-----00000030
-----00000040
-----00000050
PSC RESPONSE TIME-OUT TASK                          00000060
-----00000070
-----00000080
-----00000090
          - MODULE HISTORY -                          00000100
-----00000110
          PROJECT:      DAS                            73-01709 00000120
          SJB-SYSTEM:   PSC COMMUNICATIONS            00000130
          MODULE:       PSTO                           52      00000140
-----00000150
          - MODULE ABSTRACT -                         00000160
-----00000170
          THIS TASK IS THE PSC RESPONSE TIME-OUT TASK. IF THE PSC 00000180
          HAS NOT RESPONDED WHEN THIS TASK TIMES OUT, THE INPJT 00000190
          REQUEST OF TASK "PSIO" IS TERMINATED. WHEN A COMPLETION 00000200
          CODE IS SET BY TASK "PSIO", THIS TASK WILL POST THAT CODE 00000210
          BACK TO THE PSC CONTROL TASK "PSOC".          00000220
          RESET TIME-OUT COUNT                          00000230
21PS     WHILE COMPLETION CODE NOT POSTED AND TIME-OUT COUNT NOT ZERO 00000240
          DELAY FOR ?? MS.                             00000250
          SUBTRACT 1 FROM THE TIME-OUT COUNT           00000260
          ENDDO                                         00000270
          IF COMPLETION CODE NOT POSTED THEN          00000280
21PS     KILL THE PSC READ FUNCTION                   00000290
21PS     WAIT FOR COMPLETION CODE TO BE POSTED      00000300
          (ELSE)                                       00000310
          ENDIF                                        00000320
21PS     POST COMPLETION CODE TO PSC CONTROL TASK "PSOC" 00000330
17PS     EXIT                                         00000340
          -----00000350
          -----00000360
          -----00000370
          -----00000380
          -----00000390

```

PSC OUTPUT MESSAGE FORMATTING

- MODULE HISTORY -

PROJECT: DAS 75-01709
 SUB-SYSTEM: PSC COMMUNICATIONS
 MODULE: PSMF 50

- MODULE ABSTRACT -

THIS SUBROUTINE FORMATES THE PSC DJTPT MESSAGE INTO
 A FORM THAT CAN BE TRANSMITTED OVER THE TTY LINE TO THE
 PSC. IF THE INPJT MESSAGE TYPE IS NOT ONE OF THE VALID
 TYPES, THIS ROUTINE WILL RETURN TO THE CALLER WITH
 AN ERROR CODE.

```

16PS GET INPUTS 03000270
16PS SET RETURN CODE = ERROR 03000280
12PS PLACE THE CURRENT DJTPT SEQUENCE BYTE IN DJTPT BUFFER 03000290
16PS CASEENTRY (MESSAGE CODE) 03000300
    CASE MESSAGE CODE = IDLE 03000310
18PS     BUILD IDLE MESSAGE IN DJTPT BUFFER 03000320
16PS     SET RETURN CODE = SUCCESSFUL 03000330
    CASE MESSAGE CODE = OPEN LINK 03000340
19PS     BUILD OPEN LINK MESSAGE IN DJTPT BUFFER 03000350
16PS     SET RETURN CODE = SUCCESSFUL 03000360
    CASE MESSAGE CODE = INDUCT PACKAGE 03000370
20PS     BUILD INDUCT PACKAGE MESSAGE IN DJTPT BUFFER 03000380
    IF NO CONVERSION ERRORS FOUND AND DATA IN RANGE THEN 03000390
16PS     SET RETURN CODE = SUCCESSFUL 03000400
    (ELSE) 03000410
    ENDIF 03000420
ENDCASE 03000430
IF RETURN CODE = SUCCESSFUL THEN 03000440
27PS     COMPUTE MESSAGE CHECKSUM 03000450
18PS     PLACE CHECKSUM IN DJTPT BUFFER 03000460
19PS 03000470
20PS 03000480
(ELSE) 03000490
ENDIF 03000500
16PS RETURN TO CALLER 03000510

```

INDUCTION SCANNER #2 TASK

- MODULE HISTORY -

PROJECT: DAS 75-01709
 SUB-SYSTEM: SCANNER INPUT
 MODULE: SCSR2 10

03000010
 03000020
 03000030
 03000040
 03000050
 03000060
 03000070
 03000080
 03000090
 03000100
 03000110
 03000120
 03000130
 03000140
 03000150
 03000160
 03000170
 03000180
 03000190
 03000200
 03000210
 03000220
 03000230
 03000240
 03000250
 03000260
 03000270
 03000280
 03000290
 03000300
 03000310
 03000320
 03000330
 03000340
 03000350
 03000360
 03000370
 03000380
 03000390
 03000400
 03000410
 03000420
 03000430
 03000440
 03000450
 03000460
 03000470
 03000480
 03000490
 03000500
 03000510
 03000010
 03000020
 03000030
 03000040
 03000050
 03000060
 03000070
 03000080
 03000090
 03000100
 03000110
 03000120
 03000130
 03000140
 03000150
 03000160
 03000170
 03000180
 03000190

- MODULE ABSTRACT -

```

THIS TASK RUNS UNTIL THE SCANNER STOP FLAG IS SET (-1). IT
FIRST READS A TRANSACTION NUMBER FROM INDUCTION SCANNER #2, THEN
CONVERTS THE NUMBER TO BINARY. AN ERROR FLAG IS SENT TO THE PACK-
AGE PROCESSING TASK (SCIND) TO INDICATE THE STATUS OF THE READ.
THE SCAN INFORMATION IS PUT ON THE INDUCTION QUEUE AND THE PACK-
AGE PROCESSING TASK EVENT IS POSTED.

```

	ENQUEUE THE SCANNER	00000200
		00000210
		00000220
		00000230
		00000240
		00000250
		00000260
		00000270
		00000280
		00000290
		00000300
01SC	DO UNTIL THE SCANNER STOP FLAG IS SET	00000310
04SC	READ THE TRANSACTION NUMBER (READTEXT)	00000320
	IF THE BUFFER IS NOT COMPLETELY FULL, THEN	00000330
20SC	GET AN INDUCTION QUEUE BUFFER	00000340
20SC	IF ONE RECEIVED, THEN	00000350
06SC	PUT IN SCANNER ID	00000360
	CONVERT TRANSACTION NUMBER FROM EBCDIC TO BINARY	00000370
11SC	IF ANY ERRORS (CONVERSION, < 5-DIGITS), THEN	00000380
	SET TRANSACTION NUMBER TO ZERO AND PUT IN BUFFER	00000390
14SC	SET STOP ON NOREAD FLAG FROM SYSCOM	00000400
11SC	MOVE NO READ FLAG TO ERROR FLAG	00000410
	ELSE	00000420
	SET ERROR FLAG = -1	00000430
	ENDIF	00000440
11SC	PUT ERROR FLAG IN BUFFER	00000450
21SC	PUT ADDRESS OF BUFFER ON INDUCTION QUEUE	00000460
21SC	IF THERE WAS NO ROOM, THEN	00000470
01SC	SET SCANNER STOP FLAG	00000480
07AL	PANIC	00000490
	ELSE	00000500
18SC	POST PACKAGE PROCESSING TASK EVENT	00000510
	ENDIF	00000520
	ELSE	00000530
01SC	SET SCANNER STOP FLAG	00000540
07AL	PANIC	00000550
	ENDIF	00000560
	(ELSE)	00000570
	ENDIF	00000580
	ENDDO	00000590
	DEQUEUE THE SCANNER	00000600
	ENDTASK	00000610

INDUCTION SCANNER #1 TASK

- MODULE HISTORY -

```

PROJECT:      DAS          75-01709
SUB-SYSTEM:   SCANNER INPT

```

- MODULE ABSTRACT -

```

THIS TASK RUNS UNTIL THE SCANNER STOP FLAG IS SET (-1). IT
FIRST READS A TRANSACTION NUMBER FROM INDUCTION SCANNER #1, THEN
CONVERTS THE NUMBER TO BINARY. AN ERROR FLAG IS SENT TO THE PACK-
AGE PROCESSING TASK (SCIND) TO INDICATE THE STATUS OF THE READ.
THE SCAN INFORMATION IS PUT ON THE INDUCTION QUEUE AND THE PACK-
AGE PROCESSING TASK EVENT IS POSTED.

```

```

-----
ENQUEUE THE SCANNER ----- 00000270
01SC DD UNTIL THE SCANNER STOP FLAG IS SET 00000280
04SC READ THE TRANSACTION NUMBER (READTEXT) 00000290
      IF THE BUFFER IS NOT COMPLETELY FULL THEN 00000300
20SC GET AN INDUCTION QUEUE BUFFER 00000310
20SC IF ONE RECEIVED, THEN 00000320
06SC PUT IN SCANNER ID 00000330
      CONVERT TRANSACTION NUMBER FROM EBCDIC TO BINARY 00000340
11SC IF ANY ERRORS (CONVERSION, < 5-DIGITS), THEN 00000350
      SET TRANSACTION NUMBER TO ZERO AND PUT IN BUFFER 00000360
14SC SET STOP ON NOREAD FLAG FROM SYSCOM 00000370
11SC MOVE NO READ FLAG TO ERROR FLAG 00000380
      ELSE 00000390
      SET ERROR FLAG = -1 00000400
      ENDIF 00000410
11SC PUT ERROR FLAG IN BUFFER 00000420
21SC PUT ADDRESS OF BUFFER ON INDUCTION QUEUE 00000430
21SC IF THERE WAS NO ROOM, THEN 00000440
01SC SET SCANNER STOP FLAG 00000450
07AL PANIC 00000460
      ELSE 00000470
18SC POST PACKAGE PROCESSING TASK EVENT 00000480
      ENDIF 00000490
      ELSE 00000500
01SC SET SCANNER STOP FLAG 00000510
07AL PANIC 00000520
      ENDIF 00000530
      (ELSE) 00000540
      ENDIF 00000550
      END DD 00000560
DEQUEUE THE SCANNER ----- 00000570
ENDTASK 00000580
----- 00000590

```

SCANNER INITIALIZATION PROGRAM

- MODULE HISTORY -

```

PROJECT:      OAS              73-01709
SUB-SYSTEM:   SCANNER INPUT
MODULE:       SCINIT          5

```

- MODULE ABSTRACT -

```

THIS PROGRAM WILL BE LOADED INTO MEMORY BY THE INITIAL SYSTEM PROGRAM. ONCE LOADED IT WILL SET THE SCANNER STOP FLAG TO NOSTOP. AFTER ATTACHING ALL OF THE SCANNER ROUTINES, IT WILL GO INTO A LOOP AND POST THE PACKAGE PROCESSING TASK EVENT AND THE STATUS UPDATE TASK EVENT ONCE EVERY 5 SECONDS. IF THE SCANNER STOP FLAG IS EVER SET, IT WILL DROP OUT OF THE LOOP AND WAIT FOR AN EVENT THAT WILL NEVER BE POSTED JUST TO KEEP IT IN MEMORY.

```

```

02SC INITIALIZATION ----- 00000310
01SC INITIALIZE SCANNER STOP FLAG TO NOSTOP (0) 00000320
39FM GET NUMBER OF SCANNERS AND TYPES FROM SYSCOM 00000330
      ATTACH THOSE TASKS 00000340
      ATTACH SCTSU 00000350
      ATTACH SCINDT 00000360
      DD UNTIL THE SCANNER STOP FLAG IS SET 00000370
      DELAY 5 SECONDS 00000380
18SC POST THE PACKAGE PROCESSING TASK EVENT 00000390
----- 00000400

```



```

09FM GET A FILE MANAGER INPUT BUFFER 00000680
09FM IF ONE RECEIVED THEN 00000690
10FM PUT IN CHANGE STATUS CALL CODE 00000700
17FM PUT IN DATA 00000710
04FM 00000720
08FM CALL THE FILE MANAGER 00000730
17FM GET THE RETURN CODE 00000740
11FM IF RETURN CODE NOT = SUCCESSFUL THEN 00000750
11FM IF RETURN CODE = RECORD NOT FOUND THEN 00000760
02AL GET AN ALARM MESSAGE BUFFER 00000770
02AL IF ONE RECEIVED THEN 00000780
03AL PUT IN MESSAGE # 00000790
PUT IN TRANSACTION # 00000800
PUT IN SCANNER ID 00000810
01AL PUT ADDRESS OF BUFFER ON ALARM QUEUE 00000820
01AL IF ALREADY FULL THEN 00000830
08AL SET ERROR CODE = FATAL, ALARM QJE FULL 00000840
(ELSE) 00000850
ENDIF 00000860
ELSE 00000870
08AL SET ERROR CODE = FATAL, ALARM BUFQUE EMPTY 00000880
ENDIF 00000890
ELSE 00000900
08AL SET ERROR CODE = FATAL, BAD RTURN FROM FM 00000910
ENDIF 00000920
(ELSE) 00000930
ENDIF 00000940
08FM RETURN FILE MANAGER BUFFER 00000950
08FM IF NO ROOM THEN 00000960
08AL SET ERROR CODE = FATAL, FM BUFFER QJE FULL 00000970
(ELSE) 00000980
ENDIF 00000990
ELSE 00001000
08AL SET ERROR CODE = FATAL, NO FILE MANAGER BUFFER 00001010
ENDIF 00001020
05SC RETURN STATUS UPDATE BUFFER 00001030
05SC IF THERE IS NO ROOM THEN 00001040
08AL SET ERROR CODE = FATAL, STATUS UPDATE BUFQUE FULL 00001050
(ELSE) 00001060
ENDIF 00001070
IF ERROR CODE = FATAL THEN 00001080
08AL USE ERROR CODE FOR PANIC CODE 00001090
01SC SET SCANNER STOP FLAG 00001100
07AL PANIC 00001110
(ELSE) 00001120
ENDIF 00001130
ENDDO 00001140
ENDTASK 00001150

```

INDUCTION SCANNER TASK

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
PROJECT: DAS 75-01709 00000110
00000120
SUB-SYSTEM: SCANNER INPUT 00000130
00000140
MODULE: SCIND 13 00000150
00000160
00000170
- MODULE ABSTRACT - 00000180
00000190
00000200

```

```

THIS TASK IS ATTACHED BY THE SCANNER INITIALIZATION PROGRAM. 00000210
IT THEN WAITS ON THE PACKAGE PROCESSING EVENT TO BE POSTED. 00000220
WHEN POSTED IT WILL GET AN ENTRY FROM THE INDUCTION QUEUE AND 00000230
GET THE DATA FROM THE BUFFER. THE DATA WILL THEN BE PROCESSED. 00000240

```

		03000250
02SC	INITIALIZATION	03000260
01SC	DO UNTIL THE SCANNER STOP FLAG IS SET	03000270
	RESET FLAGS (STATJS,PSC)	03000280
	SET MSG# = 3	03000290
	SET PANIC CODE = 0	03000300
09SC	SET DESTINATION = RECIRCULATION	03000310
21SC	DO UNTIL ELEMENT RECEIVED FROM INDUCTION QUEUE	03000320
18SC	WAIT FOR PACKAGE PROCESSING TASK EVENT	03000330
21SC	GET AN ELEMENT FROM THE QUEUE	03000340
21SC	IF NONE RECEIVED, THEN	03000350
18SC	RESET EVENT	03000360
	(ELSE)	03000370
	ENDIF	03000380
	ENDDO	03000390
20SC	GET THE DATA FROM THE BUFFER	03000400
20SC	PUT THE BUFFER ON THE BUFFER QUEUE	03000410
20SC	IF THERE WAS ROOM, THEN	03000420
07PS	GET THE PSC COMMUNICATIONS FLAG	03000430
07PS	IF THE LINK IS OPEN, THEN	03000440
15SC	IF THE SLOW DOWN FLAG IS SET, THEN	03000450
16SC	SET UP TO USE THE ALTERNATE QUEUES	03000460
17SC		03000470
	ELSE	03000480
05SC	SET UP TO USE THE PRIMARY QUEUES	03000490
08SC		03000500
	ENDIF	03000510
12SC	GET THE ERROR FLAG	03000520
12SC	IF THE ERROR FLAG = -1 THEN	03000530
13SC	IF TRANSACTION# NOT = 0 THEN	03000540
23SC		03000550
03FM	COMPUTE RELATIVE RECORD NUMBER	03000560
03FM	READ RECORD FROM TRANSACTION DIRECTORY	03000570
	IF THERE WERE NO READ ERRORS THEN	03000580
	SET PSC FLAG	03000590
03FM	COMPUTE INDEX TO BUFFER	03000600
03FM	GET RELATIVE TRANSACTION NUMBER	03000610
23SC	IF ON FILE THEN	03000620
03FM	GET RELATIVE BATCH AND DESTINATION	03000630
02FM	GET BATCH STATJS FROM RESIDENT BATCH TABLE	03000640
23SC		03000650
02FM	IF STATJS = PENDING, ACTIVE, OR CMPLTE THEN	03000660
04FM	COMPUTE RELATIVE RCRD FOR THE TRANS FILE	03000670
04FM	READ RECORD FROM TRANSACTION FILE	03000680
	IF THERE WERE NO ERRORS, THEN	03000690
04FM	COMPUTE INDEX TO BUFFER	03000700
04FM	GET THE TRANSACTION NUMBER	03000710
23SC	IF THIS TRANS# EQ OUR TRANS#, THEN	03000720
	SET STATUS FLAG	03000730
23SC	IF BATCH STATJS = ACTIVE THEN	03000740
04FM	PICK UP DESTINATION	03000750
	ELSE	03000760
03AL	SET MSG# = TRANS OUT OF BATCH	03000770
23SC	IF THE BATCH IS COMPLETE, THEN	03000780
13SC	SET DESTINATION = ERROR	03000790
	(ELSE)	03000800
	ENDIF	03000810
	ENDIF	03000820
	ELSE	03000830
03AL	SET MSG# = TRANS NOT ON FILE	03000840
13SC	SET DESTINATION = ERROR	03000850
	ENDIF	03000860
	ELSE	03000870
	SET PANIC CODE = READ ERROR	03000880
	ENDIF	03000890
	ELSE	03000900
03AL	SET MSG# = TRANS NOT ON FILE	03000910
13SC	SET DESTINATION = ERROR	03000920
	ENDIF	03000930
	ELSE	03000940
03AL	SET MSG# = TRANS NOT ON FILE	03000950
13SC	SET DESTINATION = ERROR	03000960
	ENDIF	03000970
	ELSE	03000980

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08AL          SET PANIC CODE = READ ERROR          00000990
              ENDIF                                00001000
              ELSE                                  00001010
                SET PSC FLAG                        00001020
13SC          GET ERROR DESTINATION                00001030
13SC          SET DESTINATION = ERROR              00001040
              ENDIF                                00001050
              ELSE                                  00001060
40FM          INCREMENT N0READ COUNT FOR THIS SCANNER IN SYSCOM 00001070
03AL          SET MSG# = SCANNER READ ERROR        00001080
12SC          IF ERROR FLAG = 0 THEN                00001090
              SET PSC FLAG                          00001100
              (ELSE)                                00001110
              ENDIF                                  00001120
            ENDIF                                  00001130
          (ELSE)                                    00001140
          ENDIF                                     00001150
        ELSE                                       00001160
08AL          SET PANIC CODE = INDUCTION BUFFER QUEUE FULL 00001170
            ENDIF                                  00001180
            IF PSC FLAG SET THEN                    00001190
01PS          GET A BUFFER FROM THE PSC BUFFER QUEUE 00001200
01PS          IF ONE RECEIVED THEN                  00001210
03PS          PUT IN CODE FOR INDUCTION PACKAGE     00001220
04PS          PUT IN TRANSACTION NUMBER             00001230
04PS          PUT IN SCANNER ID                     00001240
              IF THE DESTINATION IS NOT RECIRC OR THE ERR CHUTE, THEN 00001250
06FM          GET DEST. FROM SRT TRNSLTN TBL        00001260
06FM          IF TRNSLTN TBL DEST = 0, THEN          00001270
06FM          DEST = DIRECTORY DESTINATION           00001280
              ELSE                                  00001290
06FM          DEST = TRNSLTION TBL DEST              00001300
              ENDIF                                  00001310
            ELSE                                    00001320
              IF THE DESTINATION IS RECIRC, THEN    00001330
40FM          INCREMENT THE RECIRC CNT IN SYSCOM    00001340
              ENDIF                                  00001350
            ENDIF                                  00001360
04PS          PUT IN DESTINATION                     00001370
02PS          PUT ADDRESS OF BUFFER ON PSC OUTPUT QUEUE 00001380
02PS          IF THE PSC OUTPUT QUEUE IS FULL THEN  00001390
08AL          SET PANIC CODE = PSC OUTPUT QUEUE FULL 00001400
            ELSE                                    00001410
02PS          GET ADDRESS OF PSC EVENT               00001420
02PS          POST PSC EVENT                         00001430
            ENDIF                                  00001440
          ELSE                                       00001450
08AL          SET PANIC CODE = PSC BUFFER QUEUE EMPTY 00001460
            ENDIF                                  00001470
          (ELSE)                                    00001480
          ENDIF                                     00001490
        IF MESSAGE NUMBER NOT EQUAL TO ZERO THEN    00001500
02AL          GET AN ALARM MESSAGE OUTPUT BUFFER    00001510
03AL          IF ONE RECEIVED THEN                  00001520
01AL          PUT IN MESSAGE#, SCANNER ID, AND TRANS # 00001530
01AL          PUT ADDRESS OF BUFFER ON ALARM MESSAGE QUEUE 00001540
01AL          IF THE ALARM QUEUE IS FULL THEN      00001550
08AL          SET PANIC CODE = ALARM QUEUE FULL     00001560
              (ELSE)                                00001570
              ENDIF                                  00001580
            ELSE                                    00001590
08AL          SET PANIC CODE = ALARM BUFFER QUEUE EMPTY 00001600
            ENDIF                                  00001610
          (ELSE)                                    00001620
          ENDIF                                     00001630
        IF STATUS FLAG SET THEN                    00001640
          ENQUEUE THE QUEUE RESOURCE                00001650
          DOUNTIL A BUFFER IS RECEIVED FROM THE BUFFER QUEUE 00001660
          GET A BUFFER                               00001670
          (ELSE)                                    00001680
          IF ONE RECEIVED THEN                      00001690
04FM          PUT IN STATUS, SCANNER ID, TRANS#, AND DIRCTRY ENTRY 00001700
08SC          PUT ADDRESS OF BUFFER ON STATUS UPDATE QUEUE 00001710
08SC          IF THE QUEUE IS ALREADY FULL THEN    00001720

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08AL      SET PANIC CODE = STATUS UPDATE QUEUE FULL      00001730
          ELSE                                           00001740
19SC      POST STATUS UPDATE TASK EVENT                  00001750
          ENDIF                                          00001760
          ELSE                                           00001770
15SC      GET THE SLOW DOWN FLAG                        00001780
15SC      IF IT IS NOT SET, THEN                        00001790
16SC      SET JP TO USE ALTERNATE QUEUES                00001800
17SC      SET THE SLOW DOWN FLAG                        00001810
15SC      SET THE SLOW DOWN FLAG                        00001820
          ELSE                                           00001830
08AL      PANIC                                          00001840
          ENDIF                                          00001850
          ENDIF                                          00001860
          ENDDO                                         00001870
          DEQUEUE THE QUEUE RESOURCE                    00001880
          (ELSE)                                         00001890
          ENDIF                                          00001900
          IF PANIC CODE NOT EQUAL TO ZERO THEN          00001910
08AL      GET PANIC CODE                                00001920
01SC      SET SCANNER STOP FLAG                        00001930
07AL      PANIC                                         00001940
          (ELSE)                                         00001950
          ENDIF                                          00001960
          ENDDO                                         00001970
          ENDTASK                                       00001980

          00000010
          00000020
          00000030
          00000040
          00000050
          00000060
          00000070
          00000080
          00000090
          00000100
          00000110
          00000120
          PROJECT:      DAS          75-01709          00000130
          00000140
          SUB-SYSTEM:   SCANNER INPUT          00000150
          00000160
          MODULE:      SCST0          6          00000170
          00000180
          00000190
          00000200
          00000210
          00000220
          00000230
          THIS TASK RUNS UNTIL THE SCANNER STOP FLAG IS SET (-1). IT
          FIRST READS A TRANSACTION NUMBER FROM THE STOCK OUT SCANNER,
          THEN CONVERTS THE NUMBER TO BINARY. IF THERE WERE ANY ERRORS AN
          ALARM MESSAGE IS PRINTED TO TELL THE OPERATOR. OTHERWISE A
          STATUS UPDATE BUFFER IS OBTAINED. IN IT ARE PUT THE NEW STATUS
          AND TRANSACTION NUMBER. THE ADDRESS OF THE BUFFER IS THEN PUT
          ON THE STATUS UPDATE QUEUE AND THE STATUS UPDATE TASK EVENT IS
          POSTED. THEN THE PROCESS STARTS ALL OVER WAITING FOR ANOTHER READ.
          00000300
          00000310
          ENQUEUE THE STOCK OUT SCANNER                00000320
02SC      INITIALIZATION                                00000330
01SC      DOUNTIL SCANNER STOP FLAG SET                00000340
          RESET THE FATAL FLAG                          00000350
04SC      READ THE TRANSACTION NUMBER (READTEXT)        00000360
          IF WE RECEIVED 7 BYTES (6-CHAR AND LINEFEED) 00000370
          CONVERT FROM EBCDIC TO BINARY                 00000380
          (ELSE)                                         00000390
          ENDIF                                          00000400
04SC      IF NO ERRORS (CONVERSION = 0 < 6 DIGITS) THEN 00000410
05SC      GET A BUFFER FROM THE STATUS UPDATE BUFFER QUEUE 00000420
05SC      IF ONE RECEIVED THEN                          00000430
06SC      PUT IN SCANNER ID                             00000440
          PUT IN TRANSACTION NUMBER.                    00000450
04FM      PUT IN STATUS                                 00000460
    
```

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085C      PUT ADDRESS OF BUFFER ON STATUS UPDATE QUEUE      00000470
085C      IF STATUS UPDATE QUEUE ALREADY FULL THEN          00000480
           SET FATAL FLAG                                  00000490
08AL      SET PANIC CODE = STATUS UPDATE QUEUE FULL        00000500
           ELSE                                            00000510
175C      POST THE STATUS UPDATE TASK EVENT                 00000520
           ENDIF                                          00000530
           ELSE                                           00000540
           SET FATAL FLAG                                  00000550
08AL      SET PANIC CODE = STATUS UPDATE QUEUE EMPTY        00000560
           ENDIF                                          00000570
           ELSE                                           00000580
           IF THE READ BUFFER WAS NOT FULL, THEN           00000590
02AL      SET A BUFFER FROM THE ALARM MESSAGE BUFFER QUEUE 00000600
02AL      IF ONE RECEIVED THEN                             00000610
03AL      PUT IN MESSAGE NUMBER = NREAD                   00000620
055C      PUT IN SCANNER ID                               00000630
01AL      PUT ADDRESS OF BUFFER ON ALARM QUEUE             00000640
01AL      IF THERE IS NO ROOM THEN                         00000650
           SET FATAL FLAG                                  00000660
08AL      SET PANIC CODE = ALRM QUE FULL                   00000670
           (ELSE)                                          00000680
           ENDIF                                          00000690
           ELSE                                           00000700
           SET FATAL FLAG                                  00000710
08AL      SET PANIC CODE = ALRM BUFQUE EMPTY               00000720
           ENDIF                                          00000730
           (ELSE)                                          00000740
           ENDIF                                          00000750
           ENDIF                                          00000760
           IF THE FATAL FLAG IS SET THEN                   00000770
015C      SET SCANNER STOP FLAG                            00000780
08AL      GET PANIC CODE                                    00000790
07AL      PANIC                                            00000800
           (ELSE)                                          00000810
           ENDIF                                          00000820
ENDDDO                                          00000830
DEQUEUE THE SCANNER                                  00000840
ENDTASK                                             00000850

00000010
00000020
00000030
00000040
00000050
EXCESS REPACK SCANNER TASK                          00000060
00000070
00000080
00000090
- MODULE HISTORY -                                  00000100
00000110
PROJECT:      DAS                    75-01709        00000120
00000130
SUB-SYSTEM:   SCANNER INPUT          00000140
00000150
MODULE:       SCXRP                   8              00000160
00000170
00000180
- MODULE ABSTRACT -                                00000190
00000200
00000210
00000220
THIS TASK RJNS UNTIL THE SCANNER STOP FLAG IS SET (-1). IT 00000230
FIRST READS A TRANSACTION NUMBER FROM THE EXCESS REPACK SCANNER, 00000240
THEN CONVERTS THE NUMBER TO BINARY. IF THERE WERE ANY ERRORS AN 00000250
ALARM MESSAGE IS PRINTED TO TELL THE OPERATOR. OTHERWISE A 00000260
STATUS UPDATE BUFFER IS OBTAINED. IN IT ARE PUT THE NEW STATUS 00000270
AND TRANSACTION NUMBER. THE ADDRESS OF THE BUFFER IS THEN PUT 00000280
ON THE STATUS UPDATE QUEUE AND THE STATUS UPDATE TASK EVENT IS 00000290
POSTED. THEN THE PROCESS STARTS ALL OVER WAITING FOR ANOTHER READ. 00000300

```

	ENQUEUE THE EXCESS REPACK SCANNER	00000310
		00000320
02SC	INITIALIZATION	00000330
01SC	DDUNTIL SCANNER STOP FLAG SET	00000340
	RESET THE FATAL FLAG	00000350
04SC	READ THE TRANSACTION NUMBER (READTEXT)	00000360
	IF WE RECEIVED 7 BYTES (6-CHAR AND LINEFEED)	00000370
	CONVERT FROM EBCDIC TO BINARY	00000380
	(ELSE)	00000390
	ENDIF	00000400
04SC	IF NO ERRORS (CONVERSION,=0, <6 DIGITS) THEN	00000410
05SC	GET A BUFFER FROM THE STATUS JDATE BUFFER QUEUE	00000420
05SC	IF ONE RECEIVED THEN	00000430
06SC	PUT IN SCANNER ID	00000440
	PUT IN TRANSACTION NUMBER	00000450
04FM	PUT IN STATJS	00000460
08SC	PUT ADDRESS OF BUFFER ON STATUS UPDATE QUEUE	00000470
08SC	IF STATJS JDATE QUEUE ALREADY FULL THEN	00000480
	SET FATAL FLAG	00000490
08AL	SET PANIC CODE = STATUS JDATE QUEUE FULL	00000500
	ELSE	00000510
19SC	POST THE STATUS JDATE TASK EVENT	00000520
	ENDIF	00000530
	ELSE	00000540
	SET FATAL FLAG	00000550
08AL	SET PANIC CODE = STATUS UPDT BUFQUE EMPTY	00000560
	ENDIF	00000570
	ELSE	00000580
	IF THE READ BUFFER WAS NOT FULL, THEN	00000590
02AL	GET A BUFFER FROM THE ALARM MESSAGE BUFFER QUEUE	00000600
02AL	IF ONE RECEIVED THEN	00000610
03AL	PUT IN MESSAGE NUMBER = NREAD	00000620
06SC	PUT IN SCANNER ID	00000630
01AL	PUT ADDRESS OF BUFFER ON ALARM QUEUE	00000640
01AL	IF THERE IS NO ROOM THEN	00000650
	SET FATAL FLAG	00000660
08AL	SET PANIC CODE = ALRM QUE FULL	00000670
	(ELSE)	00000680
	ENDIF	00000690
	ELSE	00000700
	SET FATAL FLAG	00000710
08AL	SET PANIC CODE = ALRM BUFQUE EMPTY	00000720
	ENDIF	00000730
	(ELSE)	00000740
	ENDIF	00000750
	ENDIF	00000760
	IF THE FATAL FLAG IS SET THEN	00000770
01SC	SET SCANNER STOP FLAG	00000780
08AL	GET PANIC CODE	00000790
07AL	PANIC	00000800
	(ELSE)	00000810
	ENDIF	00000820
	ENDDD	00000830
	DEQUEUE THE SCANNER	00000840
	ENDTASK	00000850
		00000010
		00000020
		00000030
		00000040
		00000050
	SHIPMENT SCANNER TASK	00000060
		00000070
		00000080
		00000090
	- MODULE HISTORY -	00000100
		00000110
		00000120
PROJECT:	DAS	73-01709
		00000130
		00000140
SUB-SYSTEM:	SCANNER INPUT	00000150
		00000160
MODULE:	SCSHP	7
		00000170
		00000180
		00000190

- MODJLE ABSTRACT -

		00000200
		00000210
		00000220
	THIS TASK RJNS UNTIL THE SCANNER STOP FLAG IS SET (-1). IT	00000230
	FIRST READS A TRANSACTION NUMBER FROM THE SHIPMENT SCANNER,	00000240
	THEN CONVERTS THE NUMBER TO BINARY. IF THERE WERE ANY ERRORS AN	00000250
	ALARM MESSAGE IS PRINTED TO TELL THE OPERATOR. OTHERWISE A	00000260
	STATUS UPDATE BUFFER IS OBTAINED. IN IT ARE PUT THE NEW STATUS	00000270
	AND TRANSACTION NUMBER. THE ADDRESS OF THE BUFFER IS THEN PUT	00000280
	ON THE STATUS UPDATE QUEUE AND THE STATUS UPDATE TASK EVENT IS	00000290
	POSTED. THEN THE PROCESS STARTS ALL OVER WAITING FOR ANOTHER READ.	00000300
		00000310
	ENQUEUE THE SHIPMENT SCANNER	00000320
02SC	INITIALIZATION	00000330
01SC	UNTIL SCANNER STOP FLAG SET	00000340
	RESET THE FATAL FLAG	00000350
04SC	READ THE TRANSACTION NUMBER (READTEXT)	00000360
	IF WE RECEIVED 7 BYTES (6-CHAR AND LINEFEED)	00000370
	CONVERT FROM EBCDIC TO BINARY	00000380
	(ELSE)	00000390
	ENDIF	00000400
04SC	IF NO ERRORS (CONVERSION=0, <6 DIGITS) THEN	00000410
05SC	GET A BUFFER FROM THE STATUS UPDATE BUFFER QUEUE	00000420
05SC	IF ONE RECEIVED THEN	00000430
06SC	PUT IN SCANNER ID	00000440
	PUT IN TRANSACTION NUMBER	00000450
04FM	PUT IN STATUS	00000460
08SC	PUT ADDRESS OF BUFFER ON STATUS UPDATE QUEUE	00000470
08SC	IF STATUS UPDATE QUEUE ALREADY FULL THEN	00000480
	SET FATAL FLAG	00000490
08AL	SET PANIC CODE = STATUS UPDATE QUEUE FULL	00000500
	ELSE	00000510
19SC	POST THE STATUS UPDATE TASK EVENT	00000520
	ENDIF	00000530
	ELSE	00000540
	SET FATAL FLAG	00000550
08AL	SET PANIC CODE = STATUS UPDATE QUEUE EMPTY	00000560
	ENDIF	00000570
	ELSE	00000580
	IF THE READ BUFFER WAS NOT FULL, THEN	00000590
02AL	GET A BUFFER FROM THE ALARM MESSAGE BUFFER QUEUE	00000600
02AL	IF ONE RECEIVED THEN	00000610
03AL	PUT IN MESSAGE NUMBER = NREAD	00000620
06SC	PUT IN SCANNER ID	00000630
01AL	PUT ADDRESS OF BUFFER ON ALARM QUEUE	00000640
01AL	IF THERE IS NO ROOM THEN	00000650
	SET FATAL FLAG	00000660
08AL	SET PANIC CODE = ALARM QUEUE FULL	00000670
	(ELSE)	00000680
	ENDIF	00000690
	ELSE	00000700
	SET FATAL FLAG	00000710
08AL	SET PANIC CODE = ALARM QUEUE EMPTY	00000720
	ENDIF	00000730
	(ELSE)	00000740
	ENDIF	00000750
	ENDIF	00000760
	IF THE FATAL FLAG IS SET THEN	00000770
01SC	SET SCANNER STOP FLAG	00000780
08AL	GET PANIC CODE	00000790
07AL	PANIC	00000800
	(ELSE)	00000810
	ENDIF	00000820
	ENDDO	00000830
	DEQUEUE THE SCANNER	00000840
	ENDTASK	00000850
		00000010
		00000020
		00000030
		00000040
		00000050

INDUCTION SCANNER #3 TASK

			0000060
			0000070
			0000080
		- MODULE HISTORY -	0000090
			0000100
			0000110
	PROJECT:	DAS	0000120
		75-01709	0000130
			0000140
	SUB-SYSTEM:	SCANNER INPUT	0000150
			0000160
			0000170
		- MODULE ABSTRACT -	0000180
			0000190
			0000200
		THIS TASK RUNS UNTIL THE SCANNER STOP FLAG IS SET (-1). IT	0000210
		FIRST READS A TRANSACTION NUMBER FROM INDUCTION SCANNER #3. THEN	0000220
		CONVERTS THE NUMBER TO BINARY. AN ERROR FLAG IS SENT TO THE PACK-	0000230
		AGE PROCESSING TASK (SCIND) TO INDICATE THE STATUS OF THE READ.	0000240
		THE SCAN INFORMATION IS PUT ON THE INDUCTION QUEUE AND THE PACK-	0000250
		AGE PROCESSING TASK EVENT IS POSTED.	0000260
			0000270
		ENQUEUE THE SCANNER	0000280
01SC	DO	UNTIL THE SCANNER STOP FLAG IS SET	0000290
04SC		READ THE TRANSACTION NUMBER (READTEXT)	0000300
		IF THE BUFFER IS NOT COMPLETELY FULL, THEN	0000310
20SC		GET AN INDUCTION QUEUE BUFFER	0000320
20SC		IF ONE RECEIVED, THEN	0000330
06SC		PUT IN SCANNER ID	0000340
		CONVERT TRANSACTION NUMBER FROM EBCDIC TO BINARY	0000350
11SC		IF ANY ERRORS (CONVERSION, < 5-DIGITS), THEN	0000360
		SET TRANSACTION NUMBER TO ZERO AND PUT IN BUFFER	0000370
14SC		GET STOP ON NOREAD FLAG FROM SYSCOM	0000380
11SC		MOVE NO READ FLAG TO ERROR FLAG	0000390
		ELSE	0000400
		SET ERROR FLAG = -1	0000410
		ENDIF	0000420
11SC		PUT ERROR FLAG IN BUFFER	0000430
21SC		PUT ADDRESS OF BUFFER ON INDUCTION QUEUE	0000440
21SC		IF THERE WAS NO ROOM, THEN	0000450
01SC		SET SCANNER STOP FLAG	0000460
07AL		PANIC	0000470
		ELSE	0000480
18SC		POST PACKAGE PROCESSING TASK EVENT	0000490
		ENDIF	0000500
		ELSE	0000510
01SC		SET SCANNER STOP FLAG	0000520
07AL		PANIC	0000530
		ENDIF	0000540
		(ELSE)	0000550
		ENDIF	0000560
		ENDDO	0000570
		DEQUEUE THE SCANNER	0000580
		ENDTASK	0000590
			0000010
			0000020
			0000030
			0000040
			0000050
		MIS-SORT REPORT	0000060
			0000070
			0000080
			0000090
		- MODULE HISTORY -	0000100
			0000110
			0000120
	PROJECT:	DAS	0000130
		75-01709	0000140
			0000150
	SUB-SYSTEM:	OPERATOR INTERFACE	0000160
			0000170
	MODULE:	QIMSR	0000180
		34	0000190

- MODULE ABSTRACT -

		00000200
		00000210
		00000220
	THIS MODULE REQUESTS A BATCH ID. THE FILE MANAGER IS	00000230
	CALLED TO READ THE BATCH FILE INFORMATION.	00000240
	IF THE BATCH INFORMATION IS FOUND, THE USER IS ASKED IF	00000250
	ALL ORDERS ARE TO BE LISTED. IF ALL ORDERS ARE NOT TO BE	00000260
	LISTED, A SPECIFIC ORDER ID IS REQUESTED, AND THE FILE	00000270
	MANAGER IS CALLED TO READ THE ORDER FILE INFORMATION.	00000280
	IF HARDCOPY IS REQUIRED, THEN THE TERMINAL IS DEQUEUED	00000290
	DEQUEUED AND THE SYSTEM PRINTER IS ENQUEUED. THE BATCH ID,	00000300
	HEADER AND STATUS ARE THEN OUTPUT TO THE ENQUEUED DEVICE.	00000310
	IF ALL ORDERS TO BE LISTED, THE FILE MANAGER IS CALLED TO	00000320
	SEQUENTIALLY READ EACH ORDER RECORD. FOR EACH ORDER RECORD,	00000330
	OR THE SPECIFICALLY REQUESTED ORDER RECORD, THE ORDER ID, ALL	00000340
	QUANTITIES AND ORDER PERCENT COMPLETE ARE OUTPUT.	00000350
	THE FILE MANAGER IS CALLED TO SEQUENTIALLY READ EACH TRAN-	00000360
	SACTION BLOCK. EACH TRANSACTION RECORD ASSOCIATED WITH THE	00000370
	ORDER IS EXAMINED. IF THE TRANSACTION STATUS IS A MIS-SORT,	00000380
	THE TRANSACTION INFORMATION IS OUTPUT TO THE ENQUEUED DEVICE.	00000390
	THE TERMINAL/PRINTER IS DEQUEUED BEFORE THE PROGRAM ENDS.	00000400
		00000410
		00000420
	ENQUEUE USER TERMINAL	00000430
0401	DISPLAY ON USER TERMINAL "MIS-SORT REPORT"	00000440
0301	ASK USER FOR BATCH ID	00000450
1201		00000460
4801	DEQUEUE AN IPM BUFFER (GET.IPM)	00000470
4901	DEQUEUE RECORD BUFFER (GET.SRB)	00000480
6001	INITIALIZE PANIC CODE	00000490
1601	ACCESS BATCH FILE (GET.BATCH)	00000500
1001	IF BATCH RECORD FOUND (OK SET), THEN	00000510
0301	! ASK USER IF ALL ORDERS TO BE LISTED	00000520
6001	! IF NOT ALL ORDERS, THEN	00000530
0301	! ASK USER FOR ORDER ID	00000540
1301	!	00000550
6001	! INITIALIZE PANIC CODE	00000560
1801	! ACCESS ORDER FILE (GET.ORDER)	00000570
1001	! IF ORDER RECORD FOUND (OK SET), THEN	00000580
1001	! SET SPECIFIC ORDER	00000590
	! (ELSE)	00000600
	! ENDF	00000610
	! ELSE	00000620
1001	! SET OK	00000630
1001	! CLEAR SPECIFIC ORDER	00000640
	! ENDF	00000650
1001	! IF ORDER FOUND (OK SET), THEN	00000660
1001	! ! SET ORDER SEQUENCE	00000670
1001	! ! SET EXCEPTIONS	00000680
0301	! ! ASK USER IF HARDCOPY OUTPUT IS REQUIRED	00000690
	! ! IF HARDCOPY REQUIRED, THEN	00000700
1001	! ! SET HARDCOPY	00000710
	! ! DEQUEUE USER TERMINAL	00000720
	! ! ENQUEUE PRINTER	00000730
	! ! ELSE	00000740
1001	! ! CLEAR HARDCOPY	00000750
	! ! ENDF	00000760
6001	! ! INITIALIZE DATE, TIME, PAGE COUNT AND LINE COUNT	00000770
6101	! ! PRINT/DISPLAY BATCH HEADINGS	00000780
6101	! ! PRINT/DISPLAY BATCH INFORMATION	00000790
6101	! ! PRINT/DISPLAY ORDER HEADINGS	00000800
6001	! ! INITIALIZE PANIC CODE	00000810
1301	! ! INITIALIZE ORDER COUNT	00000820
1301	! ! DO UNTIL ORDER COUNT GREATER THAN MAXIMUM	00000830
1001	! ! ! IF SPECIFIC ORDER SET, THEN	00000840
1301	! ! ! ! UPDATE ORDER COUNT BEYOND MAXIMUM	00000850
1201	! ! ! !	00000860
	! ! ! ! ELSE	00000870
1701	! ! ! ! FIND THE FIRST/NEXT ORDER IN THIS BATCH (FIND.ORDER)	00000880
	! ! ! ! ENDF	00000890

100I	!	!	!	!	!	!	IF ORDER RECORD FOUND (OK SET), THEN	00000900
130I	!	!	!	!	!	!	COMPUTE ORDER PERCENT COMPLETE	00000910
	!	!	!	!	!	!	SET LIST ORDER INFORMATION	00000920
610I	!	!	!	!	!	!	PRINT/DISPLAY ORDER INFORMATION AND PERCENT COMPLETE	00000930
	!	!	!	!	!	!	--- (FORMS.CTL) ---	00000940
100I	!	!	!	!	!	!	IF ABORT SET THEN	00000950
130I	!	!	!	!	!	!	SET ORDER COUNT GREATER THAN MAX.	00000960
	!	!	!	!	!	!	ELSE	00000970
	!	!	!	!	!	!	IF MIS-SORTS IN THIS ORDER THEN	00000980
	!	!	!	!	!	!	CLEAR LIST ORDER INFORMATION	00000990
610I	!	!	!	!	!	!	PRINT/DISPLAY TRANSACTION HEADINGS.	00001000
600I	!	!	!	!	!	!		00001010
60AL	!	!	!	!	!	!	INITIALIZE PANIC CODE	00001020
600I	!	!	!	!	!	!	INITIALIZE REPORT TRANSACTION COUNT	00001030
140I	!	!	!	!	!	!	INITIALIZE TRANS. CURRENT BLOCK NUMBER	00001040
140I	!	!	!	!	!	!	INITIALIZE TRANSACTION COUNT	00001050
140I	!	!	!	!	!	!	DOUNTIL TRANSACTION COUNT > MAXIMUM --	00001060
	!	!	!	!	!	!	-- OR BATCH NO LONGER ON FILE	00001070
190I	!	!	!	!	!	!	FIND THE FIRST/NEXT MIS-SORT RECORD	00001080
	!	!	!	!	!	!	--- (FIND.TRANS) ---	00001090
	!	!	!	!	!	!	IF BATCH STILL ON FILE THEN	00001100
100I	!	!	!	!	!	!	IF EXCEPTION RECORD FOUND THEN	00001110
	!	!	!	!	!	!	IF TRANS. STATUS = MIS-SORT THEN	00001120
20FM	!	!	!	!	!	!	SET UP TO READ THE TRANS. DIR.	00001130
08FM	!	!	!	!	!	!	CALL FILE MANAGER FOR DIR. REC.	00001140
20FM	!	!	!	!	!	!	IF RETJRN CODE = ERRDR THEN	00001150
20FM	!	!	!	!	!	!	IF RETURN CODE = NOT ON FILE	00001160
	!	!	!	!	!	!	PRINT ABORT MSG.	00001170
	!	!	!	!	!	!	ELSE	00001180
540I	!	!	!	!	!	!	PANIC	00001190
	!	!	!	!	!	!	ENDIF	00001200
	!	!	!	!	!	!	ELSE	00001210
03FM	!	!	!	!	!	!	SET MIS-SORT DEST. FROM REC.	00001220
	!	!	!	!	!	!	IF MIS-SORT DEST. = -1 THEN	00001230
	!	!	!	!	!	!	SET DJTPT = '?'	00001240
	!	!	!	!	!	!	ELSE	00001250
	!	!	!	!	!	!	PLACE DEST. IN OUTPJT	00001260
	!	!	!	!	!	!	ENDIF	00001270
	!	!	!	!	!	!	ENDIF	00001280
	!	!	!	!	!	!	ELSE	00001290
	!	!	!	!	!	!	SET DJTPT TO BLANKS	00001300
	!	!	!	!	!	!	ENDIF	00001310
	!	!	!	!	!	!	IF BATCH STILL ON FILE THEN	00001320
	!	!	!	!	!	!	SET LIST TRANSACTION INFO.	00001330
610I	!	!	!	!	!	!	PRINT/DISPLAY TRANS INFORMATION	00001340
	!	!	!	!	!	!	--- (FORMS.CTL) ---	00001350
	!	!	!	!	!	!	(ELSE)	00001360
	!	!	!	!	!	!	ENDIF	00001370
100I	!	!	!	!	!	!	IF ABORT SET THEN	00001380
140I	!	!	!	!	!	!	SET TRANS. COUNT > MAX	00001390
130I	!	!	!	!	!	!	SET ORDER COUNT > MAX.	00001400
	!	!	!	!	!	!	(ELSE)	00001410
	!	!	!	!	!	!	ENDIF	00001420
	!	!	!	!	!	!	(ELSE)	00001430
	!	!	!	!	!	!	ENDIF	00001440
	!	!	!	!	!	!	ELSE	00001450
040I	!	!	!	!	!	!	OUTPJT "BATCH NO LONGER ON FILE	00001460
	!	!	!	!	!	!	--- REPORT TERMINATED" ---	00001470
100I	!	!	!	!	!	!	SET ABORT	00001480
	!	!	!	!	!	!	ENDIF	00001490
	!	!	!	!	!	!	ENDDD	00001500
	!	!	!	!	!	!	CLEAR LIST TRANSACTION INFORMATION	00001510
	!	!	!	!	!	!	ELSE	00001520
020I	!	!	!	!	!	!	PRINT/DISPLAY "NO MIS-SORTS IN THIS ORDER"	00001530
	!	!	!	!	!	!	ENDIF	00001540
	!	!	!	!	!	!	IF HARD COPY SET THEN	00001550
	!	!	!	!	!	!	UPDATE LINE COUNT > MAX.	00001560
	!	!	!	!	!	!	ELSE	00001570
	!	!	!	!	!	!	MOVE REPORT TO TOP OF SCREEN	00001580
	!	!	!	!	!	!	ENDIF	00001590

```

! ! ! ! ENDFIF                                00001600
! ! ! ! ELSE                                  00001610
0401 ! ! ! ! PRINT/DISPLAY "BATCH NO LONGER ON FILE - RPT TERM" 00001620
1001 ! ! ! ! SET ABORT                          00001630
! ! ! ! IF HARD COPY NOT SET THEN            00001640
! ! ! !     MOVE REPORT TO TOP OF SCREEN    00001650
! ! ! ! (ELSE)                              00001660
! ! ! ! ENDFIF                              00001670
! ! ! ENDFIF                                00001680
! ! ENDDO                                    00001690
! ! IF HARD COPY AND EXCEPTIONS NOT SET THEN 00001700
! !     MOVE REPORT TO TOP OF SCREEN        00001710
! ! ELSE                                     00001720
! ! ENDFIF                                  00001730
! (ELSE)                                     00001740
! ENDFIF                                    00001750
ELSE                                         00001760
0201 PRINT/DISPLAY BATCH ID, "BATCH NOT ON FILE" 00001770
ENDIF                                       00001780
5101 RELEASE RECORD BUFFER (FREE.SR?)      00001790
5001 RELEASE THE IPM BUFFER (FREE.IPM)     00001800
DEQUEUE TERMINAL/PRINTER                   00001810
ENDPROG                                    00001820

FORMS.CTL SUBROUTINE                      00001830
00001840
00001850
00001850
1001 IF HARDCOPY SET, THEN                  00001870
6001 IF LINE COUNT GREATER THAN PRINTER PAGE, THEN 00001880
1001 SET JK                                00001890
ELSE                                        00001900
1001 CLEAR JK                              00001910
ENDIF                                       00001920
ELSE                                        00001930
6001 IF LINE COUNT GREATER THAN TERMINAL PAGE, THEN 00001940
1001 SET JK                                00001950
ELSE                                        00001960
1001 CLEAR JK                              00001970
ENDIF                                       00001980
ENDIF                                       00001990
1001 IF PAGE OVERFLOW (JK SET), THEN        00002000
6001 INCREMENT PAGE COUNT                  00002010
6101 PRINT/DISPLAY BATCH HEADINGS          00002020
6101 PRINT/DISPLAY BATCH INFORMATION       00002030
6101 PRINT/DISPLAY ORDER HEADINGS         00002040
6001 INITIALIZE LINE COUNT                00002050
1001 IF LIST TRANSACTION INFORMATION SET, THEN 00002060
6101 PRINT/DISPLAY ORDER INFORMATION       00002070
6101 PRINT/DISPLAY TRANSACTION HEADINGS   00002080
6001 INCREMENT LINE COUNT                 00002090
(ELSE)                                     00002100
ENDIF                                       00002110
(ELSE)                                     00002120
ENDIF                                       00002130
1001 IF LIST ORDER INFORMATION SET, THEN    00002140
6101 PRINT/DISPLAY ORDER INFORMATION       00002150
6001 INCREMENT LINE COUNT                 00002160
(ELSE)                                     00002170
ENDIF                                       00002180
1001 IF LIST TRANSACTION INFORMATION SET, THEN 00002190

6101 PRINT/DISPLAY TRANSACTION INFORMATION 00002200
6001 INCREMENT REPORT TRANSACTION COUNT    00002210
6001 INCREMENT LINE COUNT                  00002220
(ELSE)                                     00002230
ENDIF                                       00002240
IF REPORT CANCELLED AND HARD COPY SET THEN 00002250
PRINT REPORT CANCELLED MESSAGE            00002260
(ELSE)                                     00002270
ENDIF                                       00002280
RETURN                                     00002290

```


			00000010
			00000020
			00000030
			00000040
			00000050
			00000060
			00000070
			00000080
			00000090
			00000100
			00000110
			00000120
			00000130
			00000140
			00000150
			00000160
			00000170
			00000180
			00000190
			00000200
			00000210
			00000220
			00000230
			00000240
			00000250
			00000260
			00000270
			00000280
			00000290
040I	INITIALIZATION		00000300
030I	DISPLAY 'SYSTEM ERROR REPORT'		00000310
	ASK USER IF HARDCOPY IS REQUIRED		00000320
	IF YES, THEN		00000330
590I	SET A HARDCOPY FLAG		00000340
	ENQUEUE THE SYSTEM PRINTER		00000350
	EJECT TO TOP OF FORM		00000360
	ELSE		00000370
590I	CLEAR THE HARDCOPY FLAG		00000380
	ENQUEUE THE TERMINAL		00000390
	CLEAR THE SCREEN		00000400
	ENDIF		00000410
580I	PUT TIME AND DATE IN THE HEADER		00000420
580I	INITIALIZE PAGE COUNT AND PUT IN THE HEADER		00000430
580I	PRINT/DISPLAY HEADER #1		00000440
580I	PRINT/DISPLAY HEADER #2 AND HEADER #3		00000450
580I	PRINT/DISPLAY LOST, REJECTED, AND RECIRC COUNTS		00000460
39FM	GET NUMBER OF INDUCTS FROM SYSCOM		00000470
580I	CALCULATE CENTERING FOR SCANNER HEADERS AND COUNTS		00000480
580I	PRINT/DISPLAY HEADERS AND COUNTS		00000490
590I	IF HARDCOPY FLAG IS NOT SET, THEN		00000500
	MOVE THE REPORT TO THE TOP OF THE SCREEN		00000510
030I	ASK USER 'LIST MIS-SORT COUNTS ?'		00000520
	IF THE ANSWER IS YES, THEN		00000530
	CLEAR THE SCREEN		00000540
	INITIALIZE THE LINE COUNT		00000550
580I	INCREMENT THE PAGE NUMBER AND PUT IN THE HEADER		00000560
580I	PRINT THE HEADER #1		00000570
590I	CLEAR ABORT FLAG		00000580
	ELSE		00000590
590I	SET ABORT FLAG		00000600
	ENDIF		00000610
	ELSE		00000620
590I	CLEAR THE ABORT FLAG		00000630
	SKIP A FEW LINES		00000640
	ENDIF		00000650
590I	IF ABORT FLAG IS NOT SET, THEN		00000660
040I	PRINT/DISPLAY 'M I S - S O R T S'		00000670
39FM	GET NUMBER OF SORT LINES FROM SYSCOM		00000680
	CALL ROUTINE TO PRINT MIS-SORT COUNTS		00000690
	(ELSE)		00000700
	ENDIF		00000710
590I	IF HARDCOPY FLAG NOT SET, THEN		00000720
	MOVE THE REPORT TO THE TOP OF SCREEN		00000730
030I	ASK USER 'LIST TRANSFER FAILURE COUNTS ?'		00000740
	IF THE ANSWER IS YES, THEN		
	CLEAR THE SCREEN		

```

INITIALIZE THE LINE COUNT          00000750
58JI INCREMENT THE PAGE NUMBER AND PUT IN HEADER 00000760
58JI PRINT THE HEADER #1           00000770
59JI CLEAR THE ABORT FLAG          00000780
ELSE                                00000790
59JI SET THE ABORT FLAG            00000800
ENDIF                                00000810
ELSE                                00000820
59JI CLEAR THE ABORT FLAG          00000830
SKIP A FEW LINES                   00000840
ENDIF                                00000850
59JI IF THE ABORT FLAG IS NOT SET, THEN 00000860
04JI PRINT/DISPLAY 'T R A N S F E R F A I L U R E S' 00000870
39FM GET NUMBER OF SORT LINES FROM SYSCJM 00000880
CALL THE SUBROUTINE TO OUTPUT THE TRANS. FAILURE REPORT 00000890
(ELSE)                              00000900
ENDIF                                00000910
59JI IF HARDCOPY FLAG IS SET, THEN 00000920
EJECT TO TOP OF FORM               00000930
(ELSE)                              00000940
ENDIF                                00000950
DEQUEUE WHATEVER WAS ENQUEUED      00000960
EXIT                                00000970
SUBROUTINE TO OUTPUT MIS-SORT OR TRANSFER FAILURE REPORT 00000980
00001000
58JI CALCULATE THE NUMBER OF COLUMNS NEEDED 00001010
58JI PRINT/DISPLAY THAT MANY 'LINE COUNT' 00001020
40FM PRINT THE COUNTS              00001030
ADJUST THE LINE COUNT              00001040
59JI IF THE HARDCOPY FLAG IS NOT SET, THEN 00001050
MOVE TO TOP OF SCREEN              00001060
(ELSE)                              00001070
ENDIF                                00001080
RETURN TO CALLER                   00001090

```

```

RESET ERROR COUNTS

```

```

- MODULE HISTORY -

```

```

PROJECT:      DAS                73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DIRSE              92

```

```

- MODULE ABSTRACT -

```

```

THIS MODULE RESETS THE SYSTEM ERROR COUNTS TO ZERO.
ALL ERROR COUNTS CAN BE RESET OR ANY ONE OF THE FIVE
CAN BE RESET INDIVIDUALLY.

```

```

INITIALIZATION
04JI DISPLAY 'RESET ERROR COUNTS COMMAND' 00000290
03JI ASK USER IF ALL COUNTS SHOULD BE CLEARED 00000300
IF NOT, THEN 00000310
03JI ASK USER IF MISSORTS SHOULD BE CLEARED 00000320
IF YES, THEN 00000330
40FM CLEAR MISSORT COUNTS 00000340
(ELSE) 00000350
ENDIF 00000360
00000370

```

03JI	ASK USER IF THE NO-READ COUNTS SHOULD BE CLEARED	00000380
	IF YES, THEN	00000390
40FM	CLEAR NO-READ COUNTS	00000400
	(ELSE)	00000410
	ENDIF	00000420
03JI	ASK USER IF THE REJECT COUNT SHOULD BE CLEARED	00000430
	IF YES, THEN	00000440
40FM	CLEAR REJECT COUNT	00000450
	(ELSE)	00000460
	ENDIF	00000470
03JI	ASK USER IF THE RECIRC COUNT SHOULD BE CLEARED	00000480
	IF YES, THEN	00000490
40FM	CLEAR RECIRC COUNT	00000500
	(ELSE)	00000510
	ENDIF	00000520
03JI	ASK USER IF LOST FROM TRACKING COUNT SHOULD BE CLEARED	00000530
	IF YES, THEN	00000540
40FM	CLEAR LOST FROM TRACKING COUNT	00000550
	(ELSE)	00000560
	ENDIF	00000570
03JI	ASK USER IF TRANSFER FAILURE COUNTS SHOULD BE CLEARED	00000580
	IF YES, THEN	00000590
40FM	CLEAR TRANSFER FAILURE COUNTS	00000600
	(ELSE)	00000610
	ENDIF	00000620
	ELSE	00000630
40FM	CLEAR ALL COUNTS	00000640
	ENDIF	00000650
04JI	DISPLAY 'COUNTS CLEARED'	00000660
	EXIT	00000670
		00000010
		00000020
		00000030
		00000040
		00000050
	WRITE BATCH (TO DISKETTE) COMMAND	00000060
		00000070
		00000080
		00000090
	- MODULE HISTORY -	00000100
		00000110
	PROJECT: DAS 75-01709	00000120
		00000130
	SJB-SYSTEM: OPERATOR INTERFACE	00000140
		00000150
	MODULE: DIWRB 21	00000160
		00000170
		00000180
		00000190
	- MODULE ABSTRACT -	00000200
		00000210
		00000220
	THIS MODULE REQUESTS A BATCH ID. THE FILE MANAGER IS CALLED	00000230
	TO READ THE BATCH FILE RECORD. IF FOUND, THE BATCH INFORMATION	00000240
	IS DISPLAYED FOR USER VERIFICATION. THE USER IS ASKED IF OK TO	00000250
	WRITE TO DISKETTE. IF YES, THE USER IS ASKED IF DISKETTE IS	00000260
	MOUNTED. IF YES, THE FILE MANAGER IS CALLED TO SEQUENTIALLY READ	00000270
	ALL ORDER FILE RECORDS TO BUILD A TABLE OF ORDER ID'S.	00000280
	NEXT THE FILE MANAGER IS CALLED TO SEQUENTIALLY READ ALL	00000290
	TRANSACTION RECORDS, THEN EACH IS WRITTEN TO THE DISKETTE.	00000300
	ENQUEUE USER TERMINAL	00000310
	CLEAR THE SCREEN	00000320
04JI	DISPLAY ON USER TERMINAL "WRITE BATCH COMMAND"	00000330
48JI	DEQUEUE AN IPM BUFFER (GET.IPM)	00000340
	DEQUEUE A RECORD BUFFER (GET.SR3)	00000350
03JI	ASK USER FOR BATCH ID	00000360
08AL	INITIALIZE PANIC CODE	00000370
12JI		00000380
10JI	CLEAR BATCH ABORT	00000390
16JI	ACCESS BATCH FILE RECORD (GET.BATCH)	00000400
10JI	IF BATCH RECORD FOUND (OK SET), THEN	00000410
43JI	! DISPLAY BATCH HEADINGS	00000420

```

4301 ! DISPLAY BATCH INFORMATION 00000430
0301 ! ASK IF OK TO WRITE ? 00000440
! IF YES, THEN 00000450
0301 ! ! ASK IF DISKETTE MOUNTED ? 00000460
! ! IF YES, THEN 00000470
! ! CLEAR DISKETTE NOT MOUNTED FLAG 00000480
08AL ! ! ! INITIALIZE PANIC CODE 00000490
1201 ! ! ! 00000500
5701 ! ! ! SET DISKETTE COMMAND WORD 00000510
3301 ! ! ! 00000520
2301 ! ! ! WRITE DISKETTE HEADER RECORD (WRITE.BATCH) 00000530
1001 ! ! ! IF HEADER RECORD WRITTEN (OK SET), THEN 00000540
! ! ! ! DEQUEUE TERMINAL 00000550
08AL ! ! ! ! INITIALIZE PANIC CODE 00000560
1301 ! ! ! ! 00000570
1301 ! ! ! ! INITIALIZE ORDER COUNT 00000580
1301 ! ! ! ! COUNTIL ORDER COUNT GREATER THAN MAXIMUM 00000590
1701 ! ! ! ! FIND FIRST/NEXT ORDER IN THIS BATCH (FIND.ORDER) 00000600
1001 ! ! ! ! IF ORDER RECORD FOUND (OK SET), THEN 00000610
3301 ! ! ! ! STORE ORDER ID IN TABLE 00000620
! ! ! ! ELSE 00000630
1001 ! ! ! ! SET BATCH ABORT 00000640
! ! ! ! ENDIF 00000650
! ! ! ! ENDDO 00000660
1001 ! ! ! ! IF BATCH ABORT NOT SET, THEN 00000670
08AL ! ! ! ! ! INITIALIZE PANIC CODE 00000680
1201 ! ! ! ! ! 00000690
1401 ! ! ! ! ! 00000700
1001 ! ! ! ! ! CLEAR ORDER SEQUENCE, AND EXCEPTIONS 00000710
5701 ! ! ! ! ! SET DISKETTE COMMAND WORD 00000720
3301 ! ! ! ! ! 00000730
1401 ! ! ! ! ! INITIALIZE CURRENT TRANSACTION BLOCK COUNT 00000740
1401 ! ! ! ! ! INITIALIZE TRANSACTION COUNT 00000750
1001 ! ! ! ! ! COUNTIL TRANSACTION COUNT GREATER THAN MAXIMUM 00000760
1901 ! ! ! ! ! ! FIND FIRST/NEXT TRANSACTION RECORD IN BATCH 00000770
! ! ! ! ! ! --- (FIND.TRANS) --- 00000780
1001 ! ! ! ! ! ! IF TRANSACTION RECORD FOUND (OK SET), THEN 00000790
1401 ! ! ! ! ! ! IF TRANSACTION COUNT GREATER THAN MAX, THEN 00000800
5701 ! ! ! ! ! ! SET DISKETTE COMMAND WORD 00000810
! ! ! ! ! ! (ELSE) 00000820
! ! ! ! ! ! ENDIF 00000830
3301 ! ! ! ! ! ! UPDATE ORDER ID FROM TABLE 00000840
2301 ! ! ! ! ! ! WRITE NEXT DISKETTE RECORD (WRITE.BATCH) 00000850
1001 ! ! ! ! ! ! IF DISKETTE ERROR OCCURRED (OK CLEAR), THEN 00000860
1401 ! ! ! ! ! ! UPDATE TRANSACTION COUNT BEYOND MAXIMUM 00000870
! ! ! ! ! ! ELSE 00000880
! ! ! ! ! ! ENDIF 00000890
! ! ! ! ! ! ELSE 00000900
1001 ! ! ! ! ! ! SET BATCH ABORT 00000910
! ! ! ! ! ! ENDIF 00000920
! ! ! ! ! ! ENDDO 00000930
! ! ! ! ! ! (ELSE) 00000940
! ! ! ! ! ! ENDIF 00000950
! ! ! ! ! ! (ELSE) 00000960
! ! ! ! ! ! ENDIF 00000970
! ! ! ! ! ! ELSE 00000980
1001 ! ! ! ! ! ! SET BATCH ABORT 00000990
! ! ! ! ! ! ENDIF 00010000
! ! ! ! ! ! ELSE 00010010
1001 ! ! ! ! ! ! SET BATCH ABORT 00010020
! ! ! ! ! ! ENDIF 00010030
! ! ! ! ! ! ELSE 00010040
0201 ! ! ! ! ! ! DISPLAY BATCH ID, "BATCH NOT ON FILE" 00010050
1001 ! ! ! ! ! ! SET BATCH ABORT 00010060
! ! ! ! ! ! ENDIF 00010070
5001 ! ! ! ! ! ! RELEASE IPM BUFFER (FREE.IPM) 00010080
! ! ! ! ! ! ENQUEUE TERMINAL 00010090
1001 ! ! ! ! ! ! IF BATCH ABORT SET, THEN 00010100
0401 ! ! ! ! ! ! DISPLAY "WRITE BATCH COMMAND ABORTED" 00010110
! ! ! ! ! ! (ELSE) 00010120
! ! ! ! ! ! ENDIF 00010130
! ! ! ! ! ! IF DISKETTE MOUNTED THEN 00010140

```

04JI	DISPLAY "REMOVE DISKETTE"	00001150
	(ELSE)	00001160
	ENDIF	00001170
11JI	CLEAR DISKETTE IN USE	00001180
	DEQUEUE TERMINAL	00001190
	ENDPRG	00001200
		00000010
		00000020
		00000030
		00000040
		00000050
	WRITE DISKETTE DATA BLOCK SUBROUTINE	00000060
		00000070
		00000080
		00000090
	- MODULE HISTORY -	00000100
		00000110
		00000120
	PROJECT: DAS 75-01709	00000130
		00000140
	SUB-SYSTEM: DISKETTE MANAGER	00000150
		00000160
	MODULE: DMWDB 59	00000170
		00000180
		00000190
	- MODULE ABSTRACT -	00000200
		00000210
		00000220
		00000230
	THIS SUBROUTINE WHEN CALLED WILL PUT A HEADER OR	00000240
	DETAIL RECORD ON A DATA INTERCHANGE FORMAT DISKETTE.	00000250
	WHEN THE HEADER RECORD IS PROCESSED THE DISKETTE WILL	00000260
	BE CHECKED AND BROUGHT ON-LINE TO EDC WITHOUT A \$VARYON	00000270
	COMMAND. WHEN ALL RECORDS HAVE BEEN PROCESSED THE END-	00000280
	OF-DATA (EOD) POINTER ON THE DISKETTE HDRI LABEL WILL BE	00000290
	UPDATED AND THE HEADER RECORD WILL BE UPDATED WITH THE	00000300
	RECORD COUNT.	00000310
		00000320
	NOTE: THE CALLING PROGRAM MUST HAVE DS1=\$\$ AND AN ENTRY	00000330
	STATEMENT FOR DS1	00000340
		00000350
	IF COMMAND EQ TO 'PUT HEADER', THEN	00000360
	BRING DISKETTE ON-LINE	00000370
	IF NO ERRORS, THEN	00000380
	READ THE HDRI RECORD FROM DISKETTE	00000390
	IF DISKETTE EXTENTS OK, THEN	00000400
	POINT TO FIRST RECORD	00000410
	MOVE 3 LOGICAL RECORDS TO BUFFER	00000420
	SET INITILIZE SWITCH	00000430
	MOVE 'OK' TO RETURN CODE	00000440
	ELSE	00000450
	MOVE 'I/O ERROR' TO RETURN CODE	00000460
	ENDIF	00000470
	ELSE	00000480
	MOVE 'I/O ERROR' TO RETURN CODE	00000490
	ENDIF	00000500
	ELSE	00000510
	IF COMMAND EQ TO 'PUT DETAIL RECORD' OR 'LAST DETAIL', THEN	00000520
	IF INITILIZE SWITCH SET, THEN	00000530
	MOVE NEXT 24 BYTE LOGICAL RECORD TO BUFFER	00000540
	MOVE 'OK' TO RETURN CODE	00000550
	IF END OF PHYSICAL BUFFER OR COMMAND EQ 'LAST DETAIL'	00000560
	WRITE RECORD	00000570
	IF NOT END OF FILE, THEN	00000580
	IF DISK WRITE ERROR, THEN	00000590
	MOVE 'I/O ERROR' TO RETURN CODE	00000600
	ELSE	00000610
	IF COMMAND EQ LAST RECORD THEN	00000620
	PUT NUMBER OF RECORDS IN HEADER	00000630
	WRITE HEADER RECORD	00000640
	IF DISK WRITE ERROR, THEN	00000650
	MOVE 'I/O ERROR' TO RETURN CODE	

```

        (ELSE)
        ENDIF
        CALCULATE EOD PNTR FROM NUMBER OF RECORDS
        SET EOD POINTER IN HDR1 RECORD
        WRITE HDR1 RECORD TO DISK
        IF DISK WRITE ERROR, THEN
            MOVE 'I/O ERROR' TO RETURN CODE
        (ELSE)
        ENDIF
    (ELSE)
    ENDIF
ENDIF
ELSE
    MOVE 'I/O ERROR' TO RETURN CODE
ENDIF
ELSE
    IF END OF FIRST 128 BYTES, THEN
        BUMP BUFFER POINTER TO NEXT 128 BYTES
    (ELSE)
    ENDIF
ENDIF
ELSE
    MOVE 'INVALID CALL CODE' TO RETURN CODE
ENDIF
ELSE
    MOVE 'INVALID CALL CODE' TO RETURN CODE
ENDIF
RETURN
    
```

```

00000660
00000670
00000680
00000690
00000700
00000710
00000720
00000730
00000740
00000750
00000760
00000770
00000780
00000790
00000800
00000810
00000820
00000830
00000840
00000850
00000860
00000870
00000880
00000890
00000900
00000910
00000920
00000930
00000940
    
```

BATCH STATUS REPORT

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DIBSR            15
    
```

- MODULE ABSTRACT -

THIS MODULE REQUESTS THE BATCH ID. THE FILE MANAGER IS CALLED TO READ THE BATCH FILE INFORMATION.

IF THE BATCH INFORMATION IS FOUND, THE USER IS ASKED WHETHER OR NOT HARDCOPY OUTPUT IS REQUIRED. IF HARDCOPY IS REQUIRED, THEN THE TERMINAL IS DEQUEUED AND THE SYSTEM PRINTER ENQUEUED.

THE FILE MANAGER IS CALLED TO SEQUENTIALLY READ EACH ORDER RECORD. FOR EACH ORDER RECORD, THE ORDER ID AND ALL QUANTITIES AND ORDER PERCENT COMPLETE ARE OUTPUT. THE COMPLETED QUANTITY IS ACCUMULATED.

THE BATCH PERCENT COMPLETE IS OUTPUT TO THE ENQUEUED DEVICE. THE TERMINAL/PRINTER IS DEQUEUED BEFORE THE PROGRAM ENDS.

ENQUEUE USER TERMINAL

040I DISPLAY ON USER TERMINAL "BATCH STATUS REPORT"

480I DEQUEUE AN IPM BUFFER (GET.IPM)

490I DEQUEUE RECORD BUFFER (GET.SR3)

030I ASK USER FOR BATCH ID

120I

08AL INITIALIZE PANIC CODE

280I

160I ACCESS BATCH FILE (GET.BATCH)

01FM

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
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00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
    
```

```

1001 IF BATCH RECORD FOUND (OK SET), THEN 00000450
1001 SET BATCH FOUND 00000460
0301 ASK USER IF HARDCOPY OUTPUT IS REQUIRED 00000470
1001 IF HARDCOPY REQUIRED, THEN 00000480
    SET HARDCOPY 00000490
    DEQUEUE USER TERMINAL 00000500
    ENQUEUE PRINTER 00000510
    ELSE 00000520
1001 CLEAR HARDCOPY 00000530
    ENDIF 00000540
2801 INITIALIZE DATE, TIME, PAGE COUNT AND LINE COUNT 00000550
0601 PRINT/DISPLAY BATCH HEADINGS 00000560
0601 PRINT/DISPLAY BATCH INFORMATION 00000570
0601 PRINT/DISPLAY ORDER HEADINGS 00000580
0801 INITIALIZE PANIC CODE 00000590
2801 00000600
1301 INITIALIZE QUANTITY COMPLETE 00000610
1301 INITIALIZE ORDER COUNT 00000620
1301 DOUNTIL ORDER COUNT GREATER THAN MAXIMUM 00000630
1701 FIND THE FIRST/NEXT ORDER IN THIS BATCH (FIND.ORDER) 00000640
1001 IF ORDER RECORD FOUND (OK SET), THEN 00000650
1301 COMPUTE ORDER PERCENT COMPLETE 00000660
0601 PRINT/DISPLAY ORDER INFORMATION AND PERCENT COMPLETE 00000670
    --- (FORMS.CTL) --- 00000680
1001 IF ABORT SET THEN 00000690
1301 SET ORDER COUNT GREATER THAN MAX. 00000700
1001 CLEAR BATCH FOUND 00000710
    (ELSE) 00000720
    ENDIF 00000730
1301 ACCUMULATE QUANTITY COMPLETE 00000740
    ELSE 00000750
1001 CLEAR BATCH FOUND 00000760
0401 PRINT/DISPLAY "BATCH NO LONGER ON FILE - RPT TERM." 00000770
    ENJIF 00000780
    ENDDO 00000790
    IF BATCH FOUND SET, THEN 00000800
2801 COMPUTE BATCH PERCENT COMPLETE 00000810
0601 PRINT/DISPLAY BATCH PERCENT, "BATCH PERCENT COMPLETE" 00000820
    (ELSE) 00000830
    ENDIF 00000840
    IF HARD COPY NOT SET THEN 00000850
    MOVE REPORT TO TOP OF SCREEN 00000860
    (ELSE) 00000870
    ENDIF 00000880
    ELSE 00000890
0201 DISPLAY BATCH ID, "BATCH NOT ON FILE" 00000900
    ENDIF 00000910
5101 RELEASE RECORD BUFFER (FREE.SRB) 00000920
2801 RELEASE THE IPM BUFFER (FREE.IPM) 00000930
    DEQUEUE TERMINAL/PRINTER 00000940
    ENDPROG 00000950
        FORMS.CTL SUBROUTINE 00000960
        00000970
1001 IF HARDCOPY SET, THEN 00000980
2801 IF LINE COUNT GREATER THAN PRINTER PAGE, THEN 00000990
1001 SET OK 00010000
    ELSE 00010100
1001 CLEAR OK 00010120
    ENDIF 00010130
    ELSE 00010140
2801 IF LINE COUNT GREATER THAN TERMINAL PAGE, THEN 00010150
1001 SET OK 00010160
    ELSE 00010170
1001 CLEAR OK 00010180
    ENDIF 00010190
    ENDIF 00011000
1001 IF PAGE OVERFLOW (OK SET), THEN 00011110
2801 INCREMENT PAGE COUNT 00011120
0601 PRINT/DISPLAY BATCH HEADINGS 00011130
0601 PRINT/DISPLAY BATCH INFORMATION 00011140
0601 PRINT/DISPLAY ORDER HEADINGS 00011150
2801 INITIALIZE LINE COUNT 00011160
    (ELSE) 00011170

```

```

ENDIF
060I PRINT/DISPLAY ORDER INFORMATION
280I INCREMENT LINE COUNT
      IF REPORT CANCELLED AND HARD COPY SET THEN
        PRINT REPORT CANCELLED MESSAGE
      (ELSE)
ENDIF
RETURN

```

00001180
00001190
00001200
00001210
00001220
00001230
00001240
00001250

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600
00000610
00000620
00000630
00000640

READ BATCH (FROM DISKETTE) COMMAND

- MODULE HISTORY -

```

PROJECT:      DAS              75-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DIRD              20

```

- MODULE ABSTRACT -

THIS MODULE REQUESTS THE USER TO MOUNT A DISKETTE. THE DISKETTE HEADER IS READ AND DISPLAYED BACK TO THE USER. THE USER IS ASKED IF THIS DISKETTE IS OK TO READ. IF YES, THEN THE FILE MANAGER IS CALLED TO ADD THE BATCH. THE DISKETTE IS SEQUENTIALLY READ WHILE THE FILE MANAGER IS CALLED FOR EACH DISKETTE RECORD, TO ADD EACH TRANSACTION RECORD.

```

ENQUEUE USER TERMINAL
39FM GET NUMBER OF SORT LINES
140I DISPLAY ON USER TERMINAL "READ BATCH COMMAND"
030I ASK USER IF DISKETTE MOUNTED ?
      IF YES, THEN
490I ! DEQUEUE RECORD BUFFER (GET.SRB)
080I ! INITIALIZE PANIC CODE
320I !
570I ! SET DISKETTE COMMAND WORD
320I !
220I ! GET DISKETTE HEADER RECORD (READ.BATCH)
100I ! IF HEADER RECORD FOUND (OK SET), THEN
430I ! ! DISPLAY BATCH HEADINGS
430I ! ! DISPLAY BATCH INFORMATION
100I ! ! CLEAR BATCH ABORT
030I ! ! ASK IF OK TO READ ?
      ! ! IF YES, THEN
480I ! ! ! DEQUEUE AN IPM BUFFER (GET.IPM)
084L ! ! ! INITIALIZE PANIC CODE
120I ! ! !
240I ! ! ! ADD BATCH INFORMATION (ADD.BATCH)
100I ! ! ! IF BATCH ADDED (OK SET), THEN
      ! ! ! ! DEQUEUE TERMINAL
084L ! ! ! ! INITIALIZE PANIC CODE
140I ! ! ! !
570I ! ! ! ! SET DISKETTE COMMAND WORD
320I ! ! ! !
100I ! ! ! ! CLEAR LAST RECORD
100I ! ! ! ! CLEAR DELETION
100I ! ! ! ! DOUNTIL LAST RECORD SET
220I ! ! ! ! ! SET DISKETTE TRANSACTION RECORD (READ.BATCH)
100I ! ! ! ! ! IF TRANSACTION RECORD FOUND (OK SET), THEN
      ! ! ! ! ! VALIDATE TRANSACTION RECORD INFORMATION
      ! ! ! ! ! --- (VALIDATE.TRANSACTION) ---
100I ! ! ! ! ! IF TRANSACTION DATA VALID (OK SET), THEN
250I ! ! ! ! ! ADD TRANSACTION INFORMATION (ADD.TRANS)

```



```

100I      ! ! ! ! !      IF ERROR (OK CLEAR), THEN      00000650
100I      ! ! ! ! !      SET LAST RECORD      00000660
100I      ! ! ! ! !      SET BATCH ABORT      00000670
           ! ! ! ! !      (ELSE)      00000680
           ! ! ! ! !      ENDIF      00000690
           ! ! ! ! !      ELSE      00000700
100I      ! ! ! ! !      SET LAST RECORD      00000710
100I      ! ! ! ! !      SET BATCH ABORT      00000720
100I      ! ! ! ! !      SET DELETION      00000730
           ! ! ! ! !      ENDIF      00000740
           ! ! ! ! !      ELSE      00000750
100I      ! ! ! ! !      SET LAST RECORD      00000760
100I      ! ! ! ! !      SET BATCH ABORT      00000770
100I      ! ! ! ! !      SET DELETION      00000780
           ! ! ! ! !      ENDIF      00000790
           ! ! ! ! !      ENDDO      00000800
           ! ! ! ! !      ELSE      00000810
100I      ! ! ! ! !      SET BATCH ABORT      00000820
           ! ! ! ! !      ENDIF      00000830
250I      ! ! ! ! !      RELEASE IPM BUFFER (FREE.IPM)      00000840
           ! ! ! ! !      ELSE      00000850
100I      ! ! ! ! !      SET BATCH ABORT      00000860
           ! ! ! ! !      ENDIF      00000870
           ! ! ! ! !      ELSE      00000880
100I      ! ! ! ! !      SET BATCH ABORT      00000890
           ! ! ! ! !      ENDIF      00000900
510I      ! ! ! ! !      RELEASE RECORD BUFFER (FREE.SRB)      00000910
           ! ! ! ! !      ELSE      00000920
10I      ! ! ! ! !      SET BATCH ABORT      00000930
           ! ! ! ! !      ENDIF      00000940
100I      ! ! ! ! !      IF DELETION SET, THEN      00000950
120I      ! ! ! ! !      INITIALIZE BATCH DIRECTORY ADDRESS, ENTRY SIZE, BATCH COUNT      00000960
120I      ! ! ! ! !      DOWHILE RESIDENT BATCH STATUS NOT ADDING      00000970
           ! ! ! ! !      ADD 1 TO BATCH COUNT      00000980
           ! ! ! ! !      INDEX TO NEXT REC. IN RESIDENT BATCH FILE      00000990
           ! ! ! ! !      ENDDO      00001000
02FM      ! ! ! ! !      CHANGE BATCH STATUS TO DELETING      00001010
530I      ! ! ! ! !      USE FILE MANAGER TO DELETE THIS BATCH (DELETE.BATCH)      00001020
           ! ! ! ! !      (ELSE)      00001030
           ! ! ! ! !      ENDIF      00001040
100I      ! ! ! ! !      IF BATCH ABORT SET, THEN      00001050
           ! ! ! ! !      ENQUEUE TERMINAL      00001060
040I      ! ! ! ! !      DISPLAY "READ BATCH COMMAND ABORTED"      00001070
           ! ! ! ! !      (ELSE)      00001080
           ! ! ! ! !      ENDIF      00001090
040I      ! ! ! ! !      DISPLAY "REMOVE DISKETTE"      00001100
1100      ! ! ! ! !      CLEAR DISKETTE IN USE      00001110
           ! ! ! ! !      DEQUEUE USER TERMINAL      00001120
           ! ! ! ! !      ENDPROG      00001130
           ! ! ! ! !      00001140
           ! ! ! ! !      00001150
           ! ! ! ! !      VALIDATE TRANSACTION SUBROUTINE      00001150
           ! ! ! ! !      00001170
100I      ! ! ! ! !      SET OK      00001180
07FM      ! ! ! ! !      IF TRANSACTION ID ZERO OR NOT NUMERIC, THEN      00001190
100I      ! ! ! ! !      CLEAR OK      00001200
           ! ! ! ! !      (ELSE)      00001210
           ! ! ! ! !      ENDIF      00001220
07FM      ! ! ! ! !      IF STATUS NOT VALID OR NOT NUMERIC THEN      00001230
100I      ! ! ! ! !      CLEAR OK      00001240
           ! ! ! ! !      (ELSE)      00001250
           ! ! ! ! !      ENDIF      00001260
07FM      ! ! ! ! !      IF DESTINATION GREATER THAN MAXIMUM OR NOT NUMERIC, THEN      00001270
100I      ! ! ! ! !      CLEAR OK      00001280
           ! ! ! ! !      (ELSE)      00001290
           ! ! ! ! !      ENDIF      00001300
100I      ! ! ! ! !      IF VALIDATION FAILED (OK CLEAR), THEN      00001310
020I      ! ! ! ! !      DISPLAY "TRANSACTION VALIDATION ERROR"      00001320
440I      ! ! ! ! !      DISPLAY DISKETTE TRANSACTION INFORMATION      00001330
           ! ! ! ! !      (ELSE)      00001340
           ! ! ! ! !      ENDIF      00001350
           ! ! ! ! !      RETJRN      00001350

```

DISKETTE MANAGER READ SUBROUTINE

```

- MODULE HISTORY -

PROJECT:      DAS              73-01709
SUB-SYSTEM:   DISKETTE MANAGER
MODULE:       DMRDB           59

- MODULE ABSTRACT -

THIS SUBROUTINE WHEN CALLED WILL GET A HEADER OR
DETAIL RECORD FROM A DATA INTERCHANGE FORMAT DISKETTE.
WHEN THE HEADER RECORD IS REQUESTED THE DISKETTE WILL
BE CHECKED AND BROUGHT ON-LINE TO EXX WITHOUT A $VARYON
COMMAND. WHEN ALL RECORDS HAVE BEEN PROCESSED FROM THE
DISKETTE, THE RETURN CODE WILL BE SET TO INDICATE THE
LAST RECORD HAS BEEN PROCESSED.

IF COMMAND EQ TO 'GET HEADER', THEN
  BRING DISKETTE ON-LINE
  IF NO ERRORS, THEN
    READ THE HDRI RECORD FROM DISKETTE
    IF DISKETTE EXTENTS OK, THEN
      READ FIRST DATA RECORD
      IF DISK READ OK, THEN
        SET NUMBER OF RECORDS FROM HEADER
        IF NUMBER OF RECORDS BETWEEN 4 AND 9490, THEN
          SET INITILIZE SWITCH
          MOVE 3 LOGICAL RECORDS TO BUFFER
          MOVE 'JK' TO RETURN CODE
          DECREMENT RECORD COUNT
        ELSE
          MOVE 'DISKETTE FORMAT ERROR TO RTN CODE
        ENDIF
      ELSE
        MOVE 'I/O ERROR' TO RETURN CODE
      ENDIF
    ELSE
      MOVE 'I/O ERROR' TO RETURN CODE
    ENDIF
  ELSE
    MOVE 'I/O ERROR' TO RETURN CODE
  ENDIF
ELSE
  IF COMMAND EQ TO 'GET DETAIL RECORD', THEN
    IF INITILIZE SWITCH SET, THEN
      MOVE NEXT 24 BYTE LOGICAL RECORD TO BUFFER
      MOVE 'OK' TO RETURN CODE
    IF ALL RECORDS HAVE BEEN PROCESSED, THEN
      MOVE 'LAST RECORD' TO RETURN CODE
    ELSE
      IF END OF FIRST 128 BYTE BUFFER, THEN
        BUMP BUFFER POINTER TO NEXT 128 BYTES
      ELSE
        IF END OF PHYSICAL BUFFER, THEN
          READ NEXT DISK RECORD
          IF DISK READ ERRORS, THEN
            MOVE 'I/O ERROR' TO RETURN CODE
          (ELSE)
          ENDIF
        (ELSE)
        ENDIF
      ENDIF
    ENDIF
  ENDIF

```

```

00000010
00000020
00000030
00000040
00000050
00000050
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600
00000610
00000620
00000630
00000640
00000650
00000660
00000670
00000680
00000690
00000700
00000710
00000720
00000730
00000740

```

```

      ENDIF
    ENDIF
  ELSE
    MOVE 'INVALID CALL CODE' TO RETURN CODE
    PANIC
  * ENDIF
  ELSE
    MOVE 'INVALID CALL CODE' TO RETURN CODE
  ENDIF
ENDIF
RETURN

```

```

00000750
00000750
00000770
00000780
00000790
00000800
00000810
00000820
00000830
00000840
00000850

```

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220

```

TRANSACTION STATUS REPORT

- MODULE HISTORY -

```

PROJECT:      DAS              75-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DITSR            13

```

- MODULE ABSTRACT -

```

      THIS MODULE REQUESTS A TRANSACTION NUMBER, THEN CALLS THE
      FILE MANAGER TO READ THE TRANSACTION FILE.
      IF THE TRANSACTION RECORD IS FOUND, THE FILE MANAGER IS AGAIN
      CALLED, THIS TIME TO READ THE DISK BATCH FILE.
      IF THE DISK BATCH FILE INFORMATION IS FOUND, THE USER IS
      ASKED WHETHER OR NOT HARDCOPY OUTPUT IS REQUIRED. IF HARDCOPY IS
      REQUIRED, THEN THE TERMINAL IS DEQUEUED AND THE SYSTEM PRINTER
      IS ENQUEUED.
      THE COMBINED BATCH AND TRANSACTION FILE INFORMATION IS THEN
      OUTPUT TO THE ENQUEUED DEVICE.
      THE TERMINAL/PRINTER IS DEQUEUED BEFORE THE PROGRAM ENDS.
      ENQUEUE USER TERMINAL
      DISPLAY ON USER TERMINAL "TRANSACTION STATUS REPORT"
      CLEAR THE SCREEN
040I
140I
480I  DEQUEUE AN IPM BUFFER (GET,IPM)
490I  DEQUEUE RECORD BUFFER (GET,SRB)
030I  ASK USER FOR TRANSACTION ID
      IF VALID (OK<TRANSACTION ID<=999999), THEN
08AL  ! INITIALIZE PANIC CODE
140I  !
20FM  ! JDATE IPM FOR CALL
08FM  ! CALL FILE MANGR TO TRANSACTION DIRECTORY RECORD BY TRANS. ID
20FM  ! CASENTRY (FDK, FTV, FCC, ? )
      ! CASE FDK ( SUCCESSFUL )
100I  ! SET OK
03FM  ! SAVE THE MIS-SORT DESTINATION
      ! CASE FTV ( TRANSACTION NOT FOUND )
020I  ! DISPLAY TRANSACTION ID, "TRANSACTION NOT ON FILE"
100I  ! CLEAR OK
      ! CASE FCC ( INVALID CALL CODE )
540I  ! PANIC
      ! CASE ? ( INVALID COMPLETION CODE )
540I  ! PANIC
      ! ENDCASE
100I  ! IF TRANSACTION DIRECTORY RECORD FOUND (OK SET), THEN
120I  ! ! GET BATCH ID FROM SYSCOM BATCH DIRECTORY
02FM  ! !

```

```

00000230
00000240
00000250
00000250
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600

```

```

15DI ! ! ACCESS BATCH FILE (GET.BATCH) 00000610
10DI ! ! IF BATCH RECD FOUND (JK SET), THEN 00000620
08AL ! ! ! INITIALIZE PANIC CODE 00000630
14DI ! ! ! 00000640
21DI ! ! ! ACCESS TRANSACTION FILE (GET.TRANS) 00000650
10DI ! ! ! IF TRANSACTION RECORD FOUND (JK SET), THEN 00000660
23FM ! ! ! CALL FILE MANAGER FOR ORDER ID 00000670
23FM ! ! ! IF ORDER RECORD FOUND, THEN 00000680
05FM ! ! ! SAVE ORDER ID 00000690
03DI ! ! ! ASK USER IF HARDCOPY OUTPJT IS REQUIRED 00000700
! ! ! IF HARDCOPY REQUIRED, THEN 00000710
10DI ! ! ! SET HARD COPY FLAG 00000720
! ! ! DEQUEUE USER TERMINAL 00000730
! ! ! ENQUEUE PRINTER 00000740
! ! ! ELSE 00000750
10DI ! ! ! CLEAR HARD COPY FLAG 00000760
! ! ! ENDIF 00000770
30DI ! ! ! INITIALIZE DATE, TIME AND PAGE CNT 00000780
08DI ! ! ! PRINT/DISPLAY BATCH HEADINGS 00000790
08DI ! ! ! PRINT/DISPLAY BATCH INFORMATION 00000800
08DI ! ! ! PRINT/DISPLAY TRANSACTION HEADINGS 00000810
! ! ! IF TRANS. STATUS = MIS-SORT THEN 00000820
! ! ! IF MIS-SORT DESTINATION = -1 THEN 00000830
! ! ! SET MIS-SORT DEST. = '?' 00000840
! ! ! ELSE 00000850
! ! ! SET MIS-SORT DEST. IN OUTPJT LINE 00000860
! ! ! ENDIF 00000870
! ! ! (ELSE) 00000880
! ! ! ENDIF 00000890
08DI ! ! ! PRINT/DISPLAY TRANSACTION INFORMATION 00000900
10DI ! ! ! IF HARD COPY FLAG NOT SET, THEN 00000910
! ! ! MOVE DISPLAY TO TOP OF SCREEN 00000920
! ! ! (ELSE) 00000930
! ! ! ENDIF 00000940
! ! ! ELSE 00000950
23FM ! ! ! IF BATCH WAS NOT FOUND, THEN 00000960
02DI ! ! ! DISPLAY BATCH ID "BATCH NOT ON FILE" 00000970
! ! ! ELSE 00000980
23FM ! ! ! IF THE RELATIVE ORDER NUMBER WAS INVALID, THEN 00000990
08AL ! ! ! SET UP PANIC CODE FOR INVLD REL DRD 00001000
! ! ! ELSE 00001010
08AL ! ! ! SET UP PANIC CODE FOR INVLD CALL CODE 00001020
! ! ! ENDIF 00001030
07AL ! ! ! PANIC 00001040
! ! ! ENDIF 00001050
! ! ! ELSE 00001060
! ! ! ELSE 00001070
02DI ! ! ! DISPLAY TRANSACTION ID, "TRANSACTION NOT ON FILE" 00001080
! ! ! ENDIF 00001090
! ! ! ELSE 00001100
02DI ! ! ! DISPLAY BATCH ID "BATCH NOT ON FILE" 00001110
! ! ! ENDIF 00001120
! ! ! (ELSE) 00001130
! ! ! ENDIF 00001140
ELSE 00001150
02DI ! ! ! DISPLAY INVALID TRANSACTION NO 00001160
ENDIF 00001170
51DI RELEASE RECORD BUFFER (FREE.SRB) 00001180
50DI RELEASE THE IPM BUFFER (FREE.IPM) 00001190
DEQUEUE TERMINAL/PRINTER 00001200
ENDPRG 00001210

```

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090

```

FIND DISK TRANSACTION FILE RECORD BY CRITERIA (FIND.TRANS)

- MODULE HISTORY -

PROJECT: DAS 75-01709
 SUB-SYSTEM: OPERATOR INTERFACE
 MODULE: DICCFY 78

- MODULE ABSTRACT -

EACH RELATIVE TRANSACTION RECORD IS SEQUENTIALLY
 ACCESSED ONE AT A TIME BY CALLING THE SEQ.TRANS
 SUBROUTINE FOR EACH RELATIVE ORDER NUMBER.

IF SUCCESSFUL, THE ORDER SEQUENCE INDICATOR IS
 EXAMINED. IF SET, THE TRANSACTION RELATIVE ORDER
 NUMBER IS COMPARED TO THE CURRENT RELATIVE ORDER
 NUMBER. IF MATCHED, THE EXCEPTIONS INDICATOR IS
 EXAMINED. IF SET, THE TRANSACTION STATUS IS EX-
 AMINED FOR EXCEPTION STATUS.

- THIS SUBROUTINE WILL EXIT WHEN EITHER:
 1) THE DISK TRANSACTION FILE HAS BEEN ACCESSED
 WITHOUT ERROR, AND THE (OPTIONAL) SELECTION
 CRITERIA (ORDER SEQUENCE, EXCEPTIONS) HAS BEEN
 MET, OR
 2) THE DISK TRANSACTION FILE HAS BEEN COMPLETELY
 ACCESSED.

```

1401 DDUNTIL TRANSACTION COUNT GREATER THAN MAXIMUM OR JK SET--
1201 -- OR TRANSACTION COUNT = ZERO
1001
2001 FIND FIRST/NEXT TRANSACTION IN THIS BATCH (SEQ.TRANS)
1001 IF TRANSACTION RECORD FOUND (JK SET), THEN
1401 COMPUTE TRANSACTION RECORD ENTRY ADDRESS
04FM GET TRANSACTION INFORMATION
1401
1001 IF ORDER SEQUENCE SET, THEN
1401 IF RELATIVE ORDER NUMBER MATCH, THEN
1301
1001 IF EXCEPTIONS SET, THEN
1401 IF TRANSACTION STATUS NOT AN EXCEPTION, THEN
1001 CLEAR JK
(ELSE)
ENDIF
(ELSE)
ENDIF
ELSE
1001 CLEAR JK
ENDIF
(ELSE)
ENDIF
ENDIF
IF TRANSACTION COUNT > ZERO THEN
INCREMENT TRANSACTION COUNT
(ELSE)
ENDIF
ENDDD
RETJRN

```

00000100
 00000110
 00000120
 00000130
 00000140
 00000150
 00000160
 00000170
 00000180
 00000190
 00000200
 00000210
 00000220
 00000230
 00000240
 00000250
 00000260
 00000270
 00000280
 00000290
 00000300
 00000310
 00000320
 00000330
 00000340
 00000350
 00000360
 00000370
 00000380
 00000390
 00000400
 00000410
 00000420
 00000430
 00000440
 00000450
 00000460
 00000470
 00000480
 00000490
 00000500
 00000510
 00000520
 00000530
 00000540
 00000550
 00000560
 00000570
 00000580
 00000590
 00000600
 00000610
 00000620
 00000630
 00000640
 00000650
 00000660
 00000670
 00000680

END BATCH COMMAND

- MODULE HISTORY -

PROJECT: DAS 75-01709

00000010
 00000020
 00000030
 00000040
 00000050
 00000060
 00000070
 00000080
 00000090
 00000100
 00000110
 00000120
 00000130
 00000140

```

SUB-SYSTEM:  OPERATOR INTERFACE          00000150
                                           00000160
MODULE:      DIENB                        23  00000170
                                           00000180
                                           00000190
                                           00000200
                                           00000210
                                           00000220

```

- MODULE ABSTRACT -

```

THIS MODULE REQUESTS A BATCH ID. THE FILE MANAGER IS CALLED TO READ THE CORRESPONDING BATCH INFORMATION. IF THE BATCH RECORD IS FOUND, THEN THE BATCH ID, HEADER AND STATUS ARE DISPLAYED. THE USER IS ASKED WHETHER OR NOT BATCH IS TO BE ENDED. IF YES, THEN THE FILE MANAGER IS CALLED TO CHANGE THE BATCH STATUS

```

```

ENQUEUE USER TERMINAL                      00000280
040I DISPLAY ON USER TERMINAL "END BATCH COMMAND" 00000290
480I DEQUEUE AN IPM BUFFER (GET.IPM)         00000300
490I DEQUEUE RECORD BUFFER (GET.SRB)        00000310
030I ASK USER FOR BATCH ID                  00000320
120I                                         00000330
08AL INITIALIZE PANIC CODE                  00000340
350I                                         00000350
160I ACCESS BATCH FILE (GET.BATCH)          00000360
100I IF BATCH RECORD FOUND (OK SET), THEN  00000370
430I   DISPLAY BATCH HEADER                 00000380
430I   DISPLAY BATCH INFORMATION           00000390
030I   ASK USER OK TO END BATCH            00000400
      IF YES, THEN                          00000410
08AL   INITIALIZE PANIC CODE               00000420
350I                                         00000430
01FM   INITIALIZE NEW BATCH STATUS         00000440
120I                                         00000450
450I   CHANGE BATCH STATUS (CHANGE.BATCH)  00000460
      IF OK SET THEN                        00000470
      DISPLAY "BATCH ENDED"                 00000480
520I   ALARM LOG BATCH STATUS (SEND.ALARM)  00000490
      (ELSE)                                00000500
      ENDIF                                  00000510
      ELSE                                    00000520
      DISPLAY "COMMAND ABORTED"            00000530
      ENDIF                                  00000540
      ELSE                                    00000550
020I   DISPLAY BATCH ID, "BATCH NOT ON FILE" 00000560
      DISPLAY "COMMAND ABORTED"            00000570
      ENDIF                                  00000580
510I RELEASE RECORD BUFFER (FREE.SRB)      00000590
500I RELEASE THE IPM BUFFER (FREE.IPM)     00000600
ENDPRG                                       00000610

```

SORT TRANSLATION REPORT

- MODULE HISTORY -

```

PROJECT:   DAS                            75-01709
SUB-SYSTEM: OPERATOR INTERFACE          00000110
                                           00000120
MODULE:    DISTR                          19  00000130
                                           00000140
                                           00000150
                                           00000160
                                           00000170
                                           00000180
                                           00000190
                                           00000200
                                           00000210
                                           00000220

```

- MODULE ABSTRACT -

```

THIS MODULE REQUESTS WHETHER OR NOT HARDCOPY IS REQUIRED. IF HARDCOPY IS REQUIRED, THEN THE TERMINAL IS DEQUEUED AND THE SYSTEM PRINTER IS ENQUEUED. THE RESIDENT SORT TRANSLATION TABLE IS FORMATTED AND OUTPUT TO THE ENQUEUED DEVICE. THE TERMINAL/PRINTER IS DEQUEUED BEFORE THE PROGRAM ENDS.

```

```

ENQUEUE USER TERMINAL
39FM GET NUMBER OF SORT LINES FROM SYSCOM
04JI DISPLAY ON USER TERMINAL "SORT LINE SORTING REPORT"
03JI ASK USER IF HARDCOPY OUTPUT IS REQUIRED
      IF HARDCOPY REQUIRED, THEN
        DEQUEUE USER TERMINAL
        ENQUEUE PRINTER
      (ELSE)
      ENDF
31JI INITIALIZE DATE, TIME AND PAGE COUNT
31JI INITIALIZE SORT LINE COUNT
31JI INITIALIZE POINTER TO SYSCOM SORT TRANSLATION TABLE
09JI PRINT/DISPLAY REPORT HEADINGS
31JI UNTIL END OF SYSCOM SORT TRANSLATION TABLE
09JI PRINT/DISPLAY TRANSLATION INFORMATION
31JI INCREMENT POINTER TO SORT TRANSLATION TABLE
      ENDD
      DEQUEUE TERMINAL/PRINTER
      ENDPRG

```

03000280
03000290
03000300
03000310
03000320
03000330
03000340
03000350
03000360
03000370
03000380
03000390
03000400
03000410
03000420
03000430
03000440
03000450
03000450
03000470
03000480
03000490

WORK SCHEDULE REPORT

- MODULE HISTORY -

```

PROJECT:      DAS      75-01709
SJB-SYSTEM:   OPERATOR INTERFACE
MODULE:       JIWSR    15

```

- MODULE ABSTRACT -

THIS MODULE REQUESTS WHETHER OR NOT HARDCOPY OUTPUT IS REQUIRED. IF HARDCOPY REQUIRED, THEN THE TERMINAL IS DEQUEUED AND THE SYSTEM PRINTER IS ENQUEUED. THE RESIDENT BATCH DIRECTORY IS CONSULTED FOR BATCH ENTRIES IN USE.

FOR EACH BATCH ENTRY IN USE, THE FILE MANAGER IS CALLED TO READ THE BATCH FILE RECORD. THE BATCH INFORMATION IS THEN OUTPUT TO THE ENQUEUED DEVICE.

THE TERMINAL/PRINTER IS DEQUEUED BEFORE THE PROGRAM ENDS.

```

ENQUEUE USER TERMINAL
CLEAR THE SCREEN
04JI DISPLAY ON USER TERMINAL "WORK SCHEDULE REPORT"
03JI ASK USER IF HARDCOPY OUTPUT IS REQUIRED
      IF HARDCOPY REQUIRED, THEN
        SET HARD COPY FLAG
        DEQUEUE USER TERMINAL
        ENQUEUE PRINTER
      ELSE
        CLEAR HARD COPY FLAG
      ENDF
48JI DEQUEUE AN IPM BUFFER (GET.IPM)
49JI DEQUEUE RECORD BUFFER (GET.SRB)
27JI INITIALIZE DATE, TIME AND PAGE COUNT
05JI PRINT/DISPLAY BATCH HEADINGS
10JI CLEAR BATCH FOUND
02FM INITIALIZE POINTER TO SYSCOM BATCH DIRECTORY
12JI
02FM INITIALIZE SYSCOM BATCH DIRECTORY ENTRY SIZE
12JI

```

03000010
03000020
03000030
03000040
03000050
03000060
03000070
03000080
03000090
03000100
03000110
03000120
03000130
03000140
03000150
03000160
03000170
03000180
03000190
03000200
03000210
03000220
03000230
03000240
03000250
03000260
03000270
03000280
03000290
03000300
03000310
03000320
03000330
03000340
03000350
03000360
03000370
03000380
03000390
03000400
03000410
03000420
03000430
03000440
03000450
03000450
03000470
03000480
03000490
03000500

```

08AL INITIALIZE PANIC CODE 00000510
27JI 00000520
12JI INITIALIZE BATCH COJNT 00000530
12JI DDUNTIL BATCH COJNT GREATER THAN MAXIMJM 00000540
15JI FIND AN ACTIVE, PENDING OR COMPLETE BATCH RECORD (FIND.BATCH) 00000550
10JI IF BATCH RECORD FOUND (OK SET), THEN 00000560
05JI PRINT/DISPLAY BATCH INFORMATION 00000570
27JI 00000580
10JI SET BATCH FOUND 00000590
      (ELSE) 00000600
      ENDIF 00000610
ENDDD 00000620
51JI RELEASE RECORD BUFFER (FREE.SRB) 00000630
50JI RELEASE THE IPM BUFFER (FREE.IPM) 00000640
10JI IF BATCH FOUND NOT SET, THEN 00000650
04JI PRINT/DISPLAY "NO CURRENT BATCH INFORMATION" 00000660
      (ELSE) 00000670
      ENDIF 00000680
      IF HARDCOPY FLAG NOT SET, THEN 00000690
          MOVE REPORT TO TOP OF SCREEN 00000700
      (ELSE) 00000710
      ENDIF 00000720
DEQUEUE TERMINAL/PRINTER 00000730
ENDPRG 00000740

```

SET A BLOCK BUFFER

- MODULE HISTORY -

```

PROJECT:      DAS          75-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DICCGS      63

```

- MODULE ABSTRACT -

THIS SUBROUTINE DEQUEUES A 255 BYTE BLOCK BUFFER FROM THE BUFFER QUEUE IN SYSCOM. IF THE QUEUE IS EMPTY A DELAY IS PERFORMED BEFORE TRYING AGAIN.

```

37FM SET ADDRESS OF SBB (255-BYTE) BUFFER QUEUE IN SYSCOM
DDUNTIL BUFFER RECEIVED
37FM DEQUEUE A BUFFER FROM THE QUEUE
37FM IF THE QUEUE WAS EMPTY THEN
      DELAY
      (ELSE)
      ENDIF
ENDDD
RETURN

```

START BATCH COMMAND

- MODULE HISTORY -

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100

```



```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DISTR           22

```

- MODULE ABSTRACT -

```

THIS MODULE REQUESTS A BATCH ID. THE FILE MANAGER IS CALLED
TO READ THE CORRESPONDING BATCH INFORMATION. IF THE BATCH RECORD
IS FOUND, THEN THE BATCH ID, HEADER AND STATUS ARE DISPLAYED.
THE USER IS ASKED WHETHER OR NOT BATCH IS TO BE STARTED. IF
YES, THEN THE FILE MANAGER IS CALLED TO CHANGE THE BATCH STATUS
ENQUEUE USER TERMINAL
CLEAR THE SCREEN
04JI DISPLAY ON USER TERMINAL "START BATCH COMMAND"
48JI DEQUEUE AN IPM BUFFER (GET.IPM)
49JI DEQUEUE RECORD BUFFER (GET.SRB)
03JI ASK USER FOR BATCH ID
12JI
08AL INITIALIZE PANIC CODE
34JI
16JI ACCESS BATCH FILE (GET.BATCH)
10JI IF BATCH RECORD FOUND (OK SET), THEN
43JI DISPLAY BATCH HEADER
43JI DISPLAY BATCH INFORMATION
03JI ASK USER OK TO START BATCH
IF YES, THEN
08AL INITIALIZE PANIC CODE
34JI
01F4 INITIALIZE NEW BATCH STATUS
12JI
45JI CHANGE BATCH STATUS (CHANGE.BATCH)
IF OK SET THEN
DISPLAY "BATCH STARTED"
52JI ALARM LOG BATCH STATUS (SEND.ALARM)
(ELSE)
ENDIF
ELSE
DISPLAY "COMMAND ABORTED"
ENDIF
ELSE
02JI DISPLAY BATCH ID, "BATCH NOT ON FILE"
DISPLAY "COMMAND ABORTED"
ENDIF
51JI RELEASE RECORD BUFFER (FREE.SRB)
50JI RELEASE THE IPM BUFFER (FREE.IPM)
ENDPRG

```

```

00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600
00000610
00000620

```

READ RELATIVE DISK TRANSACTION FILE RECORD (SEQ.TRANS)

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DISTR           79

```

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190

```

- MODULE ABSTRACT -

EACH RELATIVE TRANSACTION NUMBER IS USED TO CALCULATE THE RELATIVE BLOCK NUMBER AND TRANSACTION RECORD ENTRY OFFSET. IF THE TRANSACTION INFORMATION REQUESTED RESIDES IN A BLOCK OTHER THAN THE CURRENT BLOCK IN THE SYSCOM RECORD BUFFER, THE FILE MANAGER IS CALLED TO READ THE NEW BLOCK INTO THE SYSCOM RECORD BUFFER.

IF SUCCESSFUL, OK IS SET. IF AN ERROR OCCURRED, TELL THE USER, CLEAR OK AND TAKE ANY FURTHER APPROPRIATE ACTION REQUIRED.

THIS SUBROUTINE WILL EXIT WHEN EITHER:

- 1) THE TRANSACTION RECORD REQUESTED IS IN THE CURRENT SYSCOM RECORD BUFFER, OR
- 2) THE DISK TRANSACTION FILE HAS BEEN ACCESSED WITHOUT ERROR, OR
- 3) THE DISK TRANSACTION FILE HAS BEEN COMPLETELY ACCESSED.

00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600
00000610
00000620
00000630
00000640
00000650
00000660
00000670
00000680
00000690

```

14DI COMPJTE TRANSACTION BLOCK NUMBER/ENTRY OFFSET
14DI IF NEW BLOCK NUMBER, THEN
55DI DEQUEUE BLOCK BUFFER (GET.S8B)
14DI SAVE TRANSACTION CURRENT BLOCK NUMBER
21FM JDATE IPM FOR CALL
08FM CALL FILE MANAGER TO READ TRANSACTION FILE BY BLOCK NUMBER
SAVE FM RETURN DATA
56DI RELEASE BLOCK BUFFER (FREE.S8B)
21FM CASEENTRY (FJK, FBN, FIR, FCC, ? )
CASE FOK ( SUCCESSFUL )
10DI SET JK
CASE FBN ( BATCH NOT FOUND )
10DI CLEAR OK
13DI UPDATE ORDER COJNT BEYOND MAXIMJM
14DI SET TRANSACTION COUNT = ZERO
12DI
CASE FCC ( INVALID CALL CODE )
54DI PANIC
CASE FIR ( INVALID RELATIVE BLOCK )
54DI PANIC
CASE ? ( INVALID COMPLETION CODE )
54DI PANIC
ENDCASE
(ELSE)
ENDIF
RETURN

```

DELETE BATCH COMMAND

- MODULE HISTORY -

PROJECT: DAS 75-01709
SUB-SYSTEM: OPERATOR INTERFACE
MODULE: DIDL3 24

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190

- MODJLE ABSTRACT -

```

                                00000200
                                00000210
                                00000220
- THIS MODJLE ASKS THE USER IF ALL BATCHES TO BE DELETED. IF
ALL BATCHES NOT TO BE DELETED, A SPECIFIC BATCH ID IS REQUESTED,
AND THE FILE MANAGER IS CALLED TO READ THE BATCH FILE INFORM-
ATION.
                                00000230
                                00000240
                                00000250
                                00000260
                                00000270
                                00000280
                                00000290
                                00000300
                                00000310
                                00000320
                                00000330
                                00000340
                                00000350
                                00000360
                                00000370
                                00000380
                                00000390
                                00000400
                                00000410
                                00000420
                                00000430
                                00000440
                                00000450
                                00000460
                                00000470
                                00000480
                                00000490
                                00000500
                                00000510
                                00000520
                                00000530
                                00000540
                                00000550
                                00000560
                                00000570
                                00000580
                                00000590
                                00000600
                                00000610
                                00000620
                                00000630
                                00000640
                                00000650
                                00000660
                                00000670
                                00000680
                                00000690
                                00000700
                                00000710
                                00000720
                                00000730
                                00000740
                                00000750
                                00000760
                                00000770
                                00000780
                                00000790
                                00000800
                                00000810
                                00000820
                                00000830
                                00000840
                                00000850
                                00000860
                                00000870
                                00000880
                                00000890
                                00000900
                                00000910
040I  DISPLAY ON USER TERMINAL "DELETE BATCH COMMAND"
480I  DEQUEUE AN IPM BUFFER (#1) (GET.IPM)
480I  DEQUEUE AN IPM BUFFER (#2) (GET.IPM)
490I  DEQUEUE RECORD BUFFER (GET.SR3)
08AL  INITIALIZE PANIC CODE
360I
120I  INITIALIZE POINTER TO SYSCOM BATCH DIRECTORY
120I  INITIALIZE SYSCOM BATCH DIRECTORY ENTRY SIZE
08AL  INITIALIZE PANIC CODE
030I  ASK USER IF ALL BATCHES TO BE DELETED
360I  IF NOT ALL BATCHES, THEN
030I    ASK USER FOR BATCH ID
120I
160I    ACCESS BATCH FILE (GET.BATCH)
        IF THE BATCH IS ACTIVE, THEN
            SET THE ACTIVE FLAG
        ELSE
            IF THE BATCH IS ACTIVE OR DELETING, THEN
                CLEAR OK FLAG
            (ELSE)
            ENDIF
        ENDIF
100I    IF BATCH RECORD FOUND (OK SET), THEN
100I        SET SPECIFIC BATCH
        ELSE
020I    DISPLAY BATCH ID, "BATCH NOT ON FILE"
        ENDIF
    ELSE
        CLEAR THE USERS SCREEN
100I    SET OK
100I    CLEAR SPECIFIC BATCH
    ENDIF
100I    IF BATCH FOUND (OK SET), THEN
100I    ! CLEAR BATCH FOUND
120I    ! INITIALIZE BATCH COUNT
120I    ! DOUNTIL BATCH COUNT GREATER THAN MAXIMUM
100I    ! ! IF SPECIFIC BATCH SET, THEN
02FM    ! !   SAVE RELATIVE BATCH ID
100I    ! !   SET BATCH FOUND
    ! ! ELSE
160I    ! !   FIND AN ACTIVE, PENDING OR COMPLETE BATCH RECORD
    ! !   --- (FIND.BATCH) ---
    ! !   SAVE RELATIVE BATCH
    ! !   IF THE BATCH IS ACTIVE, THEN
    ! !       SET THE ACTIVE FLAG
    ! !       CLEAR OK FLAG
    ! !       SET BATCH FOUND FLAG
    ! !   (ELSE)
    ! !   ENDIF
    ! ! ENDIF
100I    ! ! IF BATCH RECORD FOUND (OK SET), THEN
100I    ! !   SET BATCH FOUND
    ! !   IF THIS IS THE FIRST BATCH, THEN
    ! !       DISPLAY BATCH HEADINGS
    ! !   (ELSE)
    ! !   ENDIF
430I    ! !   DISPLAY BATCH INFORMATION

```

```

! ! IF NOT ACTIVE OR ALL BATCHES, THEN 00000920
! ! IF NOT ACTIVE, THEN 00000930
030I ! ! ASK USER OK TO DELETE BATCH 00000940
! ! IF YES, THEN 00000950
100I ! ! SET DELETE FOUND 00000960
14PM ! ! STORE RELATIVE BATCH NUMBER IN IPM BUFFER (#2) 00000970
! ! (ELSE) 00000980
! ! ENDDIF 00000990
! ! (ELSE) 00010000
! ! ENDDIF 00010100
! ! ELSE 00010200
020I ! ! DISPLAY 'BATCH MUST BE PENDING OR COMPLETE' 00010300
! ! ENDDIF 00010400
! ! (ELSE) 00010500
! ! ENDDIF 00010600
100I ! ! IF SPECIFIC BATCH, THEN 00010700
120I ! ! UPDATE BATCH COUNT BEYOND MAXIMUM 00010800
! ! (ELSE) 00010900
! ! ENDDIF 00011000
! ! CLEAR THE ACTIVE BATCH FLAG 00011100
! ENDD0 00011200
100I ! IF DELETE FOUND, THEN 00011300
530I ! ! USE FILE MANAGER TO DELETE BATCH(ES) (DELETE.BATCH) 00011400
! ELSE 00011500
100I ! ! IF BATCH FOUND NOT SET, THEN 00011500
040I ! ! DISPLAY 'NO CURRENT BATCH INFORMATION' 00011700
! ! ELSE 00011800
! ! IF ALL OF THE BATCHES WERE ACTIVE, THEN 00011900
020I ! ! DISPLAY 'ALL BATCHES ARE ACTIVE' 00012000
! ! (ELSE) 00012100
! ! ENDDIF 00012200
! ! ENDDIF 00012300
! ENDDIF 00012400
(ELSE) 00012500
ENDDIF 00012600
510I RELEASE RECORD BUFFER (FREE.SR3) 00012700
500I RELEASE THE IPM BUFFER (#1) (FREE.IPM) 00012800
500I RELEASE THE IPM BUFFER (#2) (FREE.IPM) 00012900
DISPLAY 'BATCH DELETION ENDED' 00013000
DEQUEUE USER TERMINAL 00013100
ENDPRG 00013200

00000010
00000020
00000030
00000040
00000050
READ DISKETTE BATCH/TRANSACTION RECORD ( READ.BATCH ) 00000060
00000070
00000080
00000090
- MODULE HISTORY - 00000100
00000110
00000120
PROJECT: DAS 75-01709 00000130
00000140
SUB-SYSTEM: OPERATOR INTERFACE 00000150
00000160
MODULE: DICCRB 73 00000170
00000180
00000190
- MODULE ABSTRACT - 00000200
00000210
00000220
CALL THE DISKETTE SUBROUTINE TO READ A BATCH OR 00000230
TRANSACTION RECORD. 00000240
THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL 00000250
WHETHER DISKETTE ACCESS WAS SUCCESSFUL OR NOT. 00000260
CALL READ DISKETTE DATA BLOCK 00000270
CASENTRY (DOK, DID, JBD, DLR, DCC, ? ) 00000280
CASE DOK ( SUCCESSFUL ) 00000290
100I SET OK 00000300
CASE DID ( I/O ERROR OCCURRED ) 00000310
ENQUEUE USER TERMINAL 00000320

```

```

02JI      DISPLAY "DISKETTE I/O ERROR OCCURRED"
10JI      SET LAST RECORD
10JI      SET BATCH ABORT
10JI      CLEAR OK
          CASE DBD (INVALID RECORD COUNT)
            ENQUEUE USER TERMINAL
            DISPLAY "DISKETTE FORMAT ERROR"
            CLEAR OK FLAG
          CASE DLR (LAST RECORD - SUCCESSFUL )
            SET LAST RECORD
            SET OK
          CASE OCC (INVALID CALL CODE )
            PANIC
54JI      CASE ? (INVALID COMPLETION CODE )
            PANIC
54JI      ENDCASE
          RETURN
    
```

```

00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
    
```

PANIC

- MODULE HISTORY -

```

PROJECT:      DAS              75-01709
SJB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DICCPA          84
    
```

- MODULE ABSTRACT -

THIS SJBROUTINE PERFORMS THE PANIC. THE PANIC CODE IS RETRIEVED AND PUT ON THE ALARM QUEUE. A DEQUEUE TERMINAL IS PERFORMED TO GET RID OF ANY LEFT OVER ENQUEUES. THEN A WAIT FOR AN EVENT THAT WILL NEVER BE POSTED IS PERFORMED.

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
    
```

```

08AL      GET PANIC CODE
01AL      GET ADDRESS OF ALARM QUEUE IN SYSCOM
          DEQUEUE TERMINAL / PRINTER
07AL      PUT PANIC CODE ON THE ALARM QUEUE
          WAIT FOR NEVER EVENT
          RETURN
    
```

WRITE SORT TRANSLATION TABLE TO DISK

- MODULE HISTORY -

```

PROJECT:      DAS              75-01709
SJB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DICCWS          82
    
```

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
    
```

CHECKED: S.R. MARSDEN 15-JAN-79

- MODULE ABSTRACT -

THIS SUBROUTINE WRITES THE SORT TRANSLATION TABLE TO THE DISK. IT FIRST CLEARS THE 256 BYTE BUFFER AND THEN COPIES IN THE SORT TRANSLATION TABLE. THE BUFFER IS THEN WRITTEN TO THE DISK. IF ANY ERRORS ARE ENCOUNTERED A PANIC IS PERFORMED.

```
06FM   GET ADDRESS OF SORT TRANSLATION TABLE IN SYSCOM
        CLEAR WRITE BUFFER
06FM   COPY TRANSLATION TABLE IN SYSCOM TO THE BUFFER
06FM   WRITE THE BUFFER TO DISK
        IF ANY DISK ERRORS WERE ENCOUNTERED THEN
54JI   PANIC
        (ELSE)
        ENDIF
        RETURN
```

00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410

ALARM OUTPUT

- MODULE HISTORY -

```
PROJECT:   DAS           73-01709
SJB-SYSTEM: OPERATOR INTERFACE
MODULE:    DICCSA       83
```

- MODULE ABSTRACT -

THIS SUBROUTINE SENDS AN ALARM MESSAGE TO THE ALARM MESSAGE HANDLING SJB-SYSTEM. IT GETS A BUFFER, PUTS IN THE MESSAGE NUMBER AND DATA, AND PUTS THE ADDRESS OF THE BUFFER ON THE ALARM QUEUE. IF ANY ERRORS ARE ENCOUNTERED, PANICS ARE PERFORMED.

```
02AL   GET ADDRESS ALARM BUFFER QUEUE IN SYSCOM
02AL   DEQUEUE A BUFFER
02AL   IF ONE RECEIVED THEN
03AL   PUT IN MESSAGE NUMBER
03AL   PUT IN MESSAGE DATA
01AL   GET ADDRESS OF ALARM QUEUE IN SYSCOM
01AL   PUT ADDRESS OF BUFFER ON THE ALARM QUEUE
        IF THE ALARM QUEUE WAS ALREADY FULL THEN
54JI   PANIC
        (ELSE)
        ENDIF
    ELSE
        PANIC
    ENDIF
    RETURN
```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450

FREE A RECORD BUFFER

- MODULE HISTORY -

PROJECT: DAS 75-01709

SUB-SYSTEM: OPERATOR INTERFACE

MODULE: DICCFR 67

- MODULE ABSTRACT -

THIS SUBROUTINE PUTS THE 64 BYTE RECORD BUFFER BACK
ON THE BUFFER QUEUE IN SYSCOM. IF THE QUEUE IS FULL A
PANIC IS PERFORMED.

```

36FM GET ADDRESS OF SRB (64-BYTE) BUFFER QUEUE IN SYSCOM
36FM PUT ADDRESS OF BUFFER ON THE QUEUE
36FM IF THE QUEUE WAS ALREADY FULL THEN
54JI PANIC
      (ELSE)
      ENDIF
      RETURN

```

0000010
0000020
0000030
0000040
0000050
0000060
0000070
0000080
0000090
0000100
0000110
0000120
0000130
0000140
0000150
0000160
0000170
0000180
0000190
0000200
0000210
0000220
0000230
0000240
0000250

0000260
0000270
0000280
0000290
0000300
0000310
0000320
0000330
0000340
0000350

GET A RECORD BUFFER

- MODULE HISTORY -

PROJECT: DAS 75-01709

SUB-SYSTEM: OPERATOR INTERFACE

MODULE: DICCFR 66

- MODULE ABSTRACT -

THIS SUBROUTINE GETS A 64 BYTE RECORD BUFFER FROM
THE RECORD BUFFER QUEUE IN SYSCOM. IF THE QUEUE IS EMPTY
A PANIC IS PERFORMED.

```

36FM GET ADDRESS OF SRB (64-BYTE) BUFFER QUEUE IN SYSCOM
36FM DEQUEUE A BUFFER FROM THE QUEUE
36FM IF THE QUEUE WAS EMPTY THEN
54JI PANIC
      (ELSE)
      ENDIF
      RETURN

```

0000010
0000020
0000030
0000040
0000050
0000060
0000070
0000080
0000090
0000100
0000110
0000120
0000130
0000140
0000150
0000160
0000170
0000180
0000190
0000200
0000210
0000220
0000230
0000240
0000250

0000260
0000270
0000280
0000290
0000300
0000310
0000320
0000330
0000340
0000350

FREE AN INTER-PROGRAM MESSAGE BUFFER

- MODULE HISTORY -

PROJECT: DAS 73-01709
 SUB-SYSTEM: OPERATOR INTERFACE
 MODULE: DICCFI 65

- MODULE ABSTRACT -

THIS SUBROUTINE PUTS AN INTER-PROGRAM MESSAGE BUFFER
 BACK ON THE IPM QUEUE. IF THE QUEUE IS FULL A PANIC IS
 PERFORMED.

09FM GET ADDRESS OF IPM BUFFER QUEUE IN SYSCOM
 09FM PUT ADDRESS OF BUFFER ON THE QUEUE
 09FM IF THE QUEUE WAS ALREADY FULL THEN
 54JI PANIC
 (ELSE)
 ENDIF
 RETURN

GET AN INTER-PROGRAM MESSAGE BUFFER

- MODULE HISTORY -

PROJECT: DAS 73-01709
 SUB-SYSTEM: OPERATOR INTERFACE
 MODULE: DICCGI 64

- MODULE ABSTRACT -

THIS SUBROUTINE DEQUEUES AN INTER-PROGRAM MESSAGE
 BUFFER FROM THE IPM QUEUE. IF THE QUEUE IS EMPTY A
 PANIC IS PERFORMED.

09FM GET ADDRESS OF IPM BUFFER QUEUE IN SYSCOM
 09FM DEQUEUE A BUFFER FROM THE QUEUE
 09FM IF THE QUEUE WAS EMPTY THEN
 54JI PANIC
 (ELSE)
 ENDIF
 RETURN

0000010
 0000020
 0000030
 0000040
 0000050
 0000060
 0000070
 0000080
 0000090
 0000100
 0000110
 0000120
 0000130
 0000140
 0000150
 0000160
 0000170
 0000180
 0000190
 0000200
 0000210
 0000220
 0000230
 0000240
 0000250

0000260
 0000270
 0000280
 0000290
 0000300
 0000310
 0000320
 0000330
 0000340
 0000350

0000010
 0000020
 0000030
 0000040
 0000050
 0000060
 0000070
 0000080
 0000090
 0000100
 0000110
 0000120
 0000130
 0000140
 0000150
 0000160
 0000170
 0000180
 0000190
 0000200
 0000210
 0000220
 0000230
 0000240
 0000250

0000260
 0000270
 0000280
 0000290
 0000300
 0000310
 0000320
 0000330
 0000340
 0000350


```

WRITE DISKETTE BATCH/TRANSACTION RECORD ( WRITE.BATCH )
                                00000010
                                00000020
                                00000030
                                00000040
                                00000050
                                00000060
                                00000070
                                00000080
                                00000090
                                00000100
                                00000110
                                00000120
                                00000130
                                00000140
                                00000150
                                00000160
                                00000170
                                00000180
                                00000190
                                00000200
                                00000210
                                00000220
                                00000230
                                00000240
                                00000250
                                00000260

                                - MODULE HISTORY -
PROJECT:           DAS           75-01709
SUB-SYSTEM:       OPERATOR INTERFACE
MODULE:           JICCB         75
                                00000180
                                00000190
                                00000200
                                00000210
                                00000220
                                00000230
                                00000240
                                00000250
                                00000260

                                CALL THE DISKETTE SUBROUTINE TO WRITE A BATCH OR
                                TRANSACTION RECORD.
                                THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL
                                WHETHER DISKETTE ACCESS WAS SUCCESSFUL OR NOT.
                                00000270
                                00000280
                                00000290
                                00000300
                                00000310
                                00000320
                                00000330
                                00000340
                                00000350
                                00000360
                                00000370
                                00000380
                                00000390
                                00000400
                                00000410
                                00000420

CALL WRITE DISKETTE DATA BLOCK
CASENTRY (DOK, DIO, DCC, ? )
100I   CASE DOK ( SUCCESSFUL )
        SET OK
        CASE DIO ( I/O ERROR OCCURRED )
020I   ENQUEUE USER TERMINAL
140I   DISPLAY "DISKETTE I/O ERROR OCCURRED"
100I   UPDATE TRANSACTION COUNT BEYOND MAXIMUM
100I   SET BATCH ABORT
        CLEAR OK
        CASE DCC ( INVALID CALL CODE )
540I   PANIC
540I   CASE ? ( INVALID COMPLETION CODE )
        PANIC
ENDCASE
RETURN
                                00000410
                                00000420

                                00000010
                                00000020
                                00000030
                                00000040
                                00000050
                                00000060
                                00000070
                                00000080
                                00000090
                                00000100
                                00000110
                                00000120
                                00000130
                                00000140
                                00000150
                                00000160
                                00000170
                                00000180
                                00000190
                                00000200
                                00000210
                                00000220
                                00000230
                                00000240
                                00000250
                                00000260
                                00000270
                                00000280
                                00000290
                                00000300

ASSIGN TRANSACTION STATUS COMMAND
                                00000010
                                00000020
                                00000030
                                00000040
                                00000050
                                00000060
                                00000070
                                00000080
                                00000090
                                00000100
                                00000110
                                00000120
                                00000130
                                00000140
                                00000150
                                00000160
                                00000170
                                00000180
                                00000190
                                00000200
                                00000210
                                00000220
                                00000230
                                00000240
                                00000250
                                00000260
                                00000270
                                00000280
                                00000290
                                00000300

                                - MODULE HISTORY -
PROJECT:           DAS           75-01709
SUB-SYSTEM:       OPERATOR INTERFACE
MODULE:           DIATS        28
                                00000180
                                00000190
                                00000200
                                00000210
                                00000220
                                00000230
                                00000240
                                00000250
                                00000260
                                00000270
                                00000280
                                00000290
                                00000300

                                THIS MODULE REQUESTS A BATCH ID. THE FILE MANAGER IS CALLED
                                TO READ THE BATCH FILE INFORMATION. IF THE BATCH INFORMATION IS
                                FOUND, THE USER IS ASKED FOR START AND END RANGE OF TRANSACTION
                                ID'S. EACH IS CHECKED FOR NON-ZERO VALUE,
                                AND THE END NUMBER GREATER THAN THE START NUMBER. IF TRANS-
                                ACTION VALUES ARE ACCEPTABLE, THE USER IS ASKED FOR THE NEW
                                TRANSACTION STATUS.
                                THE FILE MANAGER IS CALLED TO CHANGE TRANSACTION RECORD STATUS.

```

```

ENQJEJE USER TERMINAL
04JI DISPLAY ON USER TERMINAL "ASSIGN TRANSACTION STATUS COMMAND"
48JI DEQJEUE AN IPM BUFFER (GET.IPM)
49JI DEQJEUE RECORD BUFFER (GET.SRB)
03JI ASK USER FOR BATCH ID
12JI
08AL INITIALIZE PANIC CODE
40JI
16JI ACCESS BATCH FILE (GET.BATCH)
01FM
10JI IF BATCH RECORD FOUND (OK SET), THEN
43JI ! DISPLAY BATCH HEADER
43JI ! DISPLAY BATCH INFORMATION
03JI ! ASK USER FOR START TRANSACTION NUMBER
40JI !
40JI ! IF INPUT TRANS ID WITHIN RANGE THEN
03JI ! ! ASK USER FOR END TRANSACTION NUMBER
40JI ! ! IF INPJT TRANS ID WITHIN RANGE AND GT. OR EQ. START ID. --
! ! -- THEN
03JI ! ! ! DJ UNTIL VALID STATUS RECEIVED
40JI ! ! !
16FM ! ! ! VERIFY STATUS CODE BY RANGE CHECK
40JI ! ! ! ENDDJ
! ! ! ASK USER IF ALL PARAMETERS CORRECT
! ! ! IF YES THEN
08AL ! ! ! ! INITIALIZE PANIC CODE
! ! ! ! DISPLAY "PROCESSING STATUS CHANGES"
40JI ! ! ! !
40JI ! ! ! ! DJ UNTIL START TRANS ID GREATER THAN END TRANS ID --
! ! ! ! -- OR ERROR FOUND
! ! ! ! DELAY
16JI ! ! ! ! JDATE IPM FOR CALL
08FM ! ! ! ! CALL FILE MANAGER TO CHANGE TRANSACTION STATUS
16FM ! ! ! ! CASENTRY (FOK, FNB, FTN, FCC, FIS, ? )
! ! ! ! CASE FOK ( SUCCESSFUL )
! ! ! ! CASE FNB ( TRANSACTION NOT IN BATCH )
02JI ! ! ! ! DISPLAY "TRANS. NOT IN BATCH"
! ! ! ! SET ERROR FOUND
14JI ! ! ! ! SUB 1 FROM TRANS. ID
! ! ! ! CASE FTN ( TRANSACTION NOT ON FILE )
! ! ! ! DISPLAY "TRANS. NOT ON FILE"
! ! ! ! SET ERROR FOUND
14JI ! ! ! ! SUB 1 FROM TRANS. ID
! ! ! ! CASE FCC ( INVALID CALL CODE )
54JI ! ! ! ! PANIC
! ! ! ! CASE FIS ( INVALID STATUS )
54JI ! ! ! ! PANIC
! ! ! ! CASE ? ( INVALID COMPLETION CODE )
54JI ! ! ! ! PANIC
! ! ! ! ENDCASE
14JI ! ! ! ! INCREMENT TRANSACTION ID
! ! ! ! ENDDJ
! ! ! ! IF ANY TRANSACTIONS CHANGED THEN
! ! ! ! DISPLAY TRANS. THAT WERE CHANGED
! ! ! ! ELSE
! ! ! ! ! DISPLAY "NO CHANGE MADE"
! ! ! ! ! ENDIF
! ! ! ! (ELSE)
! ! ! ! ! ENDF
! ! ! ! ELSE
! ! ! ! ! IF END ID LESS THAN START ID THEN
! ! ! ! ! DISPLAY "END ID MUST BE GREATER THAN OR --
! ! ! ! ! -- EQUAL TO THE START ID
! ! ! ! ! ELSE
! ! ! ! ! DISPLAY "INVALID TRANS. ID"
! ! ! ! ! ENDF
! ! ! ! ENDF
! ! ! ! ELSE
! ! ! ! ! DISPLAY "INVALID TRANS. ID"
! ! ! ! ! ENDF
! ! ! ! ! DISPLAY "STATUS CHANGE COMMAND ENDED"
ELSE
02JI ! ! ! ! DISPLAY BATCH ID, "BATCH NOT ON FILE"

```

```

00000310
00000320
00000330
00000340
00000350
00000350
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600
00000610
00000620
00000630
00000640
00000650
00000660
00000670
00000680
00000690
00000700
00000710
00000720
00000730
00000740
00000750
00000760
00000770
00000780
00000790
00000800
00000810
00000820
00000830
00000840
00000850
00000860
00000870
00000880
00000890
00000900
00000910
00000920
00000930
00000940
00000950
00000960
00000970
00000980
00000990
00010000
00010100
00010200
00010300

```

```

ENDIF
51JI  RELEASE RECORD BUFFER (FREE.SRB)
50JI  RELEASE THE IPM BUFFER (FREE.IPM)
      DEQUEUE TERMINAL
      ENDPRG

```

RECIRC. ON SCANNER ERROR

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DIREC            45

```

- MODULE ABSTRACT -

THIS MODULE SETS A SWITCH VALUE IN SYSCOM WHICH WHEN INTERROGATED BY SCANNER INPUT SOFTWARE, SCANNER ERRORS WILL BE RECIRCULATED.

```

04JI  ENQUEUE USER TERMINAL
11SC  DISPLAY ON USER TERMINAL "RECIRC. ON SCANNER ERROR"
02AL  SET SYSCOM FLAG VALUE
02AL  SET AN ALARM BUFFER
02AL  IF ONE RECEIVED, THEN
03AL  PUT IN MESSAGE NUMBER
01AL  PUT ADDRESS OF THE BUFFER IN THE ALARM QUEUE
01AL  IF THERE IS NO ROOM, THEN
07AL  PANIC
      (ELSE)
      ENDIF
ELSE
07AL  PANIC
ENDIF
DEQUEUE TERMINAL
ENDPRG

```

GET DISK ORDER FILE RECORD BY ORDER ID (GET.ORDER)

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DICCSO           77

```

- MODULE ABSTRACT -

PROCESS: CALL THE FILE MANAGER TO READ THE DISK ORDER FILE RECORD. IF SUCCESSFUL, COPY THE ORDER INFORM-

```

00001040
00001050
00001060
00001070
00001080
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240

```

ATION FROM THE SYSCOM RECORD BUFFER TO THE DATA AREAS 00000250
 IN THE PROGRAM AND SET OK. IF AN ERROR OCCURRED, TELL 00000260
 THE USER, CLEAR JK AND TAKE ANY FURTHER APPROPRIATE 00000270
 ACTION REQUIRED. 00000280

THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL TO THE 00000290
 FILE MANAGER WHETHER ORDER FILE ACCESS WAS SUCCESSFUL 00000300
 OR NOT. 00000310

22FM JDATE IPM FOR CALL 00000320
 08FM CALL FILE MANAGER TO READ ORDER FILE RECORD BY ORDER ID 00000330
 22FM CASENTY (FJK, FBN, FCC, FDN, ?) 00000340
 CASE FJK (SUCCESSFUL) 00000350
 GET ORDER INFORMATION 00000360
 05FM 00000370
 13JI SET JK 00000380
 10JI CASE FBN (BATCH NOT FOUND) 00000390
 02JI PRINT/DISPLAY BATCH ID, "BATCH NO LONGER ON FILE" 00000400
 10JI CLEAR OK 00000410
 CASE FCC (INVALID CALL CODE) 00000420
 54JI PANIC 00000430
 CASE FDN (ORDER NOT ON FILE) 00000440
 02JI PRINT/DISPLAY ORDER ID, "ORDER ID NOT ON FILE" 00000450
 10JI CLEAR OK 00000460
 CASE ? (INVALID COMPLETION CODE) 00000470
 54JI PANIC 00000480
 ENDCASE 00000490
 RETURN 00000500
 00000010
 00000020
 00000030
 00000040
 00000050

FIND RELATIVE DISK ORDER FILE RECORD IN BATCH (FIND.ORDER)

- MODULE HISTORY -

PROJECT: DAS 75-01709
 SJB-SYSTEM: OPERATOR INTERFACE
 MODULE: JICCFD 76

- MODULE ABSTRACT -

EACH RELATIVE ORDER RECORD IS SEQUENTIALLY ACCESSED
 ONE AT A TIME BY CALLING THE FILE MANAGER FOR EACH
 RELATIVE ORDER NUMBER. IF SUCCESSFUL, COPY THE ORDER
 INFORMATION FROM THE SYSCOM RECORD BUFFER TO THE DATA
 AREAS IN THE PROGRAM AND SET OK. IF AN ERROR OCCURRED,
 TELL THE USER, CLEAR OK AND TAKE ANY FURTHER APPRO-
 PRIATE ACTION REQUIRED.

THIS SUBROUTINE WILL EXIT WHEN EITHER:

- 1) THE DISK ORDER FILE HAS BEEN ACCESSED WITHOUT ERROR, OR
- 2) THE DISK ORDER FILE HAS BEEN COMPLETELY ACCESSED.

13JI DDUNTIL END OF ORDER COUNT OR JK SET 00000340
 12JI 00000350
 10JI 00000360
 23FM JDATE IPM FOR CALL 00000370
 08FM CALL FILE MANAGER TO READ ORDER FILE RECORD BY ORDER COUNT 00000380
 23FM CASENTY (FJK, FBN, FCC, FIR, ?) 00000390
 CASE FJK (SUCCESSFUL) 00000400
 GET ORDER INFORMATION 00000410
 05FM 00000420
 13JI SET JK 00000430
 10JI CASE FBN (BATCH NOT FOUND) 00000440
 13JI UPDATE ORDER COUNT BEYOND MAXIMUM 00000450
 12JI 00000460

```

10JI      CLEAR OK                                00000470
          CASE FCC ( INVALID CALL CODE )         00000480
54JI      PANIC                                  00000490
          CASE FIR ( INVALID RELATIVE RECORD )  00000500
54JI      PANIC                                  00000510
          CASE ? ( INVALID COMPLETION CODE )    00000520
54JI      PANIC                                  00000530
          ENDCASE                                00000540
10JI      INCREMENT ORDER COUNT                 00000550
          ENDDJ                                  00000560
          RETURN                                 00000570

```

CHANGE DISK BATCH FILE RECORD STATUS (CHANGE.BATCH)

- MODULE HISTORY -

```

PROJECT:      DAS                                73-01709
SJB-SYSTEM:   OPERATOR INTERFACE
MODULE:       JICCCB                             72

```

- MODULE ABSTRACT -

```

CALL THE FILE MANAGER TO CHANGE THE DISK BATCH
FILE RECORD.
IF AN ERROR OCCURRED, TELL THE USER, CLEAR OK AND
TAKE ANY FURTHER APPROPRIATE ACTION REQUIRED.
THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL TO THE
FILE MANAGER WHETHER BATCH FILE ACCESS WAS SUCCESSFUL
OR NOT.

```

```

15FM      JUPDATE IPM FOR CALL                    00000300
08FM      CALL FILE MANAGER TO CHANGE BATCH STATUS 00000310
15FM      CASENTRY (FOK, FBN, FCC, FIS, ? )       00000320
          CASE FJK ( SUCCESSFUL )                 00000330
10JI      SET OK                                  00000340
          CASE FBN ( BATCH NOT FOUND )           00000350
          OUTPUT BATCH NOT ON FILE MESSAGE       00000360
          CLEAR OK                                00000370
          CASE FCC ( INVALID CALL CODE )         00000380
54JI      PANIC                                  00000390
          CASE FIS ( INVALID STATUS )           00000400
54JI      PANIC                                  00000410
          CASE ? ( INVALID COMPLETION CODE )    00000420
54JI      PANIC                                  00000430
          ENDCASE                                00000440
          RETJRN                                 00000450

```

READ DISK BATCH FILE RECORD BY BATCH ID (GET.BATCH)

- MODULE HISTORY -

```

PROJECT:      DAS                                73-01709
SJB-SYSTEM:   OPERATOR INTERFACE

```

```

00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000600
00000610
00000620
00000630
00000640
00000650
00000660
00000670
00000680
00000690
00000700
00000710
00000720
00000730
00000740
00000750
00000760
00000770
00000780
00000790
00000800
00000810
00000820
00000830
00000840
00000850
00000860
00000870
00000880
00000890
00000900
00000910
00000920
00000930
00000940
00000950
00000960
00000970
00000980
00000990

```

MODULE: D1CCGB

71

- MODULE ABSTRACT -

CALL THE FILE MANAGER TO READ THE DISK BATCH FILE RECORD. IF SUCCESSFUL, COPY THE BATCH INFORMATION FROM THE SYSCOM RECDR BUFFER TO THE DATA AREAS IN THE PROGRAM AND SET OK. IF AN ERROR OCCURRED, TELL THE USER, CLEAR OK AND TAKE ANY FURTHER APPROPRIATE ACTION REQUIRED.

THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL TO THE FILE MANAGER WHETHER BATCH FILE ACCESS WAS SUCCESSFUL OR NOT.

```

18FM  UPDATE IPM FOR CALL
08FM  CALL FILE MANAGER TO READ BATCH FILE RECDR BY BATCH ID
18FM  CASENTRY (FOK, FBN, FCC, ? )
      CASE FOK ( SUCCESSFUL )
01FM  GET BATCH INFORMATION
120I
100I  SET OK
      CASE FBN ( BATCH NOT FOUND )
100I  CLEAR OK
      CASE FCC ( INVALID CALL CODE )
540I  PANIC
      CASE ? ( INVALID COMPLETION CODE )
540I  PANIC
ENDCASE
RETURN

```

FIND CURRENT BATCH IN SYSCOM BATCH DIRECTORY (FIND.BATCH)

- MODULE HISTORY -

```

PROJECT:      DAS                70-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       D1CCFB             70

```

- MODULE ABSTRACT -

EACH SYSCOM BATCH DIRECTORY IS EXAMINED ONE AT A TIME. IF THE ENTRY CONTAINS A PENDING, ACTIVE OR COMPLETE STATUS; THE DISK BATCH FILE WILL BE ACCESSED.

THIS SUBROUTINE WILL EXIT WHEN EITHER:

- 1) THE DISK BATCH FILE HAS BEEN ACCESSED WITHOUT ERROR, OR
- 2) THE BATCH DIRECTORY HAS BEEN COMPLETELY EXAMINED.

```

120I  DDUNTIL BATCH COUNT GREATER THAN MAXIMUM OR OK SET
100I
02FM  EXAMINE SYSCOM BATCH DIRECTORY ENTRY
02FM  IF BATCH PENDING, ACTIVE OR COMPLETE, THEN
160I  ACCESS BATCH FILE (GET.BATCH)
      ELSE
100I  CLEAR OK
      ENDIF
120I  INCREMENT POINTER TO SYSCOM BATCH DIRECTORY
120I  INCREMENT BATCH COUNT
ENDDD
RETURN

```

```

00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420

```

READ DISK TRANSACTION FILE RECORD BY TRANS. ID (GET.TRANS)

- MODULE HISTORY -

PROJECT: DAS 75-01709

SUB-SYSTEM: OPERATOR INTERFACE

MODULE: JICCGT 80

- MODULE ABSTRACT -

CALL THE FILE MANAGER TO READ THE TRANSACTION FILE RECORD. IF SUCCESSFUL, COPY THE TRANSACTION DATA FROM THE SYSCOM RECORD BUFFER TO THE DATA AREAS IN THE PROGRAM AND SET OK. IF AN ERROR OCCURRED, TELL THE USER, CLEAR OK AND TAKE ANY FURTHER APPROPRIATE ACTION REQUIRED.

THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL TO THE FILE MANAGER WHETHER TRANSACTION FILE ACCESS WAS SUCCESSFUL OR NOT.

19FM JDATE IPM FOR CALL
 08FM CALL FILE MANAGER TO READ TRANSACTION FILE RECORD
 19FM CASENTRY (FOK, FTN, FCC, ?)
 CASE FOK (SUCCESSFUL)
 GET TRANSACTION INFORMATION
 04FM SET OK
 14JI CASE FTN (TRANSACTION NOT FOUND)
 10JI CLEAR OK
 10JI CASE FCC (INVALID CALL CODE)
 54OI PANIC
 54OI CASE ? (INVALID COMPLETION CODE)
 PANIC
 ENDCASE
 RETURN

HELP COMMAND (OPERATOR COMMAND LIST DISPLAY)

- MODULE HISTORY -

PROJECT: DAS 75-01709

SUB-SYSTEM: OPERATOR INTERFACE

MODULE: DIHLP 53

- MODULE ABSTRACT -

THIS MODULE DISPLAYS ALL OPERATOR INTERFACE COMMANDS ON THE USER TERMINAL.

00000010
 00000020
 00000030
 00000040
 00000050
 00000060
 00000070
 00000080
 00000090
 00000100
 00000110
 00000120
 00000130
 00000140
 00000150
 00000160
 00000170
 00000180
 00000190
 00000200
 00000210
 00000220
 00000230
 00000240
 00000250
 00000260
 00000270
 00000280
 00000290
 00000300
 00000310
 00000320
 00000330
 00000340
 00000350
 00000360
 00000370
 00000380
 00000390
 00000400
 00000410
 00000420
 00000430
 00000440
 00000450
 00000460
 00000010
 00000020
 00000030
 00000040
 00000050
 00000060
 00000070
 00000080
 00000090
 00000100
 00000110
 00000120
 00000130
 00000140
 00000150
 00000160
 00000170
 00000180
 00000190
 00000200
 00000210
 00000220
 00000230
 00000240

```

04JI ENQUEUE USER TERMINAL
47JI DISPLAY ON USER TERMINAL "HELP COMMAND"
      DISPLAY OPERATOR COMMAND LIST ON USER TERMINAL
      DEQUEUE TERMINAL
      ENDPROG

```

ADD DISK BATCH FILE RECORD (ADD.BATCH)

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DICCB           74

```

- MODULE ABSTRACT -

CALL THE FILE MANAGER TO ADD THE DISKETTE BATCH RECORD.

THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL TO THE FILE MANAGER WHETHER BATCH FILE ACCESS WAS SUCCESSFUL OR NOT.

```

00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270

```

```

12FM UPDATE IPM FOR CALL
08FM CALL FILE MANAGER TO ADD BATCH FILE RECORD
12FM CASENTRY ( FOK, FBX, FBA, FAD, FCC, ? )
      CASE FOK ( SUCCESSFUL )
10JI   SET OK
      CASE FBX ( BATCH FILE SPACE EXHAUSTED )
02JI   DISPLAY "BATCH FILE SPACE EXHAUSTED"
10JI   CLEAR OK
      CASE FBA ( BATCH ALREADY ON FILE )
02JI   DISPLAY BATCH ID, "BATCH ALREADY ON FILE"
10JI   CLEAR OK
      CASE FAD ( BATCH ALREADY BEING ADDED )
54JI   PANIC
      CASE FCC ( INVALID CALL CODE )
54JI   PANIC
      CASE ? ( INVALID COMPLETION CODE )
54JI   PANIC
ENDCASE
RETURN

```

```

00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460

```

RESTORE SHIPPING LINE COMMAND

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DIRSL           26

```

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170

```


- MODULE ABSTRACT -

```

THIS MODULE REQUESTS SORT LINE NUMBER TO BE RESTORED.
THE SORT LINE NUMBER IS VALIDATED BY RANGE CHECK. IF THE
SORT LINE NUMBER IS ACCEPTABLE, THE RESIDENT SORT TRANSLA-
TION TABLE IS UPDATED AND THE FILE MANAGER IS CALLED TO WRITE
THE RESIDENT SORT TRANSLATION TABLE TO DISK.
ENQUEUE USER TERMINAL
DISPLAY ON USER TERMINAL "RESTORE SORT LINE COMMAND"
ASK USER FOR OLD SORT LINE NUMBER
RANGE CHECK OLD SORT LINE NUMBER
IF NOT IN RANGE OR RECIRC OR ERROR CHUTE, THEN
  DISPLAY "INVALID SORT LINE NUMBER" OLD LINE NUMBER
  DISPLAY "COMMAND ABORTED"
ELSE
  IF SORT LINE HAS BEEN REROUTED, THEN
    UPDATE RESIDENT SORT TRANSLATION TABLE
    DISPLAY "SORT LINE ASSIGNED"
    INITIALIZE PANIC CODE
  WRITE SORT TRANSLATION TABLE TO DISK (WRITE.SORT)
ELSE
  DISPLAY "NO PRIOR ASSIGNMENT OF SORT LINE"
  DISPLAY "COMMAND ABORTED"
ENDIF
ENDIF
DEQUEUE TERMINAL
ENDPRG

```

ASSIGN SHIPPING LINE COMMAND

- MODULE HISTORY -

```

PROJECT:      DAS          73-01709
SJB-SYSTEM4:  OPERATOR INTERFACE
MODULE:       DIASL        25

```

- MODULE ABSTRACT -

```

THIS MODULE REQUESTS SORT LINE NUMBER TO BE (RE)ASSIGNED.
THE SORT LINE NUMBER IS VALIDATED BY RANGE CHECK. IF THE OLD
SORT LINE NUMBER IS ACCEPTABLE, THE USER IS ASKED FOR THE NEW
SORT LINE NUMBER. THIS ALSO IS VALIDATED BY RANGE CHECK. IF THE
NEW SORT LINE NUMBER IS ACCEPTABLE, THE RESIDENT SORT TRANSLA-
TION TABLE IS UPDATED AND THE FILE MANAGER IS CALLED TO WRITE
THE RESIDENT SORT TRANSLATION TABLE TO DISK.
ENQUEUE USER TERMINAL
CLEAR ABORT
GET SORT SYSTEM PARAMETERS
DISPLAY ON USER TERMINAL "REROUTE SORT LINE COMMAND"
ASK USER FOR OLD SORT LINE NUMBER
RANGE CHECK OLD SORT LINE NUMBER
IF NOT IN RANGE OR RECIRC OR ERROR CHUTE, THEN
  DISPLAY "INVALID SORT LINE NUMBER" OLD LINE NUMBER
  SET ABORT

```

```

0000180
0000190
0000200
0000210
0000220
0000230
0000240
0000250
0000260
0000270
0000280
0000290
0000300
0000310
0000320
0000330
0000340
0000350
0000360
0000370
0000380
0000390
0000400
0000410
0000420
0000430
0000440
0000450
0000460
0000470
0000480
0000490
0000010
0000020
0000030
0000040
0000050
0000060
0000070
0000080
0000090
0000100
0000110
0000120
0000130
0000140
0000150
0000160
0000170
0000180
0000190
0000200
0000210
0000220
0000230
0000240
0000250
0000260
0000270
0000280
0000290
0000300
0000310
0000320
0000330
0000340
0000350
0000360
0000370
0000380
0000390

```

```

ELSE
06FM   GET ENTRY FROM SORT TRANSLATION TABLE
      IF ALREADY RE-ASSIGNED, THEN
      TELL USER IT IS ALREADY ASSIGNED
      ASK IF WE SHOULD CONTINUE
      IF THE ANSWER IS "NO", THEN
        SET ABORT
      (ELSE)
      ENDIF
      (ELSE)
      ENDIF
      IF ABORT IS NOT SET, THEN
03JI   ASK USER FOR NEW SORT LINE NUMBER
37JI
37JI   RANGE CHECK NEW SORT LINE NUMBER
      IF NOT IN RANGE
        SET ABORT
02JI   DISPLAY "INVALID SORT LINE NUMBER" NEW LINE NUMBER
      ELSE
06FM   IF OLD SORT LINE EQ NEW SORT LINE, THEN
02JI   DISPLAY "ATTEMPT TO ASSIGN SORT LINE TO ITSELF"
      DISPLAY "COMMAND ABORTED"
      ELSE
37JI   UPDATE RESIDENT SORT TRANSLATION TABLE
04JI   DISPLAY "SORT LINE ASSIGNED"
08AL   INITIALIZE PANIC CODE
37JI
46JI   WRITE SORT TRANSLATION TABLE TO DISK (WRITE.SORT)
      ENDIF
      ENDIF
      ENDIF
      ELSE
      DISPLAY "COMMAND ABORTED"
      ENDIF
      DEQUEUE TERMINAL
      ENDPRG

```

```

00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600
00000610
00000620
00000630
00000640
00000650
00000660
00000670
00000680
00000690
00000700
00000710
00000720
00000730
00000740
00000750

```

OPERATOR INTERFACE CONTROL (ATTENTION LIST)

- MODULE HISTORY -

```

PROJECT:      DAS          75-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       DICTL        14

```

-MODULE ABSTRACT-

THIS MODULE CONTAINS THE ATTENTION LIST PROCESSING. THE USER WILL HIT THE ATTENTION KEY, THEN TYPE THE OPERATOR COMMAND FOLLOWED BY THE RETURN KEY. IN THE ATTENTION LIST PROCESSING THE PROGRAM WILL CALL A SUBROUTINE (TASK SELECT) TO SELECT ONE OF FOUR LOADING TASKS, DEPENDING ON WHICH ONES ARE NOT ACTIVE AT THE TIME. THE NUMBER ASSOCIATED WITH THE COMMAND TYPED WILL BE PASSED TO THE LOADING TASK. THEN AN END ATTENTION IS PERFORMED. THE ATTACHED TASK WILL THEN ENQUEUE A LOADING SUBROUTINE (LOADIT) AND PASS THE COMMAND NUMBER. THE SUBROUTINE WILL THEN LOAD THE REQUIRED PROGRAM.

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330

```

```

26JI   DDUNTIL FOREVER
      WAIT FOR USER TO ENTER COMMAND

```

```

00000340
00000350

```

```

0101  CASEENTRY (COMMAND-1,COMMAND-2,...,COMMAND-N)
      CASE COMMAND-1
        CALL TASK SELECT AND PASS COMMAND NUMBER
      CASE COMMAND-2
        CALL TASK SELECT AND PASS COMMAND NUMBER
      :
      :
      CASE COMMAND-N
        CALL TASK SELECT AND PASS COMMAND NUMBER
    ENDCASE
  ENDDO
ENDPRG

*****

* FORMAT OF SUBROUTINE TO ATTACH LOADING TASK *

IF FIRST LOADING TASK IS NOT ACTIVE, THEN
  PASS COMMAND NUMBER AND ATTACH THE TASK
ELSE
  IF SECOND LOADING TASK IS NOT ACTIVE, THEN
    PASS COMMAND NUMBER AND ATTACH THE TASK
  ELSE
    IF THIRD LOADING TASK IS NOT ACTIVE, THEN
      PASS COMMAND NUMBER AND ATTACH THE TASK
    ELSE
      PASS COMMAND NUMBER AND ATTACH THE FOURTH TASK
    ENDIF
  ENDIF
ENDIF
RETURN

*****

* GENERAL FORMAT OF TASKS TO CALL LOADING SUBROUTINE. *

ENQUEUE USER TERMINAL
ENQUEUE SUBROUTINE (LOAD IT)
CALL SUBROUTINE (LOAD IT) AND PASS COMMAND NUMBER
DEQUEUE SUBROUTINE
DEQUEUE USER TERMINAL
ENDTASK

*****

* FORMAT OF LOADING SUBROUTINE (LOAD IT) *

CASEENTRY (COMMAND-?,COMMAND-1,COMMAND-2,...,COMMAND-N)
  CASE COMMAND-?
    PANIC
  CASE COMMAND-1
    LOAD COMMAND-1 PROGRAM
  CASE COMMAND-2
    LOAD COMMAND-2 PROGRAM
  :
  :
  CASE COMMAND-N
    LOAD COMMAND-N PROGRAM
  ENDCASE
  GET LOAD CONDITION CODE
  IF CONDITION CODE NOT EQUAL TO 'NO ERRORS', THEN
    IF CONDITION CODE NOT EQUAL TO 'NO ROOM FOR PROGRAM', AND,
      CONDITION CODE NOT EQUAL TO 'NO ROOM FOR LOADER', THEN
      DISPLAY '*01* PROGRAM LOAD ERROR'
      DISPLAY ERROR CODE
    ELSE
      DISPLAY '**** RETRY COMMAND LATER'
    ENDIF
  (ELSE)
  ENDIF
  RETURN

```

07AL

```

00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600
00000610
00000620
00000630
00000640
00000650
00000660
00000670
00000680
00000690
00000700
00000710
00000720
00000730
00000740
00000750
00000760
00000770
00000780
00000790
00000800
00000810
00000820
00000830
00000840
00000850
00000860
00000870
00000880
00000890
00000900
00000910
00000920
00000930
00000940
00000950
00000960
00000970
00000980
00000990
00010000
00010100
00010200
00010300
00010400
00010500
00010600
00010700

```

NOTE - FOR COMMANDS 'READ BATCH' AND 'WRITE BATCH', THE DISKETTE IN USE FLAG MUST BE CHECKED BEFORE THE CORRESPONDING PROGRAM IS LOADED. IF THE FLAG IS ALREADY SET, 'DISKETTE ALREADY IN USE' IS DISPLAYED. OTHERWISE, IF THERE WERE NO LOAD ERRORS, THE FLAG IS THEN SET.

HOLD ON SCANNER ERROR

- MODULE HISTORY -

PROJECT: DAS 73-01709
 SUB-SYSTEM: OPERATOR INTERFACE
 MODULE: DI-ILT 46

- MODULE ABSTRACT -

THIS MODULE SETS A SWITCH VALUE IN SYSCOM WHICH WHEN INTERROGATED BY SCANNER INPUT SOFTWARE, SCANNER ERRORS WILL BE HALTED.

```

040I ENQUEUE USER TERMINAL
115C DISPLAY ON USER TERMINAL "HOLD ON SCANNER ERROR"
02AL SET SYSCOM FLAG VALUE
02AL GET AN ALARM BUFFER
02AL IF ONE RECEIVED, THEN
03AL PUT IN MESSAGE NUMBER
01AL PUT ADDRESS OF BUFFER ON THE ALARM QUEUE
01AL IF NO ROOM, THEN
07AL PANIC
      (ELSE)
      ENDIF
07AL ELSE
      PANIC
      ENDIF
DEQUEUE TERMINAL
ENDPRG
  
```

OPEN PSC COMMUNICATIONS COMMAND

- MODULE HISTORY -

PROJECT: DAS 73-01709
 SUB-SYSTEM: OPERATOR INTERFACE
 MODULE: DIPSC 27

- MODULE ABSTRACT -

00001080
 00001090
 00001100
 00001110
 00001120
 00001130
 00001140
 00000010
 00000020
 00000030
 00000040
 00000050
 00000060
 00000070
 00000080
 00000090
 00000100
 00000110
 00000120
 00000130
 00000140
 00000150
 00000160
 00000170
 00000180
 00000190
 00000200
 00000210
 00000220
 00000230
 00000240
 00000250
 00000260
 00000270
 00000280
 00000290
 00000300
 00000310
 00000320
 00000330
 00000340
 00000350
 00000360
 00000370
 00000380
 00000390
 00000400
 00000410
 00000420
 00000430
 00000440
 00000010
 00000020
 00000030
 00000040
 00000050
 00000060
 00000070
 00000080
 00000090
 00000100
 00000110
 00000120
 00000130
 00000140
 00000150
 00000160
 00000170
 00000180
 00000190
 00000200

THIS MODULE INITIATES PSC COMMUNICATIONS BY ENQUEUING A RESTART CODE ON THE PSC COMMUNICATION QUEUE. THE USER IS NOTIFIED OF INITIATION. IF THE PSC COMMUNICATIONS ARE ALREADY OPEN THE USER IS TOLD AND THE COMMAND IS ABORTED.

```

04JI ENQUEUE USER TERMINAL
07PS DISPLAY ON USER TERMINAL "OPEN PSC COMMUNICATIONS COMMAND"
      GET PSC COMMUNICATIONS FLAG FROM SYSCOM
      IF THE FLAG EQ COMMUNICATIONS CLOSED, THEN
01PS   DEQUEUE A PSC OUTPUT BUFFER
01PS   IF BUFFER FOUND, THEN
05PS     UPDATE PSC OUTPUT BUFFER
02PS     ENQUEJE PSC OUTPUT BUFFER ON PSC OUTPUT QUEUE
02PS     IF PSC OUTPUT QUEUE FULL
54JI       PANIC
           (ELSE)
           ENDIF
      ELSE
54JI       PANIC
           ENDIF
04JI       DISPLAY "PSC COMMUNICATION INITIATED"
      ELSE
04JI       DISPLAY "PSC COMMUNICATIONS ALREADY OPEN"
04JI       DISPLAY "COMMAND ABORTED"
           ENDIF
DEQUEUE TERMINAL
ENDPRDG

```

ADD DISK TRANSACTION (ORDER) FILE RECORD (ADD.TRANS)

- MODULE HISTORY -

```

PROJECT:      DAS                73-01709
SJB-SYSTEM:   OPERATOR INTERFACE
MODULE:       SICCAT             81

```

- MODULE ABSTRACT -

CALL THE FILE MANAGER TO ADD THE TRANSACTION FILE RECORD.
THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL TO THE FILE MANAGER WHETHER TRANSACTION FILE ACCESS WAS SUCCESSFUL OR NOT.

```

13FM UPDATE IPM FOR CALL
08FM CALL FILE MANAGER TO ADD TRANSACTION FILE RECORD
13FM CASENTRY (FOK, FTA, FTX, FOX, FCC, ? )
      CASE FOK ( SUCCESSFUL )
10JI   SET OK
      CASE FTA ( TRANSACTION ALREADY ON FILE )
02JI   PRINT/DISPLAY TRANS. ID, "TRANSACTION ALREADY ON FILE"
10JI   CLEAR OK
      CASE FTX ( TRANSACTION FILE SPACE EXHAUSTED )
02JI   PRINT/DISPLAY "TRANSACTION FILE SPACE EXHAUSTED"
10JI   CLEAR OK
      CASE FOX ( ORDER FILE SPACE EXHAUSTED )

```

00000210
00000220
00000230
00000240
00000250
00000260

00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270

00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390

```

02JI PRINT/DISPLAY "ORDER FILE SPACE EXHAUSTED"
10JI CLEAR OK
CASE FCC ( INVALID CALL CODE )
54JI PANIC
CASE ? ( INVALID COMPLETION CODE )
54JI PANIC
ENDCASE
RETURN

```

```

00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470

```

DELETE DISK BATCH FILE RECORD (DELETE BATCH)

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       JICCOB           85

```

- MODULE ABSTRACT -

```

CALL THE FILE MANAGER TO DELETE THE DISK BATCH
FILE RECORD(S).
IF AN ERROR OCCURRED, TELL THE USER, CLEAR OK AND
TAKE ANY FURTHER APPROPRIATE ACTION REQUIRED.
THIS SUBROUTINE WILL EXIT AFTER A SINGLE CALL TO THE
FILE MANAGER WHETHER BATCH FILE ACCESS WAS SUCCESSFUL
OR NOT.

```

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440

```

```

14FM UPDATE IPM FOR CALL
08FM CALL FILE MANAGER TO DELETE BATCH(ES)
14FM CASENTRY (FOK, FIR, FCC, ? )
CASE FOK ( SUCCESSFUL )
CASE FIR ( INVALID RELATIVE BATCH ID )
54JI PANIC
CASE FCC ( INVALID CALL CODE )
54JI PANIC
CASE ? ( INVALID COMPLETION CODE )
54JI PANIC
ENDCASE
RETURN

```

```

00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440

```

ORDER STATUS REPORT

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   OPERATOR INTERFACE
MODULE:       JIOSR           17

```

- MODULE ABSTRACT -

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200

```

THIS MODULE REQUESTS A BATCH ID. THE FILE MANAGER IS CALLED TO READ THE BATCH FILE INFORMATION.

IF THE BATCH INFORMATION IS FOUND, THE USER IS ASKED IF ALL ORDERS ARE TO BE LISTED. IF ALL ORDERS ARE NOT TO BE LISTED, A SPECIFIC ORDER ID IS REQUESTED, AND THE FILE MANAGER IS CALLED TO READ THE ORDER FILE INFORMATION. IF THE ORDER INFORMATION IS FOUND, OR ALL ORDERS TO BE LISTED, REQUEST IF EXCEPTIONS ARE TO BE LISTED AND WHETHER OR NOT HARDCOPY OUTPUT IS REQUIRED. IF HARDCOPY IS REQUIRED, THEN THE TERMINAL IS DEQUEUED AND THE SYSTEM PRINTER IS ENQUEUED.

THE BATCH ID, HEADER AND STATUS ARE THEN OUTPUT TO THE ENQUEUED DEVICE.

IF ALL ORDERS TO BE LISTED, THE FILE MANAGER IS CALLED TO SEQUENTIALLY READ EACH ORDER RECORD. FOR EACH ORDER RECORD, OR THE SPECIFICALLY REQUESTED ORDER RECORD, THE ORDER ID, ALL QUANTITIES AND ORDER PERCENT COMPLETE ARE OUTPUT.

IF AN EXCEPTION LIST WAS REQUIRED, THE FILE MANAGER IS CALLED TO SEQUENTIALLY READ EACH TRANSACTION BLOCK. EACH TRANSACTION RECORD ASSOCIATED WITH THE ORDER IS EXAMINED. IF THE TRANSACTION STATUS IS AN EXCEPTION, THE TRANSACTION INFORMATION IS OUTPUT TO THE ENQUEUED DEVICE.

THE TERMINAL/PRINTER IS DEQUEUED BEFORE THE PROGRAM ENDS.

		00000210
		00000220
		00000230
		00000240
		00000250
		00000260
		00000270
		00000280
		00000290
		00000300
		00000310
		00000320
		00000330
		00000340
		00000350
		00000360
		00000370
		00000380
		00000390
		00000400
		00000410
		00000420
		00000430
		00000440
		00000450
040I	ENQUEUE USER TERMINAL	00000460
030I	DISPLAY ON USER TERMINAL "ORDER STATUS REPORT"	00000470
120I	ASK USER FOR BATCH ID	00000480
480I	DEQUEUE AN IPM BUFFER (GET.IPM)	00000490
490I	DEQUEUE RECORD BUFFER (GET.SRB)	00000500
	INITIALIZE PANIC CODE	00000510
290I		00000520
160I	ACCESS BATCH FILE (GET.BATCH)	00000530
100I	IF BATCH RECORD FOUND (OK SET), THEN	00000540
030I	! ASK USER IF ALL ORDERS TO BE LISTED	00000550
290I	! IF NOT ALL ORDERS, THEN	00000560
030I	! ASK USER FOR ORDER ID	00000570
130I	!	00000580
	! INITIALIZE PANIC CODE	00000590
290I	!	00000600
180I	! ACCESS ORDER FILE (GET.ORDER)	00000610
100I	! IF ORDER RECORD FOUND (OK SET), THEN	00000620
100I	! SET SPECIFIC ORDER	00000630
	! (ELSE)	00000640
	! ENDIF	00000650
	! ELSE	00000660
100I	! SET JK	00000670
100I	! CLEAR SPECIFIC ORDER	00000680
	! ENDIF	00000690
100I	! IF ORDER FOUND (OK SET), THEN	00000700
100I	! ! SET ORDER SEQUENCE	00000710
030I	! ! ASK USER IF EXCEPTIONS TO BE LISTED	00000720
290I	! ! IF EXCEPTIONS TO BE LISTED, THEN	00000730
100I	! ! SET EXCEPTIONS	00000740
	! ! ELSE	00000750
100I	! ! CLEAR EXCEPTIONS	00000760
	! ! ENDIF	00000770
030I	! ! ASK USER IF HARDCOPY OUTPUT IS REQUIRED	00000780
	! ! IF HARDCOPY REQUIRED, THEN	00000790
100I	! ! SET HARDCOPY	00000800
	! ! DEQUEUE USER TERMINAL	00000810
	! ! ENQUEUE PRINTER	00000820
	! ! ELSE	00000830
100I	! ! CLEAR HARDCOPY	00000840
	! ! ENDIF	00000850
290I	! ! INITIALIZE DATE, TIME, PAGE COUNT AND LINE COUNT	00000860
070I	! ! PRINT/DISPLAY BATCH HEADINGS	00000870
070I	! ! PRINT/DISPLAY BATCH INFORMATION	00000880
070I	! ! PRINT/DISPLAY ORDER HEADINGS	00000890
	! ! INITIALIZE PANIC CODE	00000900
290I	! !	00000910
130I	! ! INITIALIZE ORDER COUNT	00000920

```

130I      ? ? ? ? ? DOUNTIL ORDER COUNT GREATER THAN MAXIMUM          00000930
100I      ? ? ? ? ? IF SPECIFIC ORDER SET, THEN                       00000940
130I      ? ? ? ? ?   UPDATE ORDER COUNT BEYOND MAXIMUM              00000950
120I      ? ? ? ? ?   ELSE                                           00000960
170I      ? ? ? ? ?     FIND THE FIRST/NEXT ORDER IN THIS BATCH (FIND.ORDER) 00000980
100I      ? ? ? ? ?     ENDIF                                         00000990
130I      ? ? ? ? ?     IF ORDER RECORD FOUND (OK SET), THEN          00001000
130I      ? ? ? ? ?     COMPUTE ORDER PERCENT COMPLETE                00001010
100I      ? ? ? ? ?     SET LIST ORDER INFORMATION                    00001020
070I      ? ? ? ? ?     PRINT/DISPLAY ORDER INFORMATION AND PERCENT COMPLETE 00001030
100I      ? ? ? ? ?     --- (FORMS.CTL) ---                             00001040
130I      ? ? ? ? ?     IF ABORT SET THEN                              00001050
100I      ? ? ? ? ?     SET ORDER COUNT GREATER THAN MAX.            00001060
100I      ? ? ? ? ?     ELSE                                           00001070
100I      ? ? ? ? ?     IF EXCEPTIONS SET, THEN                        00001080
100I      ? ? ? ? ?     IF EXCEPTIONS IN THIS ORDER THEN              00001090
100I      ? ? ? ? ?     CLEAR LIST ORDER INFORMATION                  00001100
070I      ? ? ? ? ?     PRINT/DISPLAY TRANSACTION HEADINGS           00001110
290I      ? ? ? ? ?     INITIALIZE PANIC CODE                          00001120
290I      ? ? ? ? ?     INITIALIZE REPORT TRANSACTION COUNT           00001130
290I      ? ? ? ? ?     INITIALIZE TRANS. CURRENT BLOCK NUMBER       00001140
140I      ? ? ? ? ?     INITIALIZE TRANSACTION COUNT                   00001150
140I      ? ? ? ? ?     INITIALIZE TRANSACTION COUNT                   00001160
140I      ? ? ? ? ?     DOUNTIL TRANSACTION COUNT > MAXIMUM --       00001180
140I      ? ? ? ? ?     -- OR BATCH NO LONGER ON FILE                 00001190
190I      ? ? ? ? ?     FIND THE FIRST/NEXT EXCEPTION RECORD         00001200
100I      ? ? ? ? ?     --- (FIND.TRANS) ---                           00001210
100I      ? ? ? ? ?     IF BATCH STILL ON FILE THEN                   00001220
100I      ? ? ? ? ?     IF EXCEPTION RECORD FOUND THEN                 00001230
100I      ? ? ? ? ?     IF TRANS. STATUS = MIS-SORT THEN              00001240
20FM      ? ? ? ? ?     SET UP TO READ THE TRANS. DIR.                00001250
08FM      ? ? ? ? ?     CALL FILE MANAGER FOR DIR. REC.                00001260
20FM      ? ? ? ? ?     IF RETURN CODE = ERROR THEN                    00001270
20FM      ? ? ? ? ?     IF RETURN CODE = NOT ON FILE                   00001280
100I      ? ? ? ? ?     PRINT ABORT MSG.                                00001290
540I      ? ? ? ? ?     ELSE                                           00001300
100I      ? ? ? ? ?     PANIC                                           00001310
100I      ? ? ? ? ?     ENDIF                                           00001320
03FM      ? ? ? ? ?     ELSE                                           00001330
100I      ? ? ? ? ?     SET MIS-SORT DEST. FROM REC.                  00001340
100I      ? ? ? ? ?     IF MIS-SORT DEST. = -1 THEN                    00001350
100I      ? ? ? ? ?     SET DJTPT = '?'                                00001360
100I      ? ? ? ? ?     ELSE                                           00001370
100I      ? ? ? ? ?     PLACE DEST. IN DJTPT                            00001380
100I      ? ? ? ? ?     ENDIF                                           00001390
100I      ? ? ? ? ?     ENDIF                                           00001400
100I      ? ? ? ? ?     ELSE                                           00001410
100I      ? ? ? ? ?     SET DJTPT TO BLANKS                             00001420
100I      ? ? ? ? ?     ENDIF                                           00001430
100I      ? ? ? ? ?     IF BATCH STILL ON FILE THEN                    00001440
070I      ? ? ? ? ?     SET LIST TRANSACTION INFO.                     00001450
100I      ? ? ? ? ?     PRINT/DISPLAY TRANS INFORMATION                00001460
100I      ? ? ? ? ?     --- (FORMS.CTL) ---                             00001470
100I      ? ? ? ? ?     (ELSE)                                          00001480
100I      ? ? ? ? ?     ENDIF                                           00001490
140I      ? ? ? ? ?     IF ABORT SET THEN                                00001500
130I      ? ? ? ? ?     SET TRANS. COUNT > MAX                          00001510
130I      ? ? ? ? ?     SET ORDER COUNT > MAX.                         00001520
100I      ? ? ? ? ?     (ELSE)                                          00001530
100I      ? ? ? ? ?     ENDIF                                           00001540
100I      ? ? ? ? ?     (ELSE)                                          00001550
100I      ? ? ? ? ?     ENDIF                                           00001560
040I      ? ? ? ? ?     ELSE                                           00001570
100I      ? ? ? ? ?     DJTPT "BATCH NO LONGER ON FILE"                00001580
100I      ? ? ? ? ?     --- REPORT TERMINATED" ---                     00001590
100I      ? ? ? ? ?     SET ABORT                                       00001600
100I      ? ? ? ? ?     ENDIF                                           00001610
100I      ? ? ? ? ?     ENDDO                                           00001620
100I      ? ? ? ? ?     CLEAR LIST TRANSACTION INFORMATION            00001630
040I      ? ? ? ? ?     ELSE                                           00001640
100I      ? ? ? ? ?     PRINT/DISPLAY "NO EXCEPTIONS IN THIS ORDER" 00001650

```



```

! ! ! ! ! ENDIF                                00001660
! ! ! ! ! (ELSE)                                00001670
! ! ! ! ! ENDIF                                00001680
! ! ! ! ! IF HARD COPY SET THEN                00001690
! ! ! ! ! UPDATE LINE COUNT > MAX.            00001700
! ! ! ! ! ELSE                                  00001710
! ! ! ! ! MOVE REPORT TO TOP OF SCREEN        00001720
! ! ! ! ! ENDIF                                00001730
! ! ! ! ! ENDIF                                00001740
! ! ! ! ! ELSE                                  00001750
0201 ! ! ! ! ! PRINT/DISPLAY "BATCH NO LONGER ON FILE - RPT TERM" 00001760
1001 ! ! ! ! ! SET ABORT                          00001770
! ! ! ! ! IF HARD COPY NOT SET THEN            00001780
! ! ! ! ! MOVE REPORT TO TOP OF SCREEN        00001790
! ! ! ! ! (ELSE)                                00001800
! ! ! ! ! ENDIF                                00001810
! ! ! ! ! ENDIF                                00001820
! ! ! ! ! ENDDO                                  00001830
! ! ! ! ! IF HARD COPY AND EXCEPTIONS NOT SET THEN 00001840
! ! ! ! ! MOVE REPORT TO TOP OF SCREEN        00001850
! ! ! ! ! ELSE                                  00001860
! ! ! ! ! ENDIF                                00001870
! ! ! ! ! (ELSE)                                00001880
! ! ! ! ! ENDIF                                00001890
! ! ! ! ! ELSE                                  00001900
0201 ! ! ! ! ! PRINT/DISPLAY BATCH ID, "BATCH NOT ON FILE" 00001910
! ! ! ! ! ENDIF                                00001920
5101 ! ! ! ! ! RELEASE RECORD BUFFER (FREE.SRB) 00001930
5001 ! ! ! ! ! RELEASE THE IPM BUFFER (FREE.IPM) 00001940
! ! ! ! ! DEQUEUE TERMINAL/PRINTER            00001950
! ! ! ! ! ENDPRDG                               00001960
! ! ! ! ! FORMS.CTL SUBROUTINE                 00001970
! ! ! ! !                                     00001980
1001 ! ! ! ! ! IF HARDCOPY SET, THEN            00001990
2901 ! ! ! ! ! IF LINE COUNT GREATER THAN PRINTER PAGE, THEN 00020000
1001 ! ! ! ! ! SET OK                          00020010
! ! ! ! ! ELSE                                  00020020
1001 ! ! ! ! ! CLEAR OK                        00020030
! ! ! ! ! ENDIF                                00020040
! ! ! ! ! ELSE                                  00020050
2901 ! ! ! ! ! IF LINE COUNT GREATER THAN TERMINAL PAGE, THEN 00020060
1001 ! ! ! ! ! SET OK                          00020070
! ! ! ! ! ELSE                                  00020080
1001 ! ! ! ! ! CLEAR OK                        00020090
! ! ! ! ! ENDIF                                0002100
! ! ! ! ! ENDIF                                0002110
1001 ! ! ! ! ! IF PAGE OVERFLOW (OK SET), THEN 0002120
2901 ! ! ! ! ! INCREMENT PAGE COUNT            0002130
0701 ! ! ! ! ! PRINT/DISPLAY BATCH HEADINGS    0002140
0701 ! ! ! ! ! PRINT/DISPLAY BATCH INFORMATION 0002150
0701 ! ! ! ! ! PRINT/DISPLAY ORDER HEADINGS    0002160
2901 ! ! ! ! ! INITIALIZE LINE COUNT           0002170
1001 ! ! ! ! ! IF LIST TRANSACTION INFORMATION SET, THEN 0002180
0701 ! ! ! ! ! PRINT/DISPLAY ORDER INFORMATION 0002190
0701 ! ! ! ! ! PRINT/DISPLAY TRANSACTION HEADINGS 0002200
2901 ! ! ! ! ! INCREMENT LINE COUNT           0002210
! ! ! ! ! (ELSE)                                0002220
! ! ! ! ! ENDIF                                0002230
! ! ! ! ! (ELSE)                                0002240
! ! ! ! ! ENDIF                                0002250
1001 ! ! ! ! ! IF LIST ORDER INFORMATION SET, THEN 0002260
0701 ! ! ! ! ! PRINT/DISPLAY ORDER INFORMATION 0002270
2901 ! ! ! ! ! INCREMENT LINE COUNT           0002280
! ! ! ! ! (ELSE)                                0002290
! ! ! ! ! ENDIF                                0002300
1001 ! ! ! ! ! IF LIST TRANSACTION INFORMATION SET, THEN 0002310
0701 ! ! ! ! ! PRINT/DISPLAY TRANSACTION INFORMATION 0002320
2901 ! ! ! ! ! INCREMENT REPORT TRANSACTION COUNT 0002330
2901 ! ! ! ! ! INCREMENT LINE COUNT           0002340
! ! ! ! ! (ELSE)                                0002350
! ! ! ! ! ENDIF                                0002360
! ! ! ! ! IF REPORT CANCELLED AND HARD COPY SET THEN 0002370
! ! ! ! ! PRINT REPORT CANCELLED MESSAGE      0002380

```

```
(ELSE)
ENDIF
RETJRN
```

ALARM MESSAGE PROGRAM

```
- MODULE HISTORY -
```

```
PROJECT:      DAS                75-01709
SJB-SYSTEM:   ALARM MESSAGES
MODULE:       ALMPG              1
```

```
- MODULE ABSTRACT -
```

```
THIS PROGRAM WILL DEQUEUE AN ALARM MESSAGE FROM THE ALARM
MESSAGE QUEUE. IF THE QUEUE IS EMPTY, IT WILL DELAY TWO
SECONDS AND START ALL OVER. IF THERE IS A MESSAGE, IT WILL GET
THE MESSAGE DATA AND RETURN THE DATA BUFFER TO THE ALARM MES-
SAGE BUFFER QUEUE. IF THE ADDRESS ON THE QUEUE WAS LESS THAN
THE ADDRESS OF $SYSCOM, THEREFORE NOT A VALID ADDRESS, THIS
PROGRAM WILL OUTPUT A FATAL ERROR MESSAGE WITH THE NUMBER AND
DEQUEUE ANOTHER MESSAGE. IF THE ADDRESS WAS VALID, THE MESSAGE
NUMBER IS USED TO FORMAT THE MESSAGE. THE MESSAGE IS THEN OUT-
PUT TO THE APPROPRIATE LOG DEVICE, AND THE SEQUENCE STARTS OVER.
```

```
INITIALIZATION
39FM  SET NUMBER OF LOG DEVICES FROM SYSCOM
      ATTACH THAT MANY OUTPUT TASKS
      ATTACH TIME OUT TASK
      DDUNTIL 1=0
01AL  GET AN ELEMENT FROM THE ALARM MESSAGE QUEUE (ALMQJE)
01AL  IF NO ELEMENT RECEIVED, THEN
      DELAY AWHILE
      ELSE
02AL  IF THE BUFFER ADDRESS > $SYSCOM, THEN
02AL  MOVE DATA FROM THE BUFFER TO AN INTERNAL BUFFER
02AL  RETURN BUFFER TO THE BUFFER POOL (BJFFQJE)
02AL  IF THE BUFFER POOL IS NOT FULL, THEN
11AL  SET COMPRESS FLAG
03AL  IF THE MESSAGE NUMBER IS WITHIN RANGE (1-N), THEN
      CASENTRY (MSG1,MSG2,.....MSGN)
03AL  CASE MSG1
      FORMAT MSG1
03AL  CASE MSG2
      FORMAT MSG2
      :
      :
03AL  CASE MSGN
      FORMAT MSGN
      ENDCASE
      PJT MESSAGE IN THE OUTPUT BUFFER
10AL  CALL THE DATE AND TIME SJBROUTINE
12AL  IF NOT A COMPRESSABLE MESSAGE, THEN
11AL  RESET COMPRESS FLAG
      (ELSE)
      ENDIF
11AL  IF COMPRESS FLAG SET, THEN
13AL  CALL COMPRESS SUBROUTINE
      (ELSE)
      ENDIF
04AL  GET ADDRESS OF TCW TABLE IN SYSCOM
04AL  ADD 2*(MSG#-1)
04AL  GET TCW AT THIS ADDRESS
```

```
00002390
00002400
00002410
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
0000100
0000110
0000120
0000130
0000140
0000150
0000160
0000170
0000180
0000190
0000200
0000210
0000220
0000230
0000240
0000250
0000260
0000270
0000280
0000290
0000300
0000310
0000320
0000330
0000340
0000350
0000360
0000370
0000380
0000390
0000400
0000410
0000420
0000430
0000440
0000450
0000460
0000470
0000480
0000490
0000500
0000510
0000520
0000530
0000540
0000550
0000560
0000570
0000580
0000590
0000600
0000610
0000620
0000630
0000640
0000650
0000660
0000670
0000680
0000690
0000700
```

```

05AL      SAVE TCW POINTER
04AL      GET POINTER TO CJTPUT TASK EVENT TABLE
04AL      DOUNTIL TCW=0
04AL      IF TCW<0, THEN
05AL      MOVE EVENT ADDR TO POST INSTRUCTION
05AL      POST THE TASK START EVENT
          (ELSE)
          ENDF
05AL      INCREMENT PTR TO NEXT TASK START EVENT
04AL      SHIFT TCW TO THE LEFT, 1 BIT
          ENDDO
09AL      POST TIMER START EVENT
          GET TCW
06AL      GET POINTER TO END EVENT TABLE
04AL      DOUNTIL TCW = 0
04AL      IF TCW<0, THEN
06AL      MOVE EVENT ADDR TO WAIT AND RESET INSTRUCTN
06AL      WAIT FOR EVENT TO OCCUR
06AL      RESET THE EVENT
          (ELSE)
          ENDF
05AL      INCREMENT POINTER TO NEXT EVENT
04AL      SHIFT TCW TO THE LEFT, 1 BIT
          ENDDO
09AL      SET TIMER COUNT TO ZERO
09AL      WAIT FOR TIMER END EVENT
09AL      RESET TIMER END EVENT
          ELSE
08AL      SET PANIC CODE = BAD MSG NUMBER
          SET FLAG = FATAL ERROR
          ENDF
          ELSE
08AL      SET PANIC CODE = ALARM BUFFER QUEUE FULL
          SET FLAG = FATAL ERROR
          ENDF
          ELSE
08AL      MOVE ADDRESS TO PANIC CODE
          SET FLAG = FATAL ERROR
          ENDF
          ENDF
08AL      IF A FATAL ERROR WAS ENCOJNTERED, THEN
08AL      FORMAT "FATAL ERROR" AND PANIC # TO CJTPUT BUFFER
10AL      CALL DATE AND TIME SUBROUTINE
13AL      CALL MESSAGE COMPRESSION SUBROUTINE
05AL      POST ALL OUTPUT TASKS
06AL      WAIT FOR ALL END EVENTS
          (ELSE)
          ENDF
          ENDDO
EXIT

TIME OUT TASK (TIMRTSK)

SET ADDRESSES OF ALL LOG DEVICE PRINT ECB'S
DOUNTIL 1=0
  WAIT FOR TIMER START EVENT TO BE POSTED
  RESET THE TIMER START EVENT
  SET UP TIME OUT COUNT
  DOWHILE TIME OUT COUNT > 0
    DELAY A LITTLE
    DECREMENT THE TIME OUT COUNT
  ENDDO
  POST ALL LOG DEVICE PRINT ECB'S
  POST TIMER END EVENT
ENDDO
EXIT

MESSAGE COMPRESSION SUBROUTINE (MDEVTEXT)

```

```

03000710
03000720
03000730
03000740
03000750
03000760
03000770
03000780
03000790
03000800
03000810
03000820
03000830
03000840
03000850
03000860
03000870
03000880
03000890
03000900
03000910
03000920
03000930
03000940
03000950
03000960
03000970
03000980
03000990
03001000
03001010
03001020
03001030
03001040
03001050
03001060
03001070
03001080
03001090
03001100
03001110
03001120
03001130
03001140
03001150
03001160
03001170
03001180
03001190
03001200
03001210
03001220
03001230
03001240
03001250
03001260
03001270
03001280
03001290
03001300
03001310
03001320
03001330
03001340
03001350
03001360
03001370
03001380
03001390
03001400
03001410
03001420
03001430

```

```

SAVE REGISTERS
GET TEXT ADDRESS
CALCULATE THE END ADDRESS FROM THE INDEX WORD
DO WHILE POINTER < END ADDRESS
  GET CURRENT BYTE
  IF CURRENT BYTE IS A BLANK AND NEXT BYTE IS A BLANK, THEN
    GET NUMBER OF BYTES TO END OF MESSAGE
    DO WHILE NOT AT END OF MESSAGE
      MOVE A BYTE TO THE LEFT
      INCREMENT POINTERS
    ENDDO
    ADJUST INDEX WORD
    ADJUST END ADDRESS
  ELSE
    INCREMENT POINTER
  ENDIF
ENDDO
RESTORE REGISTERS
RETURN

```

```

GET DATE AND TIME SUBROUTINE (DATES)

MOVE PUNCTUATION INTO THE OUTPUT BUFFER
ADJUST BUFFER INDEX WORD
GET TIME AND DATE
GET A POINTER TO THE TEXT BUFFER
GET A POINTER TO THE TIME AND DATE BUFFER
SET A COUNTER = 6
DO UNTIL COUNTER = 0
  CONVERT TIME AND DATE ENTRY TO EBCDIC
  PUT CONVERTED ENTRY IN THE BUFFER
  CHANGE LEADING SPACES TO LEADING ZEROS
  INCREMENT POINTERS
ENDDO
RETURN

```

```

ALARM MESSAGE TERMINAL 3 OUTPUT TASK

- MODULE HISTORY -

PROJECT:      DAS              73-01709
SJB-SYSTEM:   ALARM MESSAGES
MODULE:       ALTRM3          4

- MODULE ABSTRACT -

THIS TASK ENQUEUES LOG DEVICE 3. IT THEN WAITS FOR A
START EVENT TO BE POSTED. WHEN POSTED, IT WILL PRINT THE
CONTENTS OF THE OUTPUT BUFFER ON THE LOG DEVICE. THEN IT
RESETS ITS START EVENT AND POSTS ITS END EVENT. THEN RE-
TURNS TO WAIT FOR ITS START EVENT AGAIN.
4AL ENQUEUE LOG DEVICE 003
DO UNTIL 1=0
5AL WAIT FOR START EVENT TO BE POSTED
PRINT ALARM MESSAGE ON THE LOG DEVICE
5AL RESET THE START EVENT
6AL POST THE END EVENT
ENDDO
DEQUEUE THE LOG DEVICE
EXIT

```

```

00001440
00001450
00001450
00001470
00001480
00001490
00001500
00001510
00001520
00001530
00001540
00001550
00001560
00001570
00001580
00001590
00001600
00001610
00001620
00001630
00001640
00001650
00001660
00001670
00001680
00001690
00001700
00001710
00001720
00001730
00001740
00001750
00001760
00001770
00001780
00001790
00001800
00001810
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360

```

ALARM MESSAGE TERMINAL 2 OUTPUT TASK

- MODULE HISTORY -

PROJECT: DAS 73-01709
 SUB-SYSTEM: ALARM MESSAGES
 MODULE: ALTRM2 3

- MODULE ABSTRACT -

THIS TASK ENQUEUES LOG DEVICE 2. IT THEN WAITS FOR A START EVENT TO BE POSTED. WHEN POSTED, IT WILL PRINT THE CONTENTS OF THE OUTPUT BUFFER ON THE LOG DEVICE. THEN IT RESETS ITS START EVENT AND POSTS ITS END EVENT. THEN RETURNS TO WAIT FOR ITS START EVENT AGAIN.

4AL ENQUEUE LOG DEVICE 002
 00UNTIL I=0
 5AL WAIT FOR START EVENT TO BE POSTED
 PRINT ALARM MESSAGE ON THE LOG DEVICE
 5AL RESET THE START EVENT
 6AL POST THE END EVENT
 ENDDO
 DEQUEUE THE LOG DEVICE
 EXIT

ALARM MESSAGE TERMINAL 1 OUTPUT TASK

- MODULE HISTORY -

PROJECT: DAS 73-01709
 SUB-SYSTEM: ALARM MESSAGES
 MODULE: ALTRM1 2

- MODULE ABSTRACT -

THIS TASK ENQUEUES LOG DEVICE 1. IT THEN WAITS FOR A START EVENT TO BE POSTED. WHEN POSTED, IT WILL PRINT THE CONTENTS OF THE OUTPUT BUFFER ON THE LOG DEVICE. THEN IT RESETS ITS START EVENT AND POSTS ITS END EVENT. THEN RETURNS TO WAIT FOR ITS START EVENT AGAIN.

4AL ENQUEUE LOG DEVICE 001
 00UNTIL I=0
 5AL WAIT FOR START EVENT TO BE POSTED
 PRINT ALARM MESSAGE ON THE LOG DEVICE
 5AL RESET THE START EVENT
 6AL POST THE END EVENT
 ENDDO
 DEQUEUE THE LOG DEVICE
 EXIT

00000310
 00000020
 00000030
 00000040
 00000050
 00000060
 00000070
 00000080
 00000090
 00000100
 00000110
 00000120
 00000130
 00000140
 00000150
 00000160
 00000170
 00000180
 00000190
 00000200
 00000210
 00000220
 00000230
 00000240
 00000250
 00000260
 00000270
 00000280
 00000290
 00000300
 00000310
 00000320
 00000330
 00000340
 00000350
 00000360
 00000010
 00000020
 00000030
 00000040
 00000050
 00000060
 00000070
 00000080
 00000090
 00000100
 00000110
 00000120
 00000130
 00000140
 00000150
 00000160
 00000170
 00000180
 00000190
 00000200
 00000210
 00000220
 00000230
 00000240
 00000250
 00000260
 00000270
 00000280
 00000290
 00000300
 00000310
 00000320
 00000330
 00000340
 00000350
 00000360

CHANGE TRANSACTION STATUS

- MODULE HISTORY -

PROJECT: DAS 73-01709
 SUB-SYSTEM: FILE-MANAGER
 MODULE: FMCTS 35

- MODULE ABSTRACT -

THIS SUBROUTINE PROCESSES THE CHANGE TRANSACTION STATUS
 FILE MANAGER CALL CODES. IT CHECKS THE VALIDITY OF THE INPUT
 STATUS, SEARCHES THE TRANSACTION DIRECTORY FOR THE
 INPUT TRANSACTION NUMBER, THEN CHANGES THE TRANSACTION
 STATUS ON FILE TO THE INPUT STATUS. IT ALSO UPDATES THE
 ORDER STATUS COUNTS ASSOCIATED WITH THIS TRANSACTION.

```

16FM GET INPJTS
17FM
16FM GET INPUT TRANSACTION NUMBER
17FM
16FM IF CALL CODE = CHANGE STATUS FROM SCANNERS AND NEW STATUS --
-- EQUAL IN SORTATION THEN
16FM GET DIRECTORY RECORD FROM INPUT BUFFER
ELSE
03FM COMPUTE TRANSACTION DIRECTORY SECTOR FOR THIS TRANS. NUMBER
28FM READ THE TRANSACTION DIRECTORY SECTOR
03FM INDEX INTO SECTOR FOR TRANSACTION DIRECTORY RECORD
ENDIF
03FM GET RELATIVE BATCH NUMBER FROM DIRECTORY RECORD
02FM INDEX INTO THE RESIDENT BATCH FILE BY RELATIVE BATCH NUMBER
03FM IF THE TRANSACTION DIRECTORY RECORD UNUSED OR THE BATCH -
02FM -STATUS = ADDING, DELETING OR FREE THEN
16FM SET COMPLETION CODE = TRANSACTION NOT ON FILE
17FM
ELSE
16FM SAVE INPJT STATUS AS NEW STATUS
17FM
17FM IF CALL CODE = CHANGE STATUS FROM STATUS SCANNER THEN
17FM IF NEW STATUS NOT *EXCESS REPACK, STOCK-DJT, -
- STAGED FOR SHIPMENT, IN SORTATION, OR MIS-SORT THEN
17FM SET COMPLETION CODE = INVALID STATUS
(ELSE)
ENDIF
ELSE
16FM IF NEW STATUS NOT *EXCESS REPACK, STOCK-DJT, -
- STAGED FOR SHIPMENT, IN SORTATION, -
- OR NOT PICKED THEN
16FM SET COMPLETION CODE = INVALID STATUS
ELSE
16FM IF REC. BATCH ID NOT EQUAL TO INPUT BATCH ID THEN
02FM SET COMPLETION CODE = TRANSACTION NOT IN BATCH
16FM (ELSE)
ENDIF
ENDIF
ENDIF
16FM IF COMPLETION CODE = SUCCESSFUL THEN
17FM
03FM GET RELATIVE TRANS. NO. FROM TRANS. DIR. RECORD
04FM COMPUTE TRANSACTION FILE SECTOR NUMBER
28FM READ THE TRANSACTION FILE SECTOR

```

```

00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510
00000520
00000530
00000540
00000550
00000560
00000570
00000580
00000590
00000600
00000610
00000620
00000630
00000640
00000650
00000660
00000670
00000680
00000690
00000700
00000710
00000720
00000730
00000740

```

```

16FM      IF TRANS. NO. IN RECORD = INPUT TRANSACTION NUMBER THEN      00000750
04FM      INDEX TO TRANSACTION RECORD IN SECTOR                          00000760
04FM      SAVE CURRENT TRANS. STATUS AS OLD STATUS                       00000770
04FM      SET THE TRANS. STATUS = NEW STATUS                             00000780
04FM      GET THE RELATIVE ORDER RECORD NUMBER FROM TRANS. RECORD       00000790
28FM      WRITE THE TRANSACTION FILE SECTOR                               00000800
05FM      IF NEW STATUS = 'MIS-SORT' THEN                                 00000820
28FM      READ TRANS. DIR. SECTOR                                         00000830
03FM      INDEX INTO SECTOR FOR TRANS. DIRECTORY RECORD                 00000840
03FM      SET MIS-SORT DESTINATION IN TRANS. DIRECTORY RECORD           00000850
16FM      WRITE THE TRANSACTION DIRECTORY SECTOR                          00000870
28FM      (ELSE)                                                           00000890
      ENDIF                                                                00000890
05FM      COMPUTE THE ORDER FILE SECTOR NUMBER                           00000900
28FM      READ THE ORDER FILE SECTOR                                       00000910
05FM      INDEX TO ORDER RECORD IN SECTOR                                  00000920
05FM      INDEX TO OLD STATUS COUNT IN ORDER RECORD                      00000930
05FM      SUBTRACT 1 FROM THE OLD STATUS COUNT                            00000940
05FM      INDEX TO NEW STATUS COUNT IN ORDER RECORD                      00000950
05FM      ADD 1 TO THE NEW STATUS COUNT                                    00000960
28FM      WRITE THE ORDER FILE SECTOR                                       00000970
      ELSE                                                                    00000980
16FM      SET RETURN COMPLETION CODE = TRANSACTION NOT ON FILE           00000990
      ENDIF                                                                00010000
      (ELSE)                                                                  00010100
      ENDIF                                                                00010200
16FM      RETURN TO CALLER                                                 00010300
17FM

```

FILE MANAGER PANIC PROCESSOR

- MODULE HISTORY -

```

PROJECT:      DAS                      75-01709
SUB-SYSTEM:   FILE MANAGER
MODULE:       FMPAN                     43

```

- MODULE ABSTRACT -

THIS SUBROUTINE IS CALLED FROM SEVERAL FILE MANAGER
MODULES TO SEND A PANIC CODE TO THE ALARM MESSAGE SUB-SYSTEM
TO PRINT A FATAL ERROR MESSAGE AND TO SUSPEND EXECUTION OF
THE FILE MANAGER SUB-SYSTEM.

```

35FM      GET INPJTS
07AL      PUT PANIC CODE INPUT INTO ALARM MSG QUEUE
08AL      DDUNTIL NEW IPL
      WAIT FOR IMPOSSIBLE EVENT
      ENDJO
      RETURN (NOT REALLY)

```

USER INTERFACE TO FILE MANAGER

- MODULE HISTORY -

```

PROJECT:      DAS                      75-01709

```

00000750
00000760
00000770
00000780
00000790
00000800
00000810
00000820
00000830
00000840
00000850
00000860
00000870
00000890
00000890
00000900
00000910
00000920
00000930
00000940
00000950
00000960
00000970
00000980
00000990
00010000
00010100
00010200
00010300
00010400
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000250
00000270
00000280
00000290
00000300
00000310
00000310
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130

SUB-SYSTEM: FILE MANAGER

MODULE: FM 47

- MODULE ABSTRACT -

THIS SUBROUTINE PERFORMS ALL THE FUNCTIONS NEEDED TO CALL THE FILE MANAGER. A COPY OF THIS SUBROUTINE IS TO BE INCLUDED IN EACH DAS SUBSYSTEM THAT CALLS THE FILE MANAGER.

08FM GET INPJTS
34FM ENQUEUE THE FILE MANAGER RESOURCE
27FM RESET FILE MANAGER COMPLETION EVENT
25FM POST FILE MANAGER ENTRY EVENT
27FM WAIT ON FILE MANAGER COMPLETION EVENT
34FM DEQUEUE THE FILE MANAGER RESOURCE
08FM RETURN TO CALLER

ADD BATCH COMPLETE

- MODULE HISTORY -

PROJECT: DAS 75-01709

SUB-SYSTEM: FILE MANAGER

MODULE: FMADC 90

- MODULE ABSTRACT -

THIS SUBROUTINE PROCESSES THE ADD BATCH COMPLETE FILE MANAGER CALL CODE. IT FINDS THE BATCH RECORD WITH THE STATUS OF ADDING AND CHANGES IT TO PENDING IN BOTH THE RESIDENT BATCH FILE AND THE DISK BATCH FILE.

38FM GET INPJTS
02FM GET ADDRESS OF RESIDENT BATCH FILE
SET RELATIVE BATCH NUMBER = ZERO
02FM DOUNTIL END OF RESIDENT BATCH FILE OR BATCH WITH STATUS --
-- ADDING FOUND
GET FIRST/NEXT RECORD
02FM IF RECORD STATUS NOT ADDING THEN
ADD 1 TO RELATIVE BATCH NUMBER
(ELSE)
ENDIF
ENDDD
IF BATCH FOUND THEN
05FM COMPUTE FIRST ORDER FILE SECTOR ADDRESS FROM --
-- RELATIVE BATCH NUMBER
28FM UPDATE RESIDENT ORDER FILE
01FM COMPUTE BATCH FILE SECTOR ADDRESS FROM RELATIVE BATCH NO.
28FM READ THE BATCH FILE SECTOR
01FM INDEX TO THE BATCH RECORD IN THE SECTOR
01FM SET BATCH STATUS TO PENDING
28FM WRITE THE BATCH FILE SECTOR
02FM SET RESIDENT BATCH RECORD STATUS TO PENDING
ELSE
38FM SET RETURN COMPLETION CODE = BATCH NOT ON FILE
ENDIF
38FM RETURN TO CALLER

00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000500
00000510

READ TRANSACTION DIRECTORY ENTRY VIA TRANSACTION NUMBER

- MODULE HISTORY -

PROJECT: DAS 75-01709
SUB-SYSTEM: FILE MANAGER
MODULE: FMRTD 41

- MODULE ABSTRACT -

THIS SUBROUTINE READS A TRANSACTION DIRECTORY ENTRY INTO THE CALLER'S BUFFER WHEN GIVEN A TRANSACTION NUMBER AS INPUT.
20FM SET INPJTS
03FM USING TRANSACTION NO. INPUT, COMPUTE TRANS. DIR. SECTOR NO.
28FM READ TRANSACTION DIRECTORY SECTOR
03FM INDEX INTO SECTOR TO TRANSACTION DIRECTORY ENTRY
03FM IF ENTRY IS UNUSED OR BATCH STATUS = ADDING/DELETING OR FREE THE
20FM SET RETURN COMPLETION CODE = TRANSACTION NOT ON FILE
ELSE
03FM COPY TRANS. DIR. ENTRY INTO USER BUFFER
20FM
08FM GET RELATIVE TRANSACTION NO. FROM DIRECTORY RECORD
04FM COMPUTE TRANSACTION FILE SECTOR FROM REL. TRANS. NO.
28FM READ THE TRANSACTION FILE SECTOR
04FM INDEX TO TRANSACTION RECORD
20FM IF INPUT TRANSACTION NO. NOT EQUAL TO TRANSACTION NO. --
04FM -- IN RECORD THEN
20FM SET RETURN COMPLETION CODE = TRANSACTION NOT ON FILE
(ELSE)
ENDIF
ENDIF
20FM RETURN TO CALLER

0000010
0000020
0000030
0000040
0000050
0000060
0000070
0000080
0000090
0000100
0000110
0000120
0000130
0000140
0000150
0000160
0000170
0000180
0000190
0000200
0000210
0000220
0000230
0000240
0000250
0000260
0000270
0000280
0000290
0000300
0000310
0000320
0000330
0000340
0000350
0000360
0000370
0000380
0000390
0000400
0000410
0000420
0000430
0000440
0000010
0000020
0000030
0000040
0000050
0000060
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0000080
0000090
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0000110
0000120
0000130
0000140
0000150
0000160
0000170
0000180
0000190
0000200
0000210
0000220
0000230
0000240
0000250
0000260
0000270

READ ORDER RECORD VIA ORDER I.D.

- MODULE HISTORY -

PROJECT: DAS 75-01709
SUB-SYSTEM: FILE MANAGER
MODULE: FM01D 33

- MODULE ABSTRACT -

THIS SUBROUTINE PROCESSES THE READ ORDER RECORD VIA ORDER I.D. FILE MANAGER CALL CODE. IT CHECKS FOR A VALID BATCH I.D. THEN SEARCHES THE ORDER FILE FOR THE INPUT ORDER I.D.. WHEN FOUND THIS ROUTINE PLACES THE REQUESTED ORDER RECORD INTO THE CALLERS BUFFER.

```

22FM      SET INPUTS                                00000280
33FM      CALL ROUTINE TO SEARCH RESIDENT BATCH FILE FOR INPUT BATCH I.O. 00000290
33FM      IF INPUT BATCH ID FOUND THEN              00000300
01FM      COMPUTE THE BATCH FILE SECTOR FOR THE INPUT BATCH 00000310
28FM      READ THE BATCH FILE SECTOR                00000320
01FM      INDEX TO BATCH RECORD IN BATCH FILE SECTOR 00000330
01FM      GET THE NUMBER OF ORDERS FROM THE BATCH RECORD 00000340
05FM      COMPUTE STARTING SECTOR IN ORDER FILE FOR BATCH 00000350
          DOUNTIL ALL ORDER SECTORS SEARCHED OR RECORD FOUND 00000360
28FM      READ FIRST/NEXT ORDER FILE SECTOR        00000370
05FM      DOUNTIL ALL VALID ORDER RECORDS IN SECTOR SEARCHED OR - 00000380
          - RECORD FOUND                                00000390
05FM      GET FIRST/NEXT ORDER RECORD IN SECTOR    00000400
05FM      IF RECORD ORDER I.O. = INPUT ORDER I.O. THEN 00000410
          00000420
22FM      COPY ORDER RECORD INTO USER BUFFER      00000430
05FM      00000440
22FM      (ELSE)                                    00000450
          ENDIF                                        00000460
          ENDDO                                       00000470
          ENDDO                                       00000480
          IF ORDER RECORD NOT FOUND THEN              00000490
22FM      SET COMPLETION CODE = ORDER NOT FOUND 00000500
          (ELSE)                                       00000510
          ENDIF                                        00000520
ELSE                                                00000530
22FM      SET COMPLETION CODE = BATCH NOT FOUND 00000540
          ENDIF                                        00000550
22FM      RETURN TO CALLER                          00000560
          00000010
          00000020
          00000030
          00000040
          00000050
          00000060
          00000070
          00000080
          00000090
          00000100
          00000110
          00000120
          00000130
          00000140
          00000150
          00000160
          00000170
          00000180
          00000190
          00000200
          00000210
          00000220
          00000230
          00000240
          00000250
          00000260
          00000270
          00000280
          00000290
          00000300
          00000310
          00000320
          00000330
          00000340
          00000350
          00000360
          00000370
          00000380
          00000390
          00000400
          00000410
          00000420
          00000430
          00000440
          00000450

READ TRANSACTION RECORD VIA TRANSACTION NUMBER

          - MODULE HISTORY -
          00000100
          00000110
          00000120
          00000130
          00000140
          00000150
          00000160
          00000170
          00000180
          00000190
          00000200
          00000210
          00000220
          00000230
          00000240
          00000250
          00000260
          00000270
          00000280
          00000290
          00000300
          00000310
          00000320
          00000330
          00000340
          00000350
          00000360
          00000370
          00000380
          00000390
          00000400
          00000410
          00000420
          00000430
          00000440
          00000450

          PROJECT:          DAS          73-01799
          SUB-SYSTEM:      FILE MANAGER
          MODULE:          FMTID          39

          - MODULE ABSTRACT -
          00000200
          00000210
          00000220
          00000230
          00000240
          00000250
          00000260
          00000270
          00000280
          00000290
          00000300
          00000310
          00000320
          00000330
          00000340
          00000350
          00000360
          00000370
          00000380
          00000390
          00000400
          00000410
          00000420
          00000430
          00000440
          00000450

THIS SUBROUTINE COMPUTES THE TRANSACTION DIRECTORY ENTRY
FROM THE INPUT TRANSACTION NUMBER. THEN, UPON READING THE
ASSOCIATED DIRECTORY ENTRY, THE TRANSACTION FILE RELATIVE
RECORD ADDRESS CAN BE OBTAINED AND USED TO ACCESS THE DISK-
RESIDENT TRANSACTION FILE TO OBTAIN THE RELEVANT TRANSACTION
RECORD WHICH IS THEN DUTIFULLY COPIED INTO THE CALLER'S BUFFER.
19FM      GET INPUTS                                00000280
03FM      FROM TRANSACTION NUMBER, COMPUTE TRANSACTION DIRECTORY SECTOR NO 00000300
28FM      READ TRANSACTION DIRECTORY SECTOR        00000310
03FM      INDEX INTO SECTOR FOR TRANSACTION ENTRY  00000320
03FM      IF ENTRY IS UNUSED OR BATCH STATUS = ADDING/DELETING OR FREE THE 00000330
19FM      SET RETURN COMPLETION CODE = TRANSACTION NOT ON FILE 00000340
          ELSE                                        00000350
03FM      GET RELATIVE TRANSACTION NO. FROM DIRECTORY 00000360
04FM      COMPUTE TRANSACTION FILE SECTOR NUMBER  00000370
28FM      READ TRANSACTION FILE SECTOR            00000380
04FM      INDEX INTO SECTOR FOR TRANSACTION RECORD 00000390
19FM      IF INPUT TRANSACTION NO. IS EQUAL TO TRANSACTION NO. -- 00000400
04FM      -- IN RECORD THEN                        00000410
04FM      COPY RECORD INTO CALLER'S BUFFER        00000420
          ELSE                                        00000430
19FM      SET RETURN COMPLETION CODE = TRANSACTION NOT ON FILE 00000440
          ENDIF                                        00000450

```

19FM
19FM

ENDIF
RETURN TO CALLER

READ ORDER RECORD VIA RELATIVE ORDER NUMBER

- MODJLE HISTORY -

PROJECT: DAS 73-01709
SUB-SYSTEM: FILE MANAGER
MODULE: FMROL 40

- MODJLE ABSTRACT -

THIS SUBROUTINE FIRST CHECKS THE BATCH ID INPJT FOR VALIDITY AND THEN DETERMINES THE NO. OF ORDERS IN THE SPECIFIED BATCH. WITH THE NO. OF ORDERS IN THE BATCH KNOWN, THE RELATIVE ORDER NO. CAN BE VERIFIED. HAVING A GOOD RELATIVE ORDER NO., IT'S A TRIVIAL EXERCISE TO ACCESS THE ORDER FILE ON DISK (IN THE PROPER BATCH SEGMENT, OF COURSE) TO GET THE ORDER RECORD AND COPY IT INTO THE CALLER'S BUFFER. NOW LET'S GET TO THE MEAT...

23FM
33FM
33FM
01FM
01FM
23FM
23FM
05FM
05FM
28FM
05FM
05FM
23FM

GET INPUTS
CALL ROUTINE TO SEARCH RES. BATCH FILE FOR BATCH ID INPJT
IF THE BATCH WAS FOUND THEN
READ THE BATCH FILE SECTOR WHICH CONTAINS THE INDICATED BATCH
GET THE NO. OF ORDERS IN THE BATCH FROM THE RECORD
IF THE RELATIVE ORDER NO. INPUT > NO. OF ORDERS OR < 1 THEN
SET COMPLETION CODE = INVALID REL. RECORD NO.
ELSE
COMPUTE STARTING SECTOR NO. IN ORDER FILE FOR BATCH
COMPUTE SECTOR NO. FOR ORDER RECORD
READ SECTOR FROM ORDER FILE
INDEX INTO SECTOR FOR ORDER RECORD
COPY ORDER RECORD INTO USER BUFFER

ENDIF
ELSE
23FM SET COMPLETION CODE = BATCH NOT ON FILE
ENDIF
23FM RETURN TO CALLER

READ TRANSACTION FILE BLOCK

- MODJLE HISTORY -

PROJECT: DAS 73-01709
SUB-SYSTEM: FILE MANAGER
MODULE: FMRTB 37

- MODJLE ABSTRACT -

00000460
00000470
00000480
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000480
00000490
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220

```

THIS SUBROUTINE PROCESSES THE READ TRANSACTION FILE BLOCK
FILE MANAGER CALL CODE. IT SEARCHES THE RESIDENT BATCH FILE
FOR THE INPUT BATCH I.D.. THEN INPUTS THE REQUESTED TRANS.
FILE SECTOR AND PASSES THE SECTOR TO THE CALLER.
21FM SET INPUTS
33FM CALL ROUTINE TO SEARCH RESIDENT BATCH FILE FOR INPUT BATCH I.D.
33FM IF THE BATCH RECORD FOUND THEN
01FM COMPUTE THE BATCH FILE SECTOR FOR THIS BATCH
28FM READ THE BATCH FILE SECTOR
01FM INDEX TO THE BATCH RECORD IN THE SECTOR
01FM GET THE NUMBER OF TRANSACTIONS IN THIS BATCH
04FM COMPUTE THE MAX. NO. OF TRANS. FILE SECTORS IN THIS BATCH
21FM IF THE INPUT RELATIVE BLOCK NO. > MAX. TRANS. SECTORS THEN
21FM SET COMPLETION CODE = INVALID RELATIVE REC./BLOCK NO.
ELSE
04FM COMPUTE THE TRANSACTION FILE SECTOR NUMBER
28FM READ THE TRANS. FILE SECTOR
21FM COPY DATA INTO USER'S BUFFER
ENDIF
ELSE
21FM SET COMPLETION CODE = BATCH NOT ON FILE
ENDIF
21FM RETURN TO CALLER

READ/WRITE BLOCK (SECTOR) FROM/TO DISK FILE

- MODULE HISTORY -

PROJECT: DAS 73-01709
SUB-SYSTEM: FILE MANAGER
MODULE: FMRW3 42

- MODULE ABSTRACT -

THIS MODULE READS/Writes SECTORS FROM/TO DISK FILES
INDICATED BY RELATIVE FILE NUMBER (FILE CODE).
28FM SET INPUTS
CHECK VALIDITY OF BLOCK NUMBER/FILE CODE
IF INPUTS ARE VALID THEN
IF READ SECTOR WAS REQUESTED THEN
IF BLOCK ALREADY IN MEMORY THEN
COPY BLOCK INTO CALLERS BUFFER
ELSE
READ BLOCK(BLK INPUT) FROM FILE(FILE CODE INPUT)
ENDIF
ELSE
COPY CALLERS BUFFER TO MEMORY BLOCK
WRITE BLOCK(BLK INPUT) TO FILE(FILE CODE INPUT)
ENDIF
IF DISK ERROR ON READ/WRITE THEN
35FM PANIC! (DO IT WITH CALL CODE FROM USER)
09FM
WAIT FOR IPL RELIEF...
(ELSE)
ENDIF
ELSE
35FM PANIC!
ENDIF
28FM RETURN TO CALLER

```

```

00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000010
00000020
00000030
00000040

```


	DOWHILE NUMBER OF ORDERS NOT ZERO	03000780
29FM	SET FIRST/NEXT ORDER ID FROM TABLE	03000790
13FM	IF TABLE ORDER ID = INPUT ORDER ID THEN	03000800
13FM	INDEX TO STATUS COUNTER FOR THIS ORDER	03000810
29FM	ADD 1 TO NUMBER OF TRANSACTIONS OF THIS STATUS	03000820
	(ELSE)	03000830
	ENDIF	03000840
	ENDDO	03000850
29FM	IF ORDER ID NOT FOUND IN TABLE THEN	03000860
	IF NO. OF ORDERS IN TABLE G.T. OR E.Q. MAX. THEN	03000870
13FM	SET COMPLETION CODE = ORDER FILE SPACE EXHAUSTED	03000880
	ELSE	03000890
29FM	ADD INPUT ORDER ID TO TABLE	03000900
	CLEAR ALL STATUS COUNTS	03000910
13FM	INDEX TO STATUS COUNTER FOR THIS ORDER	03000920
29FM	SET NUMBER OF TRANSACTIONS OF THIS STATUS = 1	03000930
	ADD 1 TO NUMBER OF ORDERS FOUND IN BATCH THIS FAR	03000940
	ENDIF	03000950
	(ELSE)	03000960
	ENDIF	03000970
13FM	IF COMPLETION CODE = SUCCESSFUL THEN	03000980
	SAVE RELATIVE ORDER NUMBER	03000990
04FM	PLACE NEW TRANS. RECORD INTO CURRENT SECTOR BUFFER	03001000
13FM		03001010
13FM	IF THIS IS THE LAST TRANSACTION IN THE BATCH THEN	03001020
04FM	COMPUTE THE TRANSACTION FILE SECTOR	03001030
	WRITE THE CURRENT SECTOR BUFFER TO THE TRANS. FILE	03001040
29FM	SET NUMBER OF ORDERS IN ORDER TABLE	03001050
	DOUNTIL NUMBERS OF ORDERS = ZERO	03001060
	DOUNTIL ORDER SECTOR BUILT OR NUMBER OF -	03001070
	- ORDERS = ZERO	03001080
29FM	SET NEXT DATA FOR NEXT ORDER FROM TABLE	03001090
05FM	BUILD ORDER RECORD IN CURRENT SECTOR BUFFER	03001100
29FM		03001110
31FM	SUBTRACT 1 FROM NUMBER OF ORDERS	03001120
	ENDDO	03001130
05FM	COMPUTE ORDER FILE SECTOR FOR CURRENT -	03001140
	- SECTOR BUFFER	03001150
	WRITE THE CURRENT SECTOR BUFF. TO THE ORDER FILE	03001160
	ENDDO	03001170
03FM	WRITE THE BUFFERED TRANS. DIR. SECTOR TO DISK	03001180
01FM	COMPUTE THE BATCH FILE SECTOR	03001190
	READ THE BATCH FILE SECTOR	03001200
01FM	INDEX INTO SECTOR FOR BATCH RECORD	03001210
01FM	PLACE THE NUMBER OF TRANS. IN THE BATCH RECORD	03001220
01FM	PLACE THE NUMBER OF ORDERS IN THE BATCH RECORD	03001230
	WRITE THE BATCH SECTOR	03001240
09FM	GET A F.M. CALL BUFFER	03001250
	IF ONE RECEIVED THEN	03001260
38FM	SET UP BUFFER FOR ADD COMPLETE CALL	03001270
08FM	CALL THE FILE MANAGER	03001280
08FM	IF F.M. RETURN CODE = SUCCESSFUL THEN	03001290
09FM	RETURN CALL BUFFER TO FREE POOL	03001300
	IF FREE POOL FULL THEN	03001310
54JI	PANIC	03001320
	(ELSE)	03001330
	ENDIF	03001340
	ELSE	03001350
54JI	PANIC	03001360
	ENDIF	03001370
	ELSE	03001380
54JI	PANIC	03001390
	ENDIF	03001400
	ELSE	03001410
04FM	IF THE CURRENT TRANS. FILE SECTOR BUFFER FULL THEN	03001420
04FM	COMPUTE THE TRANSACTION FILE SECTOR	03001430
	WRITE THE CURRENT SECTOR BUFFER TO THE -	03001440
	- TRANSACTION FILE	03001450
	(ELSE)	03001460
	ENDIF	03001470
	ENDIF	03001480
	(ELSE)	03001490
	ENDIF	03001500

```

      (ELSE)
      ENDIF
    ENDIF
13FM  IF COMPLETION CODE = TRANSACTION ALREADY ON FILE, -
      - TRANSACTION FILE SPACE EXHAUSTED, OR ORDER FILE -
      - SPACE EXHAUSTED THEN
02FM  INDEX INTO RESIDENT BATCH FILE BY RELATIVE BATCH NO.
02FM  SET RESIDENT BATCH RECORD STATUS = DELETING
09FM  GET A.F.M. CALL BUFFER
      IF ONE RECEIVED THEN
14FM  SET RELATIVE BATCH NO. IN BATCH DELETION CALL BUFFER
08FM  CALL FILE MANAGER TO DELETE THE BATCH
08FM  IF ERROR RETURNED FROM CALL THEN
54JI  PANIC
      ELSE
09FM  RETURN CALL BUFFER TO FREE POOL
      IF FREE POOL FULL THEN
54JI  PANIC
      (ELSE)
      ENDIF
    ENDIF
      ELSE
54JI  PANIC
    ENDIF
  (ELSE)
  ENDIF
13FM  RETURN TO CALLER

READ BATCH RECORD VIA BATCH I.D.

      - MODULE HISTORY -

PROJECT:      DAS              75-01709
SUB-SYSTEM:   FILE MANAGER
MODULE:       FMBID            34

      - MODULE ABSTRACT -

THIS SUBROUTINE PROCESSES THE READ BATCH FILE RECORD
VIA BATCH ID FILE MANAGER CALL CODE. IT SEARCHES THE
RESIDENT BATCH FILE TO CHECK THE VALIDITY AND EXISTANCE
OF THE BATCH ID INPUT, RESOLVES THE RELATIVE BATCH
NUMBER, AND ACCESSES THE BATCH FILE ON DISK TO READ
THE INDICATED BATCH RECORD AND COPY THE DATA INTO THE
CALLER'S BUFFER.
18FM  SET INPUTS
33FM  CALL ROUTINE TO SEARCH RESIDENT BATCH FILE FOR INPUT BATCH I.D.
33FM  IF BATCH FOUND IN RESIDENT BATCH FILE THEN
33FM  GET RELATIVE BATCH NUMBER
28FM  READ BATCH FILE SECTOR CONTAINING INDICATED BATCH RECORD
01FM  COPY BATCH RECORD INTO CALLER'S BUFFER
18FM
      ELSE
18FM  SET COMPLETION CODE = BATCH NOT ON FILE
      ENDIF
18FM  RETURN TO CALLER

SET UP BATCH DELETION QUEUE FOR FMDLJ

```

```

00001510
00001520
00001530
00001540
00001550
00001560
00001570
00001580
00001590
00001600
00001610
00001620
00001630
00001640
00001650
00001660
00001670
00001680
00001690
00001700
00001710
00001720
00001730
00001740
00001750
00001760
00001770
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080

```

```

- MODULE HISTORY -
PROJECT:      DAS              75-01709
SUB-SYSTEM:   FILE MANAGER
MODULE:       FMDL            32

- MODULE ABSTRACT -

THIS MODULE RECEIVES ONE OR MORE RELATIVE BATCH NUMBERS
FOR BATCHES TO BE DELETED FROM THE SYSTEM. IT THEN PERFORMS
VALIDITY CHECKS ON THE INPUTS, MAKES THE NECESSARY BATCH
STATJS ADJUSTMENTS, AND PUTS THE BATCH NUMBERS INTO A
WORK QUEUE FOR THE BATCH DELETION TASK (FMDL).
14FM GET RELATIVE BATCH NUMBER(S) INPUT
14FM DDUNTIL ALL INPUTS VERIFIED OR INVALID INPUT DETECTED
14FM GET FIRST(NEXT) REL. BATCH NO. IN INPJT LIST
IF RELATIVE BATCH NUMBER IN RANGE THEN
02FM GET BATCH RECORD IN RESIDENT BATCH FILE
02FM IF RECORD STATJS = UNUSED, ACTIVE, OR ADDING THEN
14FM SET RETURN COMPLETION CODE TO "INVALID REC. NO."
(ELSE)
ENDIF
ELSE
14FM SET RETURN COMPLETION CODE TO "INVALID REC. NO."
ENDIF
ENDDD
14FM IF ALL INPUTS WERE FOUND TO BE VALID THEN
14FM DDUNTIL ALL BATCHES IN INPJT LIST ARE PROCESSED
14FM GET FIRST(NEXT) REL. BATCH NO. FROM LIST
02FM GET BATCH RECORD IN RESIDENT BATCH FILE
02FM SET STATUS IN RECORD TO "DELETING"
32FM PUT RELATIVE BATCH NO. INTO QUEUE FOR DELETION TASK
IF QUEUE IS FULL THEN
35FM PANIC!
(ELSE)
ENDIF
28FM READ SECTOR FROM BATCH FILE FOR INDICATED BATCH
01FM INDEX TO INDICATED BATCH RECORD IN SECTOR
01FM SET STATUS IN RECORD TO "DELETING"
28FM WRITE MODIFIED SECTOR BACK TO BATCH FILE ON DISK
ENDDD
(ELSE)
ENDIF
14FM RETURN

CHANGE BATCH STATUS

- MODULE HISTORY -
PROJECT:      DAS              75-01709
SUB-SYSTEM:   FILE MANAGER
MODULE:       FMCBS           35

- MODULE ABSTRACT -

```

```

0000090
0000100
0000110
0000120
0000130
0000140
0000150
0000160
0000170
0000180
0000190
0000200
0000210
0000220
0000230
0000240
0000250
0000260
0000270
0000280
0000290
0000300
0000310
0000320
0000330
0000340
0000350
0000360
0000370
0000380
0000390
0000400
0000410
0000420
0000430
0000440
0000450
0000460
0000470
0000480
0000490
0000500
0000510
0000520
0000530
0000540
0000550
0000560
0000570
0000580
0000590
0000010
0000020
0000030
0000040
0000050
0000060
0000070
0000080
0000090
0000100
0000110
0000120
0000130
0000140
0000150
0000160
0000170
0000180
0000190
0000200
0000210
0000220
0000230

```



```

THIS SUBROUTINE PROCESSES THE CHANGE BATCH STATUS FILE
MANAGER CALL CODE. IT SEARCHES THE BATCH FILE FOR THE
INPJT BATCH I.D. AND IF FOUND CHANGES THE BATCH STATUS
TO THE INPUT STATUS.
00000240
00000250
00000260
00000270
15FM   GET INPJTS                                00000280
33FM   CALL ROUTINE TO SEARCH RESIDENT BATCH FILE FOR INPUT BATCH I.D. 00000290
33FM   IF BATCH NOT FOUND THEN                   00000300
15FM       SET COMPLETION CODE = BATCH NOT ON FILE 00000310
ELSE                                         00000320
15FM     IF INPJT BATCH STATUS INVALID THEN      00000330
15FM       SET COMPLETION CODE = INVALID STATUS  00000340
ELSE                                         00000350
01FM     COMPUTE BATCH FILE SECTOR FROM RELATIVE BATCH NUMBER 00000360
28FM     READ THE BATCH FILE SECTOR            00000370
01FM     INDEX INTO SECTOR FOR BATCH RECORD    00000380
01FM     SET RECORD STATUS = INPUT STATUS      00000390
28FM     WRITE THE BATCH FILE SECTOR          00000400
02FM     INDEX INTO RESIDENT BATCH FILE BY RELATIVE BATCH NUMBER 00000410
02FM     SET RESIDENT BATCH RECORD STATUS = INPUT STATUS 00000420
ENDIF                                       00000430
ENDIF                                       00000440
15FM   RETURN TO CALLER                       00000450
00000010
00000020
00000030
00000040
00000050
00000060
00000070
00000080
00000090
00000100
00000110
00000120
00000130
00000140
00000150
00000160
00000170
00000180
00000190
00000200
00000210
00000220
00000230
00000240
00000250
00000260
00000270
00000280
00000290
00000300
00000310
00000320
00000330
00000340
00000350
00000360
00000370
00000380
00000390
00000400
00000410
00000420
00000430
00000440
00000450
00000460
00000470
00000010
00000020
00000030
00000040
00000050

```

TASK TO PROCESS BATCH DELETIONS

- MODULE HISTORY -

```

PROJECT:      DAS                73-01709
SUB-SYSTEM:   FILE MANAGER
MODULE:       FMDLD              44

```

- MODULE ABSTRACT -

THIS MODULE IS A TASK WHICH OBTAINS RELATIVE BATCH NUMBERS FROM A QUEUE WHICH INDICATE THE BATCHES TO BE DELETED. IT THEN WRITES AN INITIALIZED SECTOR TO ALL SECTORS IN THE TRANS. FILE THAT ARE ASSOCIATED WITH THE BATCH BEING DELETED. FINALLY THIS TASK MAKES THE APPROPRIATE STATUS ADJUSTMENT ENTRIES IN THE DISK AND RESIDENT BATCH FILES.

```

DO UNTIL CPJ STOPPED
  32FM   DOWHILE DATA STILL AVAILABLE IN QJEJE 00000290
  32FM     GET NEXT RELATIVE BATCH NUMBER FROM QJEJE 00000300
  04FM     COMPUTE FIRST TRANSACTION SECTOR FOR THIS BATCH 00000310
  04FM     DOWHILE ALL TRANSACTION SECTORS INITIALIZED 00000320
  34FM       ENQJEJE FILE MANAGER RESOURCE 00000330
  04FM       WRITE INITIALIZED TRANSACTION SECTOR 00000340
  34FM       DEQJEJE FILE MANAGER RESOURCE 00000350
  ENDDO 00000360
  34FM   ENQJEJE FILE MANAGER RESOURCE 00000370
  GET RELATIVE BATCH NUMBER 00000380
  READ BATCH FILE SECTOR FOR INDICATED BATCH 00000390
  SET STATUS TO UNUSED IN INDICATED BATCH RECORD 00000400
  WRITE MODIFIED SECTOR BACK TO BATCH FILE 00000410
  SET STATUS TO UNUSED IN RESIDENT BATCH FILE RECORD 00000420
  DEQUEJE FILE MANAGER RESOURCE 00000430
  ENDDO 00000440
  DELAY FOR A WHILE 00000450
ENDDO 00000460
00000470
00000010
00000020
00000030
00000040
00000050

```

ADD BATCH RECORD TO FILES

- MODULE HISTORY -

PROJECT: DAS 75-01709
 SJB-SYSTEM: FILE MANAGER
 MODULE: FMABR 31

- MODULE ABSTRACT -

THIS SUBROUTINE PROCESSES THE ADD BATCH RECORD FILE MANAGER
 CALL CODE. IT FINDS A FREE BATCH RECORD AND ADDS THE NEW
 BATCH RECORD TO THE FILES.

```

12FM  SET INPJTS
      SET RECORD POINTER = ZERO
      SET RECORD USE COUNT = ZERO
02FM  DDUNTIL END OF RESIDENT BATCH FILE OR ERROR FOUND
      GET FIRST/NEXT RESIDENT BATCH RECORD
02FM  IF RECORD STATUS NOT UNUSED OR DELETING THEN
      ADD 1 TO RECORD USE COUNT
      (ELSE)
      ENDIF
12FM  IF INPUT BATCH I.D. = RECORD BATCH I.D. THEN
02FM
02FM  IF RECORD STATUS NOT DELETING OR UNUSED THEN
02FM  SET COMPLETION CODE = BATCH ALREADY ON FILE
      ELSE
02FM  IF RECORD STATUS = UNUSED THEN
      SAVE RECORD NO. IN RECORD POINTER
      (ELSE)
      ENDIF
      ENDIF
      ELSE
02FM  IF RECORD STATUS = ADDING THEN
12FM  SET COMPLETION CODE = BATCH ALREADY BEING ADDED
      ELSE
02FM  IF RECORD STATUS = UNUSED THEN
      SAVE RECORD NO. IN RECORD POINTER
      (ELSE)
      ENDIF
      ENDIF
      ENDIF
      ENDIF
12FM  IF COMPLETION CODE = SUCCESSFUL THEN
      IF RECORD COUNT G.T. OR E.Q. 10 THEN
      SET RECORD POINTER = ZERO
      (ELSE)
      ENDIF
      IF RECORD POINTER NOT ZERO THEN
01FM  COMPUTE BATCH FILE SECTOR ADDRESS FROM RECORD POINTER
28FM  READ THE BATCH FILE SECTOR
01FM  PLACE NEW BATCH RECORD INTO SECTOR
01FM  SET NEW BATCH RECORD STATUS = ADDING
01FM  SET NEW BATCH RECORD COUNTS = ZERO
28FM  WRITE THE UPDATED BATCH SECTOR
02FM  BUILD RESIDENT BATCH RECORD WITH STATUS = ADDING
      ELSE
12FM  SET COMPLETION CODE = BATCH FILE SPACE EXHAUSTED
      ENDIF
      (ELSE)
      ENDIF
12FM  RETURN TO CALLER

```

SEARCH RESIDENT BATCH FILE SUBROUTINE

```

0000060
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0000060

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                                00000070
                                00000080
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                                00000100
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                                00000100
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                                00000370
                                00000380
                                00000390
                                00000400
                                00000410

                                - MODULE HISTORY -

PROJECT:      DAS              75-01709
SUB-SYSTEM:   FILE MANAGER
MODULE:       FMSRB           43

                                - MODULE ABSTRACT -

THIS SUBROUTINE SEARCHES THE RESIDENT BATCH FILE IN
SSYSCOM FOR THE BATCH I.D. INPJT BY THE CALLER.
IT WILL RETURN A "0" IF THE BATCH I.D. IS NOT IN FILE,
OR THE RELATIVE BATCH NO. IF THE BATCH IS FOUND.
33FM  SET INPJTS
      SET RETURN CODE = 0
      DO UNTIL END OF DATA OR BATCH I.D. FOUND
02FM  GET FIRST/LAST RECORD FROM RESIDENT BATCH FILE
02FM  IF BATCH I.D. FROM RECORD = BATCH I.D. INPUT THEN
33FM      IF BATCH STATUS NOT ADDED, DELETING, OR DELETING THEN
33FM          SET RETURN CODE = RELATIVE BATCH NUMBER
          (ELSE)
          ENDIF
      (ELSE)
      ENDIF
      ENDDO
33FM  RETURN TO CALLER

COLD START DATABASE INITIALIZATION

                                - MODULE HISTORY -

PROJECT:      DAS              75-01709
SUB-SYSTEM:   FILE MANAGER
MODULE:       FMINT           33

                                - MODULE ABSTRACT -

THIS MODULE IS CALLED BY FMGR DURING SYSTEM COLD START TO
START THE DELETION PROCESS FOR ANY BATCHES THAT HAVE A STATUS
OF "ADDING OR DELETING". THIS MODULE ALSO BUILDS THE RESIDENT
BATCH FILE IN $SYSCOM FROM CURRENT DATA IN THE DISK FILE "BATCH"
AND COPIES THE INFORMATION FROM THE DISK FILE "XLATE" INTO
THE RESIDENT TRANSLATION FILE, ALSO IN $SYSCOM, AND COPIES
INFORMATION FROM THE DISK FILE "ORDER" INTO THE RESIDENT
ORDER FILE FOUND IN THE FILE MANAGER SUB-SYSTEM.
24FM  GET INPJTS
28FM  BUILD THE RESIDENT ORDER FILE
01FM  DO UNTIL ALL SECTORS IN BATCH FILE ARE PROCESSED
28FM      READ FIRST(NEXT) BATCH FILE SECTOR
01FM      DO UNTIL ALL RECORDS IN SECTOR ARE PROCESSED
01FM          GET FIRST(NEXT) RECORD IN SECTOR
01FM          IF STATUS IN RECORD = ADDING OR DELETING THEN
01FM              SET STATUS IN RECORD TO DELETING
32FM              PUT RELATIVE BATCH NUMBER IN QUEUE FOR DELETION TASK
          (ELSE)
          ENDIF

```

```

02FM      BJILD RESIDENT BATCH FILE RECDR IN $SYSCOM
          ENDDO
28FM      WRITE SECTOR BACK TO BATCH FILE
          ENDDO
28FM      READ SORT TRANSLATION FILE SECTOR INTO $SYSCOM
06FM
24FM      RETURN

```

```

00000420
00000430
00000440
00000450
00000460
00000470
00000480
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00000100
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00000460
00000470
00000480

```

FILE MANAGER CONTROL

- MODULE HISTORY -

```

PROJECT:      DAS              73-01709
SUB-SYSTEM:   FILE MANAGER
MODULE:       FMGR              29

```

- MODULE ABSTRACT -

THIS MODULE CONTROLS ALL INPUT TO THE FILE MANAGER. IT WAITS ON THE FILE MANAGER ENTRY EVENT TO BE POSTED. IT THEN CHECKS FOR A VALID CALL CODE FOUND IN THE BUFFER PASSED TO IT BY THE CALLER. IF THE CALL CODE IS VALID THE SUBROUTINE THAT PROCESSES THAT CODE IS CALLED. WHEN THE SUBROUTINE RETURNS, THIS MODULE WILL POST THE COMPLETION EVENT BACK TO THE CALLER. THE INFORMATION REQUESTED BY THE CALLER IS HELD IN THE BUFFER THAT WAS PASSED TO THE FILE MANAGER UPON ENTRY. ATTACH BATCH DELETION WORK TASK (FMOLD) UNTIL CPU STOPPED

```

25FM      WAIT FOR THE FILE MANAGER ENTRY EVENT TO BE POSTED
09FM      IF THE CALL CODE FOUND IN THE BUFFER IS WITHIN RANGE THEN
26FM      GET ADDRESS OF CALL CODE DISPATCH TABLE
          INDEX INTO DISPATCH TABLE BY CONTROL CODE
          GET WORK SUBROUTINE ADDRESS FROM DISPATCH TABLE
          PLACE ADDRESS INTO SUBROUTINE CALL
          SET FIRST PARAMETER = INPUT BUFFER ADDRESS
09FM      SET COMPLETION CODE IN BUFFER = SUCCESSFUL
          CALL THE WORK SUBROUTINE
          ELSE
09FM      SET COMPLETION CODE = INVALID CALL CODE
          ENDIF
27FM      POST ADDRESS OF INPUT BUFFER BACK TO THE CALLER
          ENDDO
          PRGSTOP

```

```

0026      ;
0027      ; L E G E N D
0028      ;# NUMBER, NUMBER OF
0029      ;AD ADDRESS CH CHANNEL CHAR CHARACTER CM COMMAND
0030      ;COMM COMMUNICATION COUT COMMAND OUT
0031      ;CP COMPUTE CR CARRIAGE RETURN CT CONTROL
0032      ;DC DECODE, DIVERT-COMplete PHOTO
0033      ;DEC DECREMENT DIN DATA IN DOUT DATA OUT
0034      ;DT DATA DV DIVERT
0035      ;EC ENCODE EOF END OF FILE ER ERROR
0036      ;ID INDUCT PHOTO INC INCREMENT INT INTERRUPT
0037      ;JDV DIVERT NUMBER JINDU INDUCT PHOTO NUMBER
0038      ;JKB KEYBOARD NUMBER
0039      ;JPPI PULSE POSITION INDICATOR NUMBER, =0 MEANS REAL TIME CLOCK
0040      ;JUD UPDATE PHOTO NUMBER
0041      ;KB KEYBOARD
0042      ;LAG LOOK-AHEAD QUEUE
0043      ;LE LENGTH LF LINE FEED, LANE FULL PHOTO
0044      ;M MEMORY MG MERGE PHOTO MS MOST SIGNIFICANT
0045      ;MUZ MOVING UPDATABLE ZONE
0046      ;NL NEW LINE
0047      ;O OPTION

```

```

0048 ;PC PROGRAM COUNTER, PROCESS PCR PROCESSOR
0049 ;PLE PACKAGE LENGTH
0050 ;PP PACKAGE PRESENT PHOTO PPI PULSE POSITION INDICATOR
0051 ;PR PRINT
0052 ;RC RECORD RDC SILENT 700 REMOTE DEVICE CT RE READ
0053 ;RLS RELEASE MECHANISM
0054 ;RTC REAL TIME CLOCK
0055 ;SEC SECOND SIN STATUS IN
0056 ;SP STACK POINTER
0057 ;T TIME TE TRAILING EDGE PHOTO
0058 ;TR TRAINING PHOTO
0059 ;TTY TELETYPE
0060 ;UD UPDATE PHOTO
0061 ;WR WRITE XM TRANSMIT XQ EXECUTE
0062 ;
0063 ; SUBROUTINES (MODULES)
0064 ; ALL FLAGS ARE ALWAYS EFFECTED UNLESS OTHERWISE STATED.
0065 ; REGISTERS A, B, C, D, E, H, L
0066 ; EFFECTED, OR SP, PC EFFECTED UNEXPECTEDLY, WILL BE MENTIONED.
0067 ;
0068 ; COMMENTS
0069 ; COMMENTS ARE GROUPED BY SEMICOLONS. CONTINUATION OF THE SAME
0070 ; COMMENT WILL BE INDENTED 2 COLUMNS.
0071 ;
0072 ;
0073 ; REFERENCES
0074 ;
0075 ; (1) "APPLICATION MANUAL FOR PROGRAMMER SORT CONTROLLER"
0076 ;
0077 ; (2) PROGRAMMABLE SORT CONTROLLER SOFTWARE ORGANIZATION CHARTS
0078 ; AND SOFTWARE MANUAL
0079 ;
0080 ; *****
0081 ; ASSEMBLER PRESET VALUES
0082 4000 RAMAD EQU 4000H ;RAM STARTING ADDRESS
0083 0001 CVY01 EQU 1 ;FLAG FOR CONVEYOR RELAY CH
0084 4300 DCBPC EQU 4300H ;DIVERT-COMplete PHOTO-BLOCKING PROCESSOR
0085 03DD EXEC EQU 03DDH
0086 03EB DEBUG EQU 03EBH
0087 0966 COLDS EQU 0966H
0088 0C19 DEVII EQU 0C19H
0089 0C85 APEND EQU 0C85H
0090 0CA2 CASSW EQU 0CA2H
0091 0CB8 CLAKB EQU 0CB8H
0092 0CC4 CLRBP EQU 0CC4H
0093 0CC9 CLRKB EQU 0CC9H
0094 0CDF CLDP EQU 0CDFH
0095 0D17 CPKBP EQU 0D17H
0096 0D1B AD16A EQU 0D1BH
0097 0D25 CPLFM EQU 0D25H
0098 0D28 CPMSK EQU 0D28H
0099 0D3E DC2BD EQU 0D3EH
0100 0D51 DCBCD EQU 0D51H
0101 0D81 DCDV EQU 0D81H
0102 0D9E DIGPC EQU 0D9EH
0103 0DEA DSPLA EQU 0DEAH
0104 0DF4 DSPLM EQU 0DF4H
0105 3A08 DELAY EQU 3A08H
0106 3A89 DPHL EQU 3A89H
0107 3A70 DPBD EQU 3A70H
0108 0E1C ERMPC EQU 0E1CH
0109 0E01 ECDIG EQU 0E01H
0110 0F08 ERPC EQU 0F08H
0111 0F13 ERRPC EQU 0F13H
0112 0F76 FLASH EQU 0F76H
0113 0F91 GDTR E EQU 0F91H
0114 0FA2 GITRE EQU 0FA2H
0115 0FAC GUTRE EQU 0FACH
0116 0FAD GUTR1 EQU 0FADH
0117 0FBA GTREP EQU 0FBAH
0118 0FCB GETFS EQU 0FCBH
0119 100E GETKB EQU 100EH
0120 1065 GETPH EQU 1065H
0121 1097 LNKFS EQU 1097H
0122 10C4 MOVE EQU 10C4H
0123 10CF MOVEB EQU 10CFH
0124 10D8 MTKBB EQU 10D8H
0125 1109 OFCVY EQU 1109H
0126 110F OPUT EQU 110FH
0127 113F PHPPB EQU 113FH
0128 114F PUTFS EQU 114FH
0129 116C PUTKB EQU 116CH
0130 ;
0131 ;
0132 ; PROM DATA
0133 ;
0134 ;
0135 ; FACTORY PRESET PARAMETERS.

```

```

0136
0137 0000          ORG 3D00H
0138 3D00 004B    RAMT1: DW 4B00H      ;COLD START SP VALUE, = RAM TOP + 1
0139 3D02 02      NPPI:  DB 2          ;# PPI'S
0140 3D03 02      NINDU: DB 2          ;# INDUCTION STATIONS
0141 3D04 06      NUPDA: DB 6          ;# UPDATE PHOTOS
0142 3D05 05      NDVT:  DB 5          ;# DIVERTS
0143 3D06 03      PS080: DB 3          ;STARTING ADDRESS OF 1ST PM5080 BOARD;
0144 3D07 FF      DB 0FFH        ; MUST BE 0FFH (FOR COUT OR SIN);
0145 3D08 02      NKB:   DB 2          ; TOTAL # KB'S
0146 3D09 01      MAKB:  DB 1          ;MASTER KB # FOR SYSTEM OPERATIONS & CASSETTES
0147 3D0A 06      NDPLA: DB 6          ;# DISPLAY DIGITS, MUST BE 2, 4 OR 6
0148 3D0B 00      NXLTB: DB 0          ;# TRANSLATE TABLES, <= 16,
0149
0150 3D0C 01      TBSIZ: DB 1          ; 0 MEANS NO TRANSLATION OPTION
0151
0152 3D0D          COMAX: DS 3          ;TRANSLATE TABLE SIZE IN BYTES / 256,
0153 3D10 FFFF      ERO:   DW 0FFFFH      ; MUST BE 1, 2, 4, 8, 10H, 20H, OR 40H
0154 3D12          DS 6          ;ENTRY CODE UPPER BOUND IN BCD.
0155 3D18 FFFF      DW 0FFFFH      ;ERROR MESSAGES
0156 3D1A          DS 6
0157 3D20 FFFF      DW 0FFFFH      ;DIAG MODE ERROR MESSAGE MASK
0158 3D22 00      ERMKI: DB 00H        ;TRAIN MODE ERROR MESSAGE MASK
0159 3D23 00      DB 00H        ;PSC INPUT MESSAGE MASK BYTE 0
0160 3D24 08      DB 08H        ;BYTE 1
0161 3D25 09      DB 09H        ;BYTE 2
0162 3D26 00      DB 00H        ;BYTE 3
0163 3D27 00      DB 00H        ;BYTE 4
0164 3D28 00      ERMKO: DB 00H        ;PSC OUTPUT MESSAGE MASK BYTE 0
0165 3D29 03      DB 03H        ;BYTE 1
0166 3D2A 01      DB 01H        ;BYTE 2
0167 3D2B 11      DB 11H        ;BYTE 3
0168 3D2C 14      DB 14H        ;BYTE 4
0169 3D2D 00      DB 00H        ;BYTE 5
0170 3D2E FFFF      DW 0FFFFH      ;MUST BE ALL 1'S FOR FATAL ERRORS
0171 3D30 00      ERCHU: DB 0          ;ER CHUTE #
0172
0173 3D31 00      SIDE0: DB 0          ; = 0 MEANS CONVEYOR END
0174
0175
0176 3D32 30      SEND0: DB 30H        ;SIDE INDUCT OPTION:
0177
0178
0179
0180
0181 3D33 01      NICC:  DB 1          ; BIT 1 FOR AUTO-REINDUCT AFTER RECIRCULATION.
0182 3D34 01      DQLIM: DB 1          ; BIT 4 FOR SIDE-INDUCT MERGING
0183 3D35 0000    DW 0          ; "SEND" OPTION CODE.
0184 3D37 02      QLIMI: DB 2          ; BITS 0 TO 3 (LS NIBBLE) FOR LENGTH OF FIXED-
0185 3D38 00      DVTRS: DB 0          ; LENGTH CODE WITHOUT "SEND" (0 MEANS OFF),
0186
0187
0188
0189
0190 3D39 00      SCNRO: DB 0          ; BIT 4 FOR ENTER-CODE-BY-"SEND".
0191
0192
0193
0194
0195 3D3A 00      SERDV: DB 0          ; BIT 5 FOR REPEAT-BY-"SEND".
0196
0197 3D3B 00      DCO:   DB 0          ;# INTER-COMPUTER COMMUNICATION CHANNELS
0198
0199
0200
0201
0202 3D3C 02      LFO:   DB 2          ;LAQ SIZE LIMIT, MUST BE 1 TO 8
0203
0204
0205
0206 3D3D 00      LFEDV: DB 0          ;DIVERT RESET OPTION:
0207 3D3E 0000    SYS9:  DW 0          ; 0 FOR MAINTAINED DV & NO RESET UNTIL
0208 3D40 00      SYSTM: DB 0          ; NEXT PACKAGE
0209 3D41 01      DB 1          ; 1 FOR MEASURED RESET ACCORDING TO PACKAGE LE
0210 3D42 00      DB 0          ; 2 FOR HARDWARE RESET INITIATED BY PULSE
0211 3D43 03      DB 3          ;SCANNER OPTIONS:
0212 3D44 00      DB 0          ; BIT 0 TO ACTIVATE DIVER # DEFAULT.
0213 3D45          DS 1          ; BIT 1 TO ACTIVATE TRANSLATE DEFAULT;
0214
0215 3D46          ORG 3DC0H      ; BIT 6 TO ACTIVATE PARITY CHECKING.
0216 3DC0 0080    XLTBS: DW 8000H      ; BIT 7 FOR ODD PARITY
0217 3DC2 0090    DW 9000H      ;AUTO JDV FOR SCANNER FAIL
0218 3DC4 00A0    DW 0A000H     ; =0 MEANS RECIRCULATION
0219
0220 3DC6          ORG 3DF0H      ;DV-COMplete OPTION:
0221 3DF0 004C    TRSDB: DW 4C00H     ; -0 MEANS NO DIVERT-COMplete
0222 3DF2 FF4F    TLAST: DW 4FFFH     ; =1 MEANS 1 DIVERT-COMplete PHOTO PER DIVERT
0223

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0224      ; R A M   D A T A
0225      ; (PUT IN FRONT FOR CONDITIONAL MACRO EXPANSIONS)
0226      ;
0227      ;
0228      ;FIELD SET PARAMETERS FOR RUNNING MODE.
0229      ;
0230      3DF4      ORG   RAMAD
0231      4000 0000  FLAGS:  DW   0      ;BIT 0 (LS BIT) TO ENABLE ALL ERROR MESSAGES;
0232      ;                               ;BIT 7 (128) FOR AUTOMATIC INDUCTION
0233      4002 0000  SLIP:   DW   0      ;INERTIAL SLIPPAGE ALLOWANCE FOR
0234      ;                               ; HOT START
0235      4004 0000  PLEMX:  DW   0      ;T FROM RLS TO TECZ
0236      ;                               ; = MAX PLE + TE DELAY + SAFETY FACTOR
0237      4006 0000  HDSEN:  DW   0      ;HALF HOST MODE ENABLED
0238      4008 0000  DVH:    DW   0      ;FOR DVTRS=1, = MAX DIVERT HOLD TIME - PLE,
0239      ;                               ; FOR DVTRS=2, = PULSE LENGTH
0240      400A 0000  PPD:    DW   0      ;PP DWELL FOR PUTTING IN GAP
0241      400C 0000  GAP:    DW   0      ;ADDITIONAL RELEASE DELAY FOR A PACKAGE NOT
0242      ;                               ; FULLY STOPPED
0243      DCTT:      ;                               ;OVERALL MAX TIME FROM DIVERT SET
0244      ;                               ; TO DIVERT-COMplete PHOTO BLOCKED
0245      DCZSZ:    ;                               ;DC ZONE SIZE
0246      400E 0000  DCPSZ:  DW   0      ;MIN PACKAGE SIZE AT DC
0247      4010      TTBP:   DS   01E0H  ;FOR TIMING TABLES
0248      ;
0249      ;
0250      41F0      ORG   RAMAD+0200H
0251      ;
0252      ;BREAKPOINT VARIABLES.
0253      ;
0254      4200 0000  RS1AD:  DW   0      ;RESTART
0255      4202 EB03  RS2AD:  DW   DEBUG  ; MODULE
0256      4204 0000  RS3AD:  DW   0      ; AD'S
0257      4206 0000  RS4AD:  DW   0      ;
0258      4208 0000  RSSAD:  DW   0      ;
0259      420A 0000  RSGAD:  DW   0      ;
0260      420C 00   BPM:    DB   0      ;BREAKPOINT M CONTENT (TO BE RESTORED)
0261      420D 0000  ENTAD:  DW   0      ;"GO" COMMAND ENTRANCE AD
0262      420F 0000  REGAF:  DW   0      ;REGISTERS A,F VALUES
0263      4211 0000  REGBC:  DW   0      ;REGISTERS B,C VALUES
0264      4213 0000  REGDE:  DW   0      ;REGISTERS D,E VALUES
0265      4215 0000  REGHL:  DW   0      ;REGISTERS H,L VALUES
0266      4217 0000  BP:     DW   0      ;BREAKPOINT
0267      4219 00   CVCHN:  DB   0      ;COMPOSITE CHN # FOR CONVEYOR RELAY
0268      ;
0269      ;DEBUGGER AND CASSETTE VARIABLES.
0270      ;
0271      421A 00   CMCH:   DB   0      ;COMMAND CHARACTER
0272      421B 00   JRC:    DB   0      ;RECORD INDEX
0273      421C 00   RCLE:   DB   0      ;RECORD LENGTH
0274      421D 0000  DUMP0:  DW   0      ;MEMORY DUMP START AD
0275      421F 0000  DUMP9:  DW   0      ;MEMORY DUMP END AD
0276      4221 0000  LOFFS:  DW   0      ;LOAD M OFFSET
0277      4223 0000  MISMP:  DW   0      ;POINTS TO MOST RECENT MISMATCH BYTE
0278      ;
0279      ;VALUES NOT EFFECTED BY COLD START.
0280      ;
0281      4225 3C   MODE:   DB   '<'  ; '<' FOR RUN, '=' FOR DIAGNOSTIC,
0282      ;                               ; '>' FOR TRAINING, ':' FOR INTER-COMPUTER
0283      4226 08   MODEE:  DB   8      ;NOT=0 WILL ENABLE MODE TRANSITION
0284      4227 0000  FSP0:  DW   0      ;BASE AD OF FIELD-SET PARAMETERS
0285      4229 35   RS:     DB   35H   ;ASCII 5
0286      422A 35   LMR:    DB   35H   ;ASCII 5
0287      422B 00   ICHO:   DB   0      ;INTERCOMPUTER MODE ACTIVE FLAG
0288      422C 0200  TDLTM:  DW   0002H  ;MESSAGE TIMEOUT IN 0.1 SEC INCREMENTS
0289      422E 0F   VICLG:  DB   0FH   ;I/C LIGHT REFRESH TIME IN 0.1 SEC INCREMENTS
0290      ;
0291      ;TABLE POINTERS TO BE SET BY COLD START & NEVER CHANGED.
0292      ;
0293      CSRAM:
0294      422F 0000  PRET:   DW   0      ;PRE-INDUCT TIME TABLE: IF STRAIGHT MERGE
0295      ;                               ; THEN MAX T FROM RLS TO ID,
0296      ;                               ; ELSE UDBPC TO BMG
0297      4231 0000  TET:    DW   0      ;TE DELAY TABLE
0298      4233 0000  IDT:    DW   0      ;MIN T FROM ID-BLOCKING TO UD1-BLOCKING TABLE;
0299      ;                               ; IF SIDE-INDUCT, RELEASE TO UDB1 TIME
0300      4235 0000  ZNT:    DW   0      ;UDB1 TO UDB2 TIMING TABLE
0301      4237 0000  UDT:    DW   0      ;UPDATE PHOTO TIMING TABLE
0302      4239 0000  DVT:    DW   0      ;DIVERT TIMING TABLE
0303      423B 0000  DCT:    DW   0      ;MIN T LAST UD-BLOCKING TO DC-BLOCKING TABLE
0304      ;
0305      423D 0000  LDVDB:  DW   0      ;LDV DATABASE IN PROM
0306      423F 0000  FIDDB:  DW   0      ;SIDE INDUCT # FOLLOWING UPDATE
0307      4241 0000  ICDB:   DW   0      ;INTER-COMPUTER DATABASE
0308      4243 0000  PPDDB:  DW   0      ;PACKAGE-PRESENT DATABASE
0309      4245 0000  PPIDB:  DW   0      ;PPI DATABASE FOR TREE PARAMETERS
0310      ;                               ; PPI # 0 MEANS RTC

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0311 4247 0000  KBDB:  DW  0      ;KB DATA BASE.
0312 4249 0000  KBIBF: DW  0      ; KB INPUT BUFFERS.
0313 424B 0000  KB0BF: DW  0      ; KB ECHO BUFFERS
0314 424D 0000  XLDB:  DW  0      ;TRANSLATE DATABASE
0315 424F 0000  RLSDB: DW  0      ;RLS DATABASE
0316 4251 0000  UDDB:  DW  0      ;UPDA PHOTO DATABASE
0317 4253 0000  Lfdb:  DW  0      ;LF PHOTO DATABASE
0318 4255 0000  DVDB:  DW  0      ;DIVERT DATABASE FOR DC
0319 4257 0000  DCDB:  DW  0      ;DV-COMLETE DATABASE FOR DC
0320 4259 0000  OUTDB: DW  0      ;PM5007 OUTPUT BOARD IMAGES
0321 425B 0000  FS:    DW  0      ;FREE STORAGE FOR TCB'S IN RUNNING MODE
0322      ;
0323      ;VALUES TO BE SENT IN COLD START & NEVER CHANGED.
0324      ;
0325 425D AA55  RAM01: DW  55AAH   ;FOR TESTING VALIDITY OF RAM
0326 425F AA55  RAM02: DW  55AAH   ; ON POWER-UP
0327 4261      HTSL:  DS  5      ;FOR PM5008 INITIALIZATION
0328 4266 04FF  MKBCO: DW  0FF04H  ;MASTER KB COUNT AD
0329 4268 00    SL700: DB  0      ;NOT=0 MEANS MASTER KB IS SILENT 700
0330 4269 0000  K5080: DW  0      ;PM5080 SIN AD FOR 1ST KB
0331 426B 00    N5080: DB  0      ;$ PM5080 BOARDS
0332 426C 00    P5008: DB  0      ;1ST PM5008 BOARD AD
0333 426D 00    P5007: DB  0      ;1ST PM5007 BOARD AD
0334 426E 00    N5007: DB  0      ;$ PM5007 BOARDS
0335 426F 00    PPCHN: DB  0      ;1ST PP PM5008 COMPOSITE CHANNEL $
0336 4270 00    IDCHN: DB  0      ;1ST ID CH $
0337 4271 00    UDCHN: DB  0      ;1ST UD CH $
0338 4272 00    DCCHN: DB  0      ;1ST DC CH $
0339 4273 00    LFCHN: DB  0      ;1ST LF CH $
0340 4274 00    TRCHN: DB  0      ;1ST TR CH $
0341 4275 00    ILCHN: DB  0      ;OUTPUT CHANNEL $ FOR INTER-COMPUTER LIGHT
0342      ;
0343      ;
0344      ;SCRATCH VARIABLES FOR ALL MODES. INITIALLY ALL 0.
0345      ;
0346 4276 00    XYZ9:  DB  0      ;KB
0347 4277 00    XYZ8:  DB  0      ; DISPLAY
0348 4278 00    XYZ7:  DB  0      ; DIGITS
0349 4279 00    XYZ6:  DB  0      ;
0350 427A 00    XYZ5:  DB  0      ;
0351 427B 00    XYZ4:  DB  0      ;
0352 427C 00    XYZ3:  DB  0      ;
0353 427D 00    XYZ2:  DB  0      ;
0354 427E 00    XYZ1:  DB  0      ;
0355 427F 00    XYZ0:  DB  0      ;(LS ASCII DIGIT)
0356 4280 00    XYZ:   DB  0      ;MUST BE 0 (SEE TYPEN)
0357 4281 0000  PFSP:  DW  0      ;SP FOR POWER FAIL HOT START
0358 4283 0000  NNOIS: DW  0      ;$ NOISE INT'S PER 8 RTC INT'S
0359 4285 00    PCRON: DB  0      ;FLAG TO FLASH PCRON LIGHT
0360 4286 0000  FLSHR: DW  0      ;SPARE FLASHER FLAG
0361 4288 0000  RGAF:  DW  0      ;ERRPC
0362 428A 0000  RGBC:  DW  0      ; INPUT
0363 428C 0000  RGDE:  DW  0      ; REGISTERS'
0364 428E 0000  RGHL:  DW  0      ; VALUES
0365 4290 0000  RGPC:  DW  0      ;
0366 4292 00    INDEX: DB  0      ;DEVICE INDEX
0367 4293 00    DEVIC: DB  0      ;DEVICE CODE
0368 4294 00    JINDU: DB  0      ;INDUCTION STATION $
0369 4295 00    JUPDA: DB  0      ;CURRENT UPDA PHOTO $
0370 4296 00    JDV:   DB  0      ;CURRENT DIVERT $ AND
0371 4297 0000  IDENT: DW  0      ; PACKAGE IDENTIFICATION CODE
0372 4299 0000  UDBT:  DW  0      ;UPDA BLOCKING TIME
0373 429B 0000  DVST:  DW  0      ;DIVERT SET TIME
0374 429D 0000  VALUE: DW  0      ;D4BCD INPUT
0375 429F 0000  CODE:  DW  0      ;PACKAGE CODE LS 2 BYTES; OR PARAMETER CODE
0376 42A1 0000  CODE9: DW  0      ; PACKAGE CODE MS BYTE; OR OLD PARAMETER CODE
0377      ;
0378      ;MMNPL DATA
0379      ;
0380 42A3 00    MNPLC: DB  0      ;DIGIT COUNT ( >= 0 ) WHEN ER = 2 OR 5
0381 42A4 0000  MLOC:  DW  0      ;POINTS TO M BYTE TO BE DISPLAYED
0382 42A6 00    MBYTE: DB  0      ;M REPLACEMENT BYTE
0383 42A7 00    MTBD:  DB  0      ;M TO BE DISPLAYED
0384      ;
0385      ;RUNNING MODE SCRATCH VARIABLES. INITIALLY ALL 0.
0386      ;
0387 42A8 0000  FSQ1:  DW  0      ;START OF FREE STORAGE QUEUE
0388 42AA 0000  FSQ9:  DW  0      ; END OF FREE STORAGE QUEUE
0389 42AC 00    FSQ0:  DB  0      ;FREE STORAGE OUT FLAG TO LOCK OUT INDUCTS. &
0390 42AD 00    NTRE0: DB  0      ; COUNT DOWN OF $ EMPTY TREES
0391 42AE 0000  NNTR0: DW  0      ;COUNT DOWN OF $ TIMES TREES ALL EMPTY
0392 42B0 0000  COUNT: DW  0      ;TOTAL $ PACKAGES THRU UPDA $1
0393 42B2 0000  CTIMP: DW  0      ;POINTS TO CTIME FOR CURRENT PPI
0394 42B4 0000  TCBP:  DW  0      ;TCB POINTER
0395      ;
0396      ;TREE PARAMETERS FOR THE CURRENTLY USED PPI.
0397      ;
0398 42B6 0000  CTIME: DW  0      ;PPI CLOCK TIME

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0399 42B8 0000 HEAD: DW 0 ;TREE HEADER
0400 42BA 0000 RLINK: DW 0 ;POINTS TO TREE ROOT
0401 42BC 00 DB 0 ;WOULD-BE TREE BALANCE FACTOR
0402 42BD 0000 NXTIM: DW 0 ;NEXT EVENT DUE TIME
0403 42BF 00 JPPI: DB 0 ;CURRENT PPI #
0404 42C0 00 MUD: DB 0 ;LAST UD # FOR THIS PPI
0405 42C1 DS 5 ;UNUSED FILLER FOR PPIDB DATA TRANSFER
0406 42C6 0000 STKSV: DW 0 ;SP SAVE
0407 42C8 0000 QP: DW 0 ;TREE PROCESSOR
0408 42CA 0000 QX: DW 0 ; VARIABLES
0409 42CC 0000 QT: DW 0 ;
0410 42CE 0000 QW: DW 0 ;
0411 42D0 0000 QS: DW 0 ;
0412 42D2 0000 QR: DW 0 ;
0413 42D4 0000 QDIR: DW 0 ;
0414 42D6 0000 NEWTM: DW 0 ;NEW EVENT TIME
0415 42D8 0000 PPIDP: DW 0 ;POINTER TO CTIME IN PPIDB
0416 ;
0417 ;KEYBOARD DATA
0418 ;
0419 42DA 0000 KBDPB: DW 0 ;KB DATABASE POINTER
0420 42DC 0000 KBCM: DW 0 ;AD FOR KB COUT OR SIN
0421 42DE 00 CHAR: DB 0 ;KB INPUT CHARACTER
0422 42DF 00 NDIG: DB 0 ;# DIGITS ENTERED
0423 ;
0424 ;TRANSLATION DATA
0425 ;
0426 42E0 00 XL: DB 0 ;TRANSLATE FLAG FOR PACKAGE ENCODING
0427 42E1 0000 DELTA: DW 0 ;HASH FUNCTION OFFSET
0428 42E3 01 JXL: DB 1 ;TRANSLATE TABLE #
0429 42E4 01 MJXL: DB 1 ;MASTER KB TRANSLATE TABLE #
0430 42E5 0000 NCODE: DW 0 ;# CODE ENTRIES POSSIBLE PER TABLE
0431 42E7 0000 XLTB: DW 0 ;TRANSLATE TABLE BASE
0432 42E9 0000 XLTBP: DW 0 ;TRANSLATE TABLE ENTRY POINTER.
0433 ;
0434 ;INDUCTION DATA
0435 ;
0436 42EB 0000 PPDBP: DW 0 ;POINTS TO PPDB
0437 42ED 0000 PSEQ: DW 0 ;POINTS TO PACKAGE SEQUENCE THRU ID
0438 42EF 0000 RLDBP: DW 0 ;POINTS TO RLSDB
0439 42F1 0000 NEXTP: DW 0 ;NEXT TCB POINTER
0440 ;
0441 ;UPDATE PHOTO DATA
0442 ;
0443 42F3 0000 PMUZ: DW 0 ;SEE Uddb
0444 42F5 0000 PMUZP: DW 0 ;PMUZ POINTER
0445 42F7 0000 LDVP: DW 0 ;POINTS TO LDV ITEM
0446 ;
0447 ;DIVERT DATA
0448 ;
0449 42F9 0000 OACT: DW 0 ;CURRENT OACT IN TCB
0450 42FB 00 LFDV: DB 0 ;FOR LANE-FULL-CLEAR AUTOMATIC RE-INDUCTION
0451 42FC 0000 LFP: DW 0 ;POINTS TO LF IN LFDV
0452 42FE 00 CS007: DB 0 ;=N5007 - BOARD #
0453 42FF 0000 OUTIM: DW 0 ;TEMPORARY PM5007 OUTPUT IMAGE
0454 4301 0000 DCDBP: DW 0 ;DCDB ENTRY POINTER
0455 ;
0456 ;
0457 ;DIAGNOSTIC & TRAINING MODE SCRATCH VARIABLES.
0458 ; INITIALLY ALL 0.
0459 ;
0460 4303 00 CMCHR: DB 0 ;CM CHAR
0461 4304 00 XC: DB 0 ;EXERCISER FLAG
0462 4305 0000 SOURC: DW 0 ;SOURCE PHOTO DEVICE CODE (MS BYTE) & INDEX
0463 4307 0000 DEST: DW 0 ;DESTINATION PHOTO DEVICE CODE & INDEX
0464 4309 0000 BASE: DW 0 ;TRAINING TABLE BASE ADDRESS
0465 ;
0466 ;DIAGNOSTIC MODE & TRAINING MODE INITIALLY NONZERO VARIABLES.
0467 ;
0468 430B 01 SUBMO: DB 1 ;SUBMODE #
0469 430C 4000 XCITV: DW 64 ;EXERCISE INTERVAL IN PPI'S
0470 ; MUST BE A POWER OF 2
0471 430E 01 KPPI: DB 1 ;CURRENTLY ACTIVE PPI #
0472 430F 02 DVXYZ: DB 2 ;CURRENT DIVERT STATE.
0473 ; =1 MEANS SET, =2 MEANS RESET
0474 ;
0475 ;INTER-COMPUTER MODE DATA
0476 ;
0477 4310 00 XMRBS: DB 0 ;TRANSMITTER BUSY FLAG
0478 4311 00 DB 0 ;
0479 4312 0000 DW 0 ;
0480 4314 0000 DW 0 ;
0481 4316 0000 DW 0 ;
0482 4318 BYT16: DS 16 ;MESSAGE TEMPORARY BUFFER
0483 4328 DASDB: DS 16 ;PSC INPUT MESSAGE BUFFER
0484 4338 MOBUF: DS 16 ;PSC OUTPUT MESSAGE BUFFER
0485 4348 00 CSEQ: DB 0 ;COUNTER INDICATING BAD SEQ BYTE

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0486 4349 00 CSUM: DB 0 ;COUNTER INDICATING BAD CHECKSUM
0487 434A 00 CTOUT: DB 0 ;COUNTER INDICATING MESSAGE TIMEOUT
0488 434B 0000 DASIP: DW 0 ;POINTER IN PSC INPUT MESSAGE BUFFER
0489 434D 00 DTYBS: DB 0 ;FLAG INDICATING PSC IS OUTPUTTING TO I/C 5080
0490 434E 00 FDV: DB 0 ;FLAG INDICATING A DVT# OF 0
0491 434F 00 FLG9: DB 0 ;FLAG USED TO INDICATE MESSAGE 9
0492 4350 00 FLINK: DB 0 ;FLAG INDICATING I/C LINK STATUS
0493 4351 00 FMBAD: DB 0 ;FLAG INDICATING BAD MESSAGE RECEIVED BY PSC
0494 4352 00 FULBF: DB 0 ;FLAG INDICATING FULL PSC INPUT MESSAGE BUFFER
0495 4353 00 ICLGT: DB 0 ;I/C LIGHT COUNTER = VICLG
0496 4354 0000 LMP: DW 0 ;LOAD MESSAGE POINTER IN FS
0497 4356 0000 MBUFF: DW 0 ;POINTER IN PSC OUTPUT MESSAGE BUFFER
0498 4358 0000 SMP: DW 0 ;SEND MESSAGE POINTER IN FS
0499 435A 00 TCONT: DB 0 ;COUNTER USED IN TRANSACTION TABLE LOADING
0500 435B 0000 TMOUT: DW 0 ;MESSAGE TIMEOUT = TDLTM
0501 435D 0000 TRSDP: DW 0 ;POINTER IN TRANSACTION TABLE
0502
0503 ;GENERAL PURPOSE TEMPORARY VARIABLES
0504
0505 435F 0000 T2: DW 0 ;TEMPORARY
0506 4361 DS 27 ;(FILLER)
0507
0508 ;PHOTO INTERRUPT BUFFER DATA
0509 ; N.B.: BUFFER OF LENGTH 64 MUST START AT ADDRESS DIVISIBLE BY 64. ****
0510
0511 437C 8043 PH1: DW PHBF ;PHOTO IN-POINTER
0512 437E 8043 PH2: DW PHBF ;PHOTO OUT-POINTER
0513 4380 PHBF: DS 64 ;CIRCULAR BUFFER FOR 31 OR LESS INTERRUPTS
0514
0515 ;VARIABLE SIZE TABLES AND FREE STORAGE.
0516 ; N.B.: MUST START AT MULTIPLE OF 32. ****
0517
0518 43C0 TBS: DS 100 ;(DYNAMIC TABLES START HERE)
0519
0520 ;
0521 ;MISCELLANEOUS MACROS
0522 ;
0523 ;
0524 ;JUMP TO NXBC IF B.C NOT= AD OF XBC.
0525 ;REGISTER A,D, E EFFECTED.
0526 ;
0527 JBCN MACRO XBC,NXBC
0528 LXI D,XBC
0529 MOV A,C
0530 CMP E
0531 JNZ NXBC
0532 MOV A,B
0533 CMP D
0534 JNZ NXBC
0535 ENDM
0536 ;
0537 ;
0538 ;JUMP TO MEQ IF M CONTENTS AT H.L & H.L+2 ARE EQUAL, ELSE FALL THRU.
0539 ;
0540 ;OUTPUT: B.C = 1ST CONTENT, D,E = 2ND CONTENT, H.L POINTS TO MS BYTE.
0541 ;
0542 JMEQ MACRO MEQ
0543 MOV C,M ;1ST CONTENT IN B.C
0544 INX H
0545 MOV B,M
0546 INX H ;2ND CONTENT IN D.E
0547 MOV E,M
0548 INX H
0549 MOV D,M
0550 MOV A,E
0551 CMP C
0552 JNZ S+8
0553 MOV A,D
0554 CMP B
0555 JZ MEQ
0556 ENDM
0557 ;
0558 ;
0559 ;H.L = H.L - XXX. A, E, E EFFECTED. FLAG C SET IF <0.
0560 ;
0561 SUBX MACRO XXX
0562 XCHG ;D,E=SUBTRAHEND
0563 LHLD XXX ;H.L=XXX
0564 MOV A,E
0565 SUB L
0566 MOV L,A
0567 MOV A,D
0568 SBB H
0569 MOV H,A
0570 ENDM
0571 ;
0572 ;

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0573          ;BALANCED BINARY TREE MACRO DEFINITION
0574          ;
0575          ;
0576          ;GET VALUE OF LEFT LINK
0577          ;
0578          GTLLK   MACRO SRCNO,DEST
0579                  IF SRCNO          ;IS NODE ADDRESS IN (H.L)?
0580                  LHL D SRCNO        ;NO-NODE POINTER TO (H.L)
0581                  ENDIF
0582                  MOV  E,M           ;MOVE THE LEFT
0583                  INX  H             ;LINK INTO THE
0584                  MOV  D,M           ;(D,E) REGISTER
0585                  IF  DEST          ;IS IT SAVE LINK MODE?
0586                  XCHG              ;YES - LINK TO (H.L)
0587                  SHLD DEST         ;LINK TO MEMORY
0588                  ENDIF
0589                  ENDM
0590          ;
0591          ;GET VALUE OF RIGHT LINK
0592          ;
0593          GTRLK   MACRO SRCNO,DEST
0594                  IF SRCNO          ;IS NODE ADDRESS IN (H.L)?
0595                  LHL D SRCNO        ;NO - NODE POINTER TO (H.L)
0596                  ENDIF
0597                  INX  H             ;STEP POINTER TO
0598                  INX  H             ;THE RIGHT LINK
0599                  MOV  E,M           ;MOVE THE RIGHT
0600                  INX  H             ;LINK INTO THE
0601                  MOV  D,M           ;(D,E) REGISTER
0602                  IF  DEST          ;IS IT SAVE LINK MODE?
0603                  XCHG              ;YES - LINK TO (H.L)
0604                  SHLD DEST         ;LINK TO MEMORY
0605                  ENDIF
0606                  ENDM
0607          ;
0608          ;SET VALUE OF LEFT LINK
0609          ;
0610          PTLK    MACRO DSTNO,SOURC
0611          MCTEM   SET 1              ;ASSUME ADDRESS IN (H.L)
0612                  IF DSTNO          ;TEST ASSUMPTION
0613          MCTEM   SET 0              ;FALSE ASSUMPTION
0614                  ENDIF
0615                  IF SOURC          ;IS NEW LINK IN (D,E)?
0616                  IF MCTEM          ;NO - BUT IS POINTER IN (H.L)?
0617                  XCHG              ;YES - SAVE NODE POINTER
0618                  ENDIF
0619                  LHL D SOURC        ;NEW LINK TO (H.L)
0620                  XCHG              ;NEW LINK TO (D,E)
0621                  ENDIF
0622                  IF DSTNO          ;IS NODE ADDRESS IN (H.L)?
0623                  LHL D DSTNO        ;NO - NODE POINTER TO (H.L)
0624                  ENDIF
0625                  MOV  M,E           ;MOVE THE NEW
0626                  INX  H             ;LEFT LINK VALUE
0627                  MOV  M,D           ;INTO THE NODE
0628                  ENDM
0629          ;
0630          ;SET VALUE OF RIGHT LINK
0631          ;
0632          PTRLK   MACRO DSTNO,SOURC
0633          MCTEM   SET 1              ;ASSUME ADDRESS IN (H.L)
0634                  IF DSTNO          ;TEST ASSUMPTION
0635          MCTEM   SET 0              ;FALSE ASSUMPTION
0636                  ENDIF
0637                  IF SOURC          ;IS NEW LINK IN (D,E)?
0638                  IF MCTEM          ;NO - BUT IS POINTER IN (H.L)?
0639                  XCHG              ;YES - SAVE NODE POINTER
0640                  ENDIF
0641                  LHL D SOURC        ;NEW LINK TO (H.L)
0642                  XCHG              ;NEW LINK TO (D,E)
0643                  ENDIF
0644                  IF DSTNO          ;IS NODE ADDRESS IN (H.L)?
0645                  LHL D DSTNO        ;NO - NODE POINTER TO (H.L)
0646                  ENDIF
0647                  INX  H             ;NOW STEP TO THE
0648                  INX  H             ;RIGHT LINK BYTES
0649                  MOV  M,E           ;MOVE THE NEW
0650                  INX  H             ;RIGHT LINK VALUE
0651                  MOV  M,D           ;INTO THE NODE
0652                  ENDM
0653          ;
0654          ;GET BALANCE FACTOR OF NODE
0655          ;
0656          GTBAL   MACRO NODPT,DEST
0657                  IF NODPT          ;IS NODE ADDRESS IN (H.L)?
0658                  LHL D NODPT        ;NO - NODE POINTER TO (H.L)
0659                  ENDIF

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0660          INX   H           ;STEP POINTER TO
0661          INX   H           ;THE BALANCE BYTE
0662          INX   H
0663          INX   H
0664          MOV   A,M         ;BALANCE CODE TO A-REG
0665          IF   DEST         ;TEST IF SAVE NODE
0666          STA  DEST         ;YES - CODE TO MEMORY
0667          ENDIF
0668          ORA   A           ;SET CONDITION CODES
0669          ENDM
0670
0671          ;PUT BALANCE FACTOR INTO A NODE
0672
0673          PTBAL  MACRO NODPT,SOURC
0674          IF   SOURC         ;IS FACTOR IN A-REG?
0675          LDA  SOURC         ;NO - BALANCE TO A-REG
0676          ENDIF
0677          IF   NODPT         ;IS NODE ADDRESS IN (H,L)
0678          LHL  NODPT         ;NO - NODE POINTER TO (H,L)
0679          ENDIF
0680          INX   H           ;STEP POINTER TO
0681          INX   H           ;THE BALANCE BYTE
0682          INX   H
0683          INX   H
0684          MOV   M,A         ;BALANCE CODE TO NODE
0685          ENDM
0686
0687
0688          4424          ORG   1400H ;RUN MODE EXEC
0689
0690          ;RUNNING MODE EXECUTIVE.
0691
0692          ;INPUT: INTERRUPTS DISABLED, MODE DATA, COLD-START DATA.
0693          ;ALL REGISTERS & MEMORY EFFECTED.
0694
0695          1400 CDAC1A  REXEC:  CALL  RINI         ;:INI FOR RUNNING MODE
0696
0697          ;CLEAR KB'S
0698
0699          1403 FB      EI           ;:ENABLE INTERRUPTS
0700          1404 CDB80C  CALL  CLAKB        ;:CLEAR ALL KB'S
0701          1407 C30D14  JMP   REXEC1       ;:ENTER MAIN LOOP
0702
0703          ;PROCESS EVERYTHING IN PSCBF
0704
0705          140A CD2F1B  RPHPC:  CALL  PHPC         ;:PROCESS ALL PHOTOS
0706          140D CD6510  REXEC1: CALL  GETPH        ;:GET B=DEVIC.
0707          1410 C20A14  JNZ   RPHPC         ;: C=INDEX
0708
0709          ;FLASH WATCH-DOG LIGHT
0710
0711          1413 CD760F  CALL  FLASH        ;:FLASH PSC-ON LIGHT
0712
0713          ;IF FREE STORAGE IS OUT THEN LOCK OUT INDUCTS
0714
0715          1416 3AAC42  LDA   FSQ0         ;:IF FREE STORAGE
0716          1419 3D      DCR   A           ;: IS
0717          141A FA3614  JM    RKBPC         ;: OUT
0718          141D CA3C14  JZ    REVNT         ;:IF INDUCTS HAVE NOT YET
0719          1420 32AC42  STA   FSQ0         ;: BEEN LOCKED OUT
0720
0721          1423 3A033D  LDA   NINDU        ;:C= INDUCT $.
0722          1426 4F      MOV   C,A         ;: = NINDU INITIALLY
0723          1427 CD1422  RMCID:  CALL  MCLPC        ;:MASTER CLEAR INDUCT
0724          142A 3E0F    MVI   A,15        ;:ERROR
0725          142C CD130F  CALL  ERRPC        ;: 15
0726          142F 0D      DCR   C           ;:DUNTIL ALL
0727          1430 C22714  JNZ   RMCID        ;: CLEARED
0728          1433 C33C14  JMP   REVNT        ;:IGNORE KB INPUTS
0729
0730          ;PROCESS A CHAR IN KBIBF
0731
0732          1436 CD0E10  RKBPC:  CALL  GETKB        ;:A = INPUT CHAR, C = KB $.
0733          1439 C4501B  CNZ   KBPC         ;: KBDBP = HL POINTS TO KBDB BYTE 4
0734          ;:PROCESS ANY KB INPUT CHAR
0735
0736          ;PROCESS EVENTS DUE OR OVERDUE IN ALL AVL TREES
0737          ;: TREE $ 0 IS FOR RTC, TREES $ 1 TO $ NPPI ARE FOR PPI'S
0738          ;: (BEWARE OF CRITICAL SECTIONS);
0739          ;: RTC TREE IS USED FOR LOW PRIORITY DELAYED TASKS
0740
0741          143C 3A023D  REVNT:  LDA   NPPI         ;:A= $ PPI'S = $ TREES - 1
0742          143F 32AD42  STA   NTRE0        ;:COUNT DOWN $ EMPTY TREES
0743          1442 2A4542  LHL  PPIDB        ;:POINT TO CTIME IN PPIDB
0744          1445 22D842  SHLD PPIDP        ;:SAVE PPIDB POINTER
0745          1448 E5      PUSH  H           ;:SAVE AGAIN (GTREP MAY CHANGE PPIDP)
0746          1449 F5      PUSH  PSW         ;:SAVE TREE COUNT
0747          144A 11B642  LXI  D,CTIME      ;:D.E POINTS TO CTIME

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0748 144D 0610      MVI B,10H      ;GET TREE
0749 144F F3        DI              ; PARAMETERS
0750 1450 CDCF10    CALL MOVEB     ; (CRITICAL SECTION WITH
0751 1453 FB        EI              ; RTCIN & PPIIN)
0752 1454 2ABA42   LHL RLINK     ; IF THE TREE
0753 1457 7D       MOV A,L        ; IS NOT
0754 1458 B4       ORA H          ; EMPTY
0755 1459 CA8E14   JZ TREE0      ;
0756 145C 2ABD42   LHL NXTIM     ;H,L=NXTIM
0757 145F 3AB642   LDA CTIME     ; IF CTIME-NXTIM >= 0,
0758 1462 95       SUB L          ; I.E.,
0759 1463 3AB742   LDA CTIME+1   ; IF EVENT
0760 1466 9C       SBB H          ; IS OVERDUE
0761 1467 FA9214   JM NXTRE      ; OR DUE
0762                ;
0763                ; ;REMOVE EVENT FROM TREE & EXECUTE IT
0764                ;
0765 146A CD4E27    CALL DEQUE     ;H,L POINTS TO EVENT TCB
0766 146D 22B442   SHLD TCBP    ;TCBP POINTS TO TCB
0767 1470 2AD842   LHL PPIDP    ;POINT D,E TO
0768 1473 EB       XCHG         ; TREE HEADER IN PPIDB
0769 1474 13       INX D         ;POINT TO TREE PARAMETERS
0770 1475 13       INX D         ; IN PPIDB
0771 1476 21B842   LXI H,HEAD   ;POINT H,L TO CTIME
0772 1479 060E     MVI B,14     ;UPDATE
0773 147B CDCF10    CALL MOVEB     ; PPIDB
0774 147E 219214   LXI H,NXTRE  ;SIMULATE A SUBROUTINE
0775 1481 E5       PUSH H       ; RETURN ADDRESS NXTRE
0776 1482 2AB442   LHL TCBP    ;POINT H,L TO
0777 1485 110700   LXI D,7      ; BYTE 7
0778 1488 19       DAD D        ; IN TCB
0779 1489 5E       MOV E,M     ;POINT
0780 148A 23       INX H       ; D,E TO
0781 148B 56       MOV D,M     ; EVENT
0782 148C EB       XCHG         ;H,L=ENTRY, D,E POINTS TO
0783                ; ; BYTE 8 OF TCB
0784 148D E9       PCHL        ;XQ EVENT WITH SIMULATED CALL
0785                ;
0786                ; ;EXAMINE NEXT EVENT TREE
0787                ;
0788 148E 21AD42   TREE0: LXI H,NTRE0 ;COUNT DOWN
0789 1491 35       DCR M          ; $ EMPTY TREES
0790 1492 F1       NXTRE: POP PSW     ;A= $TREES LEFT TO BE EXAMINED
0791 1493 E1       POP H          ;POINT TO PPIDB
0792 1494 011000   LXI B,10H    ;POINT TO NEXT CTIME
0793 1497 09       DAD B          ; IN PPIDB
0794 1498 3D       DCR A          ;DO FOR ALL
0795 1499 F24514   JP NWTRE     ; TREES
0796                ;
0797                ; ;IF FSQ0 NOT= 0 & TREE EMPTY FOR SOME TIME THEN COLD START
0798                ; ; (A WAIT LOOP COUNT NNTR0 IS USED RATHER THAN THE RTC
0799                ; ; SO THAT THE LARGER THE SYSTEM, THE LONGER THE WAIT)
0800                ;
0801 149C 2AAC42    LHL FSQ0     ;L=FSQ0, H=NTRE0
0802 149F 2D       DCR L          ;IF FSQ0 SET, I.E.,
0803 14A0 FA0D14   JM RXEC1     ; FREE STORAGE OUT
0804 14A3 24       INR H          ; IF ALL TREES
0805 14A4 C2B914   JNZ ENNTR    ; EMPTY
0806 14A7 2AAE42   LHL NNTR0    ;WAIT UNTIL
0807 14AA 2B       DCX H          ; ALL TREES
0808 14AB 22AE42   SHLD NNTR0   ; HAVE BEEN
0809 14AE 7D       MOV A,L      ; EMPTY FOR
0810 14AF B4       ORA H          ; SOME
0811 14B0 C20D14   JNZ RXEC1    ; TIME
0812                ;
0813 14B3 CD6609    CALL COLDS   ;COLD START & RESTART
0814 14B6 C30014   JMP REXEC    ; RUNNING MODE
0815                ;
0816 14B9 210020   ENNTR: LXI H,2000H ;IF ANY TREE NONEMPTY THEN
0817 14BC 22AE42   SHLD NNTR0   ; INITIALIZE WAIT COUNT
0818 14BF C30D14   JMP RXEC1    ;DO EXEC FOREVER
0819                ; ; INTERCOMPUTE(I/C) RUNNING MODE EXECUTIVE
0820                ;
0821                ; ; INPUT: INTERRUPTS DISABLED, MODE DATA, COLD-START DATA
0822                ; ; ALL REGISTERS AND MEMORY AFFECTED
0823                ;
0824                ;
0825 14C2 CDAC1A   ICXC: CALL RINI ;INI FOR I/C RUNNING MODE
0826                ;
0827                ; ;CLEAR KB'S AND ENABLE INTERRUPTS
0828 14C5 FB       EI              ;ENABLE INTERRUPTS
0829 14C6 CDB80C   CALL CLAKB   ;CLEAR ALL KB'S
0830 14C9 C3CF14   JMP ICX1     ;ENTER MAIN LOOP
0831                ;
0832                ; ;PROCESS EVERYTHING IN PSCBF
0833                ;
0834 14CC CD2F1B   IPHC: CALL PHPC ;PROCESS ALL PHOTOS

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0835 14CF CD6510 ICEX1: CALL GETPH ;GET B=DEVIC.
0836 14D2 C2CC14          JNZ  IPHC  ; C=INDEX
0837
0838 ;FLASH WATCH-DOG LIGHT
0839
0840 14D5 CD760F          CALL FLASH ;FLASH PSC-ON LIGHT
0841
0842 ;IF FREE STORAGE RUNNOUT.LOCKOUT KB INPUTS BUT PROCESS DAS MESSAGES
0843 ;L AND I (ACCEPT BUT DON'T PROCESS D). THIS WILL ALLOW SYSTEM TO CLEAR
0844 ;PACKAGES CURRENTLY IN SYSTEM WITHOUT ACCEPTING NEW PACKAGES WHILE FS
0845 ;IS OUT. AFTER A 'SUFFICIENT' TIME, THE PSCWILL COLDSTART(WARMSTART)
0846 ;ITSELF WHEN ALL THE MESSAGES IN FS HAVE BEEN SENT TO DAS.
0847
0848 ;IF FS IS OUT THEN LOCK OUT INDUCTS
0849
0850 14D8 3AAC42          LDA  FSQ0 ;TEST FOR FS RUNOUT
0851 14DB 3D              DCR  A
0852 14DC FAF814          JM   IKBPC ;JMP IF NO FS RUNOUT
0853 14DF CAFE14          JZ   FLBCK ;JMP IF FS RUNOUT IN PROGRESS(IGNORE KB)
0854 14E2 32AC42          STA  FSQ0 ;FS RUNOUT 1ST PASS LOOP
0855 14E5 3A033D          LDA  NINDU ;C= INDUCT $ = NINDU INITIALLY
0856 14E8 4F              MOV  C,A
0857 14E9 CD1422          IMCID: CALL MCLPC ;MASTER CLEAR INDUCT
0858 14EC 3E0F            MVI  A,15 ;ERROR
0859 14EE CD130F          CALL ERRPC ; 15
0860 14F1 0D              DCR  C ;DO UNTIL ALL
0861 14F2 C2E914          JNZ  IMCID ; CLEARED
0862 14F5 C3FE14          JMP  FLBCK ;IGNORE KB INPUTS
0863
0864 ; PROCESS A CHAR IN KBIBF
0865
0866 14F8 CD0E10          IKBPC: CALL GETKB ;A=INPUT CHAR,C=KB$,
0867 ; KBDBP=HL POINTS TO KBDB BYTE 4
0868 14FB C4501B          CNZ  KBPC ;PROCESS ANY KB INPUT CHAR(WILL ONLY
0869 ;PROCESS A 'CLEAR' FOR DAS SYSTEM)
0870
0871 ;CHECK TO SEE IF A 16 BYTE MESSAGE FROM DAS HAS ARRIVED(IS MESSAGE INPUT
0872 ;BUFFER FULL?)
0873
0874 14FE 3A5243          FLBCK: LDA  FULBF
0875 1501 D601            SUI  01H
0876 1503 CA4A15          JZ   CDSCK ;JMP IF MESSAGE RECEIVED FROM DAS
0877 ;CHECK FOR MESSAGE TIMEOUT(TIMEOUT IF 1ST BYTE OF MESSAGE WAS RECEIVED
0878 ;MORE THAN TDLTM AGO)
0879 1506 3A5B43          LDA  TMOUT ;A=TMOUT
0880 1509 47              MOV  B,A ;B=TMOUT
0881 150A 3EFF            MVI  A,0FFH ;A=FF
0882 150C B8              CMP  B ;COMPARE TMOUT LSB TO FF
0883 150D CA1315          JZ   TMCAL ;CHECK FURTHER
0884 1510 C31D15          JMP  CALTM ;1ST BYTE OF MESSAGE HAS BEEN SENT
0885 1513 3E7F            TMCAL: MVI  A,7FH ;A=7F
0886 1515 47              MOV  B,A ;B=7F
0887 1516 3A5C43          LDA  TMOUT+1 ;A=TMOUT+1
0888 1519 B8              CMP  B ;COMPARE TMOUT MSB TO 7F
0889 151A CA6A15          JZ   IEVNT ;BYPASS TIMEOUT CALC IF 1ST BYTE OF
0890 ;MESSAGE HASN'T BEEN SENT YET(AS TMOUT=7FFF YET)
0891 151D 2A4542          CALTM: LHLD PPIDB ;H,L=LOC OF RTC CTIME LSB
0892 1520 5E              MOV  E,M ;E=CTIME LSB
0893 1521 23              INX  H ;H,L=CTIME MSG
0894 1522 56              MOV  D,M ;D,E=CTIME
0895 1523 2A5B43          LHLD TMOUT ;H,L=TMOUT
0896 1526 7B              MOV  A,E ;A=CTIME LSB
0897 1527 95              SUB  L ;SUB TMOUT LSB(ALSO SETS BORROW FLAG)
0898 1528 4F              MOV  C,A ;C=LSB SUBTRACTION
0899 1529 7A              MOV  A,D ;D=CTIME MSG
0900 152A 9C              SBB  H ;SUB(TMOUT MSG+BORROW FLAG SETTING)
0901 152B 47              MOV  B,A ;BIC NOW CONTAINS CTIME-TMOUT
0902 152C 2A2C42          LHLD TDLTM ;H,L=TDLTM VALUE
0903 152F 79              MOV  A,C ;A=LSB OF CTIME-TMOUT
0904 1530 95              SUB  L ;A=LSB SUBTRACTION (CTIME-TMOUT(-TDLTM
0905 1531 78              MOV  A,B ;A=MSB
0906 1532 9C              SBB  H ;A=MSB SUBTRACTION (WITH BORROW FLAG)
0907 1533 F23915          JP   MTMOT ;JMP IF (CTIME-TMOUT)=> TDLTM
0908 1536 C36A15          JMP  IEVNT ;JMP IF NO TIMEOUT
0909 ;MESSAGE TIMEOUT. RESET MESSAGE INPUT POINTER DASDB, INCREMENT COUNTER
0910 ;CTOUT FOR DEBUGGING PURPOSES
0911 1539 212843          MTMOT: LXI  H,DASDB ;DASIP= DASDB LOC
0912 153C 224B43          SHLD DASIP
0913 153F 3A4A43          LDA  CTOUT ;A=CTOUT
0914 1542 C601            ADI  1 ;CTOUT=CTOUT+1
0915 1544 324A43          STA  CTOUT
0916 1547 C36A15          JMP  IEVNT ;JMP TO IEVNT
0917
0918 ; CHECK MESSAGE FOR GOOD SEQ BYTE AND CHECKSUM
0919
0920 154A 3E00            CDSCK: MVI  A,0 ;A=0
0921 154C 325243          STA  FULBF ;RESET FULBF=0
0922 154F CDF115          CALL DSCHK ;CHECK SEQ BYTE AND CHECKSUM

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0923 1552 3A5143 LDA FMBAD ;MESSAGE GOOD ?
0924 1555 D600 SUI 00H
0925 1557 CA6715 JZ CDDEC ;JMP IF MESSAGE IS GOOD
0926
;
0927 ; MESSAGE SEQ BYTE OR CHECKSUM NO GOOD. SEND DAS LAST PSC MESSAGE(WITH
0928 ; UNTOGGLED SEQ BYTE).
0929 155A 3E00 MVI A,00 ;A= 0
0930 155C 325143 STA FMBAD ;RESET FMBAD
0931 155F 3EFF MVI A,0FFH ;A=FFH
0932 1561 CD7719 CALL DSOUT ;RESEND TO DAS THE LAST PSC MESSAGE
0933 1564 C36A15 JMP IEVNT ;JMP TO IEVNT
0934
;
0935 ; MESSAGE SEQ BYTE AND CHECKSUM IS GOOD. PROCESS THE MESSAGE AND SEND
0936 ; REPLY
0937
;
0938 1567 CD8116 CDDEC: CALL DSDEC
0939
;
0940 ; PROCESS EVENTS DUE OR OVERDUE IN ALL AVL TREES
0941 ; TREE # 0 IS FOR RTC, TREES #1 TO NPPI ARE FOR PPI'S
0942 ; (BEWARE OF CRITICAL SECTIONS);
0943 ; RTC IS USED FOR LOW PRIORITY DELAYED TASKS
0944
;
0945 156A 3A023D IEVNT: LDA NPPI ;A=$ PPI'S = $TREES-1
0946 156D 32AD42 STA NTRE0 ;COUNT DOWN $ EMPTY TREES
0947 1570 2A4542 LHL PPIDB ;POINT TO CTIME IN PPIDB
0948 1573 22D842 INWTR: SHLD PPIDP ;SAVE PPIDB POINTER
0949 1576 E5 PUSH H ;SAVE AGAIN(GTREP MAY CHANGE PPIDP)
0950 1577 F5 PUSH PSW ;SAVE TREE COUNT
0951 1578 11B642 LXI D,CTIME ;D,E POINTS TO CTIME
0952 157B 0610 MVI B,10H ;GET TREE
0953 157D F3 DI ; PARAMETERS
0954 157E CDCF10 CALL MOVEB ; (CRITICAL SECTION WITH
0955 1581 FB EI ; RTCIN AND PPIIN)
0956 1582 2ABA42 LHL RLINK ;IF THE TREE
0957 1585 7D MOV A,L ; IS NOT
0958 1586 B4 ORA H ; EMPTY
0959 1587 CAB015 JZ ITREE
0960 158A 2ABD42 LHL NXTIM ;H.L=NXTIM
0961 158D 3AB642 LDA CTIME ;IF CTIME-NXTIME>=0.
0962 1590 95 SUB L ; I.E.,
0963 1591 3AB742 LDA CTIME+1 ; IF EVENT
0964 1594 9C SBB H ;IS OVERDUE
0965 1595 FAC015 JH INXTR ;OR DUE
0966
;
0967 ; REMOVE EVENT FROM TREE AND EXECUTE IT
0968
;
0969 1598 CD4E27 CALL DEQUE ;H:L POINTS TO EVENT TCR
0970 159B 22B442 SHLD TCBP ;TCBP POINTS TO TCB
0971 159E 2AD842 LHL PPIDP ;POINT DIE TO
0972 15A1 EB XCHG ; TREE HEADER IN PPIDB
0973 15A2 13 INX D ;POINT TO TREE PARAMETERS
0974 15A3 13 INX D ; IN PPIDB
0975 15A4 21B842 LXI H,HEAD ;POINT H.L TO CTIME
0976 15A7 060E MVI B,14 ;UPDATE
0977 15A9 CDCF10 CALL MOVEB ; PPIDB
0978 15AC 21C015 LXI H,INXTR ;SIMULATE A SUBROUTINE
0979 15AF E5 PUSH H ; RETURN ADDRESS INXTR
0980 15B0 2AB442 LHL TCBP ;POINT H.L TO
0981 15B3 110700 LXI D,7 ; BYTE 7
0982 15B6 19 DAD D ; IN TCB
0983 15B7 5E MOV E,M ;POINT
0984 15B8 23 INX H ; DIE TO
0985 15B9 56 MOV D,M ;EVENT
0986 15BA EB XCHG ;H.L=ENTRY,D.E POINTS TO
0987 ; BYTE 8 OF TCB
0988 15BB E9 PCHL ;XQ EVENT WITH SIMULATED CALL
0989
;
0990 ; EXAMINE NEXT EVENT TREE
0991
;
0992 15BC 21AD42 ITREE: LXI H,NTRE0 ;COUNT DOWN
0993 15BF 35 DCR M ; $ EMPTY TREES
0994 15C0 F1 INXTR: POP PSW ;A= $TREES LEFT TO BE EXAMINED
0995 15C1 E1 POP H ;POIN TO PPIDB
0996 15C2 011000 LXI B,10H ;POINT TO NEXT CTIME
0997 15C5 09 DAD B ; IN PPIDB
0998 15C6 3D DCR A ;DO FOR ALL
0999 15C7 F27315 JP INWTR ; TREES
1000
;
1001 ; TEST TO SEE IF FS IS OUT. IF IT IS, COLDSTART(WARMSTART) AFTER A DELAY
1002 ; AND MESSAGES IN FS ARE SENT TO DAS(A WAIT LOOP COUNT NNTR0 IS USED
1003 ; RATHER THAN THE RTC SO THAT THE LARGER THE SYSTEM, THE LONGER THE WAIT
1004
;
1005 15CA 2AAC42 LHL FSQ0 ;L=FSQ0.H=NTRE0
1006 15CD 2D DCR L ;IF FSQ0 SET,I.E.
1007 15CE FACF14 JM ICX1 ; FREE STORAGE OUT
1008 15D1 24 INR H ;IF ALL TREES
1009 15D2 C2F615 JNZ IENNT ; EMPTY

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1010 15D5 2AAE42      LHL D NNTR0      ;WAIT UNTIL
1011 15D8 2B         DCX H            ; ALL TREES
1012 15D9 22AE42     SHLD NNTR0       ; HAVE BEEN
1013 15DC 7D         MOV A,L         ; EMPTY FOR
1014 15DD B4         ORA H           ; SOME
1015 15DE C2CF14     JNZ ICEX1       ; TIME
1016                ; SEE IF MESSAGES IN FS ARE ALL SENT TO DAS
1017 15E1 2AAE42     LHL D NNTR0     ;SET NNTR0= 1 AGAIN
1018 15E4 23         INX H
1019 15E5 22AE42     SHLD NNTR0
1020 15E8 2A5443     LHL D LMP       ;H,L=LMP VALUE
1021 15EB 7D         MOV A,L         ;A=L
1022 15EC B4         ORA H
1023 15ED C2CF14     JNZ ICEX1       ;JMP IF LMP2 NOT EQUAL ZERO.INDICATING
1024                ;MESSAGES STILL IN FS
1025 15F0 CD6609     CALL COLDS      ;COLDSTART(WARMSTART) AND RESTART
1026 15F3 C3C214     JMP ICEXC       ;I/C RUN MODE
1027 15F6 210020     IENNT: LXI H,2000H ;IF ANY TREE NONEMPTY THEN
1028 15F9 22AE42     SHLD NNTR0     ; INITIALIZE WAIT COUNT
1029 15FC C3CF14     JMP ICEX1       ;DO EXEC FOREVER
1030                ;I/C SUBROUTINE DSCHK
1031                ;
1032                ;CHECKS DAS MESSAGES FOR CORRECT SEQ BYTE AND CHECKSUM. IF SEQ BYTE
1033                ;OR CHECKSUM IS INVALID, DON'T TOGGLE SEQ BYTE AND RETURN TO ICEXC. IF
1034                ;VALID, TOGGLE SEQ BYTE AND RETURN TO ICEXC.
1035                ;
1036                ;INPUTS: DAS MESSAGE RESIDES IN BUFFER DASDB(16 BYTES)
1037                ;
1038                ;PROCESS:SEE ABOVE
1039                ;
1040                ;OUTPUTS: IF MESSAGE IS GOOD, TOGGLE RS,LMR AND SET FMBAD=0 AND RETURN
1041                ;TO ICEXC. IF MESSAGE IS BAD, DON'T TOGGLE LMR,RS, SET FMBAD=1
1042                ;AND RETURN TO ICEXC. ALSO SET COUNTERS FOR DEBUGGING IF SEQ
1043                ;BYTE NOT A '5' OR 'J' OR IF CHECKSUM NO GOOD.
1044                ;TEST TO SEE IF IT'S MESSAGE L
1045                ;
1046 15FF 212843     DSCHK: LXI H,DASDB ;H,L= MESSAGE BYTE 0 LOC
1047 1602 23         INX H           ;H,L= MESSAGE BYTE 1 LOC
1048 1603 46         MOV B,M        ;B= MESSAGE IDENTIFIER
1049 1604 3E4C     MVI A,'L'     ;A= ASCII L
1050 1606 B8         CMP B          ;COMPARE
1051 1607 C21316     JNZ NOL       ;JMP IF NOT MESSAGE L
1052                ;ITS MESSAGE L. TEST IF 1ST MESSAGE L
1053 160A 3A5043     YESL: LDA FLINK ;A= FLAG INDICATING L COUNT
1054 160D D600     SUI 0         ;
1055 160F CA1C16     JZ CDSCH      ;JMP IF THIS IS THE 1ST L MESSAGE
1056                ;ITS 2ND OR MORE MESSAGE L. THEREFORE IGNORE IT AND DO NOTHING SO
1057                ;THAT DAS TIMES OUT. THE PSC LINK CAN ONLY BE OPENED AGAIN WHEN THE
1058                ;HOST SWITCH IS OPERATED FOLLOWED BY A MESSAGE L FROM DAS
1059 1612 C9         RET          ;RET TO ICEXC AND TAKE MULTIPLE L PATH IN DSDEC
1060                ;
1061                ;ITS NOT MESSAGE L. TEST TO SEE IF L HAS BEEN SENT YET
1062                ;
1063 1613 3A5043     NOL:  LDA FLINK ;A=L COUNT
1064 1616 D600     SUI 0         ;
1065 1618 C21C16     JNZ CDSCH     ;JMP IF L HAS BEEN SENT AT LEAST ONCE ALREADY
1066                ;MESSAGE L HAS NOT YET BEEN SENT. THEREFORE DO NOTHING
1067 161B C9         RET          ;RET TO ICEXC AND TAKE 'DO NOTHING' PATH IN
1068                ;DSDEC
1069                ;TEST IF SEQ BYTE= COMPLEMENT OF LMR
1070 161C 3A2A42     CDSCH: LDA LMR  ;A= LMR
1071 161F 2F         CMA          ;A=COMPLEMENT OF LMR
1072 1620 E67F     ANI 7FH      ;MASK OFF BIT 7 FOR SAFETY
1073 1622 47         MOV B,A      ;B= COMPL OF LMR
1074 1623 212843     LXI H,DASDB ;H,L=DASDB LOC
1075 1626 7E         MOV A,M      ;A= SEQ BYTE
1076 1627 90         SUB B        ;A=0 IF SEQ BYTE=COMPLEMENT OF LMR
1077 1628 CA5816     JZ YSEQ      ;JMP IF SEQ GOOD
1078                ;SEQUENCE NO GOOD.CHECK TO SEE IF ITS A '5' OR 'J'
1079 162B 212843     LXI H,DASDB ;H,L=POINT TO SEQ BYTE
1080 162E 46         MOV B,M      ;B=SEQ BYTE
1081 162F 3E35     MVI A,'5'    ;A= ASCII 5
1082 1631 90         SUB B        ;
1083 1632 CA3F16     JZ NMSG      ;JMP IF SEQ IS A '5'
1084 1635 3E4A     MVI A,'J'    ;A= ASCII J
1085 1637 90         SUB B        ;
1086 1638 CA3F16     JZ NMSG      ;JMP IF SEQ IS A 'J'
1087                ; BUMP COUNTER TO INDICATE SEQ NEITHER A '5' NOR A 'J'
1088 163B 214843     LXI H,CSEQ   ;CSEQ=CSEQ+1
1089 163E 34         INR M
1090                ;
1091                ;MESSAGE SEQ OR CHECKSUM NO GOOD. LEAVE LMR,RS UNCHANGED. SET FMBAD=1
1092                ;AND RETURN TO ICEXC
1093 163F 3E01     NMSG:  MVI A,01H ;SET FMBAD=1
1094 1641 325143     STA FMBAD
1095 1644 212843     LXI H,DASDB ;H,L=MESSAGE BYTE 0 LOC
1096 1647 23         INX H           ;H,L=BYTE 1 LOC
1097 1648 46         MOV B,M        ;B=MESSAGE IDENTIFIER

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1098 1649 3E4C      MVI  A,'L'      ;A=ASCII L
1099 164B B8       CMP   B          ;COMPARE
1100 164C C0       RNZ          ;RET IF NO MESSAGE L (TO ICEXC)WITH FMBAD=1
1101 164D 3E00     MVI  A,0        ;SET FMBAD=0 FOR USE ON RETURN TO ICEXC
1102 164F 325143   STA  FMBAD      ;SO THAT DSDEC WILL BE CALLED
1103 1652 3EFF     MVI  A,-1       ;SET FLINK=-1 SO THAT WHEN DSDEC IS CALLED
1104 1654 325043   STA  FLINK      ;ON RETURN TO ICEXC, THE MULTIPLE L PATH
1105 1657 C9       RET           ;WILL BE TAKEN SO DAS TIMES OUT AND FUTURE
1106                                     ;L WILL BE TREATED AS 1ST L.
1107
1108 ;SEQ OK. CHECK FOR CORRECT CHECKSUM
1109 1658 212843   YSEQ: LXI  H,DASDB ;H,L=DASDB LOC
1110 165B CD211A   CALL CHKM      ;GET PRINTABLE ASCII CHECKSUM IN REG A
1111 165E 212843   LXI  H,DASDB ;H,L=DASDB BYTE 0 LOC
1112 1661 110D00   LXI  D,13
1113 1664 19       DAD  D          ;H,L=CHECKSUM BYTE 13 LOC
1114 1665 46       MOV  B,M        ;B=CHECKSUM FROM DAS
1115 1666 B8       CMP   B          ;COMPARE CALL WITH DAS CHECKSUM
1116 1667 CA7116   JZ   YMSG      ;JMP IF CHECKSUM GOOD
1117
1118 ;CHECKSUM NO GOOD. DON'T TOGGLE RS,LMR, AND SET FMBAD=1. RET TO ICEXC
1119 ;AFTER BUMPING COUNTER CSUM FOR DEBUGGING PURPOSES
1120 166A 214943   LXI  H,CSUM    ;H,L= CSUM LOC
1121 166D 34       INR  M          ;CSUM = CSUM+1
1122 166E C3F16   JMP  NMSG      ;
1123
1124 ;SEQ BYTE AND CHECKSUM BOTH GOOD. SET FMBAD=0, TOGGLE LMR,RS AND RETURN
1125 1671 212843   YMSG: LXI  H,DASDB ;H,L= SEQ BYTE LOC
1126 1674 7E       MOV  A,M        ;A=SEQ BYTE
1127 1675 322A42   STA  LMR       ;LMR= SEQ BYTE JUST RECEIVED FROM DAS
1128 1678 322942   STA  RS        ;RS= SAME
1129 167B 3E00     MVI  A,0
1130 167D 325143   STA  FMBAD     ;FMBAD=0
1131 1680 C9       RET           ;RET TO ICEXC
1132
1133 ;
1134 ; THIS SUBROUTINE DECODES THE VALID(SEQ BYTE AND CHECKSUM ARE OK) DAS
1135 ; MESSAGES AND PROCESSES THEM IF THEY ARE MASKED FOR PROCESSING.
1136 ;
1137 ; INPUTS: DAS MESSAGE RESIDING IN BUFFER DASDB(16 BYTES)
1138 ;
1139 ; PROCESS: NO MESSAGES ARE PROCESSED UNTIL MESSAGE L IS RECEIVED. WHEN
1140 ; MESSAGE L IS RECEIVED FOR THE 1ST TIME, DSOUT IS CALLED TO
1141 ; SEND DAS THE 1ST MESSAGE WAITING IN FS TO BE SENT TO DAS IF
1142 ; THE MESSAGE IN MOBUF IS NUMBER 8 OR 9(LAST MESSAGE SENT TO
1143 ; DAS IS MESSAGE 8 OR 9). IF THE LAST MESSAGE SENT TO DAS WAS
1144 ; NOT 8 OR 9, THEN THE MESSAGE IN MOBUF WILL BE RESENT TO DAS
1145 ; WHEN MESSAGE L IS RECEIVED FOR THE 1ST TIME.MULTIPLE MESSAGE
1146 ; L'S WILL CAUSE DAS TO TIMEOUT(NO PSC REPLY WILL BE SENT) AND
1147 ; A FLAG FLINK INCREMENTED FOR DEBUGGING PURPOSES,IF MESSAGE
1148 ; I IS SENT,DSOUT WILL BE CALLED TO REPLY WITH 1ST WAITING PSC
1149 ; MESSAGE. IF MESSAGE D IS SENT,IT WILL BE PROCESSED BY
1150 ; CALLING DASMD AND THEN A REPLY SENT TO DAS BY CALLING DSOUT.
1151 ; UNINTERPRETABLE MESSAGES WILL CAUSE A MESSAGE 8 REPLY WITH
1152 ; UNTOGGLED SEQ BYTE SO THAT DAS TIMES OUT.
1153 ;
1154 ; OUTPUTS: DAS MESSAGE DECODED AND REPLY SENT BACK (BY CALL DSOUT)
1155 ; CHECK FOR MESSAGE L
1156 ;
1157 1681 212943   DSDEC: LXI  H,DASDB+1 ;H,L= LOC OF MESSAGE BYTE 1
1158 1684 3E4C     MVI  A,'L'      ;A=ASCII L
1159 1686 96       SUB  M          ;A=ASCII L - MESSAGE BYTE 1
1160 1687 C2D316   JNZ  LPAST     ;IF NOT L, JMP TO LPAST
1161 ; ITS MESSAGE L
1162 ; IS IT 1ST MESSAGE L?
1163 168A 3A5043   LDA  FLINK     ;A=FLINK
1164 168D A7       ANA  A        ;SET ZERO FLAG AS PER VALUE OF FLINK
1165 168E CA9616   JZ   LPROS    ;JMP TO LPROS IF FLINK=0
1166 ; MULTIPLE MESSAGE L.
1167 1691 215043   LMULT: LXI  H,FLINK ;H,L= FLINK LOC IN MEMORY
1168 1694 34       INR  M        ;FLINK=FLINK+1
1169 1695 C9       RET           ;RETURN TO ICEXC WITHOUT REPLY TO DAS SO
1170 ; ;
1171 ; ;
1172 ; ;
1173 ; 1ST TIME MESSAGE L IS SENT
1174 1696 3E01     LPROS: MVI  A,01H ;A=1
1175 1698 325043   STA  FLINK     ;FLINK=1
1176 169B 3A7542   LDA  ILCHN    ;A=I/C LIGHT CHANNEL
1177 169E 4F       MOV  C,A      ;C=SAME
1178 169F 0601     MVI  B,1      ;B=1 FOR TURN ON
1179 16A1 CD0F11   CALL OPUT     ;TURN ON I/C LIGHT
1180 16A4 3A2E42   LDA  VICLG    ;A= VICLG
1181 16A7 325343   STA  ICLGT    ;SET I/C LIGHT REFRESH TIMER
1182 16AA 213943   LXI  H,MOBUF+1 ;H,L=LOC OF BYTE 1 OF MOBUF
1183 16AD 3E20     MVI  A,20H    ;A=ASCII SPACE
1184 16AF 96       SUB  M        ;A= 0 - BYTE 1 OF MOBUF
1185 16B0 CACD16   JZ   CDSOT    ;JMP IF BYTE 1 = 0 (IF NO MESS IN MOBUF)

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1186 16B3 3E38      MVI A,'8'          ;A=ASCII 8
1187 16B5 96        SUB M              ;A=ASCII 8 - BYTE 1 OF MOBUF
1188 16B6 CACD16     JZ CDSOT          ;JMP IF BYTE 1 = 8 (LAST MESSAGE SENT=8)
1189 16B9 3E39      MVI A,'9'          ;DO SAME FOR MESSAGE 9
1190 16BB 96         SUB M              ;
1191 16BC CACD16     JZ CDSOT          ;JMP IF MESSAGE 9 IN MOBUF
1192
1193
1194 16BF 212843     ; RESEND MOBUF TO DAS WITH NEW SEQ BYTE
1195 16C2 7E         LXI H,DASDB       ;
1196 16C3 213843     MOV A,M           ;A= SEQ BYTE FROM DAS
1197 16C6 77         LXI H,MOBUF       ;
1198 16C7 3EFF       MOV M,A           ;MOBUF BYTE 0= SEQ BYTE FROM DAS
1199 16C9 CD7719     MVI A,0FFH        ;A=FF
1200 16CC C9         CALL DSOUT        ;SEND MOBUF TO DAS
1201                RET          ;RET TO ICEXC
1202                ; SEND DAS 1ST MESSAGE IN FS (DEFAULTS TO IDLE MESSAGE)
1203 16CD 3E00        CDSOT: MVI A,00H   ;A=00
1204 16CF CD7719     CALL DSOUT        ;
1205 16D2 C9         RET          ;RETURN TO ICEXC
1206
1207
1208                ; NOT MESSAGE L
1209 16D3 3A5043     LPAST: LDA FLINK   ;A=FLINK
1210 16D6 D601        SUI 01H          ;A=FLINK-1
1211 16D8 CADC16     JZ ITEST         ;JMP TO ITEST IF FLINK=1(L HAS BEEN SENT
1212 16DB C9         RET          ;DO NOTHING RETURN TO ICEXC AND WAIT FOR
1213                ; MESSAGE L OTHERWISE
1214
1215 16DC 3A2E42     ; REFRESH I/C LIGHT
1216 16DF 325343     ITEST: LDA VICLG  ;A= VICLG
1217                STA ICLGT  ;ICLGT=VICLG
1218 16E2 212943     ; TEST FOR MESSAGE I
1219 16E5 3E49       LXI H,DASDB+1    ;H.L= ADDR OF BYTE 1
1220 16E7 96         MVI A,'I'        ;A= ASCII I
1221 16E8 C2FA16     SUB M             ;A= ASCII I - BYTE 1
1222                JNZ DTEST  ;JMP TO DTEST IF NOT MESSAGE I
1223                ; ITS MESSAGE I IS I MASKED?
1224 16EB 21253D     LXI H,ERMKI+3    ;
1225 16EE 7E         MOV A,M          ;A= BYTE 4 OF ERMKI
1226 16EF E601       ANI 01H          ;LOOK AT MESSAGE I BIT
1227 16F1 CAF116     JZ DTEST         ;JMP IF NOT MASKED
1228 16F4 3E00       IPROS: MVI A,00H ;A=00H
1229 16F6 CD7719     CALL DSOUT        ;CALL DSOUT TO SEND DAS 1ST WAITING
1230                ; MESSAGE IN FS.DEFAULT TO I MESSAGE
1231                ; RET TO ICEXC
1232
1233                ; TEST FOR MESSAGE D
1234 16FA 212943     DTEST: LXI H,DASDB+1 ;H.L= LOC OF BYTE 1 OF MESSAGE
1235 16FD 3E44       MVI A,'D'        ;
1236 16FF 96         SUB M             ;A=ASCII D - BYTE 1
1237 1700 C26517     JNZ DFALT        ;JMP IF NOT D
1238                ; ITS MESSAGE D. IS D MASKED?
1239 1703 21243D     LXI H,ERMKI+2    ;
1240 1706 7E         MOV A,M          ;A= BYTE 3 OF ERMKI
1241 1707 E608       ANI 08H          ;LOOK AT MESSAGE D BIT
1242 1709 CA6517     JZ DFALT        ;JMP IF D NOT MASKED
1243                ; TEST TO SEE IF FS RUNNOT IS IN PROGRESS
1244 170C 3AAC42     LDA FSQ0         ;A=FSQ0 FLAG
1245 170F A7         ANA A           ;SET ZERO FLAG ACCORDINGLY
1246 1710 C2FA16     JNZ IPROS        ;JMP IF FS RUNNOT IS IN PROGRESS
1247                ; TEST TO SEE LAQ OVERFLOW HAS OCCURRED. IF SO, DON'T PROCESS MESSAGE
1248                ; D TILL ERROR 6 IS CLEARED OFF KB.
1249
1250                ; ACCESS ER FLAG IN KBDB TO SEE IF ERROR 6 (LAQ OVERFLOW) IS ON KB
1251 1713 212843     LXI H,DASDB     ;H.L= MESSAGE BYTE 0
1252 1716 110800     LXI D,8         ;
1253 1719 19         DAD D           ;H.L= INDUCT BYTE 1 LOC
1254 171A 7E         MOV A,M         ;A= INDUCT MSB
1255 171B 0E0A       MVI C,10        ;C= WEIGHT OF 10
1256 171D CD8F18     CALL ASCII      ;CONVERT TO BINARY
1257 1720 5F         MOV E,A         ;E= CURRENT SUM
1258 1721 23         INX H          ;H.L= INDUCT BYTE 2 LOC
1259 1722 7E         MOV A,M         ;A= INDUCT LSB
1260 1723 0E01       MVI C,1         ;C= WEIGHT OF 1
1261 1725 CD8F18     CALL ASCII      ;A= BINARY 1'S
1262 1728 83         ADD E           ;A= BINARY INDUCT * = KB*
1263 1729 CD170D     CALL CPKBP      ;POINT TO KB11 IN KBDB FOR KB IN REG A
1264 172C 110400     LXI D,4         ;
1265 172F 19         DAD D           ;H.L= ER LOC
1266 1730 7E         MOV A,M         ;A= ER VALUE
1267 1731 D601       SUI 1           ;SUB 1
1268 1733 C24117     JNZ THDAS       ;JMP IF NO ERROR ON KB
1269                ; ER= 1. THEREFORE AN ERROR IS ON THIS KB. ACCESS ERROR CODE(TEMPORARILY
1270                ; STORED IN SND LOC) TO SEE IF ITS ERROR $6.
1271 1736 23         INX H           ;
1272 1737 23         INX H           ;H.L= SND LOC
1273 1738 7E         MOV A,M         ;A= SND VALUE = ERROR CODE

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1274 1739 D606      SUI 6      ;SUB 6
1275 1738 C24117   JNZ THDAS ;JMP IF NOT ERROR 6(LAQ OVERFLOW)
1276                ;LAQ OVERFLOW HAS OCCURRED ON THIS KB. THEREFORE DON'T PROCESS ANYMORE
1277                ;MESSAGE D'S UNTIL KB IS CLEARED(AND LAQ) BY HITTING 'CLEAR' KEY
1278                ;ON KB. THIS IS NECESSARY TO INSURE PACKAGES DON'T GET OUT OF SEQUENCE.
1279 173E C3F416   JMP IPROS  ;JMP TO IPROS
1280
1281                ;TEST IF IN HALF DAS MODE. IF ARE, DON'T PROSS MESSAGE D.
1282
1283 1741 3A0640   THDAS: LDA HDSEN ;A= HALF-DAS FLAG
1284 1744 D601     SUI 1      ;SUB 1
1285 1746 CAF416   JZ IPROS  ;JMP IF HALF-DAS
1286
1287                ;PROCESS MESSAGE D
1288 1749 CD9617   CALL DASMD ;PROCESS MESSAGE D
1289 174C 3A4F43   LDA FLG9  ;FLG9=1?
1290 174F D601     SUI 01H
1291 1751 CA5A17   JZ REGA  ;JMP IF FLG9=1
1292 1754 3E00     MVI A,0   ;REG A = 0
1293 1756 CD7719   CALL DSOUT ;CALL DSOUT WITH REGA=0
1294 1759 C9      RET      ;RET TO ICEXC
1295 175A 3E00     REGA: MVI A,0 ;A=0
1296 175C 324F43   STA FLG9  ;FLG9=0
1297 175F 3EFF     MVI A,0FFH ;A= FFH
1298 1761 CD7719   CALL DSOUT ;CALL DSOUT WITH REG A=FF
1299 1764 C9      RET      ;RET TO ICEXC
1300                ; DEFAULT TO MESSAGE 8. DON'T TOGGLE RS,LMR AS WANT LINK TO GO DOWN.
1301                ; SINCE WE HAVE PREVIOUSLY TOGGLED THEM AT YMESG,WE MUST RETOGGLE
1302                ; THEM HERE TO GET BACK TO ORIGINAL DAS STATE.
1303 1765 3A2843   DFALT: LDA DASDB ;A= SEQ BYTE FROM DAS
1304 1768 2F      CHA      ;A= COMPLEMENT OF SEQ BYTE
1305 1769 E67F     ANI 7FH   ;MASK BIT 7
1306 176B 322A42   STA LMR   ;SET LMR,RS & SEQ BYTE 0 IN MOBUF = TO
1307 176E 322942   STA RS    ;COMPLEMENT OF SEQ BYTE SENT BY DAS AS
1308 1771 323843   STA MOBUF ;DON'T WANT TO TOGGLE SEQ BYTE
1309 1774 212843   LXI H,DASDB
1310 1777 224B43   SHLD DASIP
1311
1312                ; LOAD REST OF MESSAGE 8 INTO MOBUF
1313 177A 212943   LXI H,DASDB+1 ;H,L= BYTE 1 LOC OF INPUT MESSAGE
1314 177D 113943   LXI D,MOBUF+1 ;D,E= BYTE 1 LOC OF OUTPUT MESSAGE
1315 1780 010F00   LXI B,15    ;B,C= 15 BYTES TO MOVE
1316 1783 CDC410   CALL MOVE   ;MOVE 15 BYTES FROM DASDB TO MOBUF
1317 1786 3E38     MVI A,'8'   ;A=ASCII 8
1318 1788 323943   STA MOBUF+1 ;STORE IN MOBUF
1319 178B 3E08     MVI A,8     ;SEND ERROR 8 TO MASTER
1320 178D CD080F   CALL ERPC   ;
1321 1790 3EFF     MVI A,0FFH ;A=FFFF"
1322 1792 CD7719   CALL DSOUT  ;SEND MOBUF TO DAS
1323 1795 C9      RET      ;RET TO ICEXC
1324                ; SUBROUTINE DASMD
1325
1326                ; INPUTS: MESSAGE D RESIDING IN DASDB
1327
1328                ; PROCESS: PUT PACKAGE IDENT AND DVT INTO LAQ. OTHER FUNCTIONS SIMILAR
1329                ; TO SNDPC
1330
1331                ; OUTPUTS: SIMILAR TO SNDPC SUBROUTINE
1332
1333                ; CONVERT ASCII DVT & AND INDUCT & IN MESSAGE D TO BINARY AND PUT ON
1334                ; STACK FOR FUTURE USE
1335
1336 1796 212843   DASMD: LXI H,DASDB ;H,L= LOC MESSAGE D BYTE 0
1337 1799 110A00   LXI D,10
1338 179C 19      DAD D      ;H,L= LOC OF DVT BYTE 1
1339 179D 7E      MOV A,M    ;A= DVT BYTE 1 (100'S DIGIT)
1340 179E 0E64     MVI C,100  ;C= WEIGHT OF 100
1341 17A0 CD8F18   CALL ASCII ;CONVERT TO BINARY
1342 17A3 5F      MOV E,A    ;E= CURRENT SUM
1343 17A4 23      INX H     ;H,L= DVT BYTE 2 (10'S DIGIT) LOC
1344 17A5 7E      MOV A,M    ;A= 10'S DIGIT
1345 17A6 0E0A     MVI C,10   ;C= WEIGHT OF 10
1346 17A8 CD8F18   CALL ASCII ;A= BINARY EQUIVALENT
1347 17AB 83      ADD E      ;ADD 100'S
1348 17AC 5F      MOV E,A    ;E= 100'S + 10'S
1349 17AD 23      INX H     ;H,L= DVT BYTE 3 (1'S DIGIT) LOC
1350 17AE 7E      MOV A,M    ;A= 1'S DIGIT
1351 17AF 0E01     MVI C,1    ;C= WEIGHT OF 1
1352 17B1 CD8F18   CALL ASCII ;A= BINARY 1'S
1353 17B4 83      ADD E      ;A= BINARY DIVERT &
1354 17B5 47      MOV B,A    ;B= SAME
1355 17B6 C5      PUSH B    ;STACK = SAME
1356 17B7 212843   LXI H,DASDB ;H,L= LOC OF MESSAGE D
1357 17BA 110800   LXI D,8
1358 17BD 19      DAD D     ;H,L= INDUCT BYTE 1 LOC
1359 17BE 7E      MOV A,M    ;A= INDUCT BYTE 1 LOC (10'S DIGIT)
1360 17BF 0E0A     MVI C,10   ;C= WEIGHT OF 10
1361 17C1 CD8F18   CALL ASCII ;CONVERT TO BINARY

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1362 17C4 5F      MOV  E,A      ;E=CURRENT SUM
1363 17C5 23      INX  H        ;H.L= INDUCT BYTE 2 LOC (1'S DIGIT)
1364 17C6 7E      MOV  A,M      ;A= 1'S DIGIT
1365 17C7 0E01     MVI  C,1      ;C= WEIGHT OF 1
1366 17C9 CD8F18   CALL ASCII    ;A= BINARY 1'S
1367 17CC 83      ADD  E        ;A= BINARY INDUCT #
1368 17CD C1      POP  B        ;B= JDV IN BINARY
1369 17CE 4F      MOV  C,A      ;C= JINDU IN BINARY
1370 17CF C5      PUSH B       ;STACK= JDV.JINDU IN BINARY
1371
1372 ; CHECK FOR LEGAL INDUCT # (>0 AND < NINDU)
1373 17D0 79      MOV  A,C      ;A= INDUCT #
1374 17D1 D600     SUI  0
1375 17D3 CA1E18   JZ   SNDM9    ;JMP IF IND = 0
1376 17D6 3A033D   LDA  NINDU    ;A=NINDU
1377 17D9 91      SUB  C        ;A=NINDU-JINDU
1378 17DA FA1E18   JM  SNDM9    ;JMP IF JINDU> NINDU (# OF INDUCTS)
1379
1380 ; INDUCT IS LEGAL. CHECK FOR LEGAL PACK CODE (EACH BYTE A NUMBER)
1381 17DD 0630     MVI  B,30H    ;B= ASCII 0
1382 17DF 0E39     MVI  C,39H    ;C= ASCII 9
1383 17E1 212A43   LXI  H,DASDB+2 ;H.L= PACKAGE CODE BYTE 1
1384 17E4 1606     MVI  D,6      ;D=LOOP COUNTER
1385 17E6 79      MOV  A,C      ;A= UPPER LIMIT ASCII 9
1386 17E7 96      SUB  M        ;A= ASCII 9 - BYTE
1387 17E8 FA1E18   JM  SNDM9    ;JMP IF BYTE > 9
1388 17EB 7E      MOV  A,M      ;A= PACK CODE BYTE
1389 17EC 90      SUB  B        ;A-BYTE - ASCII 0
1390 17ED FA1E18   JM  SNDM9    ;JMP IF BYTE < 0
1391 17F0 23      INX  H        ;GO TO NEXT BYTE
1392 17F1 15      DCR  D        ;DEC COUNTER
1393 17F2 C2E617   JNZ  LEGPC    ;DO FOR ALL 6 BYTES
1394 ; LEGAL PACKAGE CODE. CHECK FOR LEGAL DV # (>0 AND <NDVT)
1395 17F5 C1      POP  B        ;B,C= JDV,JINDU
1396 17F6 C5      PUSH B       ;STACK=SAME
1397 17F7 78      MOV  A,B      ;A= JDV
1398 17F8 D600     SUI  00H
1399 17FA CA0A18   JZ   DVEO    ;JMP IF DV#=0
1400 17FD 3A053D   LDA  NDVT     ;A= NDVT
1401 1800 57      MOV  D,A      ;0= SAME
1402 1801 78      MOV  A,B      ;A= DV#
1403 1802 3D      DCR  A        ;BIAS FOR 'JP' INSTRUCTION
1404 1803 92      SUB  D        ;A= DV#-NDVT
1405 1804 F20A18   JP  DVEO     ;JMP IF DV#>NDVT
1406 1807 C34A18   JMP  INHIB   ;LEGAL DV#.JMP TO INHIB
1407 ; DV # IS ILLEGAL. CHECK IF ZERO OR NON-ZERO
1408
1409 180A 78      DVEO: MOV  A,B      ;A= BINARY DV #
1410 180B D600     SUI  00H
1411 180D C21E18   JNZ  SNDM9    ;JMP IF DV # > 0
1412 ; DV# IS 0. DEFAULT JDV TO SERDV. SET FDV FOR USE IN DPLAQ
1413 1810 3A3A3D   LDA  SERDV
1414 1813 329642   STA  JDV      ;JDV=SERDV
1415 1816 3E01     MVI  A,01H
1416 1818 324E43   STA  FDV      ;FDV=1
1417 181B C34A18   JMP  INHIB
1418 ; SEND MESSAGE 9 TO DAS WITH SEQUENCE BYTE UNTOGGLED SO DAS TIMESOUT
1419 181E 3A2843   SNDM9: LDA  DASDB ;A= SEQ BYTE RECEIVED FROM DAS
1420 1821 2F      CMA
1421 1822 E67F     ANI  7FH     ;A= COMPLEMENT OF SEQ BYTE FROM DAS
1422 1824 322A42   STA  LMR     ;MASK BIT 7
1423 1827 322942   STA  RS      ;SET LMF,RS AND SEQ BYTE IN MOBUF
1424 182A 323843   STA  MOBUF   ;EQUAL TO ORIGINAL STATE (UNTOGGLED)
1425 ; ECHO INPUT MESSAGE BY LOADING INTO MOBUF AND SENDING TO DAS BY
1426 ; CALL DSOUT WITH REG A = FFH
1427 182D 212943   LXI  H,DASDB+1 ;H.L= BYTE 1 OF DASDB
1428 1830 113943   LXI  D,MOBUF+1 ;D,E= BYTE 1 OF MOBUF
1429 1833 010F00   LXI  B,15     ;B= 15 BYTES TO MOVE
1430 1836 CDC410   CALL MOVE    ;MOVE 15 BYTES FROM DASDB TO MOBUF
1431 1839 3E39     MVI  A,'9'    ;A= ASCII 9
1432 183B 323943   STA  MOBUF+1 ;MOBUF BYTE 1 = MESSAGE 9
1433 183E 3E09     MVI  A,09H    ;A=09H
1434 1840 CD080F   CALL ERPC    ;PUT ERROR 9 ON KB
1435 1843 3E01     MVI  A,01H    ;FLG9=1 SO WHEN DSOUT IS CALLED ON RET
1436 1845 324F43   STA  FLG9    ;TODSDEC MOBUF WILL BE SENT TO DAS
1437 1848 C1      POP  B
1438 1849 C9      RET         ;ADJUST STACK
1439 ; RET TO DSDEC
1440 ; CHECK FOR LANE FULL KB INHIBIT DESIRED
1441 184A C1      INHIB: POP  B      ;B,C=JDV, JINDU
1442 184B 3A3C3D   LDA  LFO     ;CHECK TO SEE IF LANE-FULL KB INHIBIT
1443 184E E601     ANI  01H     ; IS DESIRED
1444 1850 CA7018   JZ   CDPLQ   ;JMP IF NOT DESIRED
1445 ; GET LANE FULL MASK
1446 1853 78      MOV  A,B      ;A=JDV
1447 1854 C5      PUSH B       ;STACK= JDV, JINDU

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1449 1855 CD250D CALL CPLFM ;CP LF MASK IN A. H.L=MASK POINTER
1450 1858 C1 POP B ;B.C= JDV, JINDU
1451 1859 A6 ANA M ;IF LANE FULL JMP TO CALL 7
1452 185A C26018 JNZ CALL7 ;IF NOT FULL JMP TO CDPLQ
1453 185D C37018 JMP CDPLQ
1454 1860 3E07 CALL7: MVI A,7 ;LANE FULL. PUT ERROR 7 ON KB. IF AUTO
1455 1862 CD130F CALL ERRPC ;REINDUCT AFTER LANE CLEARS IS DESIRED.
1456 1865 3A3C3D LDA LFO ;SET AUTO REINDUCT FLAG ER=10
1457 1868 E610 ANI 10H
1458 186A C8 RZ ;RET TO DSDEC IF NO AUTO REINDUCT
1459 186B 2ADA42 LHLD KBDBP
1460 186E 77 MOV M,A ;SET ER=10
1461 186F C9 RET ;RET TO DSDEC
1462
1463 ; LOAD LAQ WITH JDV AND IDENT(IDENT= LOC IN TRANSACTION TABLE OF
1464 ; PACKAGE CODE) AND ATTEMPT TO SEND PACKAGE
1465 1870 C5 CDPLQ: PUSH B ;STACK= JDV, JINDU
1466 1871 CD9D18 CALL DPLAQ ;LOAD LAW WITH JDV,IDENT
1467 1874 C1 POP B ;B.C= JDV, JINDU
1468 1875 CA0B22 JZ SLAQF ;IF LAQ FULL, SEND ERROR 6 TO KB
1469 1878 2AEB42 LHLD PPDBP ;POINT H.L TO BYTE 1 IN PPDB
1470 187B 7E MOV A,M ;A=PWING
1471 187C 3D DCR A ;IF PACKAGE AT PP
1472 187D C28618 JNZ CRBSO ; WAITING TO BE CODED
1473
1474 ; SEND OUT WAITING PACKAGE IF NO MERGE CONFLICT
1475 1880 C5 PUSH B ;STACK= JDV, JINDU
1476 1881 23 INX H ;H.L= BYTE 2 OF PPDB
1477 1882 CD7022 CALL PSEQR ;QUEUE PACKAGE UP FOR RELEASE
1478 1885 C1 POP B ;B.C= JDV, JINDU
1479
1480 ; CONSIDER REPEAT-BY-SEND OPTION
1481 1886 3A323D CRBSO: LDA SENDO ;IF REPEAT NOT DESIRED
1482 1889 E620 ANI 20H ; THEN CLEAR
1483 188B CCC90C CZ CLRKB ; KB AND KBDB
1484 188E C9 RET ;RET TO DSDEC
1485
1486 ; SUBROUTINE ASCII
1487
1488 ; INPUTS: REG A = ASCII CHARACTER, REG C = VALUE ASSIGNED TO EACH
1489 ; DIGIT(1,10 OR 100)
1490
1491 ; PROCESS: CONVERTS ASCII CHAR IN REG A TO BINARY VALUE WEIGHTED
1492 ; BY REG C
1493
1494 ; OUTPUTS: WEIGHTED BINARY VALUE IN REG A
1495
1496 ASCII: ANI 0FH ;LOOK AT 4 LSB
1497 1891 CA9C18 JZ RETRN ;JMP IF ZERO
1498 1894 47 MOV B,A ;B= LOOP COUNT
1499 1895 3E00 MVI A,0 ;INITIALIZE A
1500 1897 81 ALOOP: ADD C ;A= WEIGHT
1501 1898 05 DCR B ;DEC COUNTER
1502 1899 C29718 JNZ ALOOP ;DO FOR ALL COUNTS
1503 189C C9 RETRN: RET ;RET WITH WEIGHTED BINARY $ IN REG A
1504
1505 ; SUBROUTINE DPLAQ
1506
1507 ; INPUTS: MESSAGE D RESIDING IN DASDB. STACK= JDV, JINDU IN BINARY
1508
1509 ; PROCESS: LOAD THE LAQ FOR THE INDUCT IN MESSAGE D WITH THE DIVERT
1510 ; NUMBER IN MESSAGE D(DVT$ IN BINARY). LOAD THE IDENT IN THE
1511 ; LAQ WITH THE ADDRESS IN THE TRANSACTION TABLE WHERE THE
1512 ; 6 BYTE ASCII PACKAGE CODE HAS BEEN STORED
1513
1514 ; OUTPUTS: STACK=JDV, JINDU. LAQ LOADED WITH JDV,IDENT
1515
1516 ; TEST TO SEE IF LAQ IS FULL
1517 189D 79 DPLAQ: MOV A,C ;A=INDUCT $
1518 189E CD9927 CALL CPPPP ;POINT H.L TO BYTE 1
1519 18A1 23 INX H ; IN PPDB FOR THIS INDUCT
1520 18A2 22EB42 SHLD PPDBP ;OUTPUT POINTER
1521 18A5 110500 LXI D,5
1522 18A8 19 DAD D
1523 18A9 3A343D LDA DQLIM
1524 18AC 96 SUB M
1525 18AD C8 RZ ;RET IF LAQ FULL
1526
1527 ; LAQ NOT FULL. TEST TO SEE IF JDV HAS ALREADY BEEN SET IN DASMD FOR
1528 ; THE CASE WHEN JDV=0.
1529 18AE 3A4E43 LDA FDV ;FDV=1?
1530 18B1 A7 ANA A
1531 18B2 CABE18 JZ CNVJD ;JMP IF FDV=0
1532 18B5 3E00 MVI A,0 ;SET FDV=0
1533 18B7 324E43 STA FDV
1534 18BA C5 PUSH B ;STACK= JDV, JINDU
1535 18BB C3C318 JMP PCODE ;JDV ALREADY LOADED IN MEMORY. JUMP
; LOAD JDV FROM MESSAGE D, AS BINARY INTO MEMORY LOC JDV
CNVJD: PUSH B ;STACK=JDV, JINDU

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1536 18BF 78      MOV  A,B          ;A=JDV
1537 18C0 329642 STA  JDV          ;MEMORY JDV = JDV IN BINARY
1538
1539 ;
1540 ; SEE IF PACKAGE CODE=000000 MEANING TO RECIRCULATE PACKAGE. IF SO.
1541 ; BYPASS PUTTING CODE INTO TRANSACTION TABLE AND SET IDENT=0000 TO
1542 ; IDENTIFY RECIRCULATING PACKAGE
1542 18C3 212A43 PCODE: LXI  H,DASDB+2 ;H.L= PACKAGE CODE BYTE 1
1543 18C6 0606     MVI  B,6          ;B= COUNTER
1544 18C8 7E      PCDCCK: MOV  A,M          ;A= PACK CODE BYTE
1545 18C9 D630     SUI  '0'         ;SUBTRACT ASCII 0
1546 18CB C2DE18  JNZ  SCODE       ;JMP IF EACH BYTE NOT A ASCII 0
1547 18CE 23     INX  H          ;LOOK AT NEXT BYTE
1548 18CF 05     DCR  B          ;DEC COUNTER
1549 18D0 C2C818 JNZ  PCDCCK      ;DO ALL 6 BYTES
1550
1551 ;
1552 ; PACKAGE CODE ='S 000000. THEREFORE SET IDENT=0000 SO WHEN PACKAGE
1553 ; LEAVES SYSTEM, IT WON'T BE NECESSARY TO CLEAR IT FROM TRANSACTION
1554 ; TABLE AS WE ARE ABOUT TO BYPASS PUTTING IT INTO TABLE.
1554 18D3 3E00     MVI  A,00H       ;A=0
1555 18D5 329742 STA  IDENT
1556 18D8 329842 STA  IDENT+1     ;IDENT = 0000H
1557 18DB C3E118 JMP  LDLAG       ;JMP LDLAG
1558
1559 ;
1560 ; LOAD 6 ASCII BYTES OF PACKAGE CODE INTO 1ST AVAILABLE 6 LOC IN
1561 ; TRANSACTION TABLE BY CALLING DSPAK. RET WITH IDENT= ADDRESS IN
1562 ; TRANSACTION TABLE WHERE CODE IS STORED.
1562 18DE CD0619 SCODE: CALL  DSPAK
1563
1564 ;
1564 ; LOAD MEMORY JDV AND IDENT INTO LAQ FOR THIS INDUCT
1565 18E1 C1      LDLAG: POP  B          ;B,C= JDV,JINDU
1566 18E2 79     MOV  A,C          ;A=INDUCT #
1567 18E3 CD9927 CALL  CPPPP      ;POINT H.L TO BYTE 1 IN PPDB
1568 18E6 23     INX  H
1569 18E7 22EB42 SHLD PPDBP      ;OUTPUT POINTER
1570 18EA 110500 LXI  D,5         ;POINT H.L TO LAQSZ IN PPDB
1571 18ED 19     DAD  D          ;H.L=LAQSZ=# OF PACKAGES IN LAQ
1572 18EE 7E     MOV  A,M          ;A=LAQSZ VALUE
1573 18EF 34     INR  M          ;INCREMENT LAQSZ
1574 18F0 23     INX  H          ;H.L=LAQ1 LOC
1575 18F1 86     ADD  M          ;A=RELATIVE NUMBER OF LOC USED UP IN LAQ
1576 18F2 E607   ANI  07H        ;MASK TO 3 BITS FOR WRAPAROUND
1577 18F4 5F     MOV  E,A
1578 18F5 1600   MVI  D,0         ;D,E= RELATIVE LOC OF AVAIL SPACE IN LAQ
1579 18F7 23     INX  H          ;H,L= START LOC OF LAQ IN PPDB
1580 18F8 19     DAD  D          ;POINT H.L TO ABSOLUTE LOC OF AVAILABLE
1581 18F9 19     DAD  D          ;SPACE IN
1582 18FA 19     DAD  D          ;LAQ
1583 18FB EB     XCHG          ;D,E = AVAIL LOC IN LAQ
1584 18FC 219642 LXI  H,JDV      ;POINT H.L TO MEMORY JDV,IDENT
1585 18FF 0603   MVI  B,3         ;3 BYTES TO MOVE
1586 1901 06CF10 CALL  MOVEB      ;COPY MEMORY JDV, IDENT INTO LAQ
1587 1904 04     INR  B          ;POSITIVE RETURN
1588 1905 C9     RET
1589
1590 ;
1591 ;
1592 ; INPUTS: POINTERS TRSDB,TRSDP,TLAST. 6 BYTE ASCII PACKAGE CODE
1593 ; IN DASDB BYTES 2-7. STACK= JDV,JINDU
1594 ;
1595 ;
1596 ; PROCESS: LOAD 6 BYTES OF ASCII CODE INTO TRANSACTION TABLE IN
1597 ; 1ST AVAILABLE 6 BYTES. TRANSACTION TABLE STARTS AT
1598 ; LOC TRSDB, ENDS AT LOC TLAST, WITH TRSDP POINTING TO
1599 ; LOC IN TABLE WHERE PREVIOUS 6 BYTES ENDED.
1600 ;
1601 ;
1602 ; OUTPUTS: PACKAGE CODE IN ASCII RESIDES IN TRANSACTION TABLE
1603 ; STARTING AT LOC POINTED TO BY H.L. TRSDP POINTS TO
1604 ; START OF NEXT 6 BYTE BLOCK IN TRANSACTION TABLE.STACK
1605 ; =JDV,JINDU
1606 ;
1607 ; TEST TO SEE IF VALUE IN LOC TRSDP = 0
1607 1906 2A5D43 DSPAK: LHLD  TRSDP ;H.L= LOC IN TRANS TAB WHERE LEFT OFF
1608 1909 7E     MOV  A,M          ;A= VALUE IN LOC TRSDP
1609 190A D600   SUI  0
1610 190C CA1719 JZ   BCONT      ;JMP IF VALUE=0
1611 190F 3E00   MVI  A,0         ;OTHERWISE, RESET COUNTER
1612 1911 325A43 STA  TCONT      ;TCONT=0
1613 1914 C31E19 JMP  BTRDP      ;JMP TO BTRDP
1614 1917 3A5A43 BCONT: LDA  TCONT ;SET TCONT = TCONT+1
1615 191A 3C     INR  A
1616 191B 325A43 STA  TCONT
1617 191E 2A5D43 BTRDP: LHLD  TRSDP ;INCREMENT TRSDP POINTER
1618 1921 23     INX  H          ;TRSDP=TRSDP+1
1619 1922 225D43 SHLD  TRSDP
1620
1621 ;
1622 ; TEST FOR WRAPAROUND OF TRANSACTION TABLE
1622 1925 3AF23D LDA  TLAST      ;A= TLAST LSB
1623 1928 57     MOV  D,A          ;B= SAME
1624 1929 3A5D43 LDA  TRSDP      ;A= TRSDP LSB
1625 192C 92     SUB  D          ;A= TRSDP LSB - TLAST LSB

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1623 192D C24919      JNZ  TTCNT      ;JMP IF NOT EQUAL
1624 1930 3AF33D      LDA  TLAST+1    ;A= TLAST MSB
1625 1933 57          MOV  D,A        ;D= SAME
1626 1934 3A5E43      LDA  TRSDP+1    ;A= TRSDP MSB
1627 1937 92          SUB  D          ;A= TRSDP MSB - TLAST MSB
1628 1938 C24919      JNZ  TTCNT      ;JMP IF NOT EQUAL
1629
1630                ; WRAPAROUND BECAUSE TRSDP = TLAST
1631 193B 3E00          MVI  A,0        ;A=0
1632 193D 325A43      STA  TCONT      ;TCONT=0
1633 1940 2AF03D      LHLD TRSDB     ;H,L= START LOC OF TRANSACTION TAB
1634 1943 225D43      SHLD TRSDP     ;RESET TRSDP POINTER TO START
1635 1946 C30619      JMP  DSPAK      ;START AT TOP OF TABLE TO LOOK FOR 6 BYTE
1636
1637                ; NO WRAP-AROUND. CHECK TO SEE IF 6 BYTES HAS BEEN FOUND
1638
1639 1949 3A5A43      TTCNT: LDA  TCONT      ;A= TCONT
1640 194C D606          SUI  6          ;TCONT=6?
1641 194E C20619      JNZ  DSPAK      ;JMP IF TCONT NOT=6
1642 1951 3E00          MVI  A,0        ;A=0
1643 1953 325A43      STA  TCONT      ;TCONT=0 (RESET COUNTER)
1644                ; SIX CONSECUTIVE BYTES FOUND. POINT H.L TO START OF 6 BYTES
1645 1956 2A5D43      LHLD TRSDP
1646 1959 2B          DCX  H
1647 195A 2B          DCX  H
1648 195B 2B          DCX  H
1649 195C 2B          DCX  H
1650 195D 2B          DCX  H
1651 195E 2B          DCX  H          ;H,L= START OF 6 AVAILABLE BYTES IN TABLE
1652 195F E5          PUSH H         ;STACK= SAME
1653 1960 EB          XCHG         ;D,E=SAME
1654 1961 212843      LXI  H,DASDB   ;H,L=DASDB LOC
1655 1964 010200      LXI  B,2       ;B,C=2
1656 1967 09          DAD  B         ;H,L=BYTE 2 OF DASDB= PACKAGE CODE START
1657 1968 010600      LXI  B,6       ;B,C= 6
1658 196B CDC410      CALL MOVE      ;MOVE 6 BYTES FROM H.L TO DIE
1659 196E E1          POP  H         ;H,L= LOC IN TRANSACTION TABLE OF CODE
1660 196F EB          XCHG         ;D,E=SAME
1661                ; PUT LOC OF STORED PAK CODE INTO MEMORY IDENT
1662 1970 219742      LXI  H,IDENT   ;H,L= IDENT LSB LOC
1663 1973 73          MOV  M,E       ;IDENT LSB= REG E
1664 1974 23          INX  H        ;H,L= IDENT MSB LOC
1665 1975 72          MOV  M,D       ;IDENT MSB= REG D
1666 1976 C9          RET          ;RET TO DPLAQ WITH IDENT=LOC IN TRANS TAB
1667
1668                ; SUBROUTINE DSOUT
1669
1670                ; WILL SEND DAS THE 1ST MESSAGE POINTED TO BY SMP IN FS. DEFAULTS TO
1671                ; IDLE MESSAGE IF NO MESSAGES IN FS. WILL RESEND LAST PSC MESSAGE IF
1672                ; CALLED FROM ICEXC WITH REG A=FFH. OR IF CALLED FROM DSDEC WITH
1673                ; REG A= FFH WILL SEND MESSAGE IN MOBUF(COULD BE SAME AS LAST PSC
1674                ; MESSAGE SENT)
1675
1676                ; INPUTS: IF REG A=FFH, MOBUF WILL CONTAIN DESIRED MESSAGE TO DAS.
1677                ; OTHERWISE, SMP WILL BE POINTING TO 1ST MESSAGE IN FS WAIT-
1678                ; ING TO BE SENT TO DAS (DEFAULT TO IDLE MESSAGE)
1679
1680                ; PROCESS: IF REG A=FFH ON ENTRY, SEND MOBUF TO DAS. IF THERE ARE NO
1681                ; MESSAGES IN FS, LOAD IDLE MESSAGE INTO MOBUF AND SEND TO
1682                ; DAS. IF THERE IS A MESSAGE IN FS, LOAD IT INTO MOBUF AND
1683                ; SEND TO DAS AND THEN UPDATE SMP POINTER TO NEXT MESSAGE.
1684
1685                ; OUTPUTS: DAS MESSAGE OUTPUT TO 5080 BOARD
1686
1687 1977 06FF      DSOUT: MVI  B,0FFH      ;JMP TO SMOBF IF REG A=FFH IN ORDER
1688 1979 90          SUB  B          ;TO RESEND MOBUF
1689 197A CAF119     JZ   SMOBF
1690                ; REG A MUST=00.
1691                ; CHECK TO SEE IF ANY MESSAGES IN FS.
1692 197D 2A5843     LHLD SMP          ;H,L= SMP VALUE
1693 1980 7C          MOV  A,H
1694 1981 B5          ORA  L          ;SMP=0000H?
1695 1982 C2A019     JNZ  YFS         ;JMP IF MESSAGES IN FS
1696
1697                ; NO MESSAGES IN FS. DEFAULT TO IDLE MESSAGE.
1698 1985 3A2942     LDA  RS          ;SET SEQ BYTE= TOGGLED RS
1699 1988 323843     STA  MOBUF
1700 198B 3E49       MVI  A,'I'       ;A= ASCII 'I'
1701 198D 323943     STA  MOBUF+1
1702 1990 3E20       MVI  A,20H      ;A= ASCII SPACE
1703 1992 213A43     LXI  H,MOBUF+2   ;STORE INTO MOBUF BYTES 2-12 AS SPACES
1704 1995 060B       MVI  B,11
1705 1997 77          LPMOB: MOV  M,A
1706 1998 23          INX  H
1707 1999 05          DCR  B
1708 199A C29719     JNZ  LPMOB
1709 199D C3F119     JMP  SMOBF      ;JMP WHEN ALL BYTES 2-12 LOADED AS SPACE
1710

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1711 ; LOAD MESSAGE IN FS POINTED TO BY SMP INTO MOBUF
1712 19A0 3A2942 YFS: LDA RS ;SET SEQ BYTE = TOGGLED RS
1713 19A3 323843 STA MOBUF
1714 19A6 2A5843 LHLD SMP ;H.L= BYTE 1 OF MESSAGE IN FS
1715 19A9 23 INX H
1716 19AA 113943 LXI D,MOBUF+1 ;D.E= MOBUF BYTE 1 LOC
1717 19AD 010C00 LXI B,12 ;B.C= 12 BYTES TO MOVE
1718 19B0 CDC410 CALL MOVE ;MOVE MESSAGE FROM FS TO MOBUF
1719 ;
1720 ; TEST TO SEE IF THIS IS THE LAST MESSAGE WAITING IN FS. IT'S LAST
1721 ; MESSAGE IF SMP=LMP(AND WE ALREADY KNOW LMP AND SMP ARE NON-ZERO)
1722 19B3 2A5843 LHLD SMP ;H.L= SMP VALUE
1723 19B6 EB XCHG ;D.E= SAME
1724 19B7 2A5443 LHLD LMP ;H.L= LMP VALUE
1725 19BA 7B MOV A,E ;A= SMP LSB
1726 19BB 95 SUB L ;A= SMP LSB- LMP LSB
1727 19BC C2D919 JNZ OTHFS ;JMP IF SMP LSB NOT= SMP LSB
1728 19BF 7A MOV A,D ;A= SMP MSB
1729 19C0 94 SUB H ;A= SMP MSB- LSB MSB
1730 19C1 C2D919 JNZ OTHFS ;JUMP IF NOT=
1731 ;
1732 ; THIS WAS LAST MESSAGE IN FS. RESET THE LOAD MESSAGE POINTER LMP AND
1733 ; SEND MESSAGE POINTER SMP.
1734 19C4 2A5843 LHLD SMP ;H.L= SMP VALUE
1735 19C7 EB XCHG ;D.E= ORIG SMP VALUE
1736 19C8 3E00 MVI A,0 ;A=0
1737 19CA 325843 STA SMP ;SMP=0000H
1738 19CD 325943 STA SMP+1
1739 19D0 325443 STA LMP ;LMP=0000H
1740 19D3 325543 STA LMP+1
1741 19D6 C3ED19 JMP RTTCB ;JMP TO RTTCB
1742 ;
1743 ; THIS IS NOT THE ONLY MESSAGE IN FS. POINT SMP POINTER TO NEXT ONE
1744 19D9 2A5843 OTHFS: LHLD SMP ;H.L= SMP VALUE
1745 19DC EB XCHG ;D.E= ORIG SMP
1746 19DD 2A5843 LHLD SMP ;POINT H.L TO BYTE 14 OF MESSAGE TCB
1747 19E0 010E00 LXI B,14
1748 19E3 09 DAD B ;H.L= BYTE 14 LOC
1749 19E4 7E MOV A,M ;A= BYTE 14
1750 19E5 325843 STA SMP ;SMP= LSB OF LINK
1751 19E8 23 INX H ;H.L= BYTE 15 LOC
1752 19E9 7E MOV A,M ;A= BYTE 15
1753 19EA 325943 STA SMP+1 ;SMP+1 = MSB OF LINK
1754 ;
1755 ; RETURN MESSAGE TCB TO FS
1756 19ED EB RTTCB: XCHG ;H.L= ORIG SMP
1757 19EE CD4F11 CALL PUTFS ;RETURN TCB POINTED TO BY ORIG SMP TO FS
1758 ;
1759 ;LOAD CR,LF INTO MOBUF,CALC CHECKSUM OF MOBUF AND LOAD INTO MOBUF.
1760 ; SEND MOBUF
1761 ; TO DAS BY CALLING PUTDS. RESET DASIP
1762 19F1 213843 SMOBF: LXI H,MOBUF ;POINT TO MOBUF
1763 19F4 CD211A CALL CHKSM ;A= PRINTABLE ASCII CHECKSUM
1764 19F7 324543 STA MOBUF+13 ;PUT CHKSM INTO MOBUF
1765 19FA 212843 LXI H,DASDB ;H.L= LOC OF DAS MESSAGE INPUT BUFFER
1766 19FD 224B43 SHLD DASIP ;RESET POINTER TO START OF BUFFER
1767 1A00 214643 LXI H,MOBUF+14 ;H.L= BYTE 14 LOC
1768 1A03 360A MVI M,0AH ;SET ASCII LF IN MOBUF
1769 1A05 23 INX H ;H.L= BYTE 15 LOC
1770 1A06 360D MVI M,0DH ;SET ASCII CR IN MOBUF
1771 1A08 CD0C1A CALL PUTDS ;PUTDS STARTS TRANSMISSION OF MOBUF
1772 1A0B C9 RET ;RET TO CALLER
1773 ;
1774 ;
1775 ; INPUTS: PSC TO DAS MESSAGE RESIDING IN BUFFER MOBUF. POINTER
1776 ; MBUFF POINTS TO BYTE 0 OF MOBUF.DTYBS=0
1777 ;
1778 ;
1779 ; PROCESS: DOUT BYTE 0 OF MOBUF TO I/C 5080 BOARD.
1780 ;
1781 ; OUTPUT: MOBUF BYTE 0 OUTPUT TO I/C 5080 BOARD.MBUFP POINTS TO
1782 ; BYTE 1 OF MOBUF. REST OF MESSAGE IN MOBUF WILL BE
1783 ; SENT BY INTERRUPT DRIVERS IN SUBROUTINE ICINT.ALSO
1784 ; SET DTYBS=1 FOR USE IN ICINT.
1785 1A0C 3E01 PUTDS: MVI A,01H
1786 1A0E 324D43 STA DTYBS ;DTYBS=1 FOR USE IN ICINT
1787 1A11 2A5643 LHLD MBUFF
1788 1A14 23 INX H
1789 1A15 225643 SHLD MBUFF ;MBUFF=MBUFF+1 (MOBUF BYTE 1)
1790 1A18 2A063D LHLD P5080 ;H.L= SIN ADDR FOR I/C 5080 BOARD
1791 1A1B 25 DCR H ;H=FE FOR DOUT OPERATION
1792 1A1C 3A3843 LDA MOBUF ;A= MOBUF BYTE 0
1793 1A1F 77 MOV M,A ;DOUT MOBUF BYTE 0 TO REG FEXX
1794 1A20 C9 RET ;RETURN TO DSOUT
1795 ;
1796 ; SUBROUTINE CHKSM
1797 ; INPUT: H.L POINTS TO 1ST OF 13 CONSECUTIVE BYTES TO CALCULATE

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1798 ;
1799 ;
1800 ;
1801 ;
1802 ;
1803 ;
1804 ;
1805 ;
1806 1A21 060D   CHKSM: MVI B,13 ;SET LOOP COUNTER
1807 1A23 0E00   MVI C,0 ;INITIALIZE C REG TO ZERO
1808 1A25 7E     MCKSM: MOV A,M ;A=BYTE
1809 1A26 E67F   ANI 7FH ;MASK BIT 7
1810 1A28 81     ADD C
1811 1A29 4F     MOV C,A ;C=CURRENT CHKSM
1812 1A2A 23     INX H ;POINT TO NEXT BYTE
1813 1A2B 05     DCR B ;DEC COUNTER
1814 1A2C C251A  JNZ MCKSM ;DO ALL 12 BYTES
1815 1A2F 79     MOV A,C ;A=CHECKSUM
1816 ;
1817 ; SET BIT 7 = 0
1818 1A30 E67F   ANI 7FH ;BIT 7=0
1819 1A32 47     MOV B,A ;STORE RESULT 0
1820 ;
1821 ; SET BIT 6 = COMPLEMENT OF BIT 5
1822 1A33 E620   ANI 20H ;LOOK AT BIT 5
1823 1A35 CA3C1A JZ BIT0 ;JMP IF BIT 5 = 0
1824 ; BIT 5 = 1. THEREFORE MAKE BIT 6 = 0.
1825 1A38 78     MOV A,B ;A= CHECKSUM WITH BIT 7=0
1826 1A39 E6BF   ANI 0BFH ;ZERO OUT BIT 6
1827 1A3B C9     RET ;RET WITH PRINTABLE ASCII CHECKSUM IN A
1828 ; BIT 5 = 0. THEREFORE MAKE BIT 6 = 1.
1829 1A3C 78     BIT0: MOV A,B ;A = CHECKSUM WITH BIT 7 = 0.
1830 1A3D F640   ORI 40H ;SET BIT 6=1
1831 1A3F C9     RET ;RET WITH PRINTABLE ASCII CHKSM IN REG A
1832 ;
1833 ; SUBROUTINE ZPACK
1834 ;
1835 ; INPUTS: H.L POINTS TO LOC WHERE IT IS DESIRED TO CLEAR 6 BYTES
1836 ; IN TRANSACTION TABLE
1837 ; PROCESS: IF H.L IS WITHIN TRANSACTION TABLE, CLEAR THE 6 BYTES
1838 ; STARTING AT H.L. IF H.L NOT WITHIN TABLE, DO FATAL
1839 ; ERROR # 1
1840 ;
1841 ; OUTPUT: 6 BYTES OF TRANSACTION TABLE CLEARED STARTING AT INPUT
1842 ; H.L LOCATIONS.
1843 ;
1844 ; CHECK TO SEE IF H.L WITHIN TRANSACTION TABLE LIMITS
1845 1A40 E5     ZPACK: PUSH H ;STACK= H.L INPUT
1846 1A41 2AF03D LHL TRSDB
1847 1A44 44     MOV B,H
1848 1A45 4D     MOV C,L ;B,C=START ADDR OF TRANS TABLE
1849 1A46 2AF23D LHL TLAST ;H,L= END ADDR OF TRANS TAB
1850 1A49 EB     XCHG ;D,E=SAME
1851 1A4A E1     POP H ;H,L= ADDR OF DESIRED DELETE
1852 1A4B E5     PUSH H ;STACK= SAME
1853 1A4C 7D     MOV A,L ;A= LSB OF DESIRED DELETE ADDR
1854 1A4D 91     SUB C ;16-BIT
1855 1A4E 7C     MOV A,H ;COMPARE
1856 1A4F 98     SBB B ;HERE
1857 1A50 FA661A JM FATLE ;JMP IF PASSED ADDR < TRANSLATION TABLE
1858 ;
1859 ; IF GET TO HERE, DESIRED DELETE ADDR IS GREATER THAN START OF TRANS TAB
1860 ; NOW SEE IF DELETE ADDR < TLAST OF TRANS TABLE
1861 1A53 7B     TTOP: MOV A,E ;A=LSB OF TLAST
1862 1A54 95     SUB L ;16-BIT
1863 1A55 7A     MOV A,D ;COMPARE
1864 1A56 9C     SBB H ;HERE
1865 1A57 FA661A JM FATLE ;JMP IF PASSED ADDR > TLAST
1866 ;
1867 1A5A E1     ZLOOP: POP H ;H,L= ADDR OF DESIRED DELETE
1868 1A5B 3E00   MVI A,0 ;A=0
1869 1A5D 0606   MVI B,6 ;COUNT=6
1870 1A5F 77     ZTRAN: MOV M,A ;ZERO THE LOC
1871 1A60 23     INX H ;INCREMENT LOC
1872 1A61 05     DCR B ;DEC COUNTER
1873 1A62 C25F1A JNZ ZTRAN ;DO ALL 6 LOC
1874 1A65 C9     RET ;RET TO CALLER
1875 ;
1876 ; H.L NOT WITHIN TRANSACTION TABLE. DO FATAL ERROR #1
1877 1A66 3E01   FATLE: MVI A,01H
1878 1A68 CD0C2C CALL FATAL
1879 1A6B C9     RET
1880 ;
1881 ; SUBROUTINE TRANP
1882 ;
1883 ; INPUTS:H.L POINTS TO 1ST OF 6 BYTES DESIRED TO BE MOVED
1884 ;
1885 ; PROCESS: MOVE 6 BYTES POINTED TO BY H,L INTO BUFFER BYT16
1886 ; IN BYTES 2-7 LOC

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1886
1887
1888
1889
1890 1A6C E5      TRANP:  PUSH H           ;STACK = LOC OF 6 BYTES TO MOVE
1891 1A6D 211843 LXI H,BYT16    ;H.L= BYT16 LOC
1892 1A70 110200 LXI D,2        ;
1893 1A73 19       DAD D           ;H.L= BYTE 2 OF BYT16
1894 1A74 EB       XCHG          ;D,E= SAME
1895 1A75 010600 LXI B,6        ;B,C= 6 COUNT
1896 1A78 E1       POP H          ;H,L= LOC OF 6 BYTES
1897 1A79 CDC410  CALL MOVE      ;MOVE 6 BYTES FROM H.L TO D1E LOC
1898 1A7C C9       RET           ;RETURN
1899
1900
1901
1902
1903
1904
1905
1906
1907
1908
1909
1910
1911 1A7D CDCB0F  DSLD:   CALL GETFS    ;H.L= LOC OF AVAIL FS
1912 1A80 EB       XCHG          ;D,E= LOC OF AVAIL FS
1913 1A81 211843 LXI H,BYT16    ;H.L= BYTE 0 LOC OF BYT16 BUFFER
1914 1A84 011000 LXI B,16       ;B,C= 16 BYTES TO MOVE
1915 1A87 D5       PUSH D         ;STACK= LOC OF AVAIL FS
1916 1A88 CDC410  CALL MOVE      ;COPY BYT16 INTO FS
1917
1918
1919 1A8B 2A5443    LHL D LMP      ;H.L= LMP VALUE
1920 1A8E 7C       MOV A,H        ;A= LMP MSB
1921 1A8F B5       ORA L          ;OR WITH LMP LSB
1922 1A90 C29B1A  JNZ MLINK     ;JMP IF LMP NOT ZERO
1923
1924
1925
1926
1927
1928
1929 1A9B 2A5443    MLINK:  LHL D LMP      ;H.L= LMP LOC = LOC IN FS WHERE LAST
1930
1931
1932
1933
1934
1935
1936
1937
1938
1939
1940
1941
1942
1943
1944
1945
1946
1947
1948 1A9E EB       XCHG          ;D,E= LMP LOC OF LAST MESSAGE
1949 1A9F 210E00  LXI H,14      ;H.L= BYTE 14 OF LAST MESSAGE
1950 1AA2 19       DAD D           ;D,E= LOC WHERE NEW MESSAGE JUST STORED
1951 1AA3 D1       POP D          ;BYTE 14 = LSB OF NEW LOC
1952 1AA4 73       MOV M,E        ;POINT TO BYTE 15
1953 1AA5 23       INX H          ;BYTE 15 OF PREVIOUS MSG=MSB OF NEW LOC
1954
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1956
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1970
1971
1972

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OUTPUTS: BYTES 2-7 OF BUFFER BYT16 CONTAIN THE 6 BYTES
 ORIGINALLY POINTED TO BY H.L

SUBROUTINE DSLD

INPUTS: MESSAGE TO BE LOADED INTO FS RESIDES IN BUFFER BYT16.
 LMP POINTS TO FS LOC WHERE LAST MESSAGE WAS LOADED.

OUTPUTS: MESSAGE IS IN FS POINTED TO BY LMP. MESSAGE IS LINKED
 TO OTHER MESSAGES IN FS.

PROCESS: GET 1ST AVAIL FS BLOCK AND LOAD MESSAGE BUFFER BYT16
 INTO IT. LINK THIS FS BLOCK TO PREVIOUS MESSAGE
 LOADED.

TEST TO SEE IF THIS IS THE 1ST MESSAGE IN FS

THIS IS THE 1ST MESSAGE IN FS. SET LMP AND SMP TO THIS FS LOC

THIS IS NOT THE ONLY MESSAGE IN FS. LINK IT TO PREVIOUS MESSAGE IN FS

RUNNING MODE INITIALIZATION.

INPUT: INTERRUPTS DISABLED.
 ALL REGISTERS EFFECTED.

POLL PP AND TE PHOTOS TO SEE IF PACKAGE IS PRESENT

RPLPP:

RPP1:

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1973 1AD8 3601      MVI  M,1      ;SET PACKAGE-PRESENT &
1974 1ADA 23       INX  H        ; PACKAGE-WAITING
1975 1ADB 3601      MVI  M,1      ; FLAGS
1976 1ADD C3E81A   JMP  RPPLP
1977
1978 1AE0 C1        RPP0:  POP  B      ;B= COMPOSITE CH #.
1979 1AE1 C5        PUSH B        ; C= INDUCT #
1980 1AE2 0601     MVI  B,1     ;START RELEASE
1981 1AE4 CD0F11   CALL OPUT    ; BELT
1982 1AE7 C1      POP  B        ;B= COMPOSITE CH #. C= INDUCT #
1983
1984 1AE8 04       RPPLP:  INR  B      ;B= NEXT PP CH #
1985 1AE9 04       INR  B
1986 1AEA 0C      INR  C        ;INC INDUCT #
1987 1AEB 3A033D   LDA  NINDU   ;DOUNTIL
1988 1AEE B9      CMP  C        ; INDUCT #
1989 1AEF D2C51A   JNC  RPLPP  ; > NINDU
1990
1991 ;POLL LF TO SET LFDB FLAGS
1992
1993 1AF2 3A3C3D   RPLLF:  LDA  LFO      ;IF
1994 1AF5 B7       ORA  A        ; THERE ARE
1995 1AF6 CA131B   JZ   REINI    ; LF'S
1996 1AF9 3A7442   LDA  TRCHN   ;A= PM5008 COMPOSITE CH #
1997 1AFC DE02     SBI  Z        ; FOR LAST LF
1998 1AFE 47       MOV  B,A      ;B= COMPOSITE CH #
1999 1AFF 3A053D   LDA  NDVT    ;A= # LF'S = COUNT
2000 1B02 4F      MOV  C,A      ;C= COUNT
2001 1B03 C5     RPLF:  PUSH B      ;SAVE CH #, COUNT
2002 1B04 CD171B  CALL POLL    ;POLL LF # C
2003 1B07 C1      POP  B        ;B= CH #,
2004 1B08 C5     PUSH B      ; C= DIVERT #
2005 1B09 DCE720  CC  LFBPC    ;IF LF BLOCKED THEN SET FLAG IN LFDB
2006 1B0C C1     POP  B        ;B= COMPOSITE CH #, C= DIVERT #
2007 1B0D 05     DCR  B        ;B= COMPOSITE CH #
2008 1B0E 05     DCR  B        ; FOR PREVIOUS LF
2009 1B0F 0D     DCR  C        ;DECREMENT DIVERT # (COUNT)
2010 1B10 C2031B  JNZ  RPLF    ; UNTIL =0
2011
2012 ;REINITIALIZE PM5008'S
2013
2014 1B13 CD190C   REINI:  CALL  DEVII ;REINITIALIZE
2015 1B16 C9      RET
2016
2017 ;POLL PM5008 COMPOSITE CHANNEL # B FOR HIGH-TRUE SIGNAL LEVEL.
2018 ; B>=0.
2019
2020 ;INPUT: A=0 MEANS LOW-TRUE, A=0FFH MEANS HIGH-TRUE.
2021 ;OUTPUT: FLAG C SET IF SIGNAL TRUE.
2022 ;ALL REGISTERS EXCEPT B EFFECTED.
2023
2024 1B17 3EFF     POLL:  MVI  A,0FFH ;FOR HIGH-TRUE PHOTOS
2025
2026 1B19 D301     POLL1:  OUT  1      ;SET HIGH-TRUE SELECT BITS
2027 1B1B 3A6C42   LDA  P5008   ;H.L
2028 1B1E 6F      MOV  L,A     ; = COUT AD FOR
2029 1B1F 26FF     MVI  H,0FFH ; 1ST PM5008 BOARD
2030 1B21 78      MOV  A,B     ;A= COMPOSITE CH #
2031 1B22 CD280D  CALL CPMSK   ;CP COUT AD & BIT MASK
2032 1B25 3606     MVI  M,06H   ;ENABLE BOARD & CLEAR ITS FLAGS
2033 1B27 25     DCR  H        ;H.L= DOUT AD
2034 1B28 3600     MVI  M,0     ;FOR SPURIOUS INT'S (SEE PM5008 MANUAL)
2035 1B2A 77     MOV  M,A     ;DOUT TO ENABLE CH
2036 1B2B 24     INR  H        ;H.L= SIN AD
2037 1B2C 7E     MOV  A,M     ;SIN
2038 1B2D 1F     RAR        ;GROUP SELECT BIT TO FLAG C
2039 1B2E C9      RET
2040
2041 ;RUNNING MODE PHOTO SENSOR CELL PROCESSOR.
2042
2043 ;INPUT: B=DEVIC. C=INDEX.
2044 ;ALL REGISTERS EFFECTED.
2045
2046 1B2F 78     PHPC:  MOV  A,B     ;CASENTRY DEVIC= 2 TO 11
2047 1B30 3D     DCR  A
2048 1B31 1E0A     MVI  E,10    ;CASE SWITCH
2049 1B33 CDA20C   CALL CASSW   ; ON 10 CASES
2050 1B36 651C     DW  PPBPC   ;DEVIC=2: PP-BLOCKING
2051 1B38 C51C     DW  TECPC   ;DEVIC=3: TE-CLEARING
2052 1B3A 2E1D     DW  IDBPC   ;DEVIC=4: ID-BLOCKING
2053 1B3C 4F1B     DW  PSRET   ;FOR FUTURE EXPANSION
2054 1B3E 6E1D     DW  UDBPC   ;DEVIC=6: UD-BLOCKING
2055 1B40 3120     DW  UDCPC   ;DEVIC=7: UD-CLEARING
2056 1B42 0043     DW  DCBPC   ;DEVIC=8: DC-BLOCKING
2057 1B44 4F1B     DW  PSRET   ;FOR FUTURE EXPANSION
2058 1B46 E720     DW  LFBPC   ;DEVIC=10: LF-BLOCKING
2059 1B48 EE20     DW  LFCPC   ;DEVIC=11: LF-CLEARING

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2060
2061 ; ERROR EXIT AND NO-OPERATION EXIT
2062
2063 1B4A 3E0A MVI A,10 ;DEFAULT CASE:
2064 1B4C CD080F CALL ERPC ; ERROR 10
2065 1B4F C9 PSRET: RET ;NO OPERATION
2066
2067 ;RUNNING MODE KEYBOARD INPUT CHARACTER PROCESSOR.
2068
2069 ;INPUT: C= KEYBOARD NUMBER, A= INPUT CHARACTER,
2070 ; KBDBP=HL POINTS TO BYTE 4 IN KBDB.
2071 ;OUTPUT: CHAR.
2072 ;ALL REGISTERS EFFECTED.
2073
2074 1B50 32DE42 KBPC: STA CHAR ;SAVE INPUT CHAR
2075 1B53 47 MOV B,A ;B=CHAR
2076 1B54 7E MOV A,M ;IF ER IN KBDB=0 THEN
2077 1B55 B7 ORA A ; SET FLAG Z
2078 1B56 C2731B JNZ IKERZ ;IF KB NOT IN ER MODE
2079
2080 ; TEST IF IN I/C MODE
2081 1B59 3A2B42 LDA ICMO
2082 1B5C D601 SUI 01H ;ICMO=1?
2083 1B5E C2EF1B JNZ NUMER ;JMP IF NOT I/C
2084
2085 ; TEST FOR HALF DAS
2086 1B61 3A0640 LDA HDSEN
2087 1B64 D601 SUI 01H
2088 1B66 CAEF1B JZ NUMER ;JMP IF HALF DAS
2089
2090 ;TEST IF IN 'PROGRAM' MODE, I.E., IS IT DESIRED TO PROGRAM
2091 ;TRANSLATION TABLES WHILE IN HOST MODE?
2092 1B69 3E2E MVI A,2EH ;A=ASCII + = 'SEND' KEY IN PROGRAM MODE
2093 1B6B B8 CMP B ;COMPARE '+' WITH INPUT CHAR
2094 1B6C CAEF1B JZ NUMER ;JMP IF 'SEND' KEY WAS HIT WHILE IN
2095 ; PROGRAM MODE
2096 1B6F D2EF1B JNC NUMER ;JMP IF A DIGIT WAS HIT WHILE IN
2097 ; PROGRAM MODE
2098 1B72 C9 RET ;RET-THEREFORE WON'T PROCESS ANY KB
2099 ;INPUTS EXCEPT 'SEND' OR A DIGIT WHILE
2100 ;IN PROGRAM MODE IF NO ERRORS ARE ON KB
2101 ; (WILL ALSO PROCESS '<' AND '>')
2102 ;AND SYSTEM IS IN HOST MODE BUT NOT IN
2103 ;HALF DAS MODE
2104
2105 ; TEST IF IN I/C MODE
2106 1B73 3A2B42 IKERZ: LDA ICMO
2107 1B76 D601 SUI 01H
2108 1B78 C2411C JNZ KERNZ ;JMP IF NOT I/C
2109
2110 ; TEST FOR HALF DAS
2111 1B7B 3A0640 LDA HDSEN
2112 1B7E D601 SUI 01H
2113 1B80 CA411C JZ KERNZ ;JMP IF HALF DAS
2114
2115 ; TEST IF INPUT WAS A 'CLEAR'
2116 1B83 3ADE42 LDA CHAR ;PICK UP INPUT CHAR
2117 1B86 F630 ORI 30H ;NORMAL ILED INPUT CHAR
2118 1B88 FE3A CPI ':' ;COMPARE WITH 'CLEAR' CHAR
2119 1B8A C8E1B JZ PRCLR ;JMP IF 'CLEAR' WAS HIT
2120 1B8D C9 RET ;RETURN AS SOME KEY OTHER THAN CLEAR WAS
2121 ; HIT WHILE IN I/C MODE(AND NOT HALF DAS)
2122
2123 ; PROCESS 'CLEAR' KEY
2124
2125 ; IF ERROR ON KB IS A 1,2,6,11,OR 12, DO SAME THING NORMALLY DONE FOR
2126 ; 'CLEAR' KEY. IF ANY OTHER ERROR, ONLY CLEAR ERROR OFF KB AND TURN RUN
2127 ; LIGHT ON KB BACK ON.
2128 PRCLR: MOV A,C ;A=KB#
2129 CALL CPKBP ;H.L POINTS TO KBI1 IN KBDB
2130 LXI D,6 ;POINT TO SND LOC(CODE IS TEMP STORED
2131 DAD D ;$ HERE IN ERRPC)
2132 MOV A,M ;A= ERROR CODE
2133 MOV D,A ;STORE ERROR CODE IN D
2134 SUI 1
2135 JZ KRLAQ ;JMP IF ERROR #1
2136 MOV A,D
2137 SUI 2
2138 JZ KRLAQ ;JMP IF ERROR #2
2139 MOV A,D
2140 SUI 6
2141 JZ KRLAQ ;JMP IF ERROR #6
2142 MOV A,D
2143 SUI 11
2144 JZ KRLAQ ;JMP IF ERROR #11
2145 MOV A,D
2146 SUI 12
2147 JZ KRLAQ ;JMP IF ERROR #12
2148
2149 ; THEREFORE JUST WANT TO CLEAR ERROR FROM KB AND TURN RUN LIGHT ON KB

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2148 ; BACK ON(FOR NON-INDUCT KB'S AND ERRORS 0,3,4,5,7,8,9,10,13,14)
2149 1BB5 CDC90C KLRKB: CALL CLRKB ;CLEAR KBDB, KB DISPLAY, TURN BACK ON
2150 ;MODE LIGHT
2151 1BB8 C9 RET ;RETURN
2152 ;
2153 ; BEFORE CLEARING LAQ DURING CALL MCLPC, CLEAR THE PACKAGES CURRENTLY
2154 ; IN LAW FROM TRANSACTION TABLE
2155 1BB9 3A033D KRLAQ: LDA NINDU ;A=NUMBER OF INDUCTS
2156 1BBC B9 CMP C ;COMPARE WITH KB# IN REG C
2157 1BBD DAB51B JC KLRKB ;JUMP IF IT'S A NON-INDUCT KB
2158 1BC0 79 MOV A,C ;A= INDUCT #
2159 1BC1 CD9927 CALL CPPPP ;POINT H.L TO PPDB FOR THIS INDUCT
2160 1BC4 C5 PUSH B ;C= KB# NOW ON STACK
2161 1BC5 110600 LXI D,6 ;D,E= 6
2162 1BC8 19 DAD D ;H,L= LOC OF LAQSZ IN PPDB
2163 1BC9 46 MOV B,M ;B= LAQSZ = LOOP COUNTER
2164 1BCA 23 INX H ;H,L= LOC OF LAQ1 IN PPDB
2165 1BCB 5E MOV E,M ;E= LAQ1
2166 1BCC 1600 MVI D,0 ;D,E= LAQ1
2167 1BCE 23 INX H ;H,L= LOC OF LAQ IN PPDB
2168 1BCF 19 DAD D ;POINT H.L TO LOC IN PPDB
2169 1BD0 19 DAD D ; OF 1ST PACKAGE THAT
2170 1BD1 19 DAD D ; HASN'T YET BEEN INDUCTED
2171 1BD2 05 DCR B ;COUNT = COUNT -1
2172 1BD3 FAEB1B LGLAQ: JM MGLAQ ;JMP IF NO PACKAGES LEFT TO CLEAR IN LAQ
2173 1BD6 23 INX H ;POINT TO IDENT LSB
2174 1BD7 5E MOV E,M ;E= IDENT LSB
2175 1BD8 23 INX H ;POINT TO IDENT MSB
2176 1BD9 56 MOV D,M ;D,E= IDENT
2177 1BDA E5 PUSH H ;STACK= LOC OF IDENT MSB
2178 1BDB EB XCHG ;H,L= IDENT VALUE
2179 1BDC 7C MOV A,H ;A= IDENT MSB
2180 1BDD B5 ORA L ;OR WITH IDENT LSB
2181 1BDE CAE61B JZ NGLAQ ;JMP IF IDENT = 0
2182 1BE1 C5 PUSH B ;SAVE COUNTER
2183 1BE2 CD401A CALL ZPACK ;REMOVE PACK CODE FROM TRANSACTION TABLE
2184 1BE5 C1 POP B ;B= COUNT AGAIN
2185 1BE6 E1 POP H ;H,L= LOC OF IDENT MSB
2186 1BE7 23 INX H ;H,L= LOC OF NEXT PACK JDV VALUE
2187 1BE8 C3D21B JMP LGLAQ ;JMP TO LGLAQ
2188 1BEB C1 MGLAQ: POP B ;B= CB# AGAIN
2189 1BEC C3611C JMP KCLR ;JMP TO KCLR TO CLEAR LAQ.ETC.
2190 1BEF 3ADE42 NUMER: LDA CHAR ;A= CHAR
2191 1BF2 F630 ORI 30H ;NORMALIZED INPUT CHAR
2192 1BF4 FE3A CPI ' ' ;IF CHAR IS
2193 1BF6 D2111C JNC KNDIG ; '0' TO '9'
2194 1BF9 CD9E0D CALL DIGPC ;PROCESS DIGIT
2195 ;
2196 ; ;AUTOMATIC "SEND" BY FIXED LENGTH CODE
2197 ;
2198 1BFC 3ADE42 LDA CHAR ;"SEND" PACKAGE
2199 1BFF E6F0 ANI 0F0H ; ONLY IF
2200 1C01 FE30 CPI 30H ; KEYSWITCH
2201 1C03 C0 RNZ ; IS IN "RUN" POSITION
2202 1C04 3A323D LDA SENDO ;A= LENGTH OF FIXED-LENGTH CODE
2203 1C07 E60F ANI 0FH ; WITHOUT "SEND"
2204 1C09 C8 RZ ;IF LENGTH SPECIFIED (> 0)
2205 1C0A 23 INX H ;POINT H.L TO BYTE 5 IN KBDB
2206 1C0B BE CMP M ;IF CODE ENTERED HAS REACHED
2207 1C0C C0 RNZ ; THAT LENGTH
2208 1C0D CD5B21 CALL SNPC ;"SEND" AUTOMATICALLY
2209 1C10 C9 RET
2210 ;
2211 ;NON-NUMERIC CHAR
2212 ;
2213 1C11 CA611C KNDIG: JZ KCLR ;IF NOT= ' '
2214 1C14 FE3C CPI '<' ;IF '<' THEN
2215 1C16 D21D1C JNC KNSND ; MUST BE ' '
2216 1C19 CD3D21 CALL SNDPC ;NORMAL "SEND"
2217 1C1C C9 RET
2218 ;
2219 1C1D C22B1C KNSND: JNZ KNLT ;IF = '<'
2220 1C20 3A0B3D LDA NXLTB ;USED ONLY FOR
2221 1C23 B7 ORA A ; TRANSLATION
2222 1C24 CA3B1C JZ KBERR ; OPTION
2223 1C27 CDB22F CALL RLTPC ;TRANSLATE FROM DV # TO CODE
2224 1C2A C9 RET
2225 ;
2226 1C2B FE3E KNLT: CPI '>' ;IF IT IS
2227 1C2D C23B1C JNZ KBERR ; ">" KEY
2228 1C30 3A0B3D LDA NXLTB ;USED ONLY FOR
2229 1C33 B7 ORA A ; TRANSLATION
2230 1C34 CA3B1C JZ KBERR ; OPTION
2231 1C37 CD6830 CALL RGTPC ;EXAMINE OR CHANGE DV #
2232 1C3A C9 RET
2233 ;
2234 ;ERRONEOUS CHAR
2235 ;

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2236 1C3B 3E08  KBERR: MVI  A.8      ;ERROR
2237 1C3D CD130F CALL  ERRPC    ; 8
2238 1C40 C9      RET
2239
2240
2241 ;KEYBOARD PROCESSOR WHEN FLAG ER NOT=0, I.E., KEYBOARD SPECIAL MODES
2242
2243 ;INPUT: C= KEYBOARD $, B=CHAR= INPUT CHARACTER.
2244
2245 1C41 3ADE42  KERNZ: LDA  CHAR    ;PICK UP INPUT CHARACTER
2246 1C44 F630    ORI  30H         ;NORMALIZED CHAR
2247 1C46 FE3A    CPI  ' '         ; IF NOT
2248 1C48 CA5D1C  JZ   KBCLR      ; "CLEAR"
2249 1C4B 7E      MOV  A.M         ;A= SPECIAL KB MODE FLAG ER IN KBDB
2250 1C4C D60A    SUI  10         ;ER=10 IS FOR MANIPULATING
2251 1C4E CABA30  JZ   RM0PC     ; TRANSLATE TABLE
2252 1C51 3D      DCR  A          ;ER=11 IS FOR MANIPULATING
2253 1C52 CA0000  JZ   0000H     ; COUNT TABLE
2254 1C55 3D      DCR  A          ;ER=12 IS FOR MANIPULATING
2255 1C56 CA0000  JZ   0000H     ; PACKAGE COUNTS
2256 1C59 CD1C0E CALL  ERMPCL   ;CALL SPECIAL KB MODE PROCESSOR
2257 1C5C C9      RET
2258
2259 ;KB "CLEAR" PROCESSING
2260
2261 1C5D 7E      KBCLR: MOV  A.M         ;IGNORE ALL INPUTS IF
2262 1C5E FE0F    CPI  0FH       ; KB IS IN
2263 1C60 C8      RZ           ; LOCK-OUT MODE
2264
2265 1C61 CD1422  KCLR:  CALL  MCLPC   ;"MASTER CLEAR"
2266 1C64 C9      RET
2267
2268 ;PACKAGE-PRESENT-PHOTO-BLOCKING PROCESSOR.
2269
2270 ;INPUT: C= INDUCT NUMBER, PPDB.
2271 ;OUTPUT: PPDB.
2272 ;ALL REGISTERS & PPDB EFFECTED.
2273
2274 1C65 79      PPBPC: MOV  A.C         ;A = INDUCT $
2275 1C66 CD9927  CALL  CPPPP    ;POINT H.L TO PPDB
2276 1C69 7E      MOV  A.M         ;PICK UP FLAG PP IN PPDB
2277 1C6A 3601    MVI  M.1       ;PP IN PPDB = 1
2278 1C6C B7      ORA  A          ;IF ORIGINAL PP IN PPDB
2279 1C6D C2B21C  JNZ  PPBES     ; == 0
2280
2281 1C70 E5      PUSH H         ;SAVE PPDB POINTER
2282 1C71 C5      PUSH B         ;SAVE INDUCT $
2283 1C72 79      MOV  A.C         ;A = INDUCT $
2284 1C73 2A5942  LHLD OUTDB    ;IF RELEASE IS STOPPED
2285 1C76 CD280D  CALL  CPMSK   ; THEN SET
2286 1C79 A6      ANA  M         ; FLAG Z
2287 1C7A C1      POP  B         ;RESTORE C= INDUCT $
2288 1C7B E1      POP  H         ;RESTORE POINTER TO PPDB
2289 1C7C CAB91C  JZ   PPBPE    ;IF RLS STOPPED THEN ERROR
2290
2291 ;LEGITIMATE PACKAGE BLOCKING PP
2292
2293 1C7F C5      PUSH B         ;SAVE C= INDUCT $
2294 1C80 E5      PUSH H         ;SAVE PPDB POINTER
2295 1C81 CDA20F  CALL  GITRE    ;GET TREE PARAMETERS FOR INDUCT
2296 1C84 2A0A40  LHLD  PPD      ;SET DE = $ PPI PULSES
2297 1C87 EB      XCHG        ; TO TIME OUT
2298 1C88 E1      POP  H         ;POINT TO
2299 1C89 23      INX  H         ; BYTE 2
2300 1C8A 23      INX  H         ; IN PPDB
2301 1C8B E5      PUSH H         ;SAVE POINTER
2302 1C8C 23      INX  H         ;POINT TO BYTE 4
2303 1C8D 23      INX  H         ; IN PPDB
2304 1C8E 019123  LXI  B,PPD10  ;SCHEDULE EVENT PPD10
2305 1C91 CD192D  CALL  RSCHD   ; TO TIMEOUT PP DWELL
2306 1C94 D1      POP  D         ;POINT DE TO BYTE 2 IN PPDB
2307 1C95 210400  LXI  H,4      ;POINT HL TO BYTE 6
2308 1C98 19      DAD  D         ; IN PPDB
2309 1C99 7E      MOV  A.M         ;PICK UP LAQ CURRENT SIZE
2310 1C9A EB      XCHG        ;POINT TO PPDB BYTE 2
2311 1C9B C1      POP  B         ;RESTORE C= INDUCT $
2312 1C9C B7      ORA  A         ;IF LAQ NOT
2313 1C9D CAA41C  JZ   PPWTG    ; EMPTY
2314 1CA0 CD7022  CALL  PSEQR   ;SEQUENCE PACKAGE TO BE INDUCTED
2315 1CA3 C9      RET
2316
2317 ;PACKAGE HAS TO WAIT TO BE ENCODED
2318
2319 1CA4 2B      PPWTG: DCX  H         ;POINT H.L TO BYTE 1 IN PPDB
2320 1CA5 7E      MOV  A.M         ;A= PWTNG IN PPDB
2321 1CA6 3601    MVI  M.1       ;PWTNG IN PPDB = 1
2322 1CA8 B7      ORA  A         ;IF ORIGINAL PWTNG IN PPDB
2323 1CA9 C2BF1C  JNZ  PPBFX    ; = 0

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2324 1CAC 0602      MVI B,2      ;RELEASE
2325 1CAE CD0F11    CALL OPUT     ; OFF
2326 1CB1 C9        RET
2327
2328 ;ILLEGITIMATE PP-BLOCKED INTERRUPT
2329
2330 1CB2 C5        PPBES:  PUSH B      ;STOP RELEASE MECHANISM
2331 1CB3 0602      MVI B,2      ; TO AVOID RELEASING
2332 1CB5 CD0F11    CALL OPUT     ; MULTIPLE
2333 1CB8 C1        POP B        ; PACKAGES
2334 1CB9 3E0C      PPBPE:  MVI A,12 ;INDUCT ERROR
2335 1CBB CD130F    CALL ERRPC   ; 12
2336 1CBE C9        RET
2337
2338 ;FLAG PP=0, RELEASE OFF, BUT FLAG PWTNG NOT=0
2339
2340 1CBF 3E0C      PPBFX:  MVI A,12 ;FATAL ERROR
2341 1CC1 CD0C2C    CALL FATAL   ; * 12
2342 1CC4 C9        RET
2343
2344 ;:
2345 ;TRAILING-EDGE-PHOTO-CLEARING PROCESSOR.
2346
2347 ;INPUT: C= INDUCT $, PPDB, RLSDB, OUTDB.
2348 ;OUTPUT: PPDB, RLSDB, EVENT TECMG OR CMG.
2349 ;ALL REGISTERS EFFECTED.
2350
2350 1CC5 79        TECPC:  MOV A,C      ;A = INDUCT $
2351 1CC6 CD9927    CALL CPPPP   ;POINT H.L TO PPDB
2352 1CC9 35        DCR M        ;RESET PP IN PPDB
2353 1CCA C2261D   JNZ TCPE0   ;IF NOT ALREADY RESET
2354
2355 1CCD C5        PUSH B      ;SAVE C= INDUCT $
2356 1CCE 79        MOV A,C      ;A = INDUCT $
2357 1CCF 2A5942   LHLD OUTDB  ;IF RLS IS STOPPED
2358 1CD2 CD280D   CALL CPMSK  ; THEN SET
2359 1CD5 A6        ANA M        ; FLAG Z
2360 1CD6 C1        POP B        ; C = INDUCT $
2361 1CD7 CA281D   JZ TCPEX   ;IF RLS RUNNING
2362
2363 1CDA CDA227   CALL CPRLP  ;POINT H.L TO TCBP IN RLSDB
2364 1CDD 5E        MOV E,M     ;POINT D.E TO
2365 1CDE 23        INX H      ; PACKAGE BEING
2366 1CDF 56        MOV D,M     ; RELEASED
2367 1CE0 7B        MOV A,E     ;IF TCB
2368 1CE1 B2        ORA D       ; NOT
2369 1CE2 CA201D   JZ TCPFE   ; NULL
2370
2371 ;
2372 ;D.E POINTS TO TCB FOR PACKAGE BEING RELEASED
2373
2373 1CE5 AF        XRA A      ;A=0
2374 1CE6 77        MOV M,A     ;CLEAR TCBP
2375 1CE7 2B        DCX H      ; IN
2376 1CE8 77        MOV M,A     ; RLSDB
2377 1CE9 210900   LXI H,9    ;POINT H.L TO BYTE 9
2378 1CEC 19        DAD D      ; IN TCB
2379 1CED 7E        MOV A,M     ;A= JINDU IN TCB
2380 1CEE B9        CMP C      ;IF IT MATCHES
2381 1CEF C2201D   JNZ TCPFE  ; THE INT
2382
2383 ;
2384 ;JUNK EVENT TEC2, CREATE EVENT TECMG OR CMG
2385
2385 1CF2 2B        DCX H      ;POINT H.L TO BYTE 8 IN TCB
2386 1CF3 118025   LXI D,JUNK ;CHANGE
2387 1CF6 72        MOV M,D     ; EVENT
2388 1CF7 2B        DCX H      ; TEC2
2389 1CF8 73        MOV M,E     ; TO JUNK
2390
2391 1CF9 C5        PUSH B      ;SAVE JINDU
2392 1CFA CDA20F   CALL GITRE  ;GET TREE PARAMETERS FOR INDUCT
2393 1CFD C1        POP B      ;B,C=JINDU
2394 1CFE CDCB0F   CALL GETFS  ;GET FREE BLOCK
2395 1D01 110900   LXI D,9    ;POINT H.L TO BYTE 9
2396 1D04 19        DAD D      ; IN TCB
2397 1D05 71        MOV M,C     ;SET JINDU IN TCB
2398 1D06 2B        DCX H      ;POINT H.L TO BYTE 8 IN TCB
2399 1D07 11CA23   LXI D,TECMG ;POINT D.E TO TECMG
2400 1D0A 3A313D   LDA SIDEO  ;IF
2401 1D0D E610     ANI 010H   ; SIDE INDUCT
2402 1D0F CA151D   JZ TSEVN   ; THEN
2403 1D12 112924   LXI D,CMG  ; POINT D.E TO CMG
2404 1D15 72        MOV M,D     ;SET
2405 1D16 2B        DCX H      ; EVENT
2406 1D17 73        MOV M,E     ; IN TCB
2407 1D18 EB      XCHG      ;POINT D.E TO BYTE 7 IN TCB
2408 1D19 2A3142   LHLD TET   ;SCHEDULE EVENT AT
2409 1D1C CD2F29   CALL ENTCB ; CTIME+ TET TIME .
2410 1D1F C9        RET
2411

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2412 ;BAD TCB, OR NULL TCB
2413
2414 1D20 3E02 TCPFE: MVI A.2 ;FATAL ERROR
2415 1D22 CD0C2C CALL FATAL ; 2
2416 1D25 C9 RET
2417
2418 ;TEPC WITHOUT PPBPC
2419
2420 1D26 3600 TCPE0: MVI M.0 ;ZERO PP IN PPDB
2421
2422 ;TEPC WHILE RELEASE NOT RUNNING
2423
2424 1D28 3E01 TCPEX: MVI A.1 ;ERROR
2425 1D2A CD130F CALL ERRPC ; 1
2426 1D2D C9 RET
2427
2428 ;INDUCT-PHOTO-BLOCKING PROCESSOR. FOR STRAIGHT INDUCT ONLY
2429 ; START TO TRACK PACKAGE.
2430
2431 ;INPUT: C= INDUCT $.
2432 ;ALL REGISTERS EFFECTED.
2433
2434 1D2E CDA227 IDBPC: CALL CPRLP ;POINT HL TO
2435 1D31 23 INX H ; BYTE 2
2436 1D32 23 INX H ; IN RLSDB
2437 1D33 22EF42 SHLD RLDBP ;SAVE RLSDB POINTER
2438 1D36 5E MOV E,M ;POINT DE
2439 1D37 23 INX H ; TO MERGE
2440 1D38 56 MOV D,M ; FIFO
2441 1D39 7B MOV A,E ;IF POINTER
2442 1D3A B2 ORA D ; NOT
2443 1D3B C8 RZ ; NULL
2444
2445 1D3C 210900 LXI H.9 ;POINT H.L TO BYTE 9
2446 1D3F 19 DAD D ; IN TCB
2447 1D40 7E MOV A,M ;IF JINDU IN TCB
2448 1D41 B9 CMP C ; MATCHES
2449 1D42 C0 RNZ ; THIS INTERRUPT
2450
2451 ;THE INTERRUPT IS AS EXPECTED: JUNK EVENT IDB2
2452
2453 1D43 210E00 LXI H.14 ;POINT H.L TO BYTE 14
2454 1D46 19 DAD D ; IN TCB
2455 1D47 E5 PUSH H ;SAVE POINTER
2456 1D48 4E MOV C,M ;POINT B.C TO NEXT
2457 1D49 23 INX H ; PACKAGE
2458 1D4A 46 MOV B,M ; IN MERGE FIFO
2459 1D4B 2AEF42 LHL RLDBP ;DELETE TCB
2460 1D4E 71 MOV M,C ; FROM
2461 1D4F 23 INX H ; INDUCT MERGE
2462 1D50 70 MOV M,B ; FIFO
2463 1D51 210700 LXI H.7 ;POINT H.L TO BYTE 7
2464 1D54 19 DAD D ; IN TCB
2465 1D55 118025 LXI D.JUNK ;CHANGE
2466 1D58 73 MOV M,E ; EVENT IDB2
2467 1D59 23 INX H ; TO
2468 1D5A 72 MOV M,D ; JUNK
2469
2470 ;CREATE EVENT UDB1
2471
2472 1D5B 3E01 MVI A.1 ;GET TREE PARAMETERS
2473 1D5D CDBA0F CALL GTREP ; FOR PPI $ 1
2474 1D60 C1 POP B ;POINT B.C TO BYTE 13
2475 1D61 0B DCX B ; IN OLD TCB
2476 1D62 CD892C CALL GADB1 ;CREATE TCB FOR UDB1
2477 1D65 4F MOV C,A ;C=JINDU
2478 1D66 EB XCHG ;POINT D.E TO BYTE 7 IN TCB
2479 1D67 2A3342 LHL IDT ;POINT H.L TO IDT TABLE
2480 1D6A CD2F29 CALL ENTCB ;SCHEDULE EVENT UDB1
2481 1D6D C9 RET
2482
2483 ;UPDATE-PHOTO-BLOCKING PROCESSOR.
2484
2485 ;INPUT: C= UPDATE PHOTO NUMBER > 0.
2486 ;OUTPUT: NEW SCHEDULED EVENTS, UDBT FOR SDVRS, LDVP FOR SCDVA,
2487 ; JDV FOR SCHDC.
2488 ;ALL REGISTERS AND JUPDA EFFECTED.
2489
2490 1D6E 79 UDBPC: MOV A,C ;SAVE UPDATE PHOTO
2491 1D6F 329542 STA JUPDA ; NUMBER
2492 1D72 0600 MVI B.0 ;B,C=JUPDA
2493 1D74 0B DCX B ;B,C=JUPDA-1
2494 1D75 2A5142 LHL UDDB ;POINT
2495 1D78 09 DAD B ; H.L
2496 1D79 09 DAD B ; TO
2497 1D7A 09 DAD B ; PMUZ(JUPDA)
2498 1D7B 09 DAD B ; IN UDDB
2499 1D7C 03 INX B ;B,C=JUPDA

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2500 1D7D 5E      MOV  E,M      ;D.E
2501 1D7E 70      MOV  M,B      ; = PMUZ,
2502 1D7F 23      INX  H        ; AND NEW
2503 1D80 56      MOV  D,M      ; PMUZ IN UDDB
2504 1D81 70      MOV  M,B      ; = 0
2505 1D82 E5      PUSH H        ;SAVE POINTER TO BYTE 1 IN UDDB
2506 1D83 7B      MOV  A,E      ;IF OLD
2507 1D84 B2      ORA  D        ; PMUZ
2508 1D85 C2AA1D  JNZ  UBPNO   ; NULL
2509
2510 ; UNIDENTIFIED PACKAGE DETECTED AT UD
2511 ;
2512 1D88 CD080F   CALL ERPC     ;ERROR 0 TO MASTER KB
2513 1D8B CDCB0F   CALL GETFS   ;GET FREE BLOCK FOR TCB
2514 1D8E E5      PUSH H        ;SAVE TCB POINTER
2515 1D8F 110D00  LXI  D,13    ;POINT H.L TO BYTE 13
2516 1D92 19      DAD  D        ; IN TCB
2517 1D93 72      MOV  M,D      ;IDENT
2518 1D94 2B      DCX  H        ; IN TCB
2519 1D95 72      MOV  M,D      ; = 0
2520 1D96 2B      DCX  H        ;POINT H.L TO BYTE 11 IN TCB
2521 1D97 3A303D  LDA  ERCHU   ;JDV IN TCB
2522 1D9A 77      MOV  M,A      ; =ERCHU
2523 1D9B 2B      DCX  H        ;POINT H.L TO BYTE 10 IN TCB
2524 1D9C 3A9542  LDA  JUPDA   ;A=JUPDA
2525 1D9F 77      MOV  M,A      ;SET JUPDA IN TCB
2526 1DA0 2B      DCX  H        ;JINDU IN TCB
2527 1DA1 3601    MVI  M,1     ; = 1
2528 1DA3 2B      DCX  H        ;POINT H.L TO BYTE 8 IN TCB
2529 1DA4 72      MOV  M,D      ;EVENT IN TCB
2530 1DA5 2B      DCX  H        ; = 0
2531 1DA6 72      MOV  M,D      ; FOR UDCPC
2532 1DA7 C3541E  JMP  UBPUD
2533 ;
2534 ; INT IS AS EXPECTED; PACKAGE WITHIN UD ZONE
2535 ;
2536 1DAA 210A00  UBPNO: LXI  H,10 ;POINT H.L TO BYTE 10
2537 1DAD 19      DAD  D        ; IN TCB AT PMUZ
2538 1DAE 79      MOV  A,C      ;A=JUPDA
2539 1DAF BE      CMP  M        ;IF JUPDA MATCHES
2540 1DB0 C2A71F  JNZ  UBPFX   ; JUPDA IN TCB
2541 ;
2542 1DB3 D5      PUSH D        ;SAVE TCB POINTER
2543 1DB4 2B      DCX  H        ;POINT H.L TO BYTE 8
2544 1DB5 2B      DCX  H        ; IN TCB
2545 ;
2546 ; TEST TO SEE IF EVENT IS UDB2 OR DCB2
2547 1DB6 2B      DCX  H        ;H.L= EVENT LSB
2548 1DB7 11D924  LXI  D,UDB2  ;D,E=UDB2
2549 1DBA 7B      MOV  A,E      ;A= LSB
2550 1DBB 96      SUB  M        ;
2551 1DBC 47      MOV  B,A      ;B= LSB RESULT
2552 1DBD 23      INX  H        ;
2553 1DBE 7A      MOV  A,D      ;A= MSB
2554 1DBF 96      SUB  M        ;MSB SUBTRACT
2555 1DC0 90      SUB  B        ;(UDB2-EVENT)
2556 1DC1 CA481E  JZ   UBPNI   ;JMP IF EVENT UDB2
2557 ;
2558 ; EVENT = DCB2. THEREFORE SEND DAS THE MESSAGE U(IF MASKED) AND SEND
2559 ; PACKAGE TO ERCHU.DON'T CLEAR FROM TRANSACTION TAB AS PACK STILL TRAKIN
2560 1DC4 3A2B42  LDA  ICMO    ;ICMO=1?
2561 1DC7 A7      ANA  A
2562 1DC8 CA391E  JZ   UBPNI   ;JMP IF NOT IN I/C MODE
2563 1DCB 212C3D  LXI  H,ERMK0+4 ;H.L= MASK BYTE 4
2564 1DCE 7E      MOV  A,M
2565 1DCP E610    ANI  10H     ;LOOK AT BIT 4
2566 1DD1 CA3E1E  JZ   UBPNI   ;JMP IF U NOT MASKED'
2567 ; SEND MESSAGE U TO DAS
2568 ; LOAD PACKAGE CODE INTO BUFFER BYT16 FROM TRANSACTION TABLE IF IDENT
2569 ; NOT=0. IF IDENT=0.LOAD 6 ASCII 0'S INTO BYT16.DON'T CLEAR PACKAGE
2570 ; CODE FROM TRANSACTION TABLE AS PACKAGE IS STILL BEING TRACKED.
2571 1DD4 E1      POP  H        ;H.L= TCB LOC
2572 1DD5 E5      PUSH H        ;STACK=SAME
2573 1DD6 110C00  LXI  D,12    ;
2574 1DD9 19      DAD  D        ;H.L= IDENT LSB
2575 1DDA 5E      MOV  E,M
2576 1ddb 23      INX  H        ;H.L=IDENT MSB
2577 1DDC 56      MOV  D,M      ;D,E=IDENT=LOC OF PACK CODE IN TRANS TAB
2578 1DDD 7B      MOV  A,E      ;A= IDENT LSB
2579 1DDE B2      ORA  D        ;OR WITH IDENT MSB
2580 1DDF CAE91D  JZ   UPACK   ;JMP IF IDENT = 0
2581 1DE2 EB      XCHG        ;H.L= TRANS TAB LOC OF PACK CODE
2582 1DE3 CD6C1A  CALL TRANP  ;LOAD PACK CODE INTO BYT16 BYTES 2-7
2583 1DE6 C3FD1D  JMP  PSTAK  ;JMP TO PSTAK
2584 1DE9 3E30    MVI  A,'0'   ;STORE ASCII 0 INTO BYT16 FOR PACK CODE
2585 1DEB 321A43  STA  BYT16+2 ;
2586 1DEE 321B43  STA  BYT16+3 ;
2587 1DF1 321C43  STA  BYT16+4 ;

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2588 1DF4 321D43      STA  BYT16+5      ;
2589 1DF7 321E43      STA  BYT16+6      ;
2590 1DFA 321F43      STA  BYT16+7      ;
2591 1DFD E1          PSTAK: POP  H      ;H.L+TCB LOC
2592 1DFE E5          PUSH H      ;STACK= SAME
2593                ; CONVERT BINARY JDV AND JINDU INTO ASCII AND STORE IN BYT16 BUFFER
2594                ;
2595 1DFF 110900       LXI  D.9
2596 1E02 19          DAD  D
2597 1E03 7E          MOV  A.M          ;H.L= JINDU LOC
2598 1E04 CD010E       CALL ECDIG        ;A= JINDU BINARY VALUE
2599 1E07 3A7E42       LDA  XYZ1         ;CONVERT TO ASCII
2600 1E0A 322043       STA  BYT16+8      ;STORE ASCII 10'S
2601 1E0D 3A7F42       LDA  XYZ0
2602 1E10 322143       STA  BYT16+9      ;STORE ASCII 1'S
2603 1E13 E1          POP  H
2604 1E14 E5          PUSH H          ;H.L= TCB LOC
2605 1E15 110B00       LXI  D.11        ;STACK= SAME
2606 1E18 19          DAD  D
2607 1E19 7E          MOV  A.M          ;H.L= JDV BINARY VALUE LOC
2608 1E1A CD010E       CALL ECDIG        ;A= JDV BINARY VALUE
2609 1E1D 3A7D42       LDA  XYZ2         ;CONVERT TO ASCII
2610 1E20 322243       STA  BYT16+10     ;STORE ASCII 100'S
2611 1E23 3A7E42       LDA  XYZ1
2612 1E26 322343       STA  BYT16+11     ;STORE ASCII 10'S
2613 1E29 3A7F42       LDA  XYZ0
2614 1E2C 322443       STA  BYT16+12     ;STORE ASCII 1'S
2615 1E2F 211843       LXI  H.BYT16     ;H.L= BYT16 LOC
2616 1E32 23          INX  H
2617 1E33 3E55        MVI  A,'U'        ;MAKE BYTE 1 = ASCII U
2618 1E35 77          MOV  M.A
2619 1E36 CD7D1A      CALL DSLD         ;LOAD BYT16 INTO FS
2620                ;
2621                ; PUT ERROR 5 ON KB
2622 1E39 3E05        UBPN2: MVI  A.5     ;A=5
2623 1E3B CD080F       CALL ERPC        ;DISPLAY ON KB
2624                ;
2625                ; SEND PACKAGE TO ERCHU
2626 1E3E E1          UBPN3: POP  H      ;H.L= TCB LOC
2627 1E3F E5          PUSH H          ;STACK= SAME
2628 1E40 110B00       LXI  D.11
2629 1E43 19          DAD  D
2630 1E44 3A303D       LDA  ERCHU       ;H.L= JDV LOC IN TCB
2631 1E47 77          MOV  M.A         ;A= ERCHU
2632                ;SET JDV=ERCHU
2633                ;
2634                ; CHANGE EVENT DCB2 TO UDB3
2635 1E48 E1          UBPN1: POP  H      ;H.L= TCB LOC
2636 1E49 E5          PUSH H          ;STACK= SAME
2637 1E4D 19          DAD  D
2638 1E4E 119925       LXI  D.UDB3      ;H.L= EVENT MSB
2639 1E51 72          MOV  M.D         ;D.E=UDB3
2640 1E52 2B          DCX  H
2641 1E53 73          MOV  M.E         ;D.E= UDB3= EVENT IN TCB
2642                ;
2643 1E54 D1          UBPU2: POP  D      ;POINT D.E TO TCB
2644 1E55 E1          POP  H          ;POINT H.L TO BYTE 2
2645 1E56 23          INX  H          ; IN UDDB
2646 1E57 7E          MOV  A.M         ;IF
2647 1E58 23          INX  H          ; PUD
2648 1E59 B6          ORA  M          ; IN UDDB
2649 1E5A CA681E      JZ   UPPUD       ; =0. THEN FINE
2650                ;
2651                ; FAKE AN UD-CLEAR INTERRUPT IF IT WAS LOST
2652                ;
2653 1E5D E5          PUSH H          ;SAVE
2654 1E5E D5          PUSH D          ; POINTERS
2655 1E5F 3A9542       LDA  JUPDA       ;PICK UP UPDATE
2656 1E62 4F          MOV  C.A        ; PHOTO #
2657 1E63 CD3120       CALL UDCPC       ;FAKE UD-CLEAR INTERRUPT
2658 1E66 D1          POP  D          ;RESTORE
2659 1E67 E1          POP  H          ; POINTERS
2660 1E68 72          MOV  M.D        ;POINT PUD
2661 1E69 2B          DCX  H          ; IN UDDB
2662 1E6A 73          MOV  M.E        ; TO TCB
2663 1E6B 210B00       LXI  H.11       ;POINT H.L TO BYTE 11
2664 1E6E 19          DAD  D          ; IN TCB
2665 1E6F E5          PUSH H          ;SAVE POINTER
2666 1E70 3A9542       LDA  JUPDA       ;A=JUPDA
2667 1E73 CDAD0F       CALL GUTR1      ;GET TREE PARAMETERS FOR UD # JUPDA
2668                ;
2669                ; SCHEDULE BMG FOR SIDE INDUCT
2670                ;
2671 1E76 3A313D       LDA  SIDE0       ;IF SIDE INDUCT
2672 1E79 E610        ANI  010H        ; THEN
2673 1E7B C4AE1F       CNZ  SCBMG       ; SCHEDULE BMG
2674                ;
2675                ; CONTINUE TO TRACK PACKAGE OR SCHEDULE DIVERTS

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2676
2677 1E7E 2AB642      LHL D CTIME      ;SET UDBT FOR
2678 1E81 229942      SHLD UDBT        ; SDVRS
2679 1E84 3A9542      LDA JUPDA        ;A
2680 1E87 D602         SUI 2            ; = UPDATE PHOTO # - 2
2681 1E89 2A3D42      LHL D LDVDB      ;POINT H.L TO LDV(1)
2682 1E8C 85          ADD L            ;POINT H.L
2683 1E8D 6F          MOV L,A          ; TO LDV(JUPDA-1)
2684 1E8E 22F742      SHLD LDVP        ;SET LDVP FOR SCDVA, UBNXU
2685 1E91 C1          POP B            ;POINT B.C TO BYTE 11 IN TCB
2686 1E92 0A          LDAX B           ;A=JDV
2687 1E93 329642      STA JDV          ;SAVE JDV FOR SCHDC AND SCDVA
2688 1E96 03          INX B            ;POINT B.C TO BYTE 13
2689 1E97 03          INX B            ; IN TCB
2690 1E98 B7          ORA A            ;JDV=0 MEANS TO
2691 1E99 CA4A1F      JZ UBNXU         ; RECIRCULATE
2692 1E9C 3D          DCR A            ;A=JDV-1
2693 1E9D BE          CMP M            ;IF JDV-1 >= LDV(JUPDA-1),
2694 1E9E DA4A1F      JC UBNXU         ; I.E., IF JDV > LDV(JUPDA-1)
2695 1EA1 23          INX H            ;POINT H.L TO LDV(JUPDA)
2696 1EA2 BE          CMP M            ;IF JDV-1 < LDV,
2697 1EA3 D24A1F      JNC UBNXU       ; I.E., IF JDV <= LDV
2698
2699 ; PACKAGE IS TO BE DIVERTED BEFORE NEXT UD OR END OF CONVEYOR
2700
2701 1EA6 C5            PUSH B            ;STACK = BYTE 13 OF PACKAGE TCB
2702 1EA7 3A2B42      LDA ICHO         ;
2703 1EAA D601         SUI 1            ;
2704 1EAC CA841E      JZ CKDLF        ;JMP IF I/C MODE
2705 1EAF CDD41F      CSCDV: CALL SCDVA  ;SCHEDULE DIVERT
2706 1EB2 C1          POP B            ;ADJUST STACK
2707 1EB3 C9          RET              ;
2708
2709 ;CHECK TO SEE IF DIVERT LANE IS FULL
2710
2710 1EB4 3A9642      CKDLF: LDA JDV    ;A= DVT#
2711 1EB7 47          MOV B,A          ;B= DVT#
2712 1EB8 CD250D      CALL CPLFM       ;CHECK FOR LANE FULL
2713 1EBB A6          ANA M            ;IF DVT LANE FULL, CLEAR ZERO FLAG
2714 1EBC C2D11E      JNZ CKAMK       ;JMP IF LANE FULL
2715
2716 ;DIVERT LANE NOT FULL. TEST FOR DIVERT COMPLETE OPTION
2717
2717 1EBF C1            POP B            ;B.C= BYTE 13 LOC OF TCB
2718 1EC0 C5            PUSH B           ;STACK = SAME
2719 1EC1 3A3B3D      LDA DCO          ;DCO= 2?
2720 1EC4 D602         SUI 2            ;
2721 1EC6 C2AF1E      JNZ CSCDV       ;JMP IF DCO NOT= 2
2722 1EC9 CD1E20      CALL SCHDC       ;SCHEDULE DCB1 IN ORDER TO SEND MESSAGE
2723 ; S OR U
2724 1ECC C1            POP B            ;B.C= BYTE 13 LOC OF TCB
2725 1ECD C5            PUSH B           ;STACK = SAME
2726 1ECE C3AF1E      JMP CSCDV        ;JMP TO SCHEDULE DIVERT
2727 ;DIVERT LANE IS FULL. SEND MESSAGE A AND SEND PACKAGE TO LFDV
2728 ;IF MESSAGE A IS MASKED.
2729 1ED1 212A3D      CKAMK: LXI H,ERMKO+2 ;H.L= ERMKO BYTE 2
2730 1ED4 7E          MOV A,M          ;A= BYTE 2
2731 1ED5 E601         ANI 01H          ;LOOK AT BIT 0
2732 1ED7 CA3F1F      JZ SLFDV        ;JMP IF A NOT MASKED
2733 1EDA C1            POP B            ;B.C= BYTE 13 OF PACKAGE TCB
2734 1EDB C5            PUSH B           ;STACK=SAME
2735 1EDC 69          MOV L,C          ;
2736 1EDD 60          MOV H,B          ;H.L= BYTE 13 LOC OF TCB
2737 1EDE 56          MOV D,M          ;
2738 1EDF 2B          DCX H            ;H.L= BYTE 12 LOC
2739 1EE0 5E          MOV E,M          ;D,E= IDENT VALUE
2740 ;ACCESS PACKAGE CODE FROM TRANSACTION TABLE AND LOAD INTO BYTES 2-7 OF
2741 ;BUFFER BYT16 IF IDENT NOT=0. IF IDENT=0,LOAD 6 ASCII 0'S INTO BYT16 FOR
2742 ;PACK CODE. IN EITHER CASE, DON'T REMOVE FROM TRANSACTION TABLE AS
2743 ;PACKAGE IS STILL TRACKING.
2744 1EE1 7B          MOV A,E          ;
2745 1EE2 B2          ORA D            ;
2746 1EE3 CAEF1E      JZ APACK        ;JMP IF IDENT=0
2747 1EE6 EB          XCHG            ;H.L= IDENT= LOC OF PACK CODE
2748 1EE7 E5          PUSH H           ;STACK= SAME
2749 1EE8 CD6C1A      CALL TRANP      ;TRANSFER PACK CODE TO BYT16 BYTES 2-7
2750 1EEB E1          POP H            ;ADJUST STACK
2751 1EEC C3031F      JMP LOADH       ;JMP TO LOADH
2752 1EEF 3E30      APACK: MVI A,30H ;STORE 6 ASCII 0'S AS PACK CODE
2753 1EF1 321A43      STA BYT16+2     ;
2754 1EF4 321B43      STA BYT16+3     ;
2755 1EF7 321C43      STA BYT16+4     ;
2756 1EFA 321D43      STA BYT16+5     ;
2757 1EFD 321E43      STA BYT16+6     ;
2758 1F00 321F43      STA BYT16+7     ;
2759 ;CONVERT BINARY INDUCT AND DVT# IN PACKAGE TCB INTO ASCII AND STORE IN
2760 ;BUFFER BYT16.
2761
2762 1F03 E1          LOADH: POP H     ;H.L= BYTE 13 LOC IN PACK TCB

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2763 1F04 E5      PUSH H           ;STACK= SAME
2764 1F05 11FCFF  LXI D,-4        ;
2765 1F08 19      DAD D           ;H.L= JINDU LOC
2766 1F09 7E      MOV A,M         ;A= JINDU IN BINARY
2767 1F0A CD010E  CALL ECDIG      ;CONVERT TO ASCII
2768 1F0D 3A7E42  LDA XYZ1        ;
2769 1F10 322043  STA BYT16+8     ;WSTORE ASCII 10'S DIGIT
2770 1F13 3A7F42  LDA XYZ0        ;
2771 1F16 322143  STA BYT16+9     ;STORE ASCII 1'S DIGIT
2772 1F19 E1      POP H           ;H.L= TCB BYTE 13 LOC
2773 1F1A E5      PUSH H         ;
2774 1F1B 11FEFF  LXI D,-2        ;H.L= JDV LOC
2775 1F1E 19      DAD D           ;
2776 1F1F 7E      MOV A,M         ;A= JDV IN BINARY
2777 1F20 CD010E  CALL ECDIG      ;CONVERT TO ASCII
2778 1F23 3A7D42  LDA XYZ2        ;
2779 1F26 322243  STA BYT16+10    ;STORE 100'S DIGIT
2780 1F29 3A7E42  LDA XYZ1        ;
2781 1F2C 322343  STA BYT16+11    ;STORE 10'S DIGIT
2782 1F2F 3A7F42  LDA XYZ0        ;
2783 1F32 322443  STA BYT16+12    ;STORE 1'S DIGIT
2784 1F35 211843  LXI H,BYT16     ;
2785 1F38 23      INX H           ;H.L= BYT16 BYTE 1
2786 1F39 3E41    MVI A,'A'       ;
2787 1F3B 77      MOV M,A         ;STORE ASCII A
2788 1F3C CD7D1A  CALL DSLD       ;STORE MESSAGE A IN FS
2789
2790 ;SEND PACKAGE TO LFEDV
2791
2792 1F3F C1      SLFDV: POP B           ;B,C= BYTE 13 OF TCB
2793 1F40 69      MOV L,C         ;H,L= BYTE 13 LOC OF TCB
2794 1F41 60      MOV H,B         ;
2795 1F42 11FEFF  LXI D,-2        ;
2796 1F45 19      DAD D           ;H,L= JDV LOC (BYTE 11)
2797 1F46 3A3D3D  LDA LFEDV       ;
2798 1F49 77      MOV M,A         ;JDV IN TCB NOW=LFEDV
2799
2800 ;TRANSFER PACKAGE TO NEXT UD
2801
2802 1F4A C5      UBNXU: PUSH B           ;SAVE POINTER TO BYTE 13 IN TCB
2803 1F4B 2AF742  LHLD LDVP       ;POINT H,L TO LDV(JUPDA-1)
2804 1F4E 4E      MOV C,M         ;C= LDV(JUPDA-1)
2805 1F4F 0C      INR C           ;C= 1ST DV $ FOR THIS UD
2806 1F50 23      INX H           ;A
2807 1F51 7E      MOV A,M         ; = LDV(JUPDA)
2808 1F52 CD7727  CALL SDVRS      ;SCHEDULE RESETS OF ALL DV'S FOR THIS UD
2809 1F55 C1      POP B           ;POINT B,C TO BYTE 13 IN TCB
2810
2811 ;SCHEDULE UDB1 IF NECESSARY
2812
2813 1F56 3A9542  LDA JUPDA       ;IF
2814 1F59 21043D  LXI H,NUPDA     ; JUPDA
2815 1F5C BE      CMP M           ; <
2816 1F5D D26C1F  JNC USUD1       ; NUPDA
2817 1F60 CD892C  CALL GUDB1      ;CREATE TCB FOR UDB1 FOR NEXT UD
2818 1F63 4B      MOV C,E         ;C= OLD JUPDA
2819 1F64 EB      XCHG          ;POINT D,E TO BYTE 7 IN TCB
2820 1F65 2A3742  LHLD UDT        ;POINT H,L TO UDT TABLE
2821 1F68 CD2F29  CALL ENTCB      ;SCHEDULE UDB1 FOR NEXT UD
2822 1F6B C9      RET
2823
2824 ; PACKAGE HAS JUST PASSED LAST UPDATE PHOTO. IF IN I/C MODE. AND IDENT
2825 ; NOT=0. REMOVE PACK CODE FROM TRANSACTION TABLE.
2826
2827 1F6C 3A2B42  USUD1: LDA ICMO     ;TEST TO SEE IF ARE IN I/C MODE
2828 1F6F D600    SUI 0           ;
2829 1F71 CA8F1F  JZ USUDP        ;JMP IF NOT IN I/C MODE
2830
2831 ; ARE IN I/C MODE. TEST TO SEE IF PACK CODE IS IN TRANS TABLE
2832 1F74 60      MOV H,B         ;
2833 1F75 69      MOV L,C         ;H.L= TCB BYTE 13 LOC
2834 1F76 7E      MOV A,M         ;A= IDENT MSB
2835 1F77 D600    SUI 0           ;
2836 1F79 C2861F  JNZ USUDJ       ;JMP IF IDENT MSB NOT=0
2837 1F7C 2B      DCX H           ;H.L= IDENT LSB LOC
2838 1F7D 7E      MOV A,M         ;
2839 1F7E D600    SUI 0           ;
2840 1F80 C2861F  JNZ USUDJ       ;JMP IF IDENT LSB NOT = 0
2841 1F83 C38F1F  JMP USUDP        ;JMP IF IDENT = 0
2842
2843 ; PACKAGE HAS JUST PASSED LAST UPDATE PHOTO WITH IDENT NOT=0. THEREFORE
2844 ; REMOVE PACKAGE CODE FROM TRANSACTION TABLE
2845 1F86 60      USUDJ: MOV H,B         ;
2846 1F87 69      MOV L,C         ;H.L=IDENT MSB LOC
2847 1F88 56      MOV D,M         ;
2848 1F89 2B      DCX H           ;
2849 1F8A 5E      MOV E,M         ;D.E=IDENT

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2850 1F8B EB XCHG ;H.L=SAME
2851 1F8C CD401A CALL ZPACK ;REMOVE PACK CODE FROM TRANS TAB
2852
2853 1F8F 3A313D USUDP: LDA SIDOE ;IF SIDE-INDUCT RE-CIRCULATION
2854 1F92 E602 ANI 02H ; AUTO-REINDUCT
2855 1F94 C8 RZ ; DESIRED
2856 1F95 CD892C CALL GUDB1 ;CREATE TCB FOR UDB1, E=JUPDA
2857 1F98 4B MOV C,E ;C=JUPDA, SHOULD = NUPDA
2858 1F99 EB XCHG ;POINT D,E TO BYTE 7 IN TCB
2859 1F9A 210300 LXI H,3 ;POINT H.L TO BYTE 10
2860 1F9D 19 DAD D ; IN TCB
2861 1F9E 3601 MVI M,1 ;JUPDA IN TCB = 1
2862 1FA0 2A3742 LHL D ;POINT H.L TO UDT TABLE
2863 1FA3 CD2F29 CALL ENTCB ;SCHEDULE UDB1 FOR UD # 1
2864 1FA6 C9 RET
2865
2866 ;OLD TCB FOR UDB2 FOUND TO BE SCRIBBLED; OR
2867 ; UPDATE DATA SCRIBBLED (MAY BE DUE TO LOST UD-CLEARING INTERRUPT)
2868
2869 1FA7 3E0A UBPFIX: MVI A,10 ;FATAL ERROR
2870 1FA9 CD0C2C CALL FATAL ; 10
2871 1FAC E1 POP H ;RESET SP
2872 1FAD C9 RET
2873
2874 ; CALLED BY UDBPC TO SCHEDULE EVENT BMG FOR SIDE INDUCT.
2875
2876 ; INPUT: JUPDA= UPDATE #. TREE PARAMETERS.
2877 ; OUTPUT: EVENT BMG IN TREE.
2878 ; ALL REGISTERS EFFECTED.
2879
2880 1FAE 3A9542 SCBMG: LDA JUPDA ;PICK UP
2881 1FB1 4F MOV C,A ; UPDATE
2882 1FB2 0600 MVI B,0 ; PHOTO #
2883 1FB4 2A3F42 LHL D FIDDB ;GET NUMBER OF INDUCT
2884 1FB7 09 DAD B ; FOLLOWING
2885 1FB8 2B DCX H ; IF
2886 1FB9 4E MOV C,M ; ANY
2887 1FBA 0D DCR C ;DO NOTHING IF
2888 1FBB 0C INR C ; NO INDUCT
2889 1FBC C8 RZ ; FOLLOWS
2890
2891 1FBD CDCB0F CALL GETFS ;GET FREE BLOCK
2892 1FC0 110900 LXI D,9 ;POINT H.L TO BYTE 9
2893 1FC3 19 DAD D ; IN TCB
2894 1FC4 71 MOV M,C ;SET INDUCT # IN TCB
2895 1FC5 2B DCX H ;POINT H.L TO BYTE 8 IN TCB
2896 1FC6 111A24 LXI D,BMG ;POINT EVENT
2897 1FC9 72 MOV M,D ; IN TCB
2898 1FCA 2B DCX H ; TO
2899 1FCB 73 MOV M,E ; BMG
2900 1FCC EB XCHG ;POINT D,E TO BYTE 7 IN TCB
2901 1FCD 2A2F42 LHL D PRET ;POINT H.L TO PRET TABLE
2902 1FD0 CD2F29 CALL ENTCB ;SCHEDULE EVENT BMG
2903 1FD3 C9 RET
2904
2905 ; SCHEDULE DIVERT ACTUATION EVENTS AFTER
2906 ; PACKAGE HAS HIT ITS LAST UPDATE PHOTO.
2907 ; CALLED BY UDBPC.
2908
2909 ; INPUT: B,C POINTS TO BYTE 13 IN OLD TCB, TREE PARAMETERS.
2910 ; UDBT, JDV, LDVP POINTS TO LDV(JUPDA-1).
2911 ; ALL REGISTERS EFFECTED.
2912
2913 1FD4 2AF742 SCDVA: LHL LDVP ;POINT H.L TO LDV(JUPDA-1)
2914 1FD7 C5 PUSH B ;STACK= BYTE 13 OF OLD TCB
2915 1FD8 4E MOV C,M ;C= LDV(JUPDA-1)
2916 1FD9 0C INR C ;C= 1ST DV FOR THIS UD
2917 1FDA 3A9642 LDA JDV ;A
2918 1FDD 3D DCR A ; =JDV-1
2919 1FDE CD7727 CALL SDVRS ;SCHEDULE RESETS OF DV'S BEFORE # JDV
2920
2921 ; SCHEDULE ACTUATION OF DV # JDV
2922
2923 1FE1 CDCB0F CALL GETFS ;GET FREE STORAGE
2924 1FE4 D1 POP D ;D,E= BYTE 13 OF OLD TCB LOC
2925 1FE5 E5 PUSH H ;STACK= LOC OF NEW TCB
2926 1FE6 EB XCHG ;H.L= BYTE 13 OF OLD TCB LOC
2927 1FE7 7E MOV A,M ;A= IDENT MSB
2928 1FE8 329842 STA IDENT+1 ;STORE IN WORKING MEMORY
2929 1FEB 2B DCX H ;H.L= BYTE 12 OF OLD TCB LOC
2930 1FEC 7E MOV A,M ;A= IDENT LSB
2931 1FED 329742 STA IDENT ;STORE IN WORKING MEMORY
2932 1FF0 2B DCX H
2933 1FF1 2B DCX H
2934 1FF2 2B DCX H ;H.L= BYTE 9 OF OLD TCB
2935 1FF3 7E MOV A,M ;A= JINDU
2936 1FF4 329442 STA JINDU ;STORE IN WORKING MEMORY

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2937 1FF7 E1      POP H           ;H,L= LOC OF NEW TCB
2938 1FF8 110D00  LXI D,13       ;D,E= 13
2939 1FFB 19       DAD D          ;H,L= LOC OF BYTE 13 IN NEW TCB
2940 1FFC 3A9842  LDA IDENT+1    ;A= IDENT MSB
2941 1FFF 77       MOV M,A        ;SET IN NEW TCB
2942 2000 3A9742  LDA IDENT      ;A= IDENT LSB
2943 2003 2B       DCX H          ;H,L= BYTE12 LOC IN NEW TCB
2944 2004 77       MOV M,A        ;SET IN NEW TCB
2945 2005 2B       DCX H          ;H,L= BYTE 11 LOC IN NEW TCB
2946 2006 3A9442  LDA JINDU      ;A= JINDU
2947 2009 77       MOV M,A        ;SET IN JDV LOC FOR DVACT TCB
2948 200A 2B       DCX H          ;H,L= BYTE 10 LOC
2949 200B 3601     MVI M,1        ;SET DIVERT FLAG = 1 FOR SETTING DV
2950 200D 2B       DCX H          ;SET JDV
2951 200E 3A9642  LDA JDV        ; IN
2952 2011 77       MOV M,A        ; TCB
2953 2012 2B       DCX H          ;POINT H.L TO BYTE 8 IN TCB
2954 2013 11A526  LXI D,DVACT    ;LET
2955 2016 72       MOV M,D        ; EVENT
2956 2017 2B       DCX H          ; BE
2957 2018 73       MOV M,E        ; DVACT
2958 2019 4F       MOV C,A        ;C=JDV
2959 201A CD1D29  CALL ENDVA     ;SCHEDULE DVACT
2960 201D C9       RET
2961
2962 ; SUBROUTINE SCHDC
2963
2964 ; INPUTS: B.C POINTS TO BYTE 13 OF PACKAGE TCB.TREE PARAMETERS.
2965 ; UDBT,JDV,LDVP POINTS TO LDV(JUPDA-1)(INPUTS SAME AS
2966 ; FOR SCDVA).
2967
2968 ; OUTPUTS: EVENT DCB1 CREATED AND SCHEDULED. DCB1 IS LOWER BOUND
2969 ; ON MUZ FOR NEXT UPDATE PHOTO
2970
2971 ; PROCESS: CREATE EVENT DCB1 BY CALLING GUDB1 AND THEN CHANGING
2972 ; EVENT TO DCB1. SCHEDULE DCB1 BY CALLING ENTCB
2973
2974 201E CD892C    SCHDC: CALL GUDB1      ;CREATE TCB FOR EVENT UDB1. RETURN WITH
2975 ; H,L AT BYTE 7 OF NEW TCB.E=OLD JUPDA
2976 2021 01A525   LXI B,DCB1     ;B,C= DCB1 EVENT LOC
2977 2024 23       INX H          ;H,L= BYTE 8 OF NEW TCB
2978 2025 70       MOV M,B        ;MAKE EVENT IN NEW TCB= DCB1
2979 2026 2B       DCX H          ;H,L= BYTE 7 OF NEW TCB
2980 2027 71       MOV M,C        ;EVENT IN NEW TCB NOW = DCB1
2981 2028 4B       MOV C,E        ;C= OLD JUPDA
2982 2029 EB       XCHG         ;POINT D,E TO BYTE 7 OF NEW TCB=EVENT
2983 202A 2A3742   LHL UDT        ;POINT H.L TO UDT TABLE
2984 202D CD2F29   CALL ENTCB     ;SCHEDULE DCB1 FOR NEXT UD
2985 2030 C9       RET          ;RET TO UDBPC
2986
2987 ; UPDATE-PHOTO-CLEARING PROCESSOR.
2988
2989 ; INPUT: C=JUPDA (P
2990 ; UPDATE-PHOTO-CLEARING PROCESSOR.
2991
2992 ; INPUT: C=JUPDA (PHOTO #).
2993 ; ALL REGISTERS EFFECTED.
2994
2995 2031 0600     UDCPC: MVI B,0      ;B,C=JUPDA
2996 2033 2A5142   LHL UDDB      ;POINT H.L TO UDDB TABLE
2997 2036 09       DAD B          ;POINT
2998 2037 09       DAD B          ; H.L TO
2999 2038 09       DAD B          ; PHUZ IN UDDB
3000 2039 09       DAD B          ; FOR UD # JUPDA+1
3001 203A 2B       DCX H          ;POINT H.L TO PUD MS BYTE IN UDDB
3002 203B 56       MOV D,M        ;SET
3003 203C 70       MOV M,B        ; D,E
3004 203D 2B       DCX H          ; = PUD IN UDDB
3005 203E 5E       MOV E,M        ; POINTING TO TCB
3006 203F 70       MOV M,B        ; & ZERO PUD IN UDDB
3007 2040 7B       MOV A,E        ;IF
3008 2041 B2       ORA D          ; TCB
3009 2042 CAA420   JZ UCPC       ; NONNULL
3010
3011 ; D,E POINTS TO NONNULL PACKAGE TCB
3012
3013 2045 D5       PUSH D         ;SAVE TCB POINTER
3014 2046 C5       PUSH B         ;SAVE JUPDA
3015 2047 CDAC0F   CALL GUTRE     ;GET TREE PARAMETERS
3016 204A C1       POP B          ;B,C=JUPDA
3017 204B C5       PUSH B         ;SAVE UPDATE PHOTO #
3018 204C 3A313D   LDA SIDE0     ;IF SIDE INDUCT
3019 204F E610     ANI 010H     ; THEN
3020 2051 C4AA20   CNZ SCCMG     ; SCHEDULE CMG
3021
3022 2054 C1       POP B          ;RESTORE BC = UPDATE PHOTO #
3023 2055 2A3D42   LHL LDVDB     ;POINT H.L TO LDV TABLE
3024 2058 09       DAD B          ;POINT H.L TO LDV

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3025 2059 2B          DCX  H          ; FOR THIS UD
3026 205A 7E          MOV  A,M         ;A= LDV FOR THIS UD
3027 205B D1          POP  D          ;POINT D.E TO TCB
3028 205C 210B00      LXI  H,11        ;POINT H.L TO BYTE 11
3029 205F 19          DAD  D          ; IN TCB
3030 2060 BE          CMP  M          ;IF LDV FOR THIS UD
3031 2061 DABE20      JC   UCUSE      ; >= JDV
3032
3033 ;SCHEDULE PROGRAM-MEASURED DIVERT RESET IF DESIRED
3034
3035 2064 3A383D      LDA  DVTRS      ;IF
3036 2067 3D          DCR  A          ; MEASURED
3037 2068 C28E20      JNZ  UCUSE      ; RESET
3038
3039 206B D5           PUSH D          ;SAVE TCB POINTER
3040 206C 4E          MOV  C,M         ;C=JDV
3041 206D 2AB642      LHL  CTIME      ;D.E
3042 2070 EB          XCHG          ; = CTIME
3043 2071 2A0840      LHL  DVH        ;ADD CONSTANT
3044 2074 19          DAD  D          ; DVH
3045 2075 229942      SHLD UDBT      ;SET UDBT FOR ENDVA
3046 2078 CDCB0F      CALL GETFS     ;GET FREE BLOCK
3047 207B 110A00      LXI  D,10       ;POINT H.L TO BYTE 10
3048 207E 19          DAD  D          ; IN TCB
3049 207F 3602        MVI  M,2        ;SET DV MODE = 2 FOR RESETTING DV
3050 2081 2B          DCX  H          ;SET DIVERT #
3051 2082 71          MOV  M,C         ; IN TCB
3052 2083 2B          DCX  H          ;POINT H.L TO BYTE 8 IN TCB
3053 2084 11A526      LXI  D,DVACT    ;LET
3054 2087 72          MOV  M,D        ; EVENT
3055 2088 2B          DCX  H          ; BE
3056 2089 73          MOV  M,E        ; DVACT
3057 208A CD1D29      CALL ENDVA     ;SCHEDULE DV RESET
3058 208D D1          POP  D          ;POINT D.E TO TCB
3059
3060 ;TCB HAS TO BE RETURNED TO FREE STORAGE AFTER HAVING BEEN USED
3061 ; BY BOTH UDB3 AND UDCPC, OR UDCPC AND JUNK.
3062
3063 208E 210700      UCUSE: LXI  H,7   ;POINT H.L TO
3064 2091 19          DAD  D          ; BYTE 7 IN TCB
3065 2092 7E          MOV  A,M         ;IF TCB IS FOR UNIDENTIFIED PACKAGE
3066 2093 23          INX  H          ; FROM UDBPC, OR HAS
3067 2094 B6          ORA  M          ; NOT YET BEEN USED
3068 2095 CA9F20      JZ   UCRFS     ; BY UDB3
3069
3070 2098 118025      LXI  D,JUNK     ;CHANGE
3071 209B 72          MOV  M,D        ; EVENT
3072 209C 2B          DCX  H          ; TO
3073 209D 73          MOV  M,E        ; JUNK
3074 209E C9          RET
3075
3076 209F EB          UCRFS: XCHG          ;H.L POINTS TO TCB
3077 20A0 CD4F11      CALL PUTFS     ;RETURN TO FREE STORAGE
3078 20A3 C9          RET
3079
3080 ;SPURIOUS INTERRUPT
3081
3082 20A4 3E0D        UCPCE: MVI  A,13 ;ERROR 13
3083 20A6 CD080F      CALL ERPC     ; TO MASTER KB
3084 20A9 C9          RET
3085
3086 ;SCHEDULE EVENT CMG FOR SIDE INDUCT.
3087 ; CALLED BY UDCPC.
3088
3089 ;INPUT: B,C= UPDATE #, TREE PARAMETERS.
3090 ;OUTPUT: EVENT CMG IN TREE.
3091 ;ALL REGISTERS EFFECTED.
3092
3093 20AA 2A3F42      SCCMG: LHL  FIDDB ;GET NUMBER OF INDUCT
3094 20AD 09          DAD  B          ; FOLLOWING
3095 20AE 2B          DCX  H          ; IF
3096 20AF 4E          MOV  C,M         ; ANY
3097 20B0 0D          DCR  C          ;DO NOTHING IF
3098 20B1 0C          INR  C          ; NO INDUCT
3099 20B2 C8          RZ           ; FOLLOWS
3100
3101 20B3 CDCB0F      CALL GETFS     ;GET FREE BLOCK
3102 20B6 110900      LXI  D,9        ;POINT H.L TO BYTE 9
3103 20B9 19          DAD  D          ; IN TCB
3104 20BA 71          MOV  M,C         ;SET JINDU IN TCB
3105 20BB 2B          DCX  H          ;POINT H.L TO BYTE 8 IN TCB
3106 20BC 112924      LXI  D,CMG     ;POINT EVENT
3107 20BF 72          MOV  M,D        ; IN TCB
3108 20C0 2B          DCX  H          ; TO
3109 20C1 73          MOV  M,E        ; CMG
3110 20C2 E5          PUSH H         ;SAVE POINTER TO BYTE 7 IN TCB
3111 20C3 0B          DCX  B         ;B,C=JINDU-1

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3112 20C4 2A2F42      LHL D  PRET      ;POINT H.L TO
3113 20C7 09          DAD  B           ;  PRET
3114 20C8 09          DAD  B           ;  TIME
3115 20C9 5E          MOV  E,M        ;D,E
3116 20CA 23          INX  H           ;  =
3117 20CB 56          MOV  D,M        ;  PRET TIME
3118 20CC 2A3142     LHL D  TET      ;POINT H.L TO
3119 20CF 09          DAD  B           ;  TET
3120 20D0 09          DAD  B           ;  TIME
3121 20D1 4E          MOV  C,M        ;B,C
3122 20D2 23          INX  H           ;  =
3123 20D3 46          MOV  B,M        ;  TET TIME
3124 20D4 2AB642     LHL D  CTIME    ;H.L= EVTIM
3125 20D7 19          DAD  D           ;  =CTIME+ PRET TIME
3126 20D8 09          DAD  B           ;  + TET TIME
3127 20D9 EB          XCHG          ;D,E= EVTIM
3128 20DA E1          POP  H           ;POINT H.L TO BYTE 6
3129 20DB 2B          DCX  H           ;  IN TCB
3130 20DC 72          MOV  M,D        ;SET EVTIM
3131 20DD 2B          DCX  H           ;  IN
3132 20DE 73          MOV  M,E        ;  TCB
3133 20DF 11FBFF     LXI  D,0FFFH    ;POINT H.L TO
3134 20E2 19          DAD  D           ;  TCB
3135 20E3 CD4429     CALL ENQU2      ;SCHEDULE CMG
3136 20E6 C9          RET
3137
3138 ; ;
3139 ; LANE-FULL-PHOTO-BLOCKED PROCESSOR.
3140 ; CALLED BY PHPC OR RINI.
3141 ; ;
3142 ; INPUT: C= DIVERT #.
3143 ; ALL REGISTERS EXCEPT B EFFECTED.
3144 ; ;
3144 20E7 79          LFBPC: MOV  A,C           ;A = DIVERT #
3145 20E8 CD250D     CALL CPLFM        ;H.L POINTS TO BYTE LF, A= MASK
3146 20EB B6          ORA  M           ;SET LF BIT IN LFDB
3147 20EC 77          MOV  M,A         ;
3148 20ED C9          RET
3149
3150 ; ;
3151 ; LANE-FULL-PHOTO-CLEARED PROCESSOR.
3152 ; ;
3153 ; INPUT: C= DIVERT #.
3154 ; ALL REGISTERS EFFECTED.
3155 ; ;
3156 20EE 79          LFCPC: MOV  A,C           ;A = DIVERT #
3157 20EF 32FB42     STA  LFDV        ;SAVE DV #
3158 20F2 CD250D     CALL CPLFM        ;A= LF BYTE MASK, H.L POINTS TO
3159 ; ;
3160 20F5 2F          CMA             ; BYTE LF
3161 20F6 A6          ANA  M           ;RESET LF BIT IN LFDB
3162 20F7 77          MOV  M,A         ;
3163
3164 ; ;
3165 ; CONSIDER AUTO-INDUCTION AFTER LANE CLEARS
3166 ; ;
3166 20F8 111000     LXI  D,10H       ;D,E=E=10H
3167 20FB 3A3C3D     LDA  LFO         ;IF LANE-CLEAR-AUTO-REINDUCT
3168 20FE A3          ANA  E           ;  OPTION
3169 20FF C8          RZ             ;  SPECIFIED
3170
3171 ; ;
3171 2100 2A4742     LHL D  KBDB      ;POINT H.L TO BYTE 4
3172 2103 010400     LXI  B,4         ;  IN KBDB
3173 2106 09          DAD  B           ;  INITIALLY FOR KB # 1
3174 2107 0E01     MVI  C,1         ;C = INDUCT #, = 1 INITIALLY
3175 2109 7E          LFCAI: MOV  A,M        ;A= AUTO-REINDUCT FLAG
3176 210A 93          SUB  E           ;IF FLAG IS
3177 210B C23321     JNZ  LNXXB       ;  SET
3178
3179 ; ;
3180 ; ;CHECK DIVERT NUMBER
3181 ; ;
3181 210E 22DA42     SHLD KBDBP       ;SAVE POINTER FOR SNPC
3182 2111 23          INX  H           ;POINT H.L TO BYTE 6
3183 2112 23          INX  H           ;  IN KBDB
3184 2113 5E          MOV  E,M        ;PICK UP
3185 2114 23          INX  H           ;  PACKAGE IDENT
3186 2115 56          MOV  D,M        ;  FROM
3187 2116 2B          DCX  H           ;  KBDB
3188 2117 B2          ORA  D           ;MS BYTE NONZERO MEANS
3189 2118 CA1C21     JZ   LFNXL       ;  CODE HAS BEEN TRANSLATED
3190 211B EB          XCHG          ;IDENT POINTS TO DIVERT #
3191 211C 3AFB42     LFNXL: LDA  LFDV        ;IF THIS DV LANE IS THE ONE
3192 211F 96          SUB  M           ;  DESIRED THEN SET FLAG Z
3193 2120 C22D21     JNZ  LKBDP       ;IF THIS IS THE DV WANTED
3194
3195 ; ;
3196 ; ;AUTOMATIC INDUCTION
3197 ; ;
3197 2123 2ADA42     LHL D  KBDBP       ;POINT H.L TO BYTE 4 IN KBDB
3198 2126 77          MOV  M,A         ;RESET FLAG ER IN KBDB

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3199 2127 CDDF0C CALL CLDP ;CLEAR KB DISPLAY
3200 212A CDSB21 CALL SNPC ;AUTO-REINDUCT PACKAGE
3201 ;
3202 ; ;CONSIDER NEXT INDUCT
3203 ;
3204 212D 2ADA42 LKBDP: LHL D KBDBP ;POINT H,L TO BYTE 4 IN KBDB
3205 2130 111000 LXI D,10H ;SET D,E=E=10H
3206 2133 19 LNXXB: DAD D ;POINT TO NEXT INDUCT KBDB
3207 2134 0C INR C ;NEXT JINDU-1
3208 2135 3A033D LDA NINDU ;DO FOR
3209 2138 B9 CMP C ; ALL
3210 2139 D20921 JNC LFCAI ; INDUCTS
3211 213C C9 RET
3212 ;
3213 ;RUN MODE "SEND" KEY PROCESSOR IN NORMAL KEYBOARD MODE.
3214 ; ' ' TO ENTER OR REPEAT CODE FOR PACKAGE INDUCTION;
3215 ; '+' TO INSERT, CHANGE OR DELETE ENTRY IN TRANSLATE TABLE.
3216 ;
3217 ;INPUT: CHAR, C= KB #, KBDBP=H.L POINTS TO KBDB BYTE 4.
3218 ;ALL REGISTERS EXCEPT C EFFECTED. JDV, IDENT, XYZ4 TO XYZ0.
3219 ; STACK, XLTBP, CODE, DELTA, TCBP, TREE EFFECTED.
3220 ;
3221 213D 3ADE42 SNNPC: LDA CHAR ;IF CHARACTER ' ' THEN
3222 2140 FE3B CPI ' ' ; SET FLAG Z
3223 2142 23 INX H ;PICK UP
3224 2143 7E MOV A,M ; NDIG IN KBDB
3225 2144 C2532F JNZ PSNPC ;IF CHAR WAS INDEED ' '
3226 ;
3227 ;"SEND" PACKAGE
3228 ;
3229 2147 B7 ORA A ;IF NDIG IN KBDB
3230 2148 C25521 JNZ SENTR ; =0
3231 ;
3232 ; ;PREVIOUS CHARACTER ENTERED WAS NOT DIGIT
3233 ;
3234 214B 23 INX H ;IF
3235 214C 7E MOV A,M ; SND
3236 214D 23 INX H ; IN KBDB
3237 214E B6 ORA M ; NOT
3238 214F CA0D22 JZ SILCH ; 0
3239 2152 C35B21 JMP SNPC ;REPEAT PACKAGE CODE
3240 ;
3241 ; ;NEW DIGIT STRING HAS BEEN ENTERED
3242 ;
3243 2155 3A323D SENTR: LDA SENDO ;IF
3244 2158 E610 ANI 10H ; ENTER-BY-"SEND"
3245 215A C8 RZ ; ENABLED
3246 ;
3247 ;
3248 ;AUTOMATIC "SEND" BY FIXED-LENGTH CODE, OR "REPEAT",
3249 ; OR CALLED BY LFCPC FOR LANE-CLEAR-AUTO-REINDUCT,
3250 ; OR CALLED BY RLTPC FOR DIRECT PACKAGE CODING BY DIVERT #.
3251 ;
3252 ;INPUT: C= KB #, KBDBP POINTS TO BYTE 4 IN KBDB.
3253 ; FLAG XL ACTIVATES CODE TRANSLATION.
3254 ;ALL REGISTERS EXCEPT C EFFECTED, KBDBP, JDV, IDENT EFFECTED.
3255 ;
3256 215B 3A033D SNNPC: LDA NINDU ;IF NINDU - KB # >= 0,
3257 215E B9 CMP C ; I.E., IT IS AN
3258 215F 3E00 MVI A,0 ; (FOR ERRPC)
3259 2161 DA0D22 JC SILCH ; INDUCTION KB
3260 ;
3261 ;DECODE DIVERT # OR PACKAGE CODE ENTERED
3262 ;
3263 2164 2ADA42 LHL D KBDBP ;POINT TO
3264 2167 23 INX H ; BYTE 6
3265 2168 23 INX H ; IN KBDB
3266 2169 3AE042 LDA XL ;IF TRANSLATION
3267 216C B7 ORA A ; IS
3268 216D CA8521 JZ SNNXL ; NECESSARY
3269 ;
3270 ; ;TRANSLATE PACKAGE CODE; DEFAULT IF NECESSARY
3271 ;
3272 2170 CD9732 CALL XLATE ;TRANSLATE CODE
3273 2173 DA7D21 JC SXLDF ;IF CODE IS
3274 2176 3A9642 LDA JDV ; OUT OF BOUND
3275 2179 B7 ORA A ; OR NOT IN
3276 217A C29F21 JNZ SNPC1 ; TRANSLATE TABLE
3277 217D 3A393D SXLDF: LDA SCNRO ;TEST DEFAULT
3278 2180 E602 ANI 02H ; OPTION
3279 2182 C39021 JMP SDFDV ;DEFAULT IF NECESSARY
3280 ;
3281 ; ;DECODE DIVERT NUMBER; DEFAULT IF NECESSARY
3282 ;
3283 2185 CD810D SNNXL: CALL DCDV ;DECODE DIVERT #
3284 2188 D29F21 JNC SNPC1 ;IF NOT LEGAL
3285 218B 3A393D LDA SCNRO ;IF DEFAULT

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3286 218E E601 ANI 01H ; OPTION
3287 2190 CA0E22 SDFDV: JZ SILDV ; SPECIFIED
3288 2193 3A3A3D LDA SERDV ;DEFAULT DIVERT $ = SERDV
3289 2196 6F MOV L,A ; AND PACKAGE
3290 2197 2600 MVI H,0 ; IDENTIFICATION = SERDV
3291 2199 329642 STA JDV ;SET DIVERT $ AND
3292 219C 229742 SHLD IDENT ; PACKAGE IDENTIFICATION
3293 ;
3294 ;SET KEYBOARD DATA IN KBDB
3295 ;
3296 219F 2ADA42 SNPC1: LHL D KBDBP ;CLEAR NDIG IN KBDB
3297 21A2 23 INX H ; TO DENOTE
3298 21A3 3600 MVI M,0 ; NON-DIGIT CHAR
3299 21A5 23 INX H ;POINT TO BYTE 6 IN KBDB
3300 21A6 3A9742 LDA IDENT ;PUT IDENT INTO
3301 21A9 77 MOV M,A ; SECONDARY
3302 21AA 23 INX H ; BUFFER
3303 21AB 3A9842 LDA IDENT+1 ; SND IN
3304 21AE 77 MOV M,A ; KBDB
3305 ;
3306 ;B= LEGAL DV $, CHECK LANE FULL
3307 ;
3308 21AF 3A3C3D SNPC2: LDA LFO ;IF LANE-FULL
3309 21B2 E601 ANI 01H ; KB-INHIBIT
3310 21B4 CAC121 JZ SPLAQ ; IS DESIRED
3311 21B7 78 MOV A,B ;A= JDV
3312 21B8 C5 PUSH B ;SAVE JDV,JINDU
3313 21B9 CD250D CALL CPLFM ;CP LF MASK IN A, H.L MASK POINTER
3314 21BC C1 POP B ;B = DIVERT $, C = INDUCT $
3315 21BD A6 ANA M ;IF LANE
3316 21BE C2FB21 JNZ SDVLF ; NOT FULL
3317 ;
3318 ;PUT CODE INTO LAQ & ATTEMPT TO SEND PACKAGE
3319 ;
3320 ; IF IN I/C AND HALF-DAS MODE, SET IDENT=0 BEFORE LOADING IT INTO LAQ
3321 ;
3322 21C1 3A2B42 SPLAQ: LDA ICMO ;TEST FOR I/C MODE
3323 21C4 D600 SUI 0 ;
3324 21C6 CAD921 JZ CSPLA ;JMP IF NOT I/C MODE
3325 21C9 3A0640 LDA HDSEN ;TEST FOR HALF-DAS MODE
3326 21CC D600 SUI 0 ;
3327 21CE CAD921 JZ CSPLA ;JMP IF NOT HALF-DAS
3328 ; ARE IN HALF-DAS MODE
3329 21D1 3E00 MVI A,0 ;ZERO IDENT BEFORE LOADING INTO LAQ
3330 21D3 329742 STA IDENT ;
3331 21D6 329842 STA IDENT+1 ;
3332 21D9 CDF02C CSPLA: CALL PLAQ ;PUT JDV,IDENT INTO LAQ,POINT PPDBP
3333 21DC CA0B22 JZ SLAQF ; TO BYTE 1 IN PPDB
3334 21DF 2AEB42 LHL D PPDBP ;POINT H.L TO BYTE 1 IN PPDB
3335 21E2 7E MOV A,M ;A= PWTNG IN PPDB
3336 21E3 3D DCR A ;IF PACKAGE AT PP
3337 21E4 C2ED21 JNZ SBEEP ; WAITING TO BE CODED
3338 ;
3339 ;SEND OUT WAITING PACKAGE IF NO MERGE CONFLICT
3340 ;
3341 21E7 C5 PUSH B ;SAVE INDUCT $
3342 21E8 23 INX H ;POINT H.L TO BYTE 2 IN PPDB
3343 21E9 CD7022 CALL PSEQR ;QUEUE PACKAGE UP FOR RELEASE
3344 ;
3345 21EC C1 POP B ;C= INDUCT $
3346 21ED 063B SBEEP: MVI B,';' ;BEEP
3347 21EF CD6C11 CALL PUTKB ; KB
3348 ;
3349 ;CONSIDER REPEAT-BY-"SEND" OPTION
3350 ;
3351 21F2 3A323D LDA SENDO ;IF REPEAT NOT DESIRED
3352 21F5 E620 ANI 20H ; THEN CLEAR
3353 21F7 CCC90C CZ CLRKB ; KB & KBDB
3354 21FA C9 RET
3355 ;
3356 ;DV LANE FULL
3357 ;
3358 21FB 3E07 SDVLF: MVI A,7 ;ERROR
3359 21FD CD130F CALL ERRPC ; 7
3360 2200 3A3C3D LDA LFO ;IF AUTO-REINDUCT-
3361 2203 E610 ANI 10H ; AFTER-LANE-CLEAR$
3362 2205 C8 RZ ; DESIRED
3363 2206 2ADA42 LHL D KBDBP ;POINT H.L TO BYTE 4 IN KBDB
3364 2209 77 MOV M,A ;SET AUTO-REINDUCT FLAG IN KBDB
3365 220A C9 RET
3366 ;
3367 ;ERROR HANDLING.
3368 ;
3369 220B 3D SLAQF: DCR A ;ERROR 06H MEANS LAQ FULL
3370 220C 3D DCR A ;ERROR 07H MEANS DV LANE FULL
3371 220D 3D DCR A ;ERROR 08H MEANS ILLEGAL COMMAND
3372 220E C609 SILDM: ADI 09H ;ERROR 09H MEANS ILLEGAL DV $

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3373 2210 CD130F CALL ERRPC ;PROCESS RUN ERROR
3374 2213 C9 RET
3375
3376 ;MASTER-CLEAR PROCESSOR FOR KEYBOARD # C.
3377 ; CLEARS KBDB, KB, PURGE LAQ IN PPDB.
3378 ; AND IF RELEASE IS OFF THEN REINITIALIZE INDUCT.
3379 ; CALLED BY KBPC.
3380
3381 ;INPUT: C= KEYBOARD #.
3382 ;ALL REGISTERS EXCEPT C, AND KBDB, PPDB EFFECTED.
3383
3384 2214 CDC90C MCLPC: CALL CLRKB ;CLEAR KBDB, KB
3385 2217 3A033D LDA NINDU ;IF IT IS A NON-INDUCT KB
3386 221A B9 CMP C ; THEN WE
3387 221B D8 RC ; ARE DONE
3388
3389 ;CLEAR INDUCT LAQ
3390
3391 221C 79 MOV A,C ;A = INDUCT #
3392 221D CD9927 CALL CPPPP ;POINT H.L TO PPDB(JINDU)
3393 2220 E5 PUSH H ;SAVE POINTER
3394 2221 110600 LXI D,6 ;POINT H.L TO BYTE 6
3395 2224 19 DAD D ; IN PPDB
3396 2225 3600 MVI M,0 ;PURGE LAQ
3397
3398 ;IF RELEASE IS OFF THEN RE-INITIALIZE INDUCT
3399
3400 2227 C5 PUSH B ;SAVE C = INDUCT #
3401 2228 79 MOV A,C ;A = INDUCT #
3402 2229 2A5942 LHLD OUTDB ;POINT H.L TO OUTIM
3403 222C CD280D CALL CPHSK ;CP A= MASK POINTED TO BY H.L
3404 222F A6 ANA M ;IF RELEASE OFF THEN SET FLAG Z
3405 2230 C1 POP B ;C=JINDU= INDUCT #
3406 2231 E1 POP H ;POINT H.L TO PPDB(JINDU)
3407 2232 C0 RNZ ;IF RELEASE OFF
3408
3409 INX H ;RETURN IF
3410 2234 23 INX H ; RELEASE
3411 2235 23 INX H ; IS
3412 2236 B6 ORA M ; BEING
3413 2237 C0 RNZ ; DELAYED
3414
3415 ;CLEAN UP PPDB MERGE DATA
3416
3417 2238 2B DCX H ;POINT TO
3418 2239 2B DCX H ; BYTE 0
3419 223A 2B DCX H ; IN PPDB
3420 223B 7E MOV A,M ;A = FLAG PP IN PPDB
3421 223C 23 INX H ;POINT H.L TO BYTE 1 IN PPDB
3422 223D 77 MOV M,A ;PWTNG IN PPDB = PP IN PPDB
3423 223E B7 ORA A ;IF PP IN PPDB
3424 223F C24B22 JNZ MICNT ; = 0
3425 2242 E5 PUSH H ;SAVE POINTER
3426 2243 C5 PUSH B ;SAVE C= INDUCT #
3427 2244 0601 MVI B,1 ;SET RELEASE
3428 2246 CD0F11 CALL OPUT ; ON
3429 2249 C1 POP B ;C= INDUCT #
3430 224A E1 POP H ;POINT H.L TO BYTE 1 IN PPDB
3431
3432 ; ;IF STRAIGHT INDUCTS THEN ADJUST COUNTS IN PPDB
3433
3434 224B 3A313D MICNT: LDA SIDE0 ;IF
3435 224E E610 ANI 010H ; STRAIGHT
3436 2250 C0 RNZ ; INDUCTION
3437
3438 2251 C5 PUSH B ;SAVE C=JINDU= INDUCT #
3439 2252 23 INX H ;B
3440 2253 46 MOV B,M ; = COUNT(JINDU) IN PPDB
3441 2254 2A4342 LHLD PPDB ;POINT H.L TO BYTE 2
3442 2257 23 INX H ; IN PPDB FOR
3443 2258 23 INX H ; INDUCT # 1
3444 2259 112000 LXI D,32 ;SET POINTER INCREMENT
3445 225C 3A033D LDA NINDU ;C IS THE LOOP
3446 225F 4F MOV C,A ; COUNT
3447 2260 78 MOV A,B ;A= COUNT(JINDU)
3448 2261 BE M9CNT: CMP M ;IF COUNT(C)
3449 2262 D26622 JNC MNXCN ; > COUNT(JINDU)
3450 2265 35 DCR M ;DEC COUNT(C)
3451 2266 19 MNXCN: DAD D ;POINT H.L TO NEXT COUNT
3452 2267 0D DCR C ;DO FOR ALL
3453 2268 C26122 JNZ M9CNT ; INDUCTS
3454
3455 226B C1 POP B ;C= INDUCT #
3456 226C CDAC2C CALL PEND ;ADJUST COUNT(JINDU)
3457 226F C9 RET
3458
3459 ;PACKAGE RELEASE SEQUENCER. MERGE CODED PACKAGES ON A 1ST-COME-

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3460 ; 1ST-SERVE BASIS; OR SIDE-INDUCT INTO GAP ON CONVEYOR.
3461 ; CALLED BY PPBPC OR SNDPC.
3462 ;
3463 ;INPUT: C= INDUCT $, PPDB, H.L POINTS TO-BYTE 2 IN PPDB.
3464 ;OUTPUT: PPDB.
3465 ;ALL REGISTERS EFFECTED.
3466 ;
3467 2270 7E PSEQR: MOV A,M ;A= COUNT IN PPDB
3468 2271 2B DCX H ;POINT H.L TO BYTE 1 IN PPDB
3469 2272 B7 ORA A ;IF COUNT IN PPDB
3470 2273 C29E22 JNZ PSTOP ; =0
3471 ;
3472 ;SPACE AVAILABLE FOR PACKAGE TO BE RELEASED
3473 ;
3474 2276 77 MOV M,A ;PWTNG IN PPDB = 0
3475 2277 23 INX H ;POINT H.L TO BYTE 2 IN PPDB
3476 2278 3A313D LDA SIDO ;IF IT IS
3477 227B E610 ANI 010H ; A STRAIGHT
3478 227D C29922 JNZ PSIDE ; MERGE
3479 ;
3480 ; ;STRAIGHT-MERGE INDUCT
3481 ;
3482 2280 E5 PUSH H ;SAVE POINTER TO BYTE 2 IN PPDB
3483 2281 2A4342 LHLD PPDB ;POINT H.L TO
3484 2284 23 INX H ; 1ST COUNT
3485 2285 23 INX H ; IN PPDB
3486 2286 3A033D LDA NINDU ;A= LOOP COUNT
3487 2289 112000 LXI D,32 ;D,E= 32
3488 228C 0601 MVI B,1 ;B=1
3489 228E 70 PICNT: MOV M,B ;EVERY COUNT IN PPDB = 1
3490 228F 19 DAD D ;POINT H.L TO NEXT COUNT IN PPDB
3491 2290 3D DCR A ;DO FOR ALL
3492 2291 C28E22 JNZ PICNT ; INDUCTS
3493 2294 E1 POP H ;POINT H.L TO BYTE 2 IN PPDB
3494 2295 CDCF22 CALL RLSPC ;PHYSICALLY RELEASE PACKAGE
3495 2298 C9 RET
3496 ;
3497 ; ;SIDE-MERGE INDUCT
3498 ;
3499 2299 34 PSIDE: INR M ;INC COUNT(JINDU) IN PPDB
3500 229A CDCF22 CALL RLSPC ;PHYSICALLY RELEASE PACKAGE
3501 229D C9 RET
3502 ;
3503 ;SPACE NOT AVAILABLE ON MAINLINE FOR PACKAGE RELEASE
3504 ;
3505 229E 3602 PSTOP: MVI M,2 ;PWTNG IN PPDB = 2
3506 22A0 5F MOV E,A ;E= COUNT(JINDU) IN PPDB
3507 22A1 3A313D LDA SIDO ;IF
3508 22A4 E610 ANI 010H ; STRAIGHT
3509 22A6 C2C922 JNZ PRSRL ; MERGE
3510 ;
3511 22A9 0601 MVI B,1 ;B= INDUCT $, = 1 INITIALLY
3512 22AB 2A4342 LHLD PPDB ;POINT H.L TO
3513 22AE 23 INX H ; 1ST COUNT
3514 22AF 23 INX H ; IN PPDB
3515 22B0 78 PICCN: MOV A,B ;A= INDUCT $
3516 22B1 B9 CMP C ;IF INDUCT $
3517 22B2 CABB22 JZ PID01 ; NOT= JINDU
3518 22B5 7E MOV A,M ;A= COUNT IN PPDB
3519 22B6 BB CMP E ;IF THIS COUNT
3520 22B7 DABB22 JC PID01 ; >= COUNT(JINDU)
3521 22BA 34 INR M ;INC COUNT IN PPDB
3522 22BB D5 PID01: PUSH D ;POINT H.L TO
3523 22BC 112000 LXI D,32 ; NEXT
3524 22BF 19 DAD D ; COUNT
3525 22C0 D1 POP D ; IN PPDB
3526 22C1 04 INR B ;INC INDUCT $
3527 22C2 3A033D LDA NINDU ;DO FOR
3528 22C5 B8 CMP B ; ALL
3529 22C6 D2B022 JNC PICCN ; INDUCTS
3530 ;
3531 22C9 0602 PRSRL: MVI B,2 ;RESET RELEASE
3532 22CB CD0F11 CALL OPUT ; $ JINDU
3533 22CE C9 RET
3534 ;
3535 ; ;PACKAGE RELEASE PROCESSOR.
3536 ; CALLED BY PSEQR, TECHG OR CMG.
3537 ;
3538 ;INPUT: C= INDUCT $, H.L POINTS TO BYTE 2 IN PPDB.
3539 ;OUTPUT: EVENT TEC2, EVENT IDB2 OR UDB1.
3540 ;ALL REGISTERS EFFECTED, JINDU, PPDBP, TCBP, RLDBP EFFECTED.
3541 ;
3542 22CF 22EB42 RLSPC: SHLD PPDBP ;SAVE POINTER
3543 22D2 CD622C CALL GLAQ ;GET B=JDV, D,E=IDENT FROM LAQ
3544 22D5 CABB23 JZ RLSFX ; IN PPDB
3545 22D8 D5 PUSH D ;SAVE IDENT
3546 22D9 C5 PUSH B ;SAVE INDUCT $, DIVERT $

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3547
3548
3549
3550 22DA CDA20F      CALL GITRE      ;GET TREE PARAMETERS FOR INDUCT
3551 22DD 2AE842      LHL D PPDBP    ;POINT TO
3552 22E0 23          INX H          ; BYTE 4
3553 22E1 23          INX H          ; IN PPDB
3554 22E2 7E          MOV A,M        ;TEST IF PP-DWELL HAS
3555 22E3 23          INX H          ; TIMED OUT
3556 22E4 B6          ORA M          ; YET
3557 22E5 C1          POP B          ;B = DIVERT #, C = INDUCT #
3558 22E6 CAFD22      JZ RTCBS      ;IF NOT YET TIMED OUT
3559 22E9 2B          DCX H          ;SET RELEASE DELAY
3560 22EA 2B          DCX H          ; FLAG FOR
3561 22EB 3601        MVI M,1        ; MCLPC
3562 22ED C5          PUSH B         ;SAVE B,C
3563 22EE 0602        MVI B,2        ;RELEASE
3564 22F0 CD0F11      CALL OPUT      ; OFF
3565 22F3 2A0C40      LHL D GAP      ;DE = PPI GAP
3566 22F6 EB          XCHG          ; TO DELAY
3567 22F7 C1          POP B          ;MUST BE CAREFUL NOT TO
3568 22F8 E1          POP H          ; LEAVE DATA IN STACK
3569 22F9 CD672D      CALL RDLPP     ;DELAY PPI PULSES
3570 22FC E5          PUSH H         ;RE-SAVE IDENT
3571 22FD 79          RTCBS: MOV A,C    ;SAVE INDUCT #
3572 22FE 329442      STA JINDU     ; IN MEMORY
3573 2301 CD9927      CALL CPPPP    ;POINT TO
3574 2304 23          INX H          ; RELEASE DELAY
3575 2305 23          INX H          ; FLAG IN
3576 2306 23          INX H          ; PPDB
3577 2307 3600        MVI M,0        ;CLEAR FLAG
3578
3579
3580
3581 2309 CDCB0F      CALL GETFS     ;GET FREE BLOCK FOR TCB
3582 230C 22B442      SHLD TCBP     ;SET TCBP FOR APEND
3583 230F 110D00      LXI D,13      ;POINT H.L TO BYTE 13
3584 2312 19          DAD D         ; IN TCB
3585 2313 D1          POP D         ;D,E=IDENT
3586 2314 72          MOV M,D       ;SET IDENT
3587 2315 2B          DCX H         ; IN
3588 2316 73          MOV M,E       ; TCB
3589 2317 2B          DCX H         ;POINT H.L TO BYTE 11 IN TCB
3590 2318 70          MOV M,B       ;SET JDV IN TCB
3591 2319 2B          DCX H         ;JUPDA IN TCB
3592 231A 3600        MVI M,0        ; = 0
3593 231C 2B          DCX H         ;SET JINDU
3594 231D 71          MOV M,C       ; IN TCB
3595 231E 2B          DCX H         ;POINT H.L TO BYTE 8 IN TCB
3596
3597
3598
3599
3600 231F E5          PUSH H         ;SAVE POINTER
3601 2320 C5          PUSH B         ;SAVE C=JINDU
3602 2321 CDA227      CALL CPRLP    ;POINT H.L TO RLSDB
3603 2324 22EF42      SHLD RLDBP    ;SAVE POINTER
3604 2327 23          INX H         ;POINT H.L TO BYTE 2
3605 2328 23          INX H         ; IN RLSDB
3606 2329 3A313D      LDA SIDE0     ;IF
3607 232C E610        ANI 010H     ; STRAIGHT
3608 232E C24623      JNZ RSIDE     ; INDUCT
3609 2331 CD850C      CALL APEND    ;APPEND TCB TO END OF MERGE FIFO
3610
3611
3612
3613 2334 C1          POP B         ;C=JINDU
3614 2335 E1          POP H         ;POINT H.L TO BYTE 8 IN TCB
3615 2336 115824      LXI D,1DB2   ;EVENT
3616 2339 72          MOV M,D       ; IN TCB
3617 233A 2B          DCX H         ; =
3618 233B 73          MOV M,E       ; 1DB2
3619 233C EB          XCHG          ;POINT DE TO BYTE 7 IN TCB
3620 233D 2A2F42      LHL D PRET    ;POINT HL TO PRET TABLE
3621 2340 CD2F29      CALL ENTCB    ;SCHEDULE EVENT 1DB2
3622 2343 C37423      JMP RTEC2
3623
3624
3625
3626 2346 C1          RLSID: POP B     ;C=JINDU
3627 2347 2A3F42      LHL D FIDDB  ;POINT TO FOLLOWING-INDUCT DATABASE
3628 234A 3A043D      LDA NUPDA     ;B IS THE LOOP
3629 234D 47          MOV B,A       ; DOWN COUNT
3630 234E 79          MOV A,C       ;A = INDUCT #
3631 234F 05          RLSID: DCR B   ;SEARCH
3632 2350 FA5F23      JM RNFID     ; FOR
3633 2353 BE          CMP M         ; INDUCT #
3634 2354 23          INX H         ; IN

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3635 2355 C24F23      JNZ  RLSID      ; FIDDB
3636 2358 3A043D      LDA  NUPDA      ;A
3637 235B 90          SUB  B          ; =
3638 235C C36023      JMP  RLUD1      ; NEXT
3639 235F AF          RNFID: XRA  A    ; UPDATE
3640 2360 3C          RLUD1: INR  A    ; PHOTO *
3641 2361 E1          POP  H          ;POINT H.L TO BYTE 8 IN TCB
3642 2362 23          INX  H          ;SET UPCOMING
3643 2363 23          INX  H          ; UPDATE
3644 2364 77          MOV  M,A        ; PHOTO
3645 2365 2B          DCX  H          ; IN
3646 2366 2B          DCX  H          ; TCB
3647 2367 119424      LXI  D,UDB1     ;POINT EVENT
3648 236A 72          MOV  M,D        ; IN TCB
3649 236B 2B          DCX  H          ; TO
3650 236C 73          MOV  M,E        ; UDB1
3651 236D EB          XCHG          ;POINT D,E TO BYTE 7 IN TCB
3652 236E 2A3342      LHL  IDT        ;POINT H.L TO IDT TABLE
3653 2371 CD2F29      CALL ENTGB      ;SCHEDULE EVENT UDB1
3654
3655                ;SCHEDULE TEC2 AND SET TCBP IN RLSDB
3656
3657 2374 2A0440      RTEC2: LHL  PLEMX ;D,E = TE TIMEOUT
3658 2377 EB          XCHG          ; PPI COUNT
3659 2378 2AEF42      LHL  RLDBP     ;POINT TO TIMEOUT POINTER
3660 237B 01A223      LXI  B,TEC2    ;SCHEDULE TE
3661 237E CD192D      CALL RSCHD     ; TIMEOUT
3662 2381 3A9442      LDA  JINDU     ;
3663 2384 4F          MOV  C,A        ; =JINDU
3664 2385 0601      MVI  B,1       ;RELEASE
3665 2387 CD0F11      CALL OPUT      ; ON
3666 238A C9          RET
3667
3668                ;EITHER LAQ EMPTY
3669
3670 238B 3E03      RLSFX: MVI  A,3 ;FATAL ERROR
3671 238D CD0C2C      CALL FATAL    ; 3
3672 2390 C9          RET
3673
3674                ;;
3675                ;PACKAGE-PRESENT DWELL TIMEOUT FOR ADDITIONAL DELAY
3676                ; BEFORE RELEASING PACKAGE.
3677
3678                ;INPUT: DE POINTS TO BYTE 8 IN TCB. TCBP POINTS TO TCB.
3679                ;ALL REGISTERS EXCEPT B, C EFFECTED.
3680
3680 2391 EB          PPDTO: XCHG          ;POINT TO
3681 2392 23          INX  H          ; BYTE 10
3682 2393 23          INX  H          ; IN TCB
3683 2394 5E          MOV  E,M        ;PICK UP
3684 2395 23          INX  H          ; BACK
3685 2396 56          MOV  D,M        ; POINTER
3686 2397 AF          XRA  A          ;CLEAR
3687 2398 12          STAX D          ; TIME
3688 2399 13          INX  D          ; OUT
3689 239A 12          STAX D          ; POINTER
3690 239B 2AB442      LHL  TCBP     ;RETURN BLOCK TO
3691 239E CD4F11      CALL PUTFS    ; FREE STORAGE
3692 23A1 C9          RET
3693
3694                ;;
3695                ;TRAILING-EDGE-PHOTO-TIME-OUT EVENT.
3696                ; CREATED BY RLSPC.
3697
3698                ;INPUT: TCBP POINTS TO TCB, D,E POINTS TO BYTE 8 IN TCB.
3699                ;ALL REGISTERS EFFECTED.
3700
3700 23A2 D5          TEC2:  PUSH  D          ;SAVE D,E
3701 23A3 EB          XCHG          ;POINT TO BYTE 9
3702 23A4 23          INX  H          ; IN TCB
3703 23A5 4E          MOV  C,M        ;PICK UP INDUCT *
3704 23A6 23          INX  H          ;POINT TO BYTE 10
3705 23A7 5E          MOV  E,M        ;PICK UP
3706 23A8 23          INX  H          ; BACK
3707 23A9 56          MOV  D,M        ; POINTER
3708 23AA AF          XRA  A          ;CLEAR
3709 23AB 12          STAX D          ; TE
3710 23AC 13          INX  D          ; TIMEOUT
3711 23AD 12          STAX D          ; POINTER
3712 23AE 3E0B      MVI  A,11      ;ERROR 11 TO
3713 23B0 CD130F      CALL ERRPC    ; INDUCT * JINDU
3714
3715                ;YIELD TO ANY NEXT PACKAGE TO BE INDUCTED
3716
3717 23B3 00          NOP
3718 23B4 0602      MVI  B,2       ;STOP RELEASE
3719 23B6 CD0F11      CALL OPUT      ; * JINDU
3720 23B9 D1          POP  D          ;POINT D,E TO BYTE 8 IN TCB
3721 23BA 3A313D      LDA  SIDE0     ;IF
3722 23BD E610      ANI  10H      ; STRAIGHT

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3723 23BF C2C623      JNZ  TE2SI      ; MERGE
3724 23C2 CDCA23      CALL  TECHG     ; CLEAR MERGE AREA & POSSIBLY RELEASE
3725 23C5 C9          RET
3726
3727 23C6 CD2924      TE2SI: CALL  CMG      ; CLEAR MERGE AREA & POSSIBLY RELEASE
3728 23C9 C9          RET
3729
3730 ; CLEAR-MERGE EVENT SCHEDULED BY TECPC FOR STRAIGHT INDUCT.
3731 ; OR CALLED BY TEC2.
3732 ; RELEASE NEXT PENDING PACKAGE IF ANY.
3733
3734 ; INPUT: TCBP, D,E POINTS TO BYTE 8 IN TCB.
3735 ; ALL REGISTERS EFFECTED, TCBP, PPDB AND RLSDB EFFECTED.
3736
3737 23CA 3A033D      TECMG: LDA  NINDU   ; B= INDUCT #
3738 23CD 47          MOV  B,A        ; = LOOP COUNT
3739 23CE 2A4342      LHLD PPDB      ; POINT H.L TO BYTE 2
3740 23D1 23          INX  H         ; IN PPDB FOR
3741 23D2 23          INX  H         ; INDUCT # 1
3742 23D3 112000     LXI  D,32      ; SET POINTER INCREMENT
3743 23D6 0E00       MVI  C,0      ; C FLAGS ANY PENDING PACKAGE TO RELEASE
3744 23D8 3E02       MVI  A,2      ; A=2
3745
3746 23DA 35          T9CNT: DCR  M        ; DEC COUNT(B)
3747 23DB FA0E24     JM   TCMFX    ; IF COUNT < 0 THEN FATAL ERROR
3748 23DE C2ED23     JNZ  TNXID    ; IF COUNT(B) = 0
3749 23E1 2B          DCX  H        ; IF
3750 23E2 BE         CMP  M        ; PWTNG(B)
3751 23E3 23          INX  H        ; =
3752 23E4 C2ED23     JNZ  TNXID    ; 2
3753
3754 23E7 0C          INR  C        ; IF C ALREADY SET THEN
3755 23E8 0D         DCR  C        ; FATAL ERROR
3756 23E9 C20E24     JNZ  TCMFX    ; ELSE
3757 23EC 48         MOV  C,B      ; C= INVERSE INDUCT # FOR PACKAGE RELEASE
3758
3759 23ED 19          TNXID: DAD  D        ; POINT H.L TO NEXT COUNT IN PPDB
3760 23EE 05         DCR  B        ; DO FOR ALL
3761 23EF C2DA23     JNZ  T9CNT    ; INDUCTS
3762
3763 ; C= INDUCT # FOR PACKAGE TO BE RELEASED IF ANY
3764
3765 23F2 2AB442      LHLD  TCBP    ; SAVE TCBP IN STACK TO AVOID
3766 23F5 E5         PUSH H       ; CONFLICT WITH RLSPC
3767 23F6 0D         DCR  C       ; C=0 MEANS NO MORE PACKAGE TO
3768 23F7 FA0924     JM   TCMGX    ; RELEASE NOW
3769 23FA 3A033D      LDA  NINDU    ; CONVERT C TO
3770 23FD 91         SUB  C       ; INDUCT #
3771 23FE 4F         MOV  C,A     ;
3772 23FF CDAC2C     CALL  PEND    ; ADJUST COUNTS IN PPDB FOR INDUCT
3773 2402 2B         DCX  H       ; POINT H.L TO BYTE 1 IN PPDB
3774 2403 3600     MVI  M,0     ; PWTNG(JINDU) = 0
3775 2405 23         INX  H       ; POINT H.L TO BYTE 2 IN PPDB
3776 2406 CDCF22     CALL  RLSPC   ; RELEASE PACKAGE
3777 2409 E1         POP  H       ; RETURN TCB TO
3778 240A CD4F11     CALL  PUTFS   ; FREE STORAGE
3779 240D C9          RET
3780
3781 ; ILLEGAL COUNT FOUND IN PPDB
3782
3783 240E 3E06       TCMFX: MVI  A,6   ; FATAL ERROR
3784 2410 CD0C2C     CALL  FATAL   ; 6
3785 2413 2AB442     LHLD  TCBP    ; RETURN TCB TO
3786 2416 CD4F11     CALL  PUTFS   ; FREE STORAGE
3787 2419 C9          RET
3788
3789 ; BLOCK-MERGE EVENT SCHEDULED BY UDBPC. FOR SIDE INDUCT ONLY.
3790
3791 ; INPUT: TCBP POINTS TO TCB, D,E POINTS TO BYTE 8 IN TCB.
3792
3793 241A 13         BMG:  INX  D        ; A
3794 241B 1A         LDAX D       ; =JINDU
3795 241C CD9927     CALL  CPPPP   ; POINT H.L TO
3796 241F 23         INX  H       ; BYTE 2
3797 2420 23         INX  H       ; IN PPDB
3798 2421 34         INR  M       ; INC COUNT IN PPDB
3799
3800 2422 2AB442     LHLD  TCBP    ; RETURN TCB TO
3801 2425 CD4F11     CALL  PUTFS   ; FREE STORAGE
3802 2428 C9          RET
3803
3804 ; CLEAR-MERGE EVENT FOR A SIDE INDUCT.
3805 ; SCHEDULED BY SCCMG OR CALLED BY UDBPC.
3806
3807 ; INPUT: TCBP, D,E POINTS TO BYTE 8 IN TCB.
3808 ; ALL REGISTERS EFFECTED, TCBP EFFECTED.
3809

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3810 2429 2AB442 CMG:  LHL D TCBP      ;SAVE TCB
3811 242C E5      PUSH H          ; POINTER
3812 242D 13      INX D          ;POINT D,E TO BYTE 9 IN TCB
3813 242E 1A      LDAX D         ;A=JINDU
3814 242F 4F      MOV C,A        ;C=JINDU
3815 2430 CD9927  CALL CPPP      ;POINT H,L TO BYTE 1
3816 2433 23      INX H          ; IN PPDB
3817 2434 7E      MOV A,M        ;A= PWTNG(JINDU)
3818 2435 23      INX H          ;POINT H,L TO BYTE 2 IN PPDB
3819 2436 35      DCR M          ;DEC COUNT(JINDU) IN PPDB
3820 2437 FA4C24  JM CMGEX      ;IF COUNT NOT< 0
3821
3822 243A C25324  JNZ CMGX      ;IF COUNT = 0
3823
3824 243D D602      SUI 2         ;IF PWTNG(JINDU) = 2 THEN
3825 243F C25324  JNZ CMGX      ; A = 0, AND
3826
3827 ;PACKAGE PENDING AT INDUCT # C
3828
3829 2442 34      INR M          ;INC COUNT(JINDU)
3830 2443 2B      DCX H          ;PWTNG(JINDU)
3831 2444 77      MOV M,A       ; = 0
3832 2445 23      INX H          ;POINT H,L TO BYTE 2 IN PPDB
3833 2446 CDCF22  CALL RLSPC    ;RELEASE PACKAGE (TCBP EFFECTED)
3834 2449 C35324  JMP CMGX
3835
3836 ;ILLEGAL COUNT IN PPDB
3837
3838 244C 3600      CMGEX: MVI M,0 ;CLEAR COUNT(JINDU)
3839 244E 3E0D      MVI A,13     ;ERROR
3840 2450 CD130F    CALL ERRC    ; 13
3841
3842 ;EVENT EXIT
3843
3844 2453 E1      CMGX:  POP H          ;RETURN TCB TO
3845 2454 CD4F11    CALL PUTFS   ; FREE STORAGE
3846 2457 C9      RET
3847
3848 ;TIME-OUT EVENT FOR IDBPC. FOR STRAIGHT INDUCT ONLY.
3849
3850 ;INPUT: TCBP, D,E POINTS TO BYTE 8 IN TCB.
3851 ;ALL REGISTERS EFFECTED.
3852
3853 2458 EB      IDB2:  XCHG         ;POINT H,L TO BYTE 9
3854 2459 23      INX H          ; IN TCB
3855 245A 4E      MOV C,M        ;C = INDUCT #
3856 245B CDA227  CALL CPRLP   ;POINT HL TO
3857 245E 23      INX H          ; BYTE 2
3858 245F 23      INX H          ; IN RLSDB
3859 2460 5E      MOV E,M        ;DE = PSEQ IN RLSDB.
3860 2461 23      INX H          ; I.E., POINT DE TO
3861 2462 56      MOV D,M        ; INDUCT FIFO
3862 2463 3AB442  LDA TCBP     ;IF THIS TCB
3863 2466 BB      CMP E          ; IS THE
3864 2467 C28E24  JNZ IB2FX   ; 1ST TCB
3865 246A 3AB542  LDA TCBP+1  ; IN
3866 246D BA      CMP D          ; INDUCT
3867 246E C28E24  JNZ IB2FX   ; FIFO
3868
3869 ;NOTHING FOUND WRONG WITH THE INDUCT FIFO
3870
3871 2471 E5      PUSH H          ;SAVE POINTER TO BYTE 3 IN RLSDB
3872 2472 210E00  LXI H,14    ;POINT HL TO BYTE 14
3873 2475 19      DAD D          ; IN TCB
3874 2476 5E      MOV E,M        ;POINT DE
3875 2477 23      INX H          ; TO NEXT
3876 2478 56      MOV D,M        ; PACKAGE
3877 2479 E1      POP H          ;POINT HL TO BYTE 3 IN RLSDB
3878 247A 72      MOV M,D        ;POINT PSEQ IN RLSDB
3879 247B 2B      DCX H          ; TO NEXT
3880 247C 73      MOV M,E        ; PACKAGE
3881
3882 247D 3E02      MVI A,2     ;ERROR
3883 247F CD130F    CALL ERRC    ; 2
3884
3885 2482 0602      MVI B,2     ;RELEASE
3886 2484 CD0F11    CALL OPUT   ; OFF
3887 2487 2AB442  LHL TCBP   ;RETURN TCB TO
3888 248A CD4F11    CALL PUTFS  ; FREE STORAGE
3889 248D C9      RET
3890
3891 ;INDUCT FIFO AT PSEQ IN RLSDB HAS BEEN SCRIBBLED
3892
3893 248E 3E09      IB2FX: MVI A,9 ;FATAL ERROR
3894 2490 CD0C2C    CALL FATAL ; 9
3895 2493 C9      RET
3896

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3897 ;EVENT UDB1 IS A PROPER LOWER BOUND FOR UPDATE PHOTO BLOCKED.
3898 ;
3899 ;INPUT: TCBP, D.E POINTS TO BYTE 8 OF TCB, TREE PARAMETERS.
3900 ;ALL REGISTERS EFFECTED.
3901 ;
3902 2494 CD8725 UDB1: CALL UDB12 ;POINT H,L TO PMUZ IN UDBB
3903 2497 5E MOV E,M ;D.E
3904 2498 23 INX H ; = PMUZ
3905 2499 56 MOV D,M ; IN UDBB
3906 249A 7B MOV A,E ; IF
3907 249B B2 ORA D ; PMUZ NOT= 0
3908 249C CABB24 JZ UB1 ;
3909 ;
3910 ;OVERLAPPING MUZ'S, LATTER OVERRIDES FORMER
3911 ;
3912 249F E5 PUSH H ;SAVE POINTER TO PMUZ MS BYTE
3913 24A0 210700 LXI H,7 ;H.L POINTS TO BYTE 7 OF TCB
3914 24A3 19 DAD D ; POINTED TO BY PMUZ IN UDBB
3915 24A4 118025 LXI D,JUNK ;JUNK
3916 24A7 73 MOV M,E ; THAT TCB
3917 24A8 23 INX H ; (ZONE POINTER
3918 24A9 72 MOV M,D ; OVERRIDEN)
3919 24AA 3E03 MVI A,3 ;ERROR 3
3920 24AC CD080F CALL ERPC ; TO MASTER KB
3921 24AF E1 POP H ;H,L POINTS TO PMUZ MS BYTE
3922 24B0 AF XRA A ;CLEAR A
3923 24B1 77 MOV M,A ;CLEAR POINTER
3924 24B2 2B DCX H ; PMUZ
3925 24B3 77 MOV M,A ; IN UDBB
3926 24B4 2AB442 LHLD TCBP ;RETURN TCB TO
3927 24B7 CD4F11 CALL PUTFS ; FREE STORAGE
3928 24BA C9 RET ;
3929 ;
3930 ;LEGITIMATE EVENT UDB1
3931 ;
3932 24BB EB UB1: XCHG ;POINT D.E
3933 24BC 2AB442 LHLD TCBP ; TO
3934 24BF EB XCHG ; TCB
3935 24C0 72 MOV M,D ;POINT PMUZ
3936 24C1 2B DCX H ; IN UDBB
3937 24C2 73 MOV M,E ; TO TCB
3938 ;
3939 24C3 210800 LXI H,8 ;POINT H,L TO BYTE 8
3940 24C6 19 DAD D ; IN TCB
3941 24C7 11D924 LXI D,UDB2 ;POINT EVENT
3942 24CA 72 MOV M,D ; IN TCB
3943 24CB 2B DCX H ; TO
3944 24CC 73 MOV M,E ; UDB2
3945 24CD EB XCHG ;POINT D.E TO BYTE 7 IN TCB
3946 24CE 2A3542 LHLD ZNT ;POINT H,L TO ZNT TABLE
3947 24D1 3A9542 LDA JUPDA ;C
3948 24D4 4F MOV C,A ; = UPDATE PHOTO *
3949 24D5 CD2F29 CALL ENTCB ;SCHEDULE EVENT UDB2
3950 24D8 C9 RET ;
3951 ;
3952 ;LEADING EDGE OF PACKAGE NOT FOUND IN MOVING UPDATABLE ZONE.
3953 ;
3954 ;INPUT: TCBP, D.E POINTS TO BYTE 8 IN TCB.
3955 ;ALL REGISTERS EFFECTED.
3956 ;
3957 24D9 CD8725 UDB2: CALL UDB12 ;POINT H,L TO PMUZ IN UDBB
3958 24DC AF XRA A ;CLEAR A
3959 24DD 77 MOV M,A ;ZERO
3960 24DE 23 INX H ; PMUZ IN UDBB
3961 24DF 77 MOV M,A ;
3962 24E0 3E04 MVI A,4 ;PUT ERROR 4 ON KB
3963 24E2 CD080F CALL ERPC ;
3964 ;
3965 ; TEST TO SEE IF UPDATE-DIVERT-COMplete OPTION SPECIFIED
3966 24E5 3A3B3D LDA DCO ;DCO=2?
3967 24E8 D602 SUI 02H
3968 24EA C28025 JNZ JUNK ;JMP IF DCO NOT = 2.
3969 ;
3970 ; TEST FOR I/C MODE
3971 24ED 3A2B42 LDA ICMO ;ICMO=1?
3972 24F0 A7 ANA A
3973 24F1 CA8025 JZ JUNK ;JMP IF NOT IN I/C MODE
3974 ;
3975 ; TEST IF MESSAGE M IS MASKED
3976 24F4 212B3D LXI H,ERMK0+3 ;H.L= MASK BYTE 4 LOC
3977 24F7 7E MOV A,M ;A= MASK BYTE 4
3978 24F8 E610 ANI 10H ;LOOK AT BIT 4
3979 24FA CA8025 JZ JUNK ;JMP IF M NOT MASKED
3980 ;
3981 ; ARE IN I/C MODE, MESSAGE M IS MASKED, AND DCO=2 IS SPECIFIED. SEND
3982 ; MESSAGE M TO DAS(BY CALLING DSLD) FIRST TEST IF IDENT=0.
3983 24FD 2AB442 LHLD TCBP ;H.L= TCB LOC
3984 2500 110C00 LXI D,12

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3985 2503 19      DAD  D           ;H,L= IDENT LSB LOC
3986 2504 7E      MOV  A,M         ;A= IDENT LSB
3987 2505 23      INX  H           ;H,L= IDENT MSB LKOC
3988 2506 46      MOV  B,M         ;B= IDENT MSB
3989 2507 B0      ORA  B
3990 2508 C21D25  JNZ  TRANS      ;JMP IF IDENT NOT = 0
3991
3992 ; IDENT IS ZERO. THEREFORE PACKAGE CODE IS ZERO (DAS WANTS TO RECIRC
3993 ; PACKAGE) AND IS NOT IN TRANSACTION TABLE. THUS PUT 6 ASCII 0'S INTO
3994 ; BUFFER BYT16 FOR PACKAGE CODE VALUE
3995 250B 211843  LXI  H,BYT16    ;H,L= BYT16 LOC
3996 250E 23      INX  H
3997 250F 23      INX  H           ;POINT TO BYTE 2
3998 2510 3E30     MVI  A,'0'       ;A=ASCII 0
3999 2512 0606     MVI  B,6         ;LOOP COUNT=6
4000 2514 77      PC000: MOV  M,A         ;BYTE = ASCII0
4001 2515 23      INX  H           ;POINT TO NEXT BYTE
4002 2516 05      DCR  B           ;DECREMENT COUNTER
4003 2517 C21425  JNZ  PC000      ;DO FOR BYTES 2-7
4004 251A C33025  JMP  CVTJD      ;JMP TO CVTJD WHEN LOADED
4005
4006 ; ACCESS PACKAGE CODE FROM TRANSACTION TABLE AND PUT INTO BYT16 BUFFER.
4007 ; CLEAR PACKAGE CODE FROM TRANSACTION TABLE
4008 251D 2AB442  TRANS: LHLD TCBP   ;H,L= TCB LOC
4009 2520 110C00  LXI  D,12
4010 2523 19      DAD  D           ;H,L= IDENT LSB LOC
4011 2524 5E      MOV  E,M         ;E= IDENT LSB
4012 2525 23      INX  H           ;H,L= IDENT MSB LOC
4013 2526 56      MOV  D,M         ;D= IDENT MSB
4014 2527 EB      XCHG          ;H,L=IDENT=LOC IN TRANS TAB OF PACK CODE
4015 2528 E5      PUSH H          ;STACK=SAME
4016 2529 CD6C1A  CALL TRANP      ;MOVE PACK CODE TO BYTES 2-7 OF BYT16
4017 252C E1      POP  H           ;H,L= PACK CODE LOC AGAIN
4018 252D CD401A  CALL ZPACK      ;REMOVE PACK CODE FROM TRANS TAB
4019
4020 ; LOAD THE ASSUMED DVT# VVV INTO BYTES 10-12 OF BYT16. THE ASSUMED
4021 ; DVT# IS LOCATED IN LDVDB IN RELATIVE LOC JUPDA-2
4022 2530 2AB442  CVTJD: LHLD TCBP   ;H,L= TCB LOC
4023 2533 110A00  LXI  D,10
4024 2536 19      DAD  D           ;H,L= JUPDA LOC
4025 2537 7E      MOV  A,M         ;A=JUPDA
4026 2538 D602     SUI  2           ;A=JUPDA-2
4027 253A F24225  JP   CVTJ1      ;JMP IF DVT# >1
4028 253D 3E00     MVI  A,0        ;A= 0
4029 253F C34A25  JMP  CVTJ2      ;JMP IF PACKAGE DISAPPEARED BEFORE DV#1
4030 2542 1600     CVTJ1: MVI  D,0      ;D= 0
4031 2544 5F      MOV  E,A         ;D,E= REL LDVDB TABLE LOC OF DV#
4032 2545 2A3D42  LHLD LDVDB     ;H,L= LDVDB LOC
4033 2548 19      DAD  D           ;H,L= LOC IN LDVDB TABLE OF DESIRED DVT#
4034 2549 7E      MOV  A,M         ;A= DVT# IN BINARY
4035 254A CD010E  CVTJ2: CALL ECDIG  ;CONVERT TO ASCII
4036 254D 3A7D42  LDA  XYZ2
4037 2550 322243  STA  BYT16+10   ;STORE ASCII 100'S
4038 2553 3A7E42  LDA  XYZ1
4039 2556 322343  STA  BYT16+11   ;STORE ASCII 10'S
4040 2559 3A7F42  LDA  XYZ0
4041 255C 322443  STA  BYT16+12   ;STORE ASCII 1'S
4042 ; CONVERT BINARY JINDU INTO ASCII AND STORE IN BYTES 8-9 OF BYT16
4043 255F 2AB442  LHLD TCBP     ;H,L = TCB LOC
4044 2562 110900  LXI  D,9
4045 2565 19      DAD  D           ;H,L= JINDU LOC
4046 2566 7E      MOV  A,M         ;A= JINDU IN BINARY
4047 2567 CD010E  CALL ECDIG    ;CONVERT TO ASCII
4048 256A 3A7E42  LDA  XYZ1
4049 256D 322043  STA  BYT16+8    ;STORE ASCII 10'S
4050 2570 3A7F42  LDA  XYZ0
4051 2573 322143  STA  BYT16+9    ;STORE ASCII 1'S
4052 2576 211843  LXI  H,BYT16   ;H,L=BYT16
4053 2579 23      INX  H           ;H,L= BYTE 1 OF BYT16
4054 257A 3E4D     MVI  A,'M'       ;A= ASCII 'M'
4055 257C 77      MOV  M,A         ;PUT ASCII M INTO BYTE 1 OF BYT16
4056 ; LOAD BYT16 INTO FS
4057 257D CD7D1A  CALL DSLD      ;PUT MESSAGE M INTO FS
4058 ; RETURN PACKAGE TCB TO FS
4059
4060 ;
4061 ; JUNK AND NO-OPERATION EVENTS DELETED FROM TREE.
4062 ; CREATED BY UDB1, UDBPC, MCLPC.
4063 ;
4064 ; INPUT: TCBP.
4065 ; REGISTERS D, E, H, L EFFECTED.
4066 ;
4067 2580 2AB442  JUNK: LHLD TCBP   ;RETURN TCB TO
4068 2583 CD4F11  CALL PUTFS     ; FREE STORAGE
4069 ;
4070 2586 C9      NOOP:  RET
4071 ;

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4072
4073
4074
4075
4076
4077
4078
4079
4080 2587 13
4081 2588 13
4082 2589 1A
4083 258A 329542
4084 258D 2A5142
4085 2590 4F
4086 2591 0600
4087 2593 0B
4088 2594 09
4089 2595 09
4090 2596 09
4091 2597 09
4092 2598 C9
4093
4094
4095
4096
4097
4098
4099
4100 2599 2AB442
4101 259C 110700
4102 259F 19
4103 25A0 AF
4104 25A1 77
4105 25A2 23
4106 25A3 77
4107 25A4 C9
4108
4109
4110
4111
4112
4113
4114
4115
4116
4117
4118
4119 25A5 CD8725
4120 25A8 5E
4121 25A9 23
4122 25AA 56
4123 25AB 7B
4124 25AC B2
4125 25AD CACC25
4126
4127
4128
4129 25B0 E5
4130 25B1 210700
4131 25B4 19
4132 25B5 118025
4133 25B8 73
4134 25B9 23
4135 25BA 72
4136 25BB 3E03
4137 25BD CD080F
4138 25C0 E1
4139 25C1 AF
4140 25C2 77
4141 25C3 2B
4142 25C4 77
4143 25C5 2AB442
4144 25C8 CD4F11
4145 25CB C9
4146
4147
4148 25CC EB
4149 25CD 2AB442
4150 25D0 EB
4151 25D1 72
4152 25D2 2B
4153 25D3 73
4154 25D4 210800
4155 25D7 19
4156 25D8 11EA25
4157 25DB 72
4158 25DC 2B

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; 1ST PART OF UDB1, UDB2. POINT H.L TO PMUZ IN UDDB.
; INPUT: DE POINTS TO BYTE 8 IN TCB.
; OUTPUT: DE POINTS TO BYTE 10 IN TCB.
; JUPDA= UPDATE PHOTO \$, BC= UPDATE \$ - 1.
; ALL REGISTERS EFFECTED.
UDB12: INX D ;D,E POINTS TO
INX D ; BYTE 10 IN TCB
LDAX D ;A = JUPDA (UPDATE PHOTO \$)
STA JUPDA ;SET UPDATE PHOTO \$ IN MEMORY
LHLD UDDB ;H,L POINTS TO UDDB TABLE
MOV C,A ;B,C
MVI B,0 ; =JUPDA
DCX B ;B,C=JUPDA-1
DAD B ;H,L
DAD B ; POINTS TO
DAD B ; PMUZ IN UDDB
DAD B ; FOR UD \$ JUPDA
RET

; CHANGED FROM UDB2 BY UDBPC. SET BYTES 7, 8 IN TCB
; ZERO TO DENOTE TCB OUT OF TREE.
; INPUT: TCBP.
; REGISTERS D, E, H, L EFFECTED.
UDB3: LHLD TCBP ;POINT H,L
LXI D,7 ; TO BYTE 7
DAD D ; IN TCB
XRA A ;A=0
MOV M,A ;ZERO
INX H ; BYTES
MOV M,A ; 7, 8
RET

SUBROUTINE DCB1
; INPUTS: TCBP,D,E POINTS TO BYTE 8 OF TCB.TREE PARAMETERS.
; OUTPUTS: EVENT DCB2 CREATED AND SCHEDULED. DCB2 IS UPPER
; BOUND ON MUZ
; PROCESS: SAME AS FOR EVENT UDB1. ONLY DIFFERENCE IS THAT WE
; ARE CREATING AND SCHEDULING DCB2 INSTEAD OF UDB2
DCB1: CALL UDB12 ;POINT H.L TO PMUZ IN UDDB
MOV E,M ;D,E
INX H ; = PMUZ
MOV D,M ; INUDDB
MOV A,E ;IF
ORA D ; PMUZ NOT= 0
JZ DB1

; OVERLAPPING MUZ'S. LATTER OVERRIDES FORMER
PUSH H ;SAVE POINTER TO PMUZ MS BYTE
LXI H,7 ;H,L POINTS TO BYTE 7 OF TCB
DAD D ; POINTED TO BY PMUZ IN UDDB
LXI D,JUNK ;JUNK
MOV M,E ; THAT TCB
INX H ; (ZONE POINTER
MOV M,D ; OVERRIDEN)
MVI A,3 ;ERROR 3
CALL ERPC ; TO MASTER KB
POP H ;H,L POINTS TO PMUZ MS BYTE
XRA A ;CLEAR A
MOV M,A ;CLEAR POINTER
DCX H ; PMUZ
MOV M,A ; IN UDDB
LHLD TCBP ;RETURN TCB TO
CALL PUTFS ; FREE STORAGE
RET

; LEGITIMATE EVENT DCB1
DB1: XCHG ;POINT D,E
LHLD TCBP ; TO
XCHG ; TCB
MOV M,D ;SET PMUZ= CURRENT-PACKAGE TCB LOC
DCX H
MOV M,E
LXI H,8 ;POINT H,L TO BYTE 8
DAD D ; IN TCB
LXI D,DCB2 ;SET EVENT = DCB2
MOV M,D
DCX H

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4159 25DD 73      MOV  M,E
4160 25DE EB      XCHG                ;D,E= BYTE 7 IN TCB
4161 25DF 2A3B42  LHL  DCT           ;H,L= 8000 SERIES TABLE
4162 25E2 3A9542  LDA  JUPDA
4163 25E5 4F      MOV  C,A           ;C= JUPDA
4164 25E6 CD2F29  CALL ENTCTB       ;SCHEDULE EVENT DCB2
4165 25E9 C9      RET
4166                ;
4167                ;
4168                ;   INPUTS: TCBP, D,E POINTS TO BYTE 8 INTCTB
4169                ;
4170                ;   OUTPUTS: PMUZ SET=0, ERROR MESSAGE
4171                ;           'S' TO DAS IF MASKED. PACKAGE TCB RETURNED TO FS
4172                ;
4173                ;   PROCESS: SAME AS FOR UDB2
4174                ;
4175 25EA CD8725  DCB2: CALL  UDB12   ;POINT H,L TO PMUZ IN UDB
4176 25ED AF      XRA  A           ;CLEA A
4177 25EE 77      MOV  M,A         ;ZERO
4178 25EF 23      INX  H           ; PMUZ IN UDB
4179 25F0 77      MOV  M,A
4180                ; CHECK TO SEE IF IN I/C MODE
4181 25F1 3A2B42  LDA  ICMO       ;ICMO=1?
4182 25F4 A7      ANA  A
4183 25F5 CA6F26  JZ   RTFS      ;JMP IF NOT I/C MODE
4184                ; CHECK TO SEE IF MESSAGE S IS MASKED
4185 25F8 21283D  LXI  H,ERMKO   ;H,L=ERROR MASK LOC
4186 25FB 23      INX  H
4187 25FC 23      INX  H           ;H,L= ERROR MASK BYTE 2 (LOC 2A)
4188 25FD 7E      MOV  A,M       ;LOOK AT MASK
4189 25FE E601    ANI  01H     ;LOOK AT BYTE 0
4190 2600 CA6F26  JZ   RTFS      ;JMP IF MESSAGE A NOT MASKED
4191                ;
4192                ; SEND MESSAGE S TO DAS. CLEAR PACKAGE CODE FROM TRANSACTION TABLE
4193                ; IF IDENT NOT = 0. SEND MESSAGE S BY CALLING DSLD.
4194 2603 2AB442  LHL  TCBP     ;H,L= PACKAGE TCB LOC
4195 2606 110C00  LXI  D,12
4196 2609 19      DAD  D           ;H,L= BYTE 12 OF TCB
4197 260A 5E      MOV  E,M     ;E= LSB OF IDENT
4198 260B 23      INX  H
4199 260C 56      MOV  D,M     ;D= MSB OF IDENT
4200 260D EB      XCHG                ;H,L= LOC IN TRANSACTION TAB OF PACK CODE
4201 260E E5      PUSH H         ;STACK= SAME
4202 260F 7C      MOV  A,H     ;A= IDENT MSB
4203 2610 B5      ORA  L           ;IDENT=0 ?
4204 2611 CA1E26  JZ   SABO     ;JMP IF IDENT=0
4205 2614 CD6C1A  CALL TRANP    ;TRANSFER PACK CODE TO BYT16 BYTES 2-7
4206 2617 E1      POP  H           ;H,L=LOC IN TRANS TAB OF PACK CODE
4207 2618 CD401A  CALL ZPACK    ;REMOVE PACK CODE FROM TABLE
4208 261B C33326  JMP  DABO     ;JMP TO DABO
4209                ; PACKAGE IDENT = 0. THEREFORE PACK NOT IN TRANS TAB. SET CODE=0
4210 261E 3E30    SABO: MVI  A,'0'   ;A= ASCII 0
4211 2620 321A43  STA  BYT16+2  ;SET PACKAGE CODE=0 IN BUFFER BYT16
4212 2623 321B43  STA  BYT16+3
4213 2626 321C43  STA  BYT16+4
4214 2629 321D43  STA  BYT16+5
4215 262C 321E43  STA  BYT16+6
4216 262F 321F43  STA  BYT16+7
4217 2632 E1      POP  H           ;ADJUST STACK
4218                ; CONVERT BINARY INDUCT AND DVT IN PACKAGE TCB INTO ASCII AND STORE
4219                ; INTO BUFFER BYT16
4220 2633 2AB442  DABO: LHL  TCBP   ;H,L= PACK TCB LOC
4221 2636 110900  LXI  D,9
4222 2639 19      DAD  D           ;H,L= JINDU LOC IN TCB
4223 263A 7E      MOV  A,M     ;A= BINARY JINDU VALUE
4224 263B CD010E  CALL ECDIG    ;CONVERT TO ASCII
4225 263E 3A7E42  LDA  XYZ1
4226 2641 322043  STA  BYT16+8  ;STORE ASCII 10'S
4227 2644 3A7F42  LDA  XYZ0
4228 2647 322143  STA  BYT16+9  ;STORE ASCII 1'S
4229 264A 2AB442  LHL  TCBP     ;H,L= PACK TCB LOC
4230 264D 110B00  LXI  D,11
4231 2650 19      DAD  D           ;H,L= JDV IN TCB LOC
4232 2651 7E      MOV  A,M     ;A= BINARY OUT#
4233 2652 CD010E  CALL ECDIG    ;CONVERT TO ASCII
4234 2655 3A7D42  LDA  XYZ2
4235 2658 322243  STA  BYT16+10 ;STORE ASCII 100'S
4236 265B 3A7E42  LDA  XYZ1
4237 265E 322343  STA  BYT16+11 ;STORE ASCII 10'S
4238 2661 3A7F42  LDA  XYZ0
4239 2664 322443  STA  BYT16+12 ;STORE ASCII 1'S
4240 2667 3E53    MVI  A,'S'     ;A=ASCII S
4241 2669 321943  STA  BYT16+1  ;STORE
4242 266C CD7D1A  CALL DSLD     ;LOAD BYT16 INTO FS
4243                ; RETURN PACKAGE TCB TO FS
4244                ;
4245 266F 2AB442  RTFS: LHL  TCBP ;H,L= TCB LOC

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4246 2672 CD4F11 CALL PUTFS ;RET TCB TO FS
4247 2675 C9 RET ;RET TO CALLER
4248
4249 ;EVENT TO RESET DIVERT AND SCHEDULE ANY FURTHER RESETS WITH
4250 ; SAME PPI IF REQUIRED.
4251
4252 ;INPUT: D.E POINTS TO BYTE 8 IN TCB, TCBP, TREE PARAMETERS.
4253 ;ALL REGISTERS EFFECTED.
4254
4255 2676 EB DVRS: XCHG ;POINT H.L TO BYTE 9
4256 2677 23 INX H ; IN TCB
4257 2678 4E MOV C,M ;C=JDV
4258 2679 34 INR M ;INC JDV IN TCB FOR NEXT DV
4259 267A 23 INX H ;POINT H.L TO BYTE 10 IN TCB
4260 267B 35 DCR M ;IF MORE DV'S
4261 267C CA9426 JZ DVORS ; TO RESET
4262
4263 ;SCHEDULE DVRS AGAIN TO RESET MORE DV'S
4264
4265 267F C5 PUSH B ;SAVE C=JDV
4266 2680 23 INX H ;POINT H.L TO BYTE 11 IN TCB
4267 2681 5E MOV E,M ;D.E
4268 2682 23 INX H ; = UDBT
4269 2683 56 MOV D,M ; IN TCB
4270 2684 EB XCHG ;SET UDBT
4271 2685 229942 SHLD UDBT ; FOR ENDVA
4272 2688 21FBFF LXI H,0FFFBH ;POINT H.L TO BYTE 12-5=7
4273 268B 19 DAD D ; IN TCB
4274 268C 0C INR C ;C= NEXT JDV
4275 268D CD1D29 CALL ENDVA ;SCHEDULE FURTHER DV RESETS
4276 2690 C1 POP B ;C=JDV
4277 2691 C39A26 JMP DORS
4278
4279 ;NO MORE DV'S TO RESET FOR THIS TASK
4280
4281 2694 2AB442 DVORS: LHLD TCBP ;RETURN TCB TO
4282 2697 CD4F11 CALL PUTFS ; FREE STORAGE
4283
4284 ;OUTPUT DV RESET SIGNAL
4285
4286 269A 3A033D DORS: LDA NINDU ;C
4287 269D 81 ADD C ; =NINDU+JDV
4288 269E 4F MOV C,A ; = DV OUTPUT CH #
4289 269F 0602 MVI B,2 ;RESET
4290 26A1 CD0F11 CALL OPUT ; DV # JDV
4291 26A4 C9 RET
4292
4293 ;DIVERT ACTUATION EVENT.
4294
4295 ;INPUT: TCBP, D.E POINTS TO BYTE 8 OF TCB, TREE PARAMETERS.
4296 ;OUTPUT: DVST.
4297 ;ALL REGISTERS EFFECTED, JUPDA EFFECTED.
4298
4299 26A5 2AB642 DVACT: LHLD CTIME ;FOR DVRCC
4300 26A8 229B42 SHLD DVST
4301 26AB 13 INX D ;POINT H.L TO BYTE 9
4302 26AC EB XCHG ; IN TCB
4303 26AD 3A033D LDA NINDU ;C
4304 26B0 86 ADD M ; =NINDU+JDV
4305 26B1 4F MOV C,A ; = DIVERT OUTPUT CH #
4306 26B2 23 INX H ;B
4307 26B3 46 MOV B,M ; = DIVERT MODE
4308 26B4 C5 PUSH B ;SAVE PARAMETERS FOR OPUT
4309 26B5 05 DCR B ;IF B=1, I.E.,
4310 26B6 C2F926 JNZ DVVNS ; DV TO BE SET
4311
4312 ;CHECK LANE-FULL CONDITIONS LOCALLY
4313
4314 26B9 3A2B42 LDA ICMO ;A= I/C FLAG
4315 26BC A7 ANA A
4316 26BD C2E026 JNZ DVLF0 ;BYPASS LANE FULL CHECK IF I/C MODE AS
4317 ;CHECK HAS ALREADY BEEN MADE AT PEVIOUS
4318 ;UPDATE PHOTO
4319 26C0 3A3C3D LDA LFO ;IF LANE-FULL-
4320 26C3 E602 ANI 02H ; DV-INHIBIT
4321 26C5 CAE026 JZ DVLF0 ; DESIRED
4322 26C8 E5 PUSH H ;SAVE POINTER TO TCB BYTE 10
4323 26C9 2B DCX H ;POINT H.L TO BYTE 9 IN TCB
4324 26CA 7E MOV A,M ;A=JDV
4325 26CB 47 MOV B,A ;SAVE JDV IN B
4326 26CC CD250D CALL CPLFM ;IF DV LANE FULL THEN
4327 26CF A6 ANA M ; CLEAR FLAG Z
4328 26D0 E1 POP H ;POINT H.L TO BYTE 10 IN TCB
4329 26D1 CAE026 JZ DVLF0 ;IF LANE FULL
4330 26D4 48 MOV C,B ;C=JDV
4331 26D5 CD0427 CALL DVRCC ;SEND PACKAGE TO LFDV(LANE FULL_ERR DVT
4332 26D8 2AB442 LHLD TCBP ;RETURN PACKAGE TCB TO FS

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4333 26DB CD4F11 CALL PUTFS
4334 26DE C1 POP B ;ADJUST STACK
4335 26DF C9 RET ;RET
4336
4337 ;IF DVTRS=2 THEN SCHEDULE RESET OF DV # JDV
4338 ;
4339 26E0 3A383D DVLFO: LDA DVTRS ;IF DVTRS=2,
4340 26E3 D602 SUI 2 ; I.E., FIXED-LE-
4341 26E5 C2F926 JNZ DDVNS ; DV-PULSE OPTION
4342 26E8 34 INR M ;DV MODE IN TCB = 2 TO RESET
4343 26E9 11FDFE LXI D,0FFFDDH ;POINT D,E TO
4344 26EC 19 DAD D ; BYTE 7
4345 26ED EB XCHG ; IN TCB
4346 26EE 0E01 MVI C,1 ;SET PARAMETERS FOR ENTCTB
4347 26F0 210340 LXI H,DVH ; ( DVH(1)=DVH )
4348 26F3 CD2F29 CALL ENTCTB ;SCHEDULE DV RESET
4349 26F6 C3FF26 JMP DOPUT
4350
4351 ;OUTPUT DV SIGNAL AND EXIT
4352 ;
4353 26F9 2AB442 DDVNS: LHLD TCBP ;RETURN FREE
4354 26FC CD4F11 CALL PUTFS ; STORAGE
4355 26FF C1 DOPUT: POP B ;ACTUATE
4356 2700 CD0F11 CALL OPUT ; DV
4357 2703 C9 RET
4358
4359 ;CALLED BY DVACT FOR PACKAGE RECIRCULATION.
4360 ;
4361 ;INPUT: C=JDV (DIVERT #), NDVT, TREE PARAMETERS, DVST= CTIME AT DIVERT.
4362 ;OUTPUT: UDBT= UD-BLOCKING TIME.
4363 ;ALL REGISTERS EFFECTED, CTIME EFFECTED.
4364 ;
4365 DVRC:
4366 ;
4367 ;SCHEDULE RESETS FOR DIVERTS # JDV TO # LAST FOR THIS UD
4368 ;
4369 2704 0D DCR C ;C=JDV-1
4370 2705 0600 MVI B,0 ;B,C=JDV-1
4371 2707 2A3942 LHLD DVT ;H,L POINTS TO TIME IN
4372 270A 09 DAD B ; DVT TABLE
4373 270B 09 DAD B ;
4374 270C 5E MOV E,M ;D,E= TIME IN DVT
4375 270D 23 INX H ;
4376 270E 56 MOV D,M ;
4377 270F 2A9B42 LHLD DVST ;H,L= CTIME AT DIVERT
4378 2712 7D MOV A,L ;H,L=CTIME- TIME IN DVT FOR
4379 2713 93 SUB E ; JDV=
4380 2714 6F MOV L,A ; LAST UPDA BLOCKING
4381 2715 7C MOV A,H ; TIME
4382 2716 9A SBB D ;
4383 2717 67 MOV H,A ;
4384 2718 229942 SHLD UDBT ;SAVE IN UDBT FOR SDVRS
4385 ;
4386 271B 3AC042 LDA MUD ;A= LAST UD # FOR THIS PPI
4387 271E 5F MOV E,A ;D,E
4388 271F 1600 MVI D,0 ; = LAST UD #
4389 2721 2A3D42 LHLD LDVDB ;POINT H,L TO LDV OF
4390 2724 19 DAD D ; LAST UD
4391 2725 2B DCX H ; FOR THIS PPI
4392 2726 79 MOV A,C ;A=JDV-1
4393 2727 2B DPVUD: DCX H ;POINT H,L TO PREVIOUS LDV
4394 2728 1D DCR E ;E = PREVIOUS UD #
4395 2729 BE CMP M ;DUNTIL JDV-1 >= LDV.
4396 272A DA2727 JC DPVUD ; I.E., LDV < JDV
4397 272D 1C INR E ;COMPUTE
4398 272E D5 PUSH D ; UPDATE PHOTO #
4399 272F 23 INX H ;A= LDV
4400 2730 7E MOV A,M ; FOR THIS UD
4401 2731 0C INR C ;C=JDV
4402 2732 CD7727 CALL SDVRS ;SCHEDULE THEIR RESETS
4403 ;
4404 ;SCHEDULE EVENT UDB1 FOR NEXT UPDATE PHOTO
4405 ;
4406 2735 C1 POP B ;C = UPDATE PHOTO #
4407 2736 3A043D LDA NUPDA ;IF ALREADY
4408 2739 A9 XRA C ; PASSED LAST
4409 273A 47 MOV B,A ; UPDATE PHOTO
4410 273B 3A313D LDA SIDE0 ; AND NO SIDE-INDUCT
4411 273E E602 ANI 02H ; AUTO
4412 2740 B0 ORA B ; RECIRCULATION
4413 2741 C8 RZ ; THEN RETURN
4414 ;
4415 2742 CDCB0F CALL GETFS ;GET FREE BLOCK
4416 2745 110D00 LXI D,13 ;POINT TO BYTE 13
4417 2748 19 DAD D ; IN TCB
4418 2749 3600 MVI M,0 ;DEFAULT
4419 274B 2B DCX H ; IDENT IN TCB

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4420 274C 3600      MVI  M.0      ; = 0
4421 274E 2B       DCX  H         ;POINT TO BYTE 11 IN TCB
4422 274F 3A3D3D   LDA  L FEDV    ;ROUTE PACKAGE TO LANE-FULL
4423 2752 77       MOV  M.A       ; ERROR DIVERT
4424 2753 2B       DCX  H         ;POINT TO BYTE 10 IN TCB
4425 2754 71       MOV  M.C       ;SET JUPDA IN TCB
4426 2755 34       INR  M         ; FOR NEXT UPDATE PHOTO
4427 2756 3A043D   LDA  NUPDA    ;UPDATE PHOTO $ NUPDA + 1
4428 2759 BE       CMP  M         ; MEANS
4429 275A D25F27   JNC  DRTCB    ; UPDATE PHOTO
4430 275D 3601     MVI  M.1      ; $ 1
4431 275F 2B       DRTCB: DCX  H      ;DEFAULT
4432 2760 3601     MVI  M.1      ; JINDU IN TCB = 1
4433 2762 2B       DCX  H         ;POINT TO TCB BYTE 8
4434 2763 119424   LXI  D,UDB1   ;LET
4435 2766 72       MOV  M.D      ; EVENT
4436 2767 2B       DCX  H         ; BE
4437 2768 73       MOV  M.E      ; UDB1
4438 2769 EB       XCHG        ;POINT D,E TO BYTE 7 IN TCB
4439 276A 2A9942   LHLD UDBT    ;SET CTIME FOR
4440 276D 22B642   SHLD CTIME   ; ENTCB
4441 2770 2A3742   LHLD UDT     ;SCHEDULE UDB1 FOR
4442 2773 CD2F29   CALL ENTCB   ; UD $ C
4443 2776 C9       RET
4444
4445 ;:
4446 ;SCHEDULE RESETS FOR DIVERTS $ C TO $ A.
4447 ; ALL FOR CURRENT PPI.
4448 ; CALLED BY UDBPC, SCHDV OR DVRCC.
4449 ;
4450 ;INPUT: UDBT= LAST UPDA BLOCKING TIME, TREE PARAMETERS.
4451 ;ALL REGISTERS EFFECTED.
4452 ;
4453 2777 91       SDVRS: SUB  C      ;RETURN
4454 2778 D8       RC          ; IF NO DIVERT TO RESET
4455 ;
4456 2779 3C       INR  A         ;A= $ DV'S TO RESET
4457 277A 47       MOV  B.A       ;B= $ DIVERTS TO BE RESET
4458 277B CDCB0F   CALL GETFS    ;GET FREE BLOCK
4459 277E 110C00   LXI  D,12     ;POINT H,L TO BYTE 12
4460 2781 19       DAD  D         ; IN TCB
4461 2782 EB       XCHG        ;D,E
4462 2783 2A9942   LHLD UDBT    ; =
4463 2786 EB       XCHG        ; UDBT
4464 2787 72       MOV  M.D      ;SET UDBT
4465 2788 2B       DCX  H         ; IN
4466 2789 73       MOV  M.E      ; TCB
4467 278A 2B       DCX  H         ;POINT H,L TO BYTE 10 IN TCB
4468 278B 70       MOV  M.B       ;STORE $ DV'S TO RESET
4469 278C 2B       DCX  H         ;SET $ OF 1ST DV
4470 278D 71       MOV  M.C       ; TO RESET
4471 278E 2B       DCX  H         ;POINT H,L TO BYTE 8 IN TCB
4472 278F 117626   LXI  D,DVRS   ;LET
4473 2792 72       MOV  M.D      ; EVENT
4474 2793 2B       DCX  H         ; BE
4475 2794 73       MOV  M.E      ; DVRS
4476 2795 CD1D29   CALL ENDVA   ;SCHEDULE DV RESETS
4477 2798 C9       RET
4478 ;:
4479 ;POINT H,L TO PPDB( INDUCT $).
4480 ;
4481 ;INPUT: A= INDUCT $
4482 ;REGISTERS A, D, E EFFECTED.
4483 ;
4484 2799 3D       CPPPP: DCR  A         ;A = INDUCT $ - 1
4485 279A 87       ADD  A         ;A = ( INDUCT $ - 1 ) * 2
4486 279B 2A4342   LHLD PPDB    ;POINT H,L TO PPDB
4487 279E CD1B0D   CALL AD16A   ;ADD ( INDUCT $ * 32 )
4488 27A1 C9       RET
4489 ;
4490 ;
4491 ;COMPUTE RSLDB POINTER.
4492 ;
4493 ;INPUT: C= INDUCT $
4494 ;OUTPUT: HL POINTS TO RLSDB( INDUCT $ ).
4495 ;REGISTER B ALSO EFFECTED.
4496 ;
4497 27A2 2A4F42   CPRLP: LHLD RLSDB ;POINT H,L TO RLSDB START
4498 27A5 0600     MVI  B.0     ;BC
4499 27A7 0D       DCR  C         ; = INDUCT $ - 1
4500 27A8 09       DAD  B         ;POINT
4501 27AA 09       DAD  B         ; H,L
4502 27AB 09       DAD  B         ; TO
4503 27AC 0C       INR  C         ; RLSDB(JINDU)
4504 27AD C9       RET
4505 ;:
4506 ;REMOVE THE NEXT ITEM FROM THE FUTURE EVENT TREE

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4507      ; AND UPDATE NXTIM.
4508
4509      ; INPUT: TREE PARAMETERS, AVL TREE.
4510      ; OUTPUT: H.L POINTS TO TCB DELETED.
4511      ; ALL REGISTERS EFFECTED.
4512
4513 27AE 2ABA42  DEQUE:  LHLD  RLINK      ;GET ADDRESS OF ROOT
4514 27B1 EB      XCHG                ;SAVE ADDRESS OF ROOT
4515 27B2 210000  LXI   H,0      ;SAVE THE SP VALUE
4516 27B5 39      DAD   SP        ;FOR STACK PURGE
4517 27B6 22C642  SHLD  STKSV
4518 27B9 21B842  LXI   H,HEAD    ;STACK (HEAD,+)
4519 27BC E5      PUSH  H
4520 27BD 2601    MVI   H,01
4521 27BF E5      TEST:  PUSH  H      ;STACK DIRECTION
4522 27C0 EB      XCHG                ;NODE ADDRESS TO (H,L)
4523          GTLLK  0,0    ;MOVE TO THE LEFT
4523          IF   00000H  ;IS NODE ADDRESS IN (H,L)?
4523          LHLD  00000H  ;NO-NODE POINTER TO (H,L)
4523          ENDIF
4523 27C1 5E      MOV   E,M      ;MOVE THE LEFT
4523 27C2 23      INX   H        ;LINK INTO THE
4523 27C3 56      MOV   D,M      ;(D,E) REGISTER
4523          IF   00000H  ;IS IT SAVE LINK MODE?
4523          XCHG                ;YES - LINK TO (H,L)
4523          SHLD  00000H  ;LINK TO MEMORY
4523          ENDIF
4524 27C4 2B      DCX   H        ;RESTORE NODE ADDRESS
4525 27C5 7A      MOV   A,D      ;HAVE WE FOUND THE
4526 27C6 B3      ORA   E        ;LEFTMOST NODE IN TREE?
4527 27C7 CAD027  JZ    DLTE     ;YES - DELETE IT
4528 27CA E5      PUSH  H        ;STACK AS PARENT ADDRESS
4529 27CB 26FF    MVI   H,0FFH   ;MARK DIRECTION LEFT
4530 27CD C3BF27  JMP   TEST     ;GO LEFT AGAIN
4531
4532      ;
4533      ; DELETE A RECORD FROM TREE
4534
4534 27D0 22C842  DLTE:  SHLD  QP      ;SAVE NODE ADDRESS
4535          GTRLK  0,0    ;GET ITS RIGHT LINK
4535          IF   00000H  ;IS NODE ADDRESS IN (H,L)?
4535          LHLD  00000H  ;NO - NODE POINTER TO (H,L)
4535          ENDIF
4535 27D3 23      INX   H        ;STEP POINTER TO
4535 27D4 23      INX   H        ;THE RIGHT LINK
4535 27D5 5E      MOV   E,M      ;MOVE THE RIGHT
4535 27D6 23      INX   H        ;LINK INTO THE
4535 27D7 56      MOV   D,M      ; (D,E) REGISTER
4535          IF   00000H  ;IS IT SAVE LINK MODE?
4535          XCHG                ;YES - LINK TO (H,L)
4535          SHLD  00000H  ;LINK TO MEMORY
4535          ENDIF
4536 27D8 E1      POP   H        ;POP AND SAVE THE
4537 27D9 7C      MOV   A,H      ;DIRECTION INTO NODE
4538 27DA E1      POP   H        ;GET THE PARENT NODE
4539 27DB 22CA42  SHLD  QX      ;SAVE ADDRESS OF PARENT
4540 27DE B7      ORA   A        ;DIRECTION TO CONDITION CODE
4541 27DF F5      PUSH  PSW     ;SAVE DIRECTION STATUS
4542 27E0 F21329  JP    QRTY    ;PARENT IS TREE HEADER
4543          PTLLK  0,0    ;RLINK(SON) TO LLINK(PARENT)
4543 0001      MCTEM  SET  1      ;ASSUME ADDRESS IN (H,L)
4543          IF   00000H  ;TEST ASSUMPTION
4543          MCTEM  SET  0      ;FALSE ASSUMPTION
4543          ENDIF
4543          IF   00000H  ;IS NEW LINK IN (D,E)?
4543          IF   MCTEM    ;NO - BUT IS POINTER IN (H,L)?
4543          XCHG                ;YES - SAVE NODE POINTER
4543          ENDIF
4543          LHLD  00000H  ;NEW LINK TO (H,L)
4543          XCHG                ;NEW LINK TO (D,E)
4543          ENDIF
4543          IF   00000H  ;IS NODE ADDRESS IN (H,L)?
4543          LHLD  00000H  ;NO - NODE POINTER TO (H,L)
4543          ENDIF
4543 27E3 73      MOV   M,E      ;MOVE THE NEW
4543 27E4 23      INX   H        ;LEFT LINK VALUE
4543 27E5 72      MOV   M,D      ;INTO THE NODE
4544
4544 27E6 7A      QRTX:  MOV   A,D      ;IF RLINK(SON) IS ZERO THE
4545 27E7 B3      ORA   E        ;PARENT IS THE POST ORDER
4546 27E8 CAF427  JZ    SETTM   ;SUCCESSOR OF TASK NODE
4547 27EB EB      FIND:  XCHG                ;FIND THE POST ORDER
4548          GTLLK  0,0    ;SUCCESSOR OF THE
4548          IF   00000H  ;IS NODE ADDRESS IN (H,L)?
4548          LHLD  00000H  ;NO-NODE POINTER TO (H,L)
4548          ENDIF
4548 27EC 5E      MOV   E,M      ;MOVE THE LEFT

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4548	27ED 23		INX H		;LINK INTO THE
4548	27EE 56		MOV D,M		; (D,E) REGISTER
4548			IF 00000H		; IS IT SAVE LINK MODE?
4548			XCHG		; YES - LINK TO (H,L)
4548			SHLD 00000H		; LINK TO MEMORY
4548			ENDIF		
4548					
4549	27EF 7A		MOV A,D		;SELECTED NODE
4550	27F0 B3		ORA E		
4551	27F1 C2EB27		JNZ FIND		
4552	27F4 110400	SETTM:	LXI D,4		;TIME OFFSET LESS ONE TO (D,E)
4553	27F7 19		DAD D		;POINT TO NEXT EVENT TIME
4554	27F8 5E		MOV E,M		;GET THE TIME OF
4555	27F9 23		INX H		;THE NEXT EVENT
4556	27FA 56		MOV D,M		
4557	27FB EB		XCHG		;TIME TO (H,L)
4558	27FC 22BD42		SHLD NXTIM		;UPDATE NEXT EVENT TIME
4559	27FF F1		POP PSW		;RECOVER ENTRY DIRECTION
4560	2800 2ACA42	QLOOP:	LHLD QX		;POINT TO PARENT NODE
4561	2803 B7	QLRP:	ORA A		;ENTRY MODE TO FLAGS
4562	2804 F2F528		JP QXIT		;PARENT IS TREE HEADER
4563			GTBAL 0,0		;GET BALANCE FACTOR OF PARENT
4563			IF 00000H		;IS NODE ADDRESS IN (H,L)?
4563			LHLD 00000H		;NO - NODE POINTER TO (H,L)
4563			ENDIF		
4563	2807 23		INX H		;STEP POINTER TO
4563	2808 23		INX H		;THE BALANCE BYTE
4563	2809 23		INX H		
4563	280A 23		INX H		
4563	280B 7E		MOV A,M		;BALANCE CODE TO A-REG
4563			IF 00000H		;TEST IF SAVE NODE
4563			STA 00000H		;YES - CODE TO MEMORY
4563			ENDIF		
4563	280C B7		ORA A		;SET CONDITION CODES
4563					
4564	280D FAC228		JM BMNUS		;PARENT WAS HEAVY TO LEFT
4565	2810 CACD28		JZ BZERO		;PARENT WAS BALANCED
4566			GTRLK QX,QT		;PARENT WAS HEAVY TO RIGHT
4566			IF QX		;IS NODE ADDRESS IN (H,L)?
4566	2813 2ACA42		LHLD QX		;NO - NODE POINTER TO (H,L)
4566			ENDIF		
4566	2816 23		INX H		;STEP POINTER TO
4566	2817 23		INX H		;THE RIGHT LINK
4566	2818 5E		MOV E,M		;MOVE THE RIGHT
4566	2819 23		INX H		;LINK INTO THE
4566	281A 56		MOV D,M		; (D,E) REGISTER
4566			IF QT		; IS IT SAVE LINK MODE?
4566	281B EB		XCHG		; YES - LINK TO (H,L)
4566	281C 22CC42		SHLD QT		;LINK TO MEMORY
4566			ENDIF		
4566					
4567			GTLK 0,0		;GET THE BALANCE FACTOR OF THE
4567			IF 00000H		;IS NODE ADDRESS IN (H,L)?
4567			LHLD 00000H		;NO-NODE POINTER TO (H,L)
4567			ENDIF		
4567	281F0 5E		MOV E,M		;MOVE THE LEFT
4567	2820 23		INX H		;LINK INTO THE
4567	2821 56		MOV D,M		; (D,E) REGISTER
4567			IF 00000H		; IS IT SAVE LINK MODE?
4567			XCHG		; YES - LINK TO (H,L)
4567			SHLD 00000H		; LINK TO MEMORY
4567			ENDIF		
4567					
4568	2822 2B		DCX H		;PARENTS RIGHT SUBTREE AFTER
4569			GTBAL 0,0		;GETTING THE LEFT LINK OF THE
4569			IF 00000H		;IS NODE ADDRESS IN (H,L)?
4569			LHLD 00000H		;NO - NODE POINTER TO (H,L)
4569			ENDIF		
4569	2823 23		INX H		;STEP POINTER TO
4569	2824 23		INX H		;THE BALANCE BYTE
4569	2825 23		INX H		
4569	2826 23		INX H		
4569	2827 7E		MOV A,M		;BALANCE CODE TO A-REG
4569			IF 00000H		;TEST IF SAVE NODE
4569			STA 00000H		;YES - CODE TO MEMORY
4569			ENDIF		
4569	2828 B7		ORA A		;SET CONDITION CODES
4569					
4570	2829 CAD228		JZ CASE3		;SUBTREE ROOT NODE
4571	282C FA5028		JM CASE2		;RIGHT SUBTREE HEAVY ON LEFT
4572	282F 3600		MVI M,0		;RIGHT SUBTREE HEAVY ON RIGHT
4573			PTRLK QX,0		;ROTATE SUBTREE INTO BALANCE
4573	0001	MCTEM	SET 1		;ASSUME ADDRESS IN (H,L)
4573			IF QX		;TEST ASSUMPTION
4573	0000	MCTEM	SET 0		;FALSE ASSUMPTION
4573			ENDIF		
4573			IF 00000H		;IS NEW LINK IN (D,E)?
4573			IF MCTEM		;NO - BUT IS POINTER IN (H,L)?

4573			XCHG		:YES - SAVE NODE POINTER
4573			ENDIF		
4573			LHLD 00000H		:NEW LINK TO (H,L)
4573			XCHG		:NEW LINK TO (D,E)
4573			ENDIF		
4573			IF QX		:IS NODE ADDRESS IN (H,L)?
4573	2831	2ACA42	LHLD QX		:NO - NODE POINTER TO (H,L)
4573			ENDIF		
4573	2834	23	INX H		:NOW STEP TO THE
4573	2835	23	INX H		:RIGHT LINK BYTES
4573	2836	73	MOV M,E		:MOVE THE NEW
4573	2837	23	INX H		:RIGHT LINK VALUE
4573	2838	72	MOV M,D		:INTO THE NODE
4573					
4574			PTLLK QT,QX		:PARENT TO LLINK(SUBTREE ROOT)
4574	0001	MCTEM	SET 1		:ASSUME ADDRESS IN (H,L)
4574			IF QT		:TEST ASSUMPTION
4574	0000	MCTEM	SET 0		:FALSE ASSUMPTION
4574			ENDIF		
4574			IF QX		:IS NEW LINK IN (D,E)?
4574			IF MCTEM		:NO - BUT IS POINTER IN (H,L)?
4574			XCHG		:YES - SAVE NODE POINTER
4574			ENDIF		
4574	2839	2ACA42	LHLD QX		:NEW LINK TO (H,L)
4574	283C	EB	XCHG		:NEW LINK TO (D,E)
4574			ENDIF		
4574			IF QT		:IS NODE ADDRESS IN (H,L)?
4574	283D	2ACC42	LHLD QT		:NO - NODE POINTER TO (H,L)
4574			ENDIF		
4574	2840	73	MOV M,E		:MOVE THE NEW
4574	2841	23	INX H		:LEFT LINK VALUE
4574	2842	72	MOV M,D		:INTO THE NODE
4574					
4575	2843	AF	XRA A		:THE ORIGINAL PARENT IS
4576	2844	EB	XCHG		:NOW BALANCED BUT THE ROOT
4577			PTBAL 0.0		:OF RIGHT SUBTREE HAS REPLACED IT
4577			IF 00000H		:IS FACTOR IN A-REG?
4577			LDA 00000H		:NO - BALANCE TO A-REG
4577			ENDIF		
4577			IF 00000H		:IS NODE ADDRESS IN (H,L)
4577			LHLD 00000H		:NO - NODE POINTER TO (H,L)
4577			ENDIF		
4577	2845	23	INX H		:STEP POINTER TO
4577	2846	23	INX H		:THE BALANCE BYTE
4577	2847	23	INX H		
4577	2848	23	INX H		
4577	2849	77	MOV M,A		:BALANCE CODE TO NODE
4577					
4578	284A	CDF928	CALL CSUBR		:SO ADJUST THE GRANDPARENT
4579	284D	C30028	JMP QLOOP		:TREE HEIGHT CHANGED, SO GO ON
4580					
4581	2850	EB	XCHG	CASE2:	:SAVE THE LEFT LINK OF
4582	2851	22CE42	SHLD QW		:THE SUBTREE ROOT NODE
4583			GTRLK 0.0		:PLACE THIS NODE INTO PARENTS
4583			IF 00000H		:IS NODE ADDRESS IN (H,L)?
4583			LHLD 00000H		:NO - NODE POINTER TO (H,L)
4583			ENDIF		
4583	2854	23	INX H		:STEP POINTER TO
4583	2855	23	INX H		:THE RIGHT LINK
4583	2856	5E	MOV E,M		:MOVE THE RIGHT
4583	2857	23	INX H		:LINK INTO THE
4583	2858	56	MOV D,M		: (D,E) REGISTER
4583			IF 00000H		:IS IT SAVE LINK NODE?
4583			XCHG		:YES - LINK TO (H,L)
4583			SHLD 00000H		:LINK TO MEMORY
4583			ENDIF		
4583					
4584			PTLLK QT,0		:TREE POSITION WITH PARENT AS
4584	0001	MCTEM	SET 1		:ASSUME ADDRESS IN (H,L)
4584			IF QT		:TEST ASSUMPTION
4584	0000	MCTEM	SET 0		:FALSE ASSUMPTION
4584			ENDIF		
4584			IF 00000H		:IS NEW LINK IN (D,E)?
4584			IF MCTEM		:NO - BUT IS POINTER IN (H,L)?
4584			XCHG		:YES - SAVE NODE POINTER
4584			ENDIF		
4584			LHLD 00000H		:NEW LINK TO (H,L)
4584			XCHG		:NEW LINK TO (D,E)
4584			ENDIF		
4584			IF QT		:IS NODE ADDRESS IN (H,L)?
4584	2859	2ACC42	LHLD QT		:NO - NODE POINTER TO (H,L)
4584			ENDIF		
4584	285C	73	MOV M,E		:MOVE THE NEW
4584	285D	23	INX H		:LEFT LINK VALUE
4584	285E	72	MOV M,D		:INTO THE NODE
4584					
4585			PTRLK QW,QT		:ITS LEFT SUBTREE AND ITS OLD LEFT

4585	0001	MCTEM	SET	1	:ASSUME ADDRESS IN (H.L.)
4585			IF	QW	:TEST ASSUMPTION
4585	0000	MCTEM	SET	0	:FALSE ASSUMPTION
4585			ENDIF		
4585			IF	QT	:IS NEW LINK IN (D,E)?
4585			IF	MCTEM	:NO - BUT IS POINTER IN (H.L)?
4585			XCHG		:YES - SAVE NODE POINTER
4585			ENDIF		
4585	285F	ZACC42	LHLD	QT	:NEW LINK TO (H.L.)
4585	2862	EB	XCHG		:NEW LINK TO (D,E)
4585			ENDIF		
4585			IF	QW	:IS NODE ADDRESS IN (H.L)?
4585	2863	ZACE42	LHLD	QW	:NO - NODE POINTER TO (H.L.)
4585			ENDIF		
4585	2866	23	INX	H	:NOW STEP TO THE
4585	2867	23	INX	H	:RIGHT LINK BYTES
4585	2868	73	MOV	M.E	:MOVE THE NEW
4585	2869	23	INX	H	:RIGHT LINK VALUE
4585	286A	72	MOV	M.D	:INTO THE NODE
4585					
4586			GTLK	QW.0	:SUBTREE AS PARENTS RIGHT SUBTREE
4586			IF	QW	:IS NODE ADDRESS IN (H.L)?
4586	286B	ZACE42	LHLD	QW	:NO-NODE POINTER TO (H.L.)
4586			ENDIF		
4586	286E	5E	MOV	E.M	:MOVE THE LEFT
4586	286F	23	INX	H	:LINK INTO THE
4586	2870	56	MOV	D.H	: (D,E) REGISTER
4586			IF	00000H	:IS IT SAVE LINK MODE?
4586			XCHG		:YES - LINK TO (H.L.)
4586			SHLD	00000H	:LINK TO MEMORY
4586			ENDIF		
4586					
4587			PTRLK	QX.0	:ITS RIGHT SUBTREE REPLACES THE
4587	0001	MCTEM	SET	1	:ASSUME ADDRESS IN (H.L.)
4587			IF	QX	:TEST ASSUMPTION
4587	0000	MCTEM	SET	0	:FALSE ASSUMPTION
4587			ENDIF		
4587			IF	00000H	:IS NEW LINK IN (D,E)?
4587			IF	MCTEM	:NO - BUT IS POINTER IN (H.L)?
4587			XCHG		:YES - SAVE NODE POINTER
4587			ENDIF		
4587			LHLD	00000H	:NEW LINK TO (H.L.)
4587			XCHG		:NEW LINK TO (D,E)
4587			ENDIF		
4587			IF	QX	:IS NODE ADDRESS IN (H.L)?
4587	2871	ZACA42	LHLD	QX	:NO - NODE POINTER TO (H.L.)
4587			ENDIF		
4587	2874	23	INX	H	:NOW STEP TO THE
4587	2875	23	INX	H	:RIGHT LINK BYTES
4587	2876	73	MOV	M.E	:MOVE THE NEW
4587	2877	23	INX	H	:RIGHT LINK VALUE
4587	2878	72	MOV	M.D	:INTO THE NODE
4587					
4588			PTLLK	QW.QX	:LEFT SUBTREE OF ORIGINAL RIGHT
4588	0001	MCTEM	SET	1	:ASSUME ADDRESS IN (H.L.)
4588			IF	QW	:TEST ASSUMPTION
4588	0000	MCTEM	SET	0	:FALSE ASSUMPTION
4588			ENDIF		
4588			IF	QX	:IS NEW LINK IN (D,E)?
4588			IF	MCTEM	:NO - BUT IS POINTER IN (H.L)?
4588			XCHG		:YES - SAVE NODE POINTER
4588			ENDIF		
4588	2879	ZACA42	LHLD	QX	:NEW LINK TO (H.L.)
4588	287C	EB	XCHG		:NEW LINK TO (D,E)
4588			ENDIF		
4588			IF	QW	:IS NODE ADDRESS IN (H.L)?
4588	287D	ZACE42	LHLD	QW	:NO - NODE POINTER TO (H.L.)
4588			ENDIF		
4588	2880	73	MOV	M.E	:MOVE THE NEW
4588	2881	23	INX	H	:LEFT LINK VALUE
4588	2882	72	MOV	M.D	:INTO THE NODE
4588					
4589	2883	2B	DCX	H	:SUBTREE OF THE PARENT
4590			GTBAL	0.0	:GET BALANCE FACTOR OF PARENTS REPLCMT
4590			IF	00000H	:IS NODE ADDRESS IN (H.L)?
4590			LHLD	00000H	:NO - NODE POINTER TO (H.L.)
4590			ENDIF		
4590	2884	23	INX	H	:STEP POINTER TO
4590	2885	23	INX	H	:THE BALANCE BYTE
4590	2886	23	INX	H	
4590	2887	23	INX	H	
4590	2888	7E	MOV	A.M	:BALANCE CODE TO A-REG
4590			IF	00000H	:TEST IF SAVE NODE
4590			STA	00000H	:YES - CODE TO MEMORY
4590			ENDIF		
4590	2889	B7	ORA	A	:SET CONDITION CODES
4590					

4591	288A CAB628	JZ	CSSZ	:BALANCED NODE
4592	288D F2B428	JP	CSSP	:HEAVY ON THE RIGHT
4593	2890 AF	XRA	A	:NEW RIGHT SUBTREE OF PARENT
4594		PTBAL	QX,0	:BRINGS IT INTO BALANCE BUT
4594		IF	00000H	:IS FACTOR IN A-REG?
4594		LDA	00000H	:NO - BALANCE TO A-REG
4594		ENDIF		
4594		IF	QX	:IS NODE ADDRESS IN (H,L)
4594	2891 2ACA42	LHLD	QX	:NO - NODE POINTER TO (H,L)
4594		ENDIF		
4594	2894 23	INX	H	:STEP POINTER TO
4594	2895 23	INX	H	:THE BALANCE BYTE
4594	2896 23	INX	H	
4594	2897 23	INX	H	
4594	2898 77	MOV	M,A	:BALANCE CODE TO NODE
4594				
4595	2899 3C	INR	A	:ROOT OF ORIGINAL RIGHT SUBTREE
4596		PTBAL	QT,0	:IS NOW HEAVY TO THE RIGHT
4596		IF	00000H	:IS FACTOR IN A-REG?
4596		LDA	00000H	:NO - BALANCE TO A-REG
4596		ENDIF		
4596		IF	QT	:IS NODE ADDRESS IN (H,L)
4596	289A 2ACC42	LHLD	QT	:NO - NODE POINTER TO (H,L)
4596		ENDIF		
4596	289D 23	INX	H	:STEP POINTER TO
4596	289E 23	INX	H	:THE BALANCE BYTE
4596	289F 23	INX	H	
4596	28A0 23	INX	H	
4596	28A1 77	MOV	M,A	:BALANCE CODE TO NODE
4596				
4597	28A2 2ACE42	LHLD	QW	:SAVE ADDRESS OF PARENTS
4598	28A5 22CC42	SHLD	QT	:REPLACEMENT NODE AND
4599	28A8 AF	XRA	A	:ALSO MARK IT AS A
4600		PTBAL	0,0	:BALANCED SUBTREE AND
4600		IF	00000H	:IS FACTOR IN A-REG?
4600		LDA	00000H	:NO - BALANCE TO A-REG
4600		ENDIF		
4600		IF	00000H	:IS NODE ADDRESS IN (H,L)
4600		LHLD	00000H	:NO - NODE POINTER TO (H,L)
4600		ENDIF		
4600	28A9 23	INX	H	:STEP POINTER TO
4600	28AA 23	INX	H	:THE BALANCE BYTE
4600	28AB 23	INX	H	
4600	28AC 23	INX	H	
4600	28AD 77	MOV	M,A	:BALANCE CODE TO NODE
4600				
4601	28AE CDF928	CALL	CSUBR	:ADJUST THE GRANDPARENT NODE LINK
4602	28B1 C30028	JMP	QLOOP	:TREE HEIGHT CHANGED - SO GO ON
4603				
4604	28B4 3EFF	MVI	A,0FFH	:PARENT WILL BE LEFT HEAVY
4605		PTBAL	QX,0	:SET PARENT BALANCE (0,-)
4605		IF	00000H	:IS FACTOR IN A-REG?
4605		LDA	00000H	:NO - BALANCE TO A-REG
4605		ENDIF		
4605		IF	QX	:IS NODE ADDRESS IN (H,L)
4605	28B6 2ACA42	LHLD	QX	:NO - NODE POINTER TO (H,L)
4605		ENDIF		
4605	28B9 23	INX	H	:STEP POINTER TO
4605	28BA 23	INX	H	:THE BALANCE BYTE
4605	28BB 23	INX	H	
4605	28BC 23	INX	H	
4605	28BD 77	MOV	M,A	:BALANCE CODE TO NODE
4605				
4606	28BE AF	XRA	A	:BUT ORIGINAL RIGHT SUBTREE
4607	28BF C39A28	JMP	CSSX	:OF PARENT WILL BE BALANCED
4608				
4609	28C2 AF	XRA	A	:REMOVED ITEM ON LEFT SO
4610	28C3 77	MOV	M,A	:PARENT IS NOW BALANCED
4611	28C4 E1	POP	H	:BUT WE MAY HAVE UPSET
4612	28C5 7C	MOV	A,H	:BALANCES ABOVE THIS NODE
4613	28C6 E1	POP	H	:SO BACKTRACK UP THE TREE
4614	28C7 22CA42	SHLD	QX	:SET NEW PARENT ADDRESS
4615	28CA C30328	JMP	QLRP	:AND GO ON
4616				
4617	28CD 3C	INR	A	:PARENT NOW HEAVY ON RIGHT
4618	28CE 77	MOV	M,A	:BUT SUBTREE HEIGHT UNCHANGED
4619	28CF C3F128	JMP	QDONE	:SO THE TOTAL TREE IS NOW OK
4620				
4621	28D2 3EFF	MVI	A,0FFH	:PARENTS REPLACEMENT WILL
4622		PTBAL	QT,0	:BE HEAVY TO THE LEFT
4622		IF	00000H	:IS FACTOR IN A-REG?
4622		LDA	00000H	:NO - BALANCE TO A-REG
4622		ENDIF		
4622		IF	QT	:IS NODE ADDRESS IN (H,L)
4622	28D4 2ACC42	LHLD	QT	:NO - NODE POINTER TO (H,L)
4622		ENDIF		
4622	28D7 23	INX	H	:STEP POINTER TO

4622	28D8 23		INX	H		;THE BALANCE BYTE
4622	28D9 23		INX	H		
4622	28DA 23		INX	H		
4622	28DB 77		MOV	M,A		;BALANCE CODE TO NODE
4622						
4623			PTRLK	QX,0		;SUBTREE ROOT LLINK TO PARENT RLINK
4623	0001	MCTEM	SET	1		;ASSUME ADDRESS IN (H,L)
4623			IF	QX		;TEST ASSUMPTION
4623	0000	MCTEM	SET	0		;FALSE ASSUMPTION
4623			ENDIF			
4623			IF	00000H		;IS NEW LINK IN (D,E)?
4623			IF	MCTEM		;NO - BUT IS POINTER IN (H,L)?
4623			XCHG			;YES - SAVE NODE POINTER
4623			ENDIF			
4623			LHLD	00000H		;NEW LINK TO (H,L)
4623			XCHG			;NEW LINK TO (D,E)
4623			ENDIF			
4623			IF	QX		;IS NODE ADDRESS IN (H,L)?
4623	28DC 2ACA42		LHLD	QX		;NO - NODE POINTER TO (H,L)
4623			ENDIF			
4623	28DF 23		INX	H		;NOW STEP TO THE
4623	28E0 23		INX	H		;RIGHT LINK BYTES
4623	28E1 73		MOV	M,E		;MOVE THE NEW
4623	28E2 23		INX	H		;RIGHT LINK VALUE
4623	28E3 72		MOV	M,D		;INTO THE NODE
4623						
4624			PTRLK	QT,QX		;ADD PARENT AS LEFT SUBTREE
4624	0001	MCTEM	SET	1		;ASSUME ADDRESS IN (H,L)
4624			IF	QT		;TEST ASSUMPTION
4624	0000	MCTEM	SET	0		;FALSE ASSUMPTION
4624			ENDIF			
4624			IF	QX		;IS NEW LINK IN (D,E)?
4624			IF	MCTEM		;NO - BUT IS POINTER IN (H,L)?
4624			XCHG			;YES - SAVE NODE POINTER
4624			ENDIF			
4624	28E4 2ACA42		LHLD	QX		;NEW LINK TO (H,L)
4624	28E7 EB		XCHG			;NEW LINK TO (D,E)
4624			ENDIF			
4624			IF	QT		;IS NODE ADDRESS IN (H,L)?
4624	28E8 2ACC42		LHLD	QT		;NO - NODE POINTER TO (H,L)
4624			ENDIF			
4624	28EB 73		MOV	M,E		;MOVE THE NEW
4624	28EC 23		INX	H		;LEFT LINK VALUE
4624	28ED 72		MOV	M,D		;INTO THE NODE
4624						
4625	28EE CDF928		CALL	CSUBR		;LINK NEW PARENT TO TREE
4626						
4627	28F1 2AC642	QDONE:	LHLD	STKSV		;PURGE THE STACK BACK
4628	28F4 F9		SPHL			;TO ITS ENTRY VALUE
4629	28F5 2AC842	QXIT:	LHLD	QP		;ADDRESS OF SELECTED NODE
4630	28F8 C9		RET			;RETURN
4631						
4632	28F9 2ACC42	CSUBR:	LHLD	QT		;GET ADDRESS OF OLD PARENT
4633	28FC EB		XCHG			;NODE REPLACEMENT AND
4634	28FD E1		POP	H		;POP THE ROUTINE RETURN THEN
4635	28FE E3		XTHL			;GET GRANDPARENTS EXIT DIRECTION
4636	28FF 7C		MOV	A,H		;SAVE IT
4637	2900 E1		POP	H		;THE RETURN AGAIN
4638	2901 E3		XTHL			;GET GRANDPARENT ADDRESS
4639	2902 22CA42		SHLD	QX		;AND SAVE IT
4640	2905 B7		ORA	A		;DIRECTION TO CONDITION CODE
4641	2906 FA0F29		JM	CS225		;WE CAME FROM THE LEFT
4642			PTRLK	0,0		;ADJUST THE RIGHT LINK
4642	0001	MCTEM	SET	1		;ASSUME ADDRESS IN (H,L)
4642			IF	00000H		;TEST ASSUMPTION
4642		MCTEM	SET	0		;FALSE ASSUMPTION
4642			ENDIF			
4642			IF	00000H		;IS NEW LINK IN (D,E)?
4642			IF	MCTEM		;NO - BUT IS POINTER IN (H,L)?
4642			XCHG			;YES - SAVE NODE POINTER
4642			ENDIF			
4642			LHLD	00000H		;NEW LINK TO (H,L)
4642			XCHG			;NEW LINK TO (D,E)
4642			ENDIF			
4642			IF	00000H		;IS NODE ADDRESS IN (H,L)?
4642			LHLD	00000H		;NO - NODE POINTER TO (H,L)
4642			ENDIF			
4642	2909 23		INX	H		;NOW STEP TO THE
4642	290A 23		INX	H		;RIGHT LINK BYTES
4642	290B 73		MOV	M,E		;MOVE THE NEW
4642	290C 23		INX	H		;RIGHT LINK VALUE
4642	290D 72		MOV	M,D		;INTO THE NODE
4642						
4643	290E C9		RET			;AND RETURN
4644						
4645		CS225:	PTRLK	0,0		;ADJUST THE LEFT LINK
4645	0001	MCTEM	SET	1		;ASSUME ADDRESS IN (H,L)

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4645          IF      00000H          ;TEST ASSUMPTION
4645 MCTEM      SET      0              ;FALSE ASSUMPTION
4645          ENDIF
4645          IF      00000H          ;IS NEW LINK IN (D,E)?
4645          IF      MCTEM          ;NO - BUT IS POINTER IN (H,L)?
4645          XCHG          ;YES - SAVE NODE POINTER
4645          ENDIF
4645          LHL      00000H          ;NEW LINK TO (H,L)
4645          XCHG          ;NEW LINK TO (D,E)
4645          ENDIF
4645          IF      00000H          ;IS NODE ADDRESS IN (H,L)?
4645          LHL      00000H          ;NO - NODE POINTER TO (H,L)
4645          ENDIF
4645 290F 73    MOV      M.E          ;MOVE THE NEW
4645 2910 23    INX      H            ;LEFT LINK VALUE
4645 2911 72    MOV      M.D          ;INTO THE NODE
4645
4646 2912 C9    RET                    ;AND RETURN
4647
4648          QRTY: PTRLK 0.0          ;ESTABLISH THE NEW ROOT
4648 0001 MCTEM  SET      1              ;ASSUME ADDRESS IN (H,L)
4648          IF      00000H          ;TEST ASSUMPTION
4648 MCTEM      SET      0              ;FALSE ASSUMPTION
4648          ENDIF
4648          IF      00000H          ;IS NEW LINK IN (D,E)?
4648          IF      MCTEM          ;NO - BUT IS POINTER IN (H,L)?
4648          XCHG          ;YES - SAVE NODE POINTER
4648          ENDIF
4648          LHL      00000H          ;NEW LINK TO (H,L)
4648          XCHG          ;NEW LINK TO (D,E)
4648          ENDIF
4648          IF      00000H          ;IS NODE ADDRESS IN (H,L)?
4648          LHL      00000H          ;NO - NODE POINTER TO (H,L)
4648          ENDIF
4648 2913 23    INX      H            ;NOW STEP TO THE
4648 2914 23    INX      H            ;RIGHT LINK BYTES
4648 2915 73    MOV      M.E          ;MOVE THE NEW
4648 2916 23    INX      H            ;RIGHT LINK VALUE
4648 2917 72    MOV      M.D          ;INTO THE NODE
4648
4649 2918 2B    DCX      H            ;NODE FOR THE TREE AND
4650 2919 2B    DCX      H            ;ADJUST THE POINTERS SO
4651 291A C3E627 JMP      QRTX          ;THAT WE MAY CONTINUE
4652
4653          ; SCHEDULE DIVERT ACTUATION EVENT AT UDBT + DVT(C).
4654          ;   CALLED BY SCDVA, DVRS, SDVRS.
4655
4656          ; INPUT: H,L POINTS TO BYTE 7 IN TCB, TREE PARAMETERS, UDBT, C=JDV.
4657          ; OUTPUT: EVENT IN TREE, UPDATED PPIDB.
4658          ; ALL REGISTERS EFFECTED.
4659
4660 291D EB    ENDVA: XCHG          ;POINT D,E TO BYTE 7 IN TCB
4661 291E 2A3942 LHL      DVT          ;POINT H,L TO DVT TABLE
4662 2921 0600   MVI      B.0          ;B,C=JDV
4663 2923 09    DAD      B            ;POINT H,L TO
4664 2924 09    DAD      B            ; DVT(JDV)
4665 2925 2B    DCX      H            ; MS BYTE
4666 2926 46    MOV      B.M          ;B,C
4667 2927 2B    DCX      H            ; =
4668 2928 4E    MOV      C.M          ; DVT(JDV)
4669 2929 2A9942 LHL      UDBT          ;H,L=UDBT
4670 292C C33A29 JMP      ENTC1
4671
4672          ; SCHEDULE EVENT AT CTIME + H,L(C).
4673
4674          ; INPUT: D,E POINTS TO BYTE 7 IN TCB, H,L POINTS TO TIME
4675          ;   TABLE, C ITS INDEX, TREE PARAMETERS.
4676          ; OUTPUT: EVENT IN TREE.
4677          ; ALL REGISTERS EFFECTED.
4678
4679 292F 0600   ENTCB: MVI      B.0          ;B,C= INDEX
4680 2931 09    DAD      B            ;POINT H,L
4681 2932 09    DAD      B            ; TO TIME
4682 2933 2B    DCX      H            ; MS BYTE
4683 2934 46    MOV      B.M          ;B,C
4684 2935 2B    DCX      H            ; = TABLE
4685 2936 4E    MOV      C.M          ; TIME
4686 2937 2AB642 LHL      CTIME          ;H,L= CTIME
4687
4688 293A 09    ENTC1: DAD      B            ;H,L= EVENT TIME
4689 293B EB    XCHG          ;D,E= EVTIM, POINT H,L TO
4690 293C 2B    DCX      H            ; BYTE 6 IN TCB
4691 293D 72    MOV      M.D          ;SET EVTIM
4692 293E 2B    DCX      H            ; IN
4693 293F 73    MOV      M.E          ; TCB
4694 2940 11FBFF LXI      D.0FFFBH      ;POINT H,L TO
4695 2943 19    DAD      D            ; TCB

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4696
4697
4698 ; INSERT TCB INTO EVENT TREE & UPDATE PPIDB.
4699
4700 ; INPUT: H.L POINTS TO TCB. TREE PARAMETERS FOR THIS PPI.
4701 ; PPIDP POINTS TO CTIME IN PPIDB (SEE EXEC. GTREP).
4702 ; ALL REGISTERS EFFECTED.
4703
4704 2944 CD5629 ENQ2: CALL ENQUE ; INSERT EVENT TCB
4705 2947 2AD842 LHL D PPIDP ; POINT D.E TO
4706 294A EB XCHG ; CTIME IN PPIDB
4707 294B 13 INX D ; SKIP PPI CLOCK TIME SINCE
4708 294C 13 INX D ; IT MAY HAVE BEEN CHANGED BY PPIIN
4709 294D 21B842 LXI H,HEAD ; POINT H.L TO TREE HEADER
4710 2950 060E MVI B,14 ; UPDATE
4711 2952 CDCF10 CALL MOVEB ; PPIDB
4712 2955 C9 RET
4713
4714 ; ADD A NEW NODE TO THE FUTURE EVENT TREE
4715
4716 2956 EB ENQUE: XCHG ; SAVE NEW NODE ADDRESS
4717 2957 2ABA42 LHL D RLINK ; TEST IF THE
4718 295A 7C MOV A,H ; TREE IS EMPTY
4719 295B B5 ORA L
4720 295C CAE22B JZ TREMT ; THE TREE IS EMPTY
4721 295F D5 PUSH D ; SAVE NODE ADDRESS
4722 2960 EB XCHG ; SAVE ADDRESS OF THE ROOT
4723 2961 21B842 LXI H,HEAD ; POINT TO THE TREE HEADER
4724 2964 22CC42 SHLD QT ; SET PARENT ADDRESS
4725 2967 EB XCHG ; GET BALANCE POINT NODE
4726 2968 22D042 SHLD QS ; AND SAVE IT
4727 296B AF XRA A ; MARK EXIT DIRECTION
4728 296C 32D442 STA QDIR ; AS RIGHT AND SAVE IT
4729 296F 22C842 SHLD QP ; SET CURRENT NODE ADDRESS
4730 2972 E1 POP H ; RECOVER THE ADDRESS
4731 2973 E5 PUSH H ; OF THE NEW NODE
4732 2974 CDEE2B CALL NDCLR ; CLEAR ITS LINKS & GET 'TIME'
4733 2977 2ABD42 LHL D NXTIM ; GET THE NEXT SCHEDULED
4734 297A EB XCHG ; EVENT TIME AND COMPARE
4735 297B 7B MOV A,E ; IT TO THE SCHEDULED TIME
4736 297C 95 SUB L ; IN THE NEW FUTURE EVENT
4737 297D 7A MOV A,D ; (NXTIME)-(NEW TIME)
4738 297E 9C SBB H
4739 ; ADD A ; DECISION BIT TO CONDITION CODE
4740 297F FA8529 JM ENQ10 ; NEW EVENT IS NOT NEXT
4741 2982 22BD42 SHLD NXTIM ; UPDATE NEXT EVENT TIME
4742 2985 22D642 ENQ10: SHLD NEWTM ; SAVE THIS EVENT TIME
4743 2988 2AC842 LHL D QP ; COMPARE THE NEW 'PPI'
4744 298B CDFD2B ENQ2: CALL CMPTM ; TIME WITH NODE VALUE
4745 298E F2CA29 JP ENQ4 ; MOVE TO RIGHT IN SUBTREE
4746
4746 ENQ3: GTLLK QP,QW ; MOVE TO THE LEFT
4746 IF QP ; IS NODE ADDRESS IN (H,L)?
4746 2991 2AC842 LHL D QP ; NO-NODE POINTER TO (H,L)
4746 ENDIF
4746 2994 5E MOV E,M ; MOVE THE LEFT
4746 2995 23 INX H ; LINK INTO THE
4746 2996 56 MOV D,M ; (D,E) REGISTER
4746 IF QW ; IS IT SAVE LINK MODE?
4746 2997 EB XCHG ; YES - LINK TO (H,L)
4746 2998 22CE42 SHLD QW ; LINK TO MEMORY
4746 ENDIF
4746
4747 299B 7C MOV A,H ; BUT TEST FOR A
4748 299C B5 ORA L ; NULL LEFT SUBTREE
4749 299D CAE929 JZ ENQ5 ; NULL - ADD NEW NODE
4750 29A0 16FF MVI D,0FFH ; MARK DIRECTION AS LEFT
4751
4751 ENQ3A: GTBAL 0,0 ; TEST THE BALANCE FACTOR
4751 IF 00000H ; IS NODE ADDRESS IN (H,L)?
4751 LHL D 00000H ; NO - NODE POINTER TO (H,L)
4751 ENDIF
4751 29A2 23 INX H ; STEP POINTER TO
4751 29A3 23 INX H ; THE BALANCE BYTE
4751 29A4 23 INX H
4751 29A5 23 INX H
4751 29A6 7E MOV A,M ; BALANCE CODE TO A-REG
4751 IF 00000H ; TEST IF SAVE NODE
4751 STA 00000H ; YES - CODE TO MEMORY
4751 ENDIF
4751 29A7 B7 ORA A ; SET CONDITION CODES
4751
4752 29A8 CAC129 JZ ENQ3B ; OF THIS NODE - BALANCED
4753 29AB 7A MOV A,D ; GET EXIT DIRECTION FROM
4754 29AC 32D442 STA QDIR ; THIS NODE'S PARENT AND SAVE
4755 29AF 2AC842 LHL D QP ; RECORD PARENT OF THE
4756 29B2 22CC42 SHLD QT ; NEW REBALANCE POINT
4757 29B5 2ACE42 LHL D QW ; MARK THE NEXT NODE
4758 29B8 22D042 SHLD QS ; AS REBALANCE POINT

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4759 29BB 22C842      SHLD  QP      ;AND AS CURRENT NODE
4760 29BE C38B29      JMP   ENQ2    ;FIND DIRECTION OF NEXT STEP
4761
;
4762 29C1 2ACE42      ENQ3B: LHL  QW      ;MARK THIS AS
4763 29C4 22C842      SHLD  QP      ; THE CURRENT NODE
4764 29C7 C38B29      JMP   ENQ2    ;FIND DIRECTION OF NEXT STEP
4765
;
4766                ENQ4:  GTRLK QP,QW      ;MOVE TO THE RIGHT
4766                IF     QP      ;IS NODE ADDRESS IN (H.L)?
4766 29CA 2AC842      LHL  QP      ;NO - NODE POINTER TO (H.L)
4766                ENDIF
4766 29CD 23          INX  H      ;STEP POINTER TO
4766 29CE 23          INX  H      ;THE RIGHT LINK
4766 29CF 5E          MOV  E,M     ;MOVE THE RIGHT
4766 29D0 23          INX  H      ;LINK INTO THE
4766 29D1 56          MOV  D,M     ;(D,E) REGISTER
4766                IF     QW      ;IS IT SAVE LINK MODE?
4766 29D2 EB          XCHG                ;YES - LINK TO (H.L)
4766 29D3 22CE42     SHLD  QW      ;LINK TO MEMORY
4766                ENDIF
4766
4767 29D6 7C          MOV  A,H     ;BUT TEST FOR A
4768 29D7 B5          ORA  L      ;NULL RIGHT SUBTREE
4769 29D8 1600        MVI  D,0     ;BUT MARK A STEP RIGHT
4770 29DA C2A229     JNZ  ENQ3A   ;CONTINUE IF NOT NULL SUBTREE
4771 29DD D1          POP  D      ;GET NEW NODE ADDRESS
4772                PTRLK QP,0     ;ATTACH AS RIGHT SUBTREE
4772 0001            MCTEM  SET  1      ;ASSUME ADDRESS IN (H.L)
4772                IF     QP      ;TEST ASSUMPTION
4772 0000            MCTEM  SET  0      ;FALSE ASSUMPTION
4772                ENDIF
4772                IF     00000H   ;IS NEW LINK IN (D,E)?
4772                IF     MCTEM    ;NO - BUT IS POINTER IN (H.L)?
4772                XCHG                ;YES - SAVE NODE POINTER
4772                ENDIF
4772                LHL  00000H   ;NEW LINK TO (H.L)
4772                XCHG                ;NEW LINK TO (D,E)
4772                ENDIF
4772                IF     QP      ;IS NODE ADDRESS IN (H.L)?
4772 29DE 2AC842     LHL  QP      ;NO - NODE POINTER TO (H.L)
4772                ENDIF
4772                INX  H      ;NOW STEP TO THE
4772 29E2 23          INX  H      ;RIGHT LINK BYTES
4772 29E3 73          MOV  M,E     ;MOVE THE NEW
4772 29E4 23          INX  H      ;RIGHT LINK VALUE
4772 29E5 72          MOV  M,D     ;INTO THE NODE
4772
4773 29E6 C3F029     JMP   ENQ6   ;IF NEEDED - BALANCE THE TREE
4774
;
4775 29E9 D1          ENQ5:  POP  D      ;GET NEW NODE ADDRESS
4776                PTLK  QP,0     ;ATTACH AS LEFT SUBTREE
4776 0001            MCTEM  SET  1      ;ASSUME ADDRESS IN (H.L)
4776                IF     QP      ;TEST ASSUMPTION
4776 0000            MCTEM  SET  0      ;FALSE ASSUMPTION
4776                ENDIF
4776                IF     00000H   ;IS NEW LINK IN (D,E)?
4776                IF     MCTEM    ;NO - BUT IS POINTER IN (H.L)?
4776                XCHG                ;YES - SAVE NODE POINTER
4776                ENDIF
4776                LHL  00000H   ;NEW LINK TO (H.L)
4776                XCHG                ;NEW LINK TO (D,E)
4776                ENDIF
4776                IF     QP      ;IS NODE ADDRESS IN (H,L)?
4776 29EA 2AC842     LHL  QP      ;NO - NODE POINTER TO (H.L)
4776                ENDIF
4776 29ED 73          MOV  M,E     ;MOVE THE NEW
4776 29EE 23          INX  H      ;LEFT LINK VALUE
4776 29EF 72          MOV  M,D     ;INTO THE NODE
4776
;
4777                ;BALANCE TREE
4778
;
4779
4780 29F0 EB          ENQ6:  XCHG                ;NOW SAVE THE ADDRESS
4781 29F1 22CE42     SHLD  QW      ;OF THE ADDED NODE
4782 29F4 2AD042     LHL  QS      ;DETERMINE IF LEFT OR
4783 29F7 CDFD2B     CALL CMPTM   ;RIGHT SUBTREE TO BE
4784 29FA F5          PUSH PSW     ;ADJUSTED - SAVE INDICATOR
4785 29FB 2AD042     LHL  QS      ;POINT TO REBALANCE POINT
4786 29FE FA092A     JM   ENGA1   ;TEST FOR LEFT SUBTREE
4787                GTRLK 0,0     ;IT IS RIGHT SUBTREE
4787                IF     00000H   ;IS NODE ADDRESS IN (H,L)?
4787                LHL  00000H   ;NO - NODE POINTER TO (H.L)
4787                ENDIF
4787 2A01 23          INX  H      ;STEP POINTER TO
4787 2A02 23          INX  H      ;THE RIGHT LINK
4787 2A03 5E          MOV  E,M     ;MOVE THE RIGHT
4787 2A04 23          INX  H      ;LINK INTO THE

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4787 2A05 56      MOV    D.M      ;(D,E) REGISTER
4787             IF    00000H      ;IS IT SAVE LINK MODE?
4787             XCHG             ;YES - LINK TO (H,L)
4787             SHLD    00000H      ;LINK TO MEMORY
4787             ENDIF
4788 2A06 C30C2A   JMP     EN6A2      ;GO TO COMMON PATH
4789
4790             ;
4790             EN6A1: GTLLK  0,0      ;IT IS LEFT SUBTREE
4790             IF    00000H      ;IS NODE ADDRESS IN (H,L)?
4790             LHL  00000H      ;NO-NODE POINTER TO (H,L)
4790             ENDIF
4790 2A09 5E      MOV    E.M      ;MOVE THE LEFT
4790 2A0A 23      INX    H        ;LINK INTO THE
4790 2A0B 56      MOV    D.M      ;(D,E) REGISTER
4790             IF    00000H      ;IS IT SAVE LINK MODE?
4790             XCHG             ;YES - LINK TO (H,L)
4790             SHLD    00000H      ;LINK TO MEMORY
4790             ENDIF
4791 2A0C EB      EN6A2: XCHG             ;GET START OF CHAIN
4792 2A0D 22C842  SHLD   QP      ;INITIALIZE THE
4793 2A10 22D242  SHLD   QR      ;PATH POINTERS
4794 2A13 EB      EN6A3: XCHG             ;SAVE NEXT NODE ADDRESS
4795 2A14 2ACE42 LHL    QW      ;BUT TEST IF IT IS
4796 2A17 7D      MOV    A,L      ;THE NEWLY ADDED NODE
4797 2A18 AB      XRA     E
4798 2A19 C2212A JNZ    EN6A4      ;NO
4799 2A1C 7C      MOV    A,H      ;MAYBE - COMPARE THE
4800 2A1D AA      XRA     D        ;SECOND BYTES
4801 2A1E CA532A JZ     ENQ7      ;OK - BALANCE FACTORS ADJUSTED
4802 2A21 EB      EN6A4: XCHG             ;NEW NODE ADDRESS TO (H,L)
4803 2A22 CDFD2B CALL   CMPTM     ;COMPARE THE KEYS
4804 2A25 2AC842 LHL    QP      ;AGAIN POINT TO THE NODE
4805 2A28 FA412A JM     EN6A5     ;WE MUST CHASE LEFTWARDS
4806 2A2B 3E01   MVI    A,01     ;SET BALANCE FACTOR TO +
4807             PTBAL  0,0      ;AND ADJUST THIS NODE
4807             IF    00000H      ;IS FACTOR IN A-REG?
4807             LDA    00000H      ;NO - BALANCE TO A-REG
4807             ENDIF
4807             IF    00000H      ;IS NODE ADDRESS IN (H,L)
4807             LHL  00000H      ;NO - NODE POINTER TO (H,L)
4807             ENDIF
4807 2A2D 23      INX    H        ;STEP POINTER TO
4807 2A2E 23      INX    H        ;THE BALANCE BYTE
4807 2A2F 23      INX    H
4807 2A30 23      INX    H
4807 2A31 77      MOV    M,A      ;BALANCE CODE TO NODE
4807
4808             GTRLK  QP,QP     ;MOVE TO THE RIGHT
4808             IF    QP      ;IS NODE ADDRESS IN (H,L)?
4808 2A32 2AC842  LHL    QP      ;NO - NODE POINTER TO (H,L)
4808             ENDIF
4808 2A35 23      INX    H        ;STEP POINTER TO
4808 2A36 23      INX    H        ;THE RIGHT LINK
4808 2A37 5E      MOV    E.M      ;MOVE THE RIGHT
4808 2A38 23      INX    H        ;LINK INTO THE
4808 2A39 56      MOV    D.M      ;(D,E) REGISTER
4808             IF    QP      ;IS IT SAVE LINK MODE?
4808             XCHG             ;YES - LINK TO (H,L)
4808 2A3A EB      SHLD   QP      ;LINK TO MEMORY
4808 2A3B 22C842  ENDIF
4808
4809 2A3E C3132A   JMP     EN6A3      ;CONTINUE
4810
4811             ;
4812             ;
4813             EN6A5: PTBAL  0,0      ;SET NEW BALANCE AS (-)
4813             IF    00000H      ;IS FACTOR IN A-REG?
4813             LDA    00000H      ;NO - BALANCE TO A-REG
4813             ENDIF
4813             IF    00000H      ;IS NODE ADDRESS IN (H,L)
4813             LHL  00000H      ;NO - NODE POINTER TO (H,L)
4813             ENDIF
4813 2A41 23      INX    H        ;STEP POINTER TO
4813 2A42 23      INX    H        ;THE BALANCE BYTE
4813 2A43 23      INX    H
4813 2A44 23      INX    H
4813 2A45 77      MOV    M,A      ;BALANCE CODE TO NODE
4813
4814             GTLLK  QP,QP     ;MOVE TO THE LEFT
4814             IF    QP      ;IS NODE ADDRESS IN (H,L)?
4814 2A46 2AC842  LHL    QP      ;NO-NODE POINTER TO (H,L)
4814             ENDIF
4814 2A49 5E      MOV    E.M      ;MOVE THE LEFT
4814 2A4A 23      INX    H        ;LINK INTO THE
4814 2A4B 56      MOV    D.M      ;(D,E) REGISTER

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4814          IF      QP          ;IS IT SAVE LINK MODE?
4814 2A4C EB      XCHG          ;YES - LINK TO (H,L)
4814 2A4D 22C842 SHLD      QP          ;LINK TO MEMORY
4814          ENDIF
4814
4815 2A50 C3132A JMP      EN6A3          ;CONTINUE
4816
4817 2A53 F1      ENQ7: POP      PSW          ;RECOVER REBALANCE DIRECTION
4818 2A54 FA072B JM       EN7AM          ;WE ARE ADJUSTING A LEFT SUBTREE
4819          GTBAL  QS.0       ;WAS REBALANCE POINT BALANCED?
4819          IF      QS          ;IS NODE ADDRESS IN (H,L)?
4819 2A57 2AD042 LHL D  QS          ;NO - NODE POINTER TO (H,L)
4819          ENDIF
4819 2A5A 23      INX      H          ;STEP POINTER TO
4819 2A5B 23      INX      H          ;THE BALANCE BYTE
4819 2A5C 23      INX      H
4819 2A5D 23      INX      H
4819 2A5E 7E      MOV      A,M          ;BALANCE CODE TO A-REG
4819          IF      00000H      ;TEST IF SAVE NODE
4819          STA     00000H      ;YES - CODE TO MEMORY
4819          ENDIF
4819 2A5F B7      ORA      A          ;SET CONDITION CODES
4819
4820 2A60 C2672A JNZ     EN7P1          ;NO
4821 2A63 3E01     MVI     A,01      ;YES - SET BALANCE FACTOR TO (+)
4822 2A65 77      MOV      M,A          ;BECAUSE THE TREE IS HIGHER
4823 2A66 C9      RET
4824          ;WE ARE DONE
4825 2A67 F26D2A EN7P1: JP      EN7P2          ;SUBTREE WAS HEAVY TO RIGHT
4826 2A6A AF      EN7X1: XRA     A          ;IT WAS HEAVY TO THE LEFT
4827 2A6B 77      MOV      M,A          ;NOW IT IS BALANCED
4828 2A6C C9      RET
4829          ;WE ARE DONE
4830          EN7P2: GTBAL  QR.0       ;TREE TOO FAR OUT OF BALANCE
4830          IF      QR          ;IS NODE ADDRESS IN (H,L)?
4830 2A6D 2AD242 LHL D  QR          ;NO - NODE POINTER TO (H,L)
4830          ENDIF
4830 2A70 23      INX      H          ;STEP POINTER TO
4830 2A71 23      INX      H          ;THE BALANCE BYTE
4830 2A72 23      INX      H
4830 2A73 23      INX      H
4830 2A74 7E      MOV      A,M          ;BALANCE CODE TO A-REG
4830          IF      00000H      ;TEST IF SAVE NODE
4830          STA     00000H      ;YES - CODE TO MEMORY
4830          ENDIF
4830 2A75 B7      ORA      A          ;SET CONDITION CODES
4830
4831 2A76 FA972A JM      EN7P9          ;DETERMINE ADJUSTMENT NEEDS
4832 2A79 2AD242 LHL D  QR          ;SAVE ADDRESS OF ROOT
4833 2A7C 22C842 SHLD      QP          ;OF THE RIGHT SUBTREE
4834          GTLLK  0.0         ;ATTACH ITS LEFT SUBTREE AS
4834          IF      00000H      ;IS NODE ADDRESS IN (H,L)?
4834          LHL D  00000H      ;NO-NODE POINTER TO (H,L)
4834          ENDIF
4834 2A7F 5E      MOV      E,M          ;MOVE THE LEFT
4834 2A80 23      INX      H          ;LINK INTO THE
4834 2A81 56      MOV      D,M          ; (D,E) REGISTER
4834          IF      00000H      ;IS IT SAVE LINK MODE?
4834          XCHG          ;YES - LINK TO (H,L)
4834          SHLD  00000H      ;LINK TO MEMORY
4834          ENDIF
4834
4835          PTRLK  QS.0       ;THE RIGHT SUBTREE OF REBALANCE POINT
4835 0001 MCTEM SET     1          ;ASSUME ADDRESS IN (H,L)
4835          IF      QS          ;TEST ASSUMPTION
4835 0000 MCTEM SET     0          ;FALSE ASSUMPTION
4835          ENDIF
4835          IF      00000H      ;IS NEW LINK IN (D,E)?
4835          IF      MCTEM      ;NO - BUT IS POINTER IN (H,L)?
4835          XCHG          ;YES - SAVE NODE POINTER
4835          ENDIF
4835          LHL D  00000H      ;NEW LINK TO (H,L)
4835          XCHG          ;NEW LINK TO (D,E)
4835          ENDIF
4835          IF      QS          ;IS NODE ADDRESS IN (H,L)?
4835 2A82 2AD042 LHL D  QS          ;NO - NODE POINTER TO (H,L)
4835          ENDIF
4835 2A85 23      INX      H          ;NOW STEP TO THE
4835 2A86 23      INX      H          ;RIGHT LINK BYTES
4835 2A87 73      MOV      M,E          ;MOVE THE NEW
4835 2A88 23      INX      H          ;RIGHT LINK VALUE
4835 2A89 72      MOV      M,D          ;INTO THE NODE
4835
4836          PTLK  QR.QS       ;BALANCE POINT BECOMES LEFT SUBTREE
4836 0001 MCTEM SET     1          ;ASSUME ADDRESS IN (H,L)
4836          IF      QR          ;TEST ASSUMPTION
4836 0000 MCTEM SET     0          ;FALSE ASSUMPTION

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4836                                     ENDF
4836 IF QS ;IS NEW LINK IN (D,E)?
4836 IF MCTEM ;NO - BUT IS POINTER IN (H,L)?
4836 XCHG ;YES - SAVE NODE POINTER
4836 ENDF
4836 2A8A 2AD042 LHLQ QS ;NEW LINK TO (H,L)
4836 2A8D EB XCHG ;NEW LINK TO (D,E)
4836 ENDF
4836 IF QR ;IS NODE ADDRESS IN (H,L)?
4836 2A8E 2AD242 LHLQ QR ;NO - NODE POINTER TO (H,L)
4836 ENDF
4836 2A91 73 MOV M,E ;MOVE THE NEW
4836 2A92 23 INX H ;LEFT LINK VALUE
4836 2A93 72 MOV M,D ;INTO THE NODE
4836
4837 2A94 C3422B JMP EN7T5 ;RESET BOTH BALANCE FACTORS
4838
4839
4840
4841 EN7P9: GTLLK QR,QP ;MOVE RIGHT THEN LEFT FROM BALANCE POINT
4841 IF QR ;IS NODE ADDRESS IN (H,L)?
4841 2A97 2AD242 LHLQ QR ;NO-NODE POINTER TO (H,L)
4841 ENDF
4841 2A9A 5E MOV E,M ;MOVE THE LEFT
4841 2A9B 23 INX H ;LINK INTO THE
4841 2A9C 56 MOV D,M ;(D,E) REGISTER
4841 IF QP ;IS IT SAVE LINK MODE?
4841 2A9D EB XCHG ;YES - LINK TO (H,L)
4841 2A9E 22C842 SHLD QP ;LINK TO MEMORY
4841 ENDF
4842
4842 GTRLK 0,0 ;THIS CODE EFFECTS A CLOCKWISE ROTATION
4842 IF 00000H ;IS NODE ADDRESS IN (H,L)?
4842 LHLQ 00000H ;NO - NODE POINTER TO (H,L)
4842 ENDF
4842 2AA1 23 INX H ;STEP POINTER TO
4842 2AA2 23 INX H ;THE RIGHT LINK
4842 2AA3 5E MOV E,M ;MOVE THE RIGHT
4842 2AA4 23 INX H ;LINK INTO THE
4842 2AA5 56 MOV D,M ;(D,E) REGISTER
4842 IF 00000H ;IS IT SAVE LINK MODE?
4842 XCHG ;YES - LINK TO (H,L)
4842 SHLD 00000H ;LINK TO MEMORY
4842 ENDF
4842
4843 PTLK QR,0 ;OF THE RIGHT SUBTREE FOLLOWED BY
4843 0001 MCTEM SET 1 ;ASSUME ADDRESS IN (H,L)
4843 IF QR ;TEST ASSUMPTION
4843 0000 MCTEM SET 0 ;FALSE ASSUMPTION
4843 ENDF
4843 IF 00000H ;IS NEW LINK IN (D,E)?
4843 IF MCTEM ;NO - BUT IS POINTER IN (H,L)?
4843 XCHG ;YES - SAVE NODE POINTER
4843 ENDF
4843 LHLQ 00000H ;NEW LINK TO (H,L)
4843 XCHG ;NEW LINK TO (D,E)
4843 ENDF
4843 IF QR ;IS NODE ADDRESS IN (H,L)?
4843 2AA6 2AD242 LHLQ QR ;NO - NODE POINTER TO (H,L)
4843 ENDF
4843 2AA9 73 MOV M,E ;MOVE THE NEW
4843 2AAA 23 INX H ;LEFT LINK VALUE
4843 2AAB 72 MOV M,D ;INTO THE NODE
4843
4844 PTRLK QP,QR ;A COUNTER CLOCKWISE ROTATION OF
4844 0001 MCTEM SET 1 ;ASSUME ADDRESS IN (H,L)
4844 IF QP ;TEST ASSUMPTION
4844 0000 MCTEM SET 0 ;FALSE ASSUMPTION
4844 ENDF
4844 IF QR ;IS NEW LINK IN (D,E)?
4844 IF MCTEM ;NO - BUT IS POINTER IN (H,L)?
4844 XCHG ;YES - SAVE NODE POINTER
4844 ENDF
4844 2AAC 2AD242 LHLQ QR ;NEW LINK TO (H,L)
4844 2AAF EB XCHG ;NEW LINK TO (D,E)
4844 ENDF
4844 IF QP ;IS NODE ADDRESS IN (H,L)?
4844 2AB0 2AC842 LHLQ QP ;NO - NODE POINTER TO (H,L)
4844 ENDF
4844 2AB3 23 INX H ;NOW STEP TO THE
4844 2AB4 23 INX H ;RIGHT LINK BYTES
4844 2AB5 73 MOV M,E ;MOVE THE NEW
4844 2AB6 23 INX H ;RIGHT LINK VALUE
4844 2AB7 72 MOV M,D ;INTO THE NODE
4844
4845 GTLLK QP,0 ;THE TREE. IN THIS CONTEXT THE
4845 IF QP ;IS NODE ADDRESS IN (H,L)?

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4845	ZAB8 2AC842	LHLD QP	;NO-NODE POINTER TO (H,L)
4845		ENDIF	
4845	2ABB 5E	MOV E,M	;MOVE THE LEFT
4845	2ABC 23	INX H	;LINK INTO THE
4845	2ABD 56	MOV D,M	; (D,E) REGISTER
4845		IF 00000H	; IS IT SAVE LINK MODE?
4845		XCHG	; YES - LINK TO (H,L)
4845		SHLD 00000H	; LINK TO MEMORY
4845		ENDIF	
4846		PTRLK QS.0	; REBALANCE POINT IS
4846	0001 MCTEM	SET 1	; ASSUME ADDRESS IN (H,L)
4846		IF QS	; TEST ASSUMPTION
4846	0000 MCTEM	SET 0	; FALSE ASSUMPTION
4846		ENDIF	
4846		IF 00000H	; IS NEW LINK IN (D,E)?
4846		IF MCTEM	; NO - BUT IS POINTER IN (H,L)?
4846		XCHG	; YES - SAVE NODE POINTER
4846		ENDIF	
4846		LHLD 00000H	; NEW LINK TO (H,L)
4846		XCHG	; NEW LINK TO (D,E)
4846		ENDIF	
4846		IF QS	; IS NODE ADDRESS IN (H,L)?
4846	2ABE 2AD042	LHLD QS	; NO - NODE POINTER TO (H,L)
4846		ENDIF	
4846	2AC1 23	INX H	; NOW STEP TO THE
4846	2AC2 23	INX H	; RIGHT LINK BYTES
4846	2AC3 73	MOV M,E	; MOVE THE NEW
4846	2AC4 23	INX H	; RIGHT LINK VALUE
4846	2AC5 72	MOV M,D	; INTO THE NODE
4846			
4847		PTRLK QP,QS	; CONSIDERED TO BE THE ROOT
4847	0001 MCTEM	SET 1	; ASSUME ADDRESS IN (H,L)
4847		IF QP	; TEST ASSUMPTION
4847	0000 MCTEM	SET 0	; FALSE ASSUMPTION
4847		ENDIF	
4847		IF QS	; IS NEW LINK IN (D,E)?
4847		IF MCTEM	; NO - BUT IS POINTER IN (H,L)?
4847		XCHG	; YES - SAVE NODE POINTER
4847		ENDIF	
4847	2AC6 2AD042	LHLD QS	; NEW LINK TO (H,L)
4847	2AC9 EB	XCHG	; NEW LINK TO (D,E)
4847		ENDIF	
4847		IF QP	; IS NODE ADDRESS IN (H,L)?
4847	2ACA 2AC842	LHLD QP	; NO - NODE POINTER TO (H,L)
4847		ENDIF	
4847	2ACD 73	MOV M,E	; MOVE THE NEW
4847	2ACE 23	INX H	; LEFT LINK VALUE
4847	2ACF 72	MOV M,D	; INTO THE NODE
4847			
4848	2AD0 23	INX H	; STEP (H,L) TO
4849	2AD1 23	INX H	; POINT AT THE BALANCE
4850	2AD2 23	INX H	; FACTOR OF THE NEW ROOT
4851	2AD3 56	MOV D,M	; GET THE ORIGINAL
4852	2AD4 AF	XRA A	; BALANCE FACTOR AND
4853	2AD5 77	MOV M,A	; MARK THE ROOT BALANCED
4854	2AD6 B2	ORA D	; TEST ORIGINAL BALANCE FACTOR
4855	2AD7 CA422B	JZ EN7T5	; BALANCED - SO SUBTREES ARE BALANCED
4856	2ADA FAF22A	JM EN7P8	; IT WAS HEAVY TO THE LEFT
4857	2ADD AF	XRA A	; IT WAS HEAVY TO THE RIGHT
4858		PTBAL QR.0	; SET BALANCE FACTORS
4858		IF 00000H	; IS FACTOR IN A-REG?
4858		LDA 00000H	; NO - BALANCE TO A-REG
4858		ENDIF	
4858		IF QR	; IS NODE ADDRESS IN (H,L)
4858	2ADE 2AD242	LHLD QR	; NO - NODE POINTER TO (H,L)
4858		ENDIF	
4858	2AE1 23	INX H	; STEP POINTER TO
4858	2AE2 23	INX H	; THE BALANCE BYTE
4858	2AE3 23	INX H	
4858	2AE4 23	INX H	
4858	2AE5 77	MOV M,A	; BALANCE CODE TO NODE
4858			
4859	2AE6 2F	CMA	; FOR THE SUBTREES THAT
4860		PTBAL QS.0	; WERE ROTATED DURING ADJUSTMENT
4860		IF 00000H	; IS FACTOR IN A-REG?
4860		LDA 00000H	; NO - BALANCE TO A-REG
4860		ENDIF	
4860		IF QS	; IS NODE ADDRESS IN (H,L)
4860	2AE7 2AD042	LHLD QS	; NO - NODE POINTER TO (H,L)
4860		ENDIF	
4860	2AEA 23	INX H	; STEP POINTER TO
4860	2AEB 23	INX H	; THE BALANCE BYTE
4860	2AEC 23	INX H	
4860	2AED 23	INX H	
4860	2AEE 77	MOV M,A	; BALANCE CODE TO NODE
4860			

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4861 2AEF C3C32B      JMP      NQA10      ;ADJUST LINKAGE TO TOTAL TREE
4862
4863 2AF2 AF          EN7P8:  XRA      A          ;SET THE BALANCE FACTORS
4864                   PTBAL   QS.0      ;FOR THE SUBTREES THAT WERE
4864                   IF      00000H    ;IS FACTOR IN A-REG?
4864                   LDA      00000H    ;NO - BALANCE TO A-REG
4864                   ENDIF
4864                   IF      QS          ;IS NODE ADDRESS IN (H,L)
4864 2AF3 2AD042      LHLD   QS          ;NO - NODE POINTER TO (H,L)
4864                   ENDIF
4864                   INX      H          ;STEP POINTER TO
4864 2AF6 23          INX      H          ;THE BALANCE BYTE
4864 2AF7 23          INX      H
4864 2AF8 23          INX      H
4864 2AF9 23          INX      H
4864 2AFA 77          MOV      M,A        ;BALANCE CODE TO NODE
4864
4865 2AFB 3C          INR      A          ;ROTATED DURING ADJUSTMENT
4866                   PTBAL   QR.0
4866                   IF      00000H    ;IS FACTOR IN A-REG?
4866                   LDA      00000H    ;NO - BALANCE TO A-REG
4866                   ENDIF
4866                   IF      QR          ;IS NODE ADDRESS IN (H,L)
4866 2AFC 2AD242      LHLD   QR          ;NO - NODE POINTER TO (H,L)
4866                   ENDIF
4866                   INX      H          ;STEP POINTER TO
4866 2AFF 23          INX      H          ;THE BALANCE BYTE
4866 2B00 23          INX      H
4866 2B01 23          INX      H
4866 2B02 23          INX      H
4866 2B03 77          MOV      M,A        ;BALANCE CODE TO NODE
4866
4867 2B04 C3C32B      JMP      NQA10      ;ADJUST LINKAGE TO TOTAL TREE
4868
4869                   EN7AM:  GTBAL   QS.0      ;WAS REBALANCE POINT BALANCED?
4869                   IF      QS          ;IS NODE ADDRESS IN (H,L)?
4869 2B07 2AD042      LHLD   QS          ;NO - NODE POINTER TO (H,L)
4869                   ENDIF
4869                   INX      H          ;STEP POINTER TO
4869 2B0A 23          INX      H          ;THE BALANCE BYTE
4869 2B0B 23          INX      H
4869 2B0C 23          INX      H
4869 2B0D 23          INX      H
4869 2B0E 7E          MOV      A,M        ;BALANCE CODE TO A-REG
4869                   IF      00000H    ;TEST IF SAVE NODE
4869                   STA      00000H    ;YES - CODE TO MEMORY
4869                   ENDIF
4869 2B0F B7          ORA      A          ;SET CONDITION CODES
4869
4870 2B10 C2162B      JNZ      EN7M1      ;NO
4871 2B13 2F          CMA          ;YES - SET BALANCE FACTOR TO (-)
4872 2B14 77          MOV      M,A        ;BECAUSE THE TREE IS HIGHER.
4873 2B15 C9          RET          ;WE ARE DONE
4874
4875 2B16 F26A2A      EN7M1:  JP      EN7X1      ;SUBTREE WAS HEAVY TO RIGHT
4876                   GTBAL   QR.0      ;TREE TOO FAR OUT OF BALANCE
4876                   IF      QR          ;IS NODE ADDRESS IN (H,L)?
4876 2B19 2AD242      LHLD   QR          ;NO - NODE POINTER TO (H,L)
4876                   ENDIF
4876                   INX      H          ;STEP POINTER TO
4876 2B1D 23          INX      H          ;THE BALANCE BYTE
4876 2B1E 23          INX      H
4876 2B1F 23          INX      H
4876 2B20 7E          MOV      A,M        ;BALANCE CODE TO A-REG
4876                   IF      00000H    ;TEST IF SAVE NODE
4876                   STA      00000H    ;YES - CODE TO MEMORY
4876                   ENDIF
4876 2B21 B7          ORA      A          ;SET CONDITION CODES
4876
4877 2B22 F2562B      JP      EN7M9      ;DETERMINE ADJUSTMENT NEEDS
4878 2B25 2AD242      LHLD   QR          ;SAVE ADDRESS OF ROOT
4879 2B28 22C842      SHLD   OP          ;OF THE LEFT SUBTREE
4880                   GTRLK   0.0      ;ATTACH ITS RIGHT SUBTREE AS
4880                   IF      00000H    ;IS NODE ADDRESS IN (H,L)?
4880                   LHLD   00000H    ;NO - NODE POINTER TO (H,L)
4880                   ENDIF
4880                   INX      H          ;STEP POINTER TO
4880 2B2B 23          INX      H          ;THE RIGHT LINK
4880 2B2C 23          INX      H
4880 2B2D 5E          MOV      E,M        ;MOVE THE RIGHT
4880 2B2E 23          INX      H          ;LINK INTO THE
4880 2B2F 56          MOV      D,M        ;(D,E) REGISTER
4880                   IF      00000H    ;IS IT SAVE LINK MODE?
4880                   XCHG          ;YES - LINK TO (H,L)
4880                   SHLD   00000H    ;LINK TO MEMORY
4880                   ENDIF
4880
4881                   PTLK   QS.0      ;THE LEFT SUBTREE OF REBALANCE POINT
4881 0001           MCTEM  SET      1          ;ASSUME ADDRESS IN (H,L)
4881                   IF      QS          ;TEST ASSUMPTION
4881 0000           MCTEM  SET      0          ;FALSE ASSUMPTION

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4881      ENDIF
4881      IF      00000H      ;IS NEW LINK IN (D,E)?
4881      IF      MCTEM      ;NO - BUT IS POINTER IN (H,L)?
4881      XCHG                    ;YES - SAVE NODE POINTER
4881      ENDIF
4881      LHL D 00000H      ;NEW LINK TO (H,L)
4881      XCHG                    ;NEW LINK TO (D,E)
4881      ENDIF
4881      IF      QS          ;IS NODE ADDRESS IN (H,L)?
4881      LHL D 2B30 2AD042 ;NO - NODE POINTER TO (H,L)
4881      ENDIF
4881      2B33 73      MOV      M,E      ;MOVE THE NEW
4881      2B34 23      INX      H          ;LEFT LINK VALUE
4881      2B35 72      MOV      M,D      ;INTO THE NODE
4881
4882      PTRLK  QR, QS      ;BALANCE POINT BECOMES RIGHT SUBTREE
4882      0001      MCTEM    SET      1          ;ASSUME ADDRESS IN (H,L)
4882      IF      QR          ;TEST ASSUMPTION
4882      0000      MCTEM    SET      0          ;FALSE ASSUMPTION
4882      ENDIF
4882      IF      QS          ;IS NEW LINK IN (D,E)?
4882      IF      MCTEM      ;NO - BUT IS POINTER IN (H,L)?
4882      XCHG                    ;YES - SAVE NODE POINTER
4882      ENDIF
4882      LHL D 2B36 2AD042 ;NEW LINK TO (H,L)
4882      2B39 EB      XCHG                    ;NEW LINK TO (D,E)
4882      ENDIF
4882      IF      QR          ;IS NODE ADDRESS IN (H,L)?
4882      LHL D 2B3A 2AD242 ;NO - NODE POINTER TO (H,L)
4882      ENDIF
4882      INX      H          ;NOW STEP TO THE
4882      2B3E 23      INX      H          ;RIGHT LINK BYTES
4882      2B3F 73      MOV      M,E      ;MOVE THE NEW
4882      2B40 23      INX      H          ;RIGHT LINK VALUE
4882      2B41 72      MOV      M,D      ;INTO THE NODE
4882
4883
4884      2B42 AF      EN7T5: XRA      A          ;TOTAL TREE IS BALANCED - SO
4885      PTBAL  QS, 0      ;RESET THE BALANCE FACTORS
4885      IF      00000H      ;IS FACTOR IN A-REG?
4885      LDA      00000H      ;NO - BALANCE TO A-REG
4885      ENDIF
4885      IF      QS          ;IS NODE ADDRESS IN (H,L)
4885      LHL D 2B43 2AD042 ;NO - NODE POINTER TO (H,L)
4885      ENDIF
4885      INX      H          ;STEP POINTER TO
4885      2B46 23      INX      H          ;THE BALANCE BYTE
4885      2B47 23      INX      H
4885      2B48 23      INX      H
4885      2B49 23      INX      H
4885      2B4A 77      MOV      M,A          ;BALANCE CODE TO NODE
4885
4886      PTBAL  QR, 0      ;IN THE AFFECTED NODES
4886      IF      00000H      ;IS FACTOR IN A-REG?
4886      LDA      00000H      ;NO - BALANCE TO A-REG
4886      ENDIF
4886      IF      QR          ;IS NODE ADDRESS IN (H,L)
4886      LHL D 2B4B 2AD242 ;NO - NODE POINTER TO (H,L)
4886      ENDIF
4886      INX      H          ;STEP POINTER TO
4886      2B4E 23      INX      H          ;THE BALANCE BYTE
4886      2B4F 23      INX      H
4886      2B50 23      INX      H
4886      2B51 23      INX      H
4886      2B52 77      MOV      M,A          ;BALANCE CODE TO NODE
4886
4887      2B53 C3C32B      JMP      NQA10      ;ADJUST LINKAGE TO TOTAL TREE
4887
4888
4889      EN7M9: GTRLK  QR, QP      ;MOVE LEFT THEN RIGHT FROM BALANCE POINT
4889      IF      QR          ;IS NODE ADDRESS IN (H,L)?
4889      LHL D 2B56 2AD242 ;NO - NODE POINTER TO (H,L)
4889      ENDIF
4889      INX      H          ;STEP POINTER TO
4889      2B59 23      INX      H          ;THE RIGHT LINK
4889      2B5A 23      INX      H
4889      2B5B 5E      MOV      E,M      ;MOVE THE RIGHT
4889      2B5C 23      INX      H          ;LINK INTO THE
4889      2B5D 56      MOV      D,M      ;(D,E) REGISTER
4889      IF      QP          ;IS IT SAVE LINK MODE?
4889      2B5E EB      XCHG                    ;YES - LINK TO (H,L)
4889      2B5F 22C842      SHLD  QP          ;LINK TO MEMORY
4889      ENDIF
4889
4890      GTLLK  0, 0          ;THIS CODE EFFECTS A COUNTER CLOCKWISE
4890      IF      00000H      ;IS NODE ADDRESS IN (H,L)?
4890      LHL D 00000H      ;NO-NODE POINTER TO (H,L)
4890      ENDIF
4890      2B62 5E      MOV      E,M      ;MOVE THE LEFT
4890      2B63 23      INX      H          ;LINK INTO THE
4890      2B64 56      MOV      D,M      ;(D,E) REGISTER

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4890          IF      00000H          ;IS IT SAVE LINK MODE?
4890          XCHG                    ;YES - LINK TO (H,L)
4890          SHLD   00000H          ;LINK TO MEMORY
4890          ENDIF
4891          PTRLK  QR.0            ;ROTATION OF THE LEFT SUBTREE FOLLOWED
4891  0001      MCTEM  SET      1            ;ASSUME ADDRESS IN (H,L)
4891          IF      QR            ;TEST ASSUMPTION
4891  0000      MCTEM  SET      0            ;FALSE ASSUMPTION
4891          ENDIF
4891          IF      00000H          ;IS NEW LINK IN (D,E)?
4891          IF      MCTEM          ;NO - BUT IS POINTER IN (H,L)?
4891          XCHG                    ;YES - SAVE NODE POINTER
4891          ENDIF
4891          LHL   00000H          ;NEW LINK TO (H,L)
4891          XCHG                    ;NEW LINK TO (D,E)
4891          ENDIF
4891          IF      QR            ;IS NODE ADDRESS IN (H,L)?
4891  2B65 2AD242 LHL   QR            ;NO - NODE POINTER TO (H,L)
4891          ENDIF
4891          INX   H                ;NOW STEP TO THE
4891  2B68 23      INX   H                ;RIGHT LINK BYTES
4891  2B69 23      INX   H
4891  2B6A 73      MOV   M,E            ;MOVE THE NEW
4891  2B6B 23      INX   H                ;RIGHT LINK VALUE
4891  2B6C 72      MOV   M,D            ;INTO THE NODE
4891
4892          PTRLK  QP.QR          ;BY A CLOCKWISE ROTATION OF THE TREE.
4892  0001      MCTEM  SET      1            ;ASSUME ADDRESS IN (H,L)
4892          IF      QP            ;TEST ASSUMPTION
4892  0000      MCTEM  SET      0            ;FALSE ASSUMPTION
4892          ENDIF
4892          IF      QR            ;IS NEW LINK IN (D,E)?
4892          IF      MCTEM          ;NO - BUT IS POINTER IN (H,L)?
4892          XCHG                    ;YES - SAVE NODE POINTER
4892          ENDIF
4892          LHL   QR            ;NEW LINK TO (H,L)
4892  2B6D 2AD242 XCHG                    ;NEW LINK TO (D,E)
4892  2B70 EB      XCHG
4892          ENDIF
4892          IF      QP            ;IS NODE ADDRESS IN (H,L)?
4892  2B71 2AC842 LHL   QP            ;NO - NODE POINTER TO (H,L)
4892          ENDIF
4892          MOV   M,E            ;MOVE THE NEW
4892  2B74 73      INX   H                ;LEFT LINK VALUE
4892  2B75 23      INX   H
4892  2B76 72      MOV   M,D            ;INTO THE NODE
4892
4893          GTRLK  QP.0            ;IN THIS CONTEXT THE REBALANCE
4893          IF      QP            ;IS NODE ADDRESS IN (H,L)?
4893  2B77 2AC842 LHL   QP            ;NO - NODE POINTER TO (H,L)
4893          ENDIF
4893          INX   H                ;STEP POINTER TO
4893  2B7A 23      INX   H                ;THE RIGHT LINK
4893  2B7B 23      INX   H
4893  2B7C 5E      MOV   E,M            ;MOVE THE RIGHT
4893  2B7D 23      INX   H                ;LINK INTO THE
4893  2B7E 56      MOV   D,M            ; (D,E) REGISTER
4893          IF      00000H          ;IS IT SAVE LINK MODE?
4893          XCHG                    ;YES - LINK TO (H,L)
4893          SHLD   00000H          ;LINK TO MEMORY
4893          ENDIF
4893
4894          PTRLK  QS.0            ;POINT IS CONSIDERED TO BE
4894  0001      MCTEM  SET      1            ;ASSUME ADDRESS IN (H,L)
4894          IF      QS            ;TEST ASSUMPTION
4894  0000      MCTEM  SET      0            ;FALSE ASSUMPTION
4894          ENDIF
4894          IF      00000H          ;IS NEW LINK IN (D,E)?
4894          IF      MCTEM          ;NO - BUT IS POINTER IN (H,L)?
4894          XCHG                    ;YES - SAVE NODE POINTER
4894          ENDIF
4894          LHL   00000H          ;NEW LINK TO (H,L)
4894          XCHG                    ;NEW LINK TO (D,E)
4894          ENDIF
4894          IF      QS            ;IS NODE ADDRESS IN (H,L)?
4894  2B7F 2AD042 LHL   QS            ;NO - NODE POINTER TO (H,L)
4894          ENDIF
4894          MOV   M,E            ;MOVE THE NEW
4894  2B82 73      INX   H                ;LEFT LINK VALUE
4894  2B83 23      INX   H
4894  2B84 72      MOV   M,D            ;INTO THE NODE
4894
4895          PTRLK  QP.QS          ;THE ROOT NODE OF THE TREE
4895  0001      MCTEM  SET      1            ;ASSUME ADDRESS IN (H,L)
4895          IF      QP            ;TEST ASSUMPTION
4895  0000      MCTEM  SET      0            ;FALSE ASSUMPTION
4895          ENDIF
4895          IF      QS            ;IS NEW LINK IN (D,E)?
4895          IF      MCTEM          ;NO - BUT IS POINTER IN (H,L)?
4895          XCHG                    ;YES - SAVE NODE POINTER
4895          ENDIF

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4895 2B85 2AD042      LHLD QS      ;NEW LINK TO (H,L)
4895 2B88 EB          XCHG         ;NEW LINK TO (D,E)
4895                ENDIF
4895                IF QP
4895 2B89 2AC842      LHLD QP      ;IS NODE ADDRESS IN (H,L)?
4895                ENDIF
4895                INX H
4895 2B8D 23          INX H        ;NOW STEP TO THE
4895 2B8E 73          MOV M.E     ;RIGHT LINK BYTES
4895 2B8F 23          INX H        ;MOVE THE NEW
4895 2B90 72          MOV M.D     ;RIGHT LINK VALUE
4895                ;INTO THE NODE
4896 2B91 23          INX H
4897 2B92 56          MOV D.M     ;STEP TO THE BALANCE FACTOR
4898 2B93 AF          XRA A       ;GET ORINAL BALANCE FACTOR
4899 2B94 77          MOV M.A     ;OF NEW SUBTREE ROOT NODE
4900 2B95 B2          ORA D       ;AND MARK IT AS BALANCED
4901 2B96 CA422B     JZ EN7T5    ;TEST ORIGINAL BALANCE FACTOR
4902 2B99 F2B12B     JP EN7M8    ;OK - ALL SUBTREES BALANCED
4903 2B9C AF          XRA A       ;IT WAS HEAVY TO THE RIGHT
4904                PTBAL QR,0 ;IT WAS HEAVY TO THE LEFT
4904                IF 00000H   ;SET BALANCE FACTORS FOR
4904                LDA 00000H   ;IS FACTOR IN A-REG?
4904                ENDIF      ;NO - BALANCE TO A-REG
4904                IF QR
4904 2B9D 2AD242      LHLD QR      ;IS NODE ADDRESS IN (H,L)
4904                ENDIF      ;NO - NODE POINTER TO (H,L)
4904 2BA0 23          INX H
4904 2BA1 23          INX H        ;STEP POINTER TO
4904 2BA2 23          INX H        ;THE BALANCE BYTE
4904 2BA3 23          INX H
4904 2BA4 77          MOV M.A     ;BALANCE CODE TO NODE
4905 2BA5 3C          INR A
4906                PTBAL QS,0 ;THE SUBTREES THAT WERE
4906                IF 00000H   ;ROTATED DURING ADJUSTMENT
4906                LDA 00000H   ;IS FACTOR IN A-REG?
4906                ENDIF      ;NO - BALANCE TO A-REG
4906                IF QS
4906 2BA6 2AD042      LHLD QS      ;IS NODE ADDRESS IN (H,L)
4906                ENDIF      ;NO - NODE POINTER TO (H,L)
4906 2BA9 23          INX H
4906 2BAA 23          INX H        ;STEP POINTER TO
4906 2BAB 23          INX H        ;THE BALANCE BYTE
4906 2BAC 23          INX H
4906 2BAD 77          MOV M.A     ;BALANCE CODE TO NODE
4907 2BAE C3C32B     JMP NQA10   ;ADJUST LINKAGE TO TOTAL TREE
4908
4909 2BB1 AF          EN7M8: XRA A
4910                PTBAL QS,0 ;SET BALANCE FACTORS
4910                IF 00000H   ;FOR THE SUBTREES THAT WERE
4910                LDA 00000H   ;IS FACTOR IN A-REG?
4910                ENDIF      ;NO - BALANCE TO A-REG
4910                IF QS
4910 2BB2 2AD042      LHLD QS      ;IS NODE ADDRESS IN (H,L)
4910                ENDIF      ;NO - NODE POINTER TO (H,L)
4910 2BB5 23          INX H
4910 2BB6 23          INX H        ;STEP POINTER TO
4910 2BB7 23          INX H        ;THE BALANCE BYTE
4910 2BB8 23          INX H
4910 2BB9 77          MOV M.A     ;BALANCE CODE TO NODE
4911 2BBA 2F          CMA
4912                PTBAL QR,0 ;ROTATED DURING ADJUSTMENT
4912                IF 00000H   ;IS FACTOR IN A-REG?
4912                LDA 00000H   ;NO - BALANCE TO A-REG
4912                ENDIF
4912                IF QR
4912 2BBB 2AD242      LHLD QR      ;IS NODE ADDRESS IN (H,L)
4912                ENDIF      ;NO - NODE POINTER TO (H,L)
4912 2BBE 23          INX H
4912 2BBF 23          INX H        ;STEP POINTER TO
4912 2BC0 23          INX H        ;THE BALANCE BYTE
4912 2BC1 23          INX H
4912 2BC2 77          MOV M.A     ;BALANCE CODE TO NODE
4913
4914 2BC3 3AD442      NQA10: LDA QDIR
4915 2BC6 B7          ORA A
4916 2BC7 FAD72B     JM ENQ11   ;GET THE ENTRY DIRECTION
4917                PTRLK QT,QP ;INTO THE BALANCE POINT
4917 0001           MCTEM SET 1       ;AND TEST IT
4917                IF QT      ;LINK ADJUSTED SUBTREE
4917 0000           MCTEM SET 0       ;ASSUME ADDRESS IN (H,L)
4917                ENDIF     ;TEST ASSUMPTION
4917                IF QP      ;FALSE ASSUMPTION
4917                ;IS NEW LINK IN (D,E)?

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4917          IF      MCTEM          ;NO - BUT IS POINTER IN (H.L)?
4917          XCHG
4917          ENDIF          ;YES - SAVE NODE POINTER
4917 2BCA 2AC842  LHL  QP          ;NEW LINK TO (H.L)
4917 2BCD EB      XCHG          ;NEW LINK TO (D,E)
4917          ENDIF
4917          IF      QT          ;IS NODE ADDRESS IN (H.L)?
4917 2BCE 2ACC42  LHL  QT          ;NO - NODE POINTER TO (H.L)
4917          ENDIF
4917 2BD1 23      INX  H          ;NOW STEP TO THE
4917 2BD2 23      INX  H          ;RIGHT LINK BYTES
4917 2BD3 73      MOV  M,E        ;MOVE THE NEW
4917 2BD4 23      INX  H          ;RIGHT LINK VALUE
4917 2BD5 72      MOV  M,D        ;INTO THE NODE
4917
4918 2BD6 C9      RET              ;AS A RIGHT SUBTREE
4919
4920          ENQ11: PTLK  QT,QP      ;LINK ADJUSTED SUBTREE
4920 0001         MCTEM  SET   1      ;ASSUME ADDRESS IN (H.L)
4920          IF      QT          ;TEST ASSUMPTION
4920 0000         MCTEM  SET   0      ;FALSE ASSUMPTION
4920          ENDIF
4920          IF      QP          ;IS NEW LINK IN (D,E)?
4920          IF      MCTEM        ;NO - BUT IS POINTER IN (H.L)?
4920          XCHG          ;YES - SAVE NODE POINTER
4920          ENDIF
4920 2BD7 2AC842  LHL  QP          ;NEW LINK TO (H.L)
4920 2BDA EB      XCHG          ;NEW LINK TO (D,E)
4920          ENDIF
4920          IF      QT          ;IS NODE ADDRESS IN (H.L)?
4920 2BDB 2ACC42  LHL  QT          ;NO - NODE POINTER TO (H.L)
4920          ENDIF
4920 2BDE 73      MOV  M,E        ;MOVE THE NEW
4920 2BDF 23      INX  H          ;LEFT LINK VALUE
4920 2BE0 72      MOV  M,D        ;INTO THE NODE
4920
4921 2BE1 C9      RET              ;AS A LEFT SUBTREE
4922
4923 2BE2 EB      TREMT: XCHG          ;RECOVER THE NODE ADDRESS
4924 2BE3 22BA42  SHLD RLINK        ;MARK TREE NONEMPTY
4925 2BE6 CDEE2B  CALL  NDCLR          ;RESET LINK AND BALANCE FACTOR
4926 2BE9 EB      XCHG          ;GET EVENT TIME FROM NODE
4927 2BEA 22BD42  SHLD NXTIM        ;SET AS NEXT EVENT TIME
4928 2BED C9      RET              ;ALL DONE
4929
4930 2BEE AF      NDCLR: XRA  A          ;RESET THE LINKS AND
4931 2BEF 77      MOV  M,A        ;THE BALANCE FACTOR OF
4932 2BF0 23      INX  H          ;THE SELECTED NODE
4933 2BF1 77      MOV  M,A
4934 2BF2 23      INX  H
4935 2BF3 77      MOV  M,A
4936 2BF4 23      INX  H
4937 2BF5 77      MOV  M,A
4938 2BF6 23      INX  H
4939 2BF7 77      MOV  M,A
4940 2BF8 23      INX  H          ;STEP TO THE EVENT TIME
4941 2BF9 5E      MOV  E,M        ;RETURN WITH THE
4942 2BFA 23      INX  H          ;EVENT TIME IN (D,E)
4943 2BFB 56      MOV  D,M
4944 2BFC C9      RET
4945
4946 2BFD 110500  CMPTM: LXI  D,05      ;POINT TO THE TIME WORD
4947 2C00 19      DAD  D          ;OF SELECTED NODE
4948 2C01 5E      MOV  E,M        ;GET LOW ORDER BITS
4949 2C02 23      INX  H          ;POINT AT HIGH BITS
4950 2C03 56      MOV  D,M        ;AND GET THEM
4951 2C04 2AD642  LHL  NEWTM        ;GET TIME OF NEW NODE
4952 2C07 7D      MOV  A,L          ;FORM
4953 2C08 93      SUB  E          ;(NEW TIME)-(NODE TIME)
4954 2C09 7C      MOV  A,H
4955 2C0A 9A      SBB  D
4956          ADD  A          ;DECISION BIT TO SIGN
4957 2C0B C9      RET              ;EXIT
4958
4959          ;FATAL ERROR HANDLER. WAIT FOR "CLEAR" FROM MASTER KEYBOARD
4960          ; AND THEN COLD START THE COMPUTER.
4961          ; IF "SEND" IS HIT INSTEAD, WILL RETURN TO CALLER.
4962
4963          ;INPUT: A= ERROR CODE.
4964          ;ALL REGISTERS RESTORED.
4965
4966 2C0C F5      FATAL: PUSH  PSW      ;SAVE PSW FOR DEBUGGER
4967 2C0D 3A6842  LDA  SL700        ;IF MASTER KEYBOARD
4968 2C10 B7      ORA  A          ; IS T.I.
4969 2C11 CA162C  JZ   FATL        ; SILENT 700
4970 2C14 F1      POP  PSW      ;RESTORE PSW
4971 2C15 D7      RST  2          ;ENTER DEBUGGER

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4972 ;
4973 ;NORMAL FATAL ERROR HANDLING
4974 ;
4975 2C16 F1   FATL:  POP   PSW       ;A = ERROR CODE
4976 2C17 E5           PUSH   H         ;SAVE
4977 2C18 D5           PUSH   D         ; ALL
4978 2C19 C5           PUSH   B         ;  REGISTERS
4979 2C1A F5           PUSH   PSW      ;
4980 2C1B F6F0        ORI    0F0H      ;FATAL ERROR CODES ARE 0F0H TO 0FFH
4981 2C1D CD080F      CALL  ERPC      ;DISPLAY MESSAGE
4982 2C20 CD0911      CALL  OFCVY     ;TURN OFF CONVEYOR
4983 2C23 3E01        MVI   A,1       ;LOCK OUT
4984 2C25 32AC42      STA   FSQ0      ; ALL OTHER KB PROCESSING
4985 ;
4986 ; BEEP AND WAIT FOR COMMAND
4987 ;
4988 2C28 063B        FALRM: MVI   B,','   ;B = "BEEP"
4989 2C2A 3A093D      LDA   MAKB      ;C
4990 2C2D 4F          MOV   C,A       ; = MASTER KB #
4991 2C2E CD6C11      CALL  PUTKB     ;ALARM OPERATOR
4992 2C31 110500      LXI   D,5       ;IDLE ABOUT
4993 2C34 CD5A2D      CALL  RDELA     ; 500 MS.
4994 2C37 C5          PUSH  B         ;SAVE MASTER KB #
4995 2C38 CD0E10      CALL  GETKB     ;C = KB #, A = CHAR
4996 2C3B 47          MOV   B,A       ;SAVE CHARACTER IN B
4997 2C3C D1          POP   D         ;E = MASTER KB #
4998 2C3D 79          MOV   A,C       ;WAIT UNTIL
4999 2C3E BB          CMP   E         ; THERE IS INPUT
5000 2C3F C2282C      JNZ  FALRM     ; FROM MASTER KB.
5001 2C42 78          MOV   A,B       ;A = CHAR
5002 2C43 FE3B        CPI   ','       ;IF ',' THEN
5003 2C45 CA4B2C      JZ   FATLX     ; RETURN
5004 2C48 C3DD03      JMP  EXEC       ;ELSE COLD START
5005 ;
5006 2C4B CD190C      FATLX: CALL  DEVII ;RE-INITIALIZE DEVICES
5007 2C4E F1          POP   PSW      ;RESTORE
5008 2C4F C1          POP   B         ; ALL
5009 2C50 D1          POP   D         ;  REGISTERS
5010 2C51 E1          POP   H         ;
5011 2C52 C9          RET
5012 ;
5013 ;GET DCDB QUEUE POINTER IN D,E.
5014 ;
5015 ;INPUT: C=JDV.
5016 ;OUTPUT: B,C=JDV, H,L POINTS TO DCDB(JDV), SO DOES DCDBP.
5017 ;
5018 2C53 2A5742      GDCQ:  LHLD  DCDB   ;POINT H,L TO DCDB TABLE
5019 2C56 0600        MVI   B,0       ;B,C=JDV
5020 2C58 09          DAD   B         ;POINT H,L TO
5021 2C59 09          DAD   B         ; DCDB(JDV)
5022 2C5A 2B          DCX  H         ; MS BYTE
5023 2C5B 56          MOV   D,M       ;D,E
5024 2C5C 2B          DCX  H         ; =
5025 2C5D 5E          MOV   E,M       ; DCDB(JDV)
5026 2C5E 220143     SHLD  DCDBP     ;SET DCDBP
5027 2C61 C9          RET
5028 ;
5029 ;
5030 ;GET JDV, IDENT FROM LAQ IN PPDB.
5031 ;
5032 ;INPUT: H,L POINTS TO BYTE 2 IN PPDB.
5033 ;OUTPUT: B=JDV, D,E=IDENT, H,L POINTS TO IDENT MS BYTE IN PPDB,
5034 ; FLAG Z SET & H,L NOT EFFECTED ON NEGATIVE RETURN.
5035 ;ALL REGISTERS EXCEPT C EFFECTED.
5036 ;
5037 2C62 110400      GLAQ:  LXI   D,4   ;POINT TO BYTE 6
5038 2C65 19          DAD   D         ; IN PPDB
5039 2C66 7E          MOV   A,M       ;IF
5040 2C67 B7          ORA   A         ; LAQ
5041 2C68 C8          RZ           ; NOT EMPTY
5042 ;
5043 ;REMOVE JDV, IDENT FROM LAQ IN PPDB
5044 ;
5045 2C69 3A0040      LDA   FLAGS     ;IF AUTO-INDUCT
5046 2C6C B7          ORA   A         ; IS
5047 2C6D F2752C      JP   GDLAG     ; DESIRED
5048 2C70 23          INX  H         ;PICK UP LAQ
5049 2C71 5E          MOV   E,H       ; ITEM WITHOUT
5050 2C72 C37D2C      JMP  GLAQ1     ; DELETING
5051 ;
5052 2C75 35          GDLAG: DCR   M         ;DEC LAQSZ IN PPDB
5053 2C76 23          INX  H         ;H,L POINTS TO LAQ1
5054 2C77 7E          MOV   A,M       ;PICK UP LAQ 1ST ITEM # LAQ1
5055 2C78 5F          MOV   E,A       ;SAVE LAQ1 IN E
5056 2C79 3C          INR  A         ;NEXT ITEM IN LAQ BECOMES 1ST ITEM
5057 2C7A E607        ANI  07H       ;WRAP AROUND IF NECESSARY
5058 2C7C 77          MOV   M,A       ;UPDATE LAQ1 IN PPDB

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5059 2C7D 2C      GLAQ1:  INR  L      ;H,L POINTS TO LAQ (ALSO RESET FLAG Z)
5060 2C7E 1600    MVI  D,0      ;D,E=LAQ1
5061 2C80 19      DAD  D      ;H,L POINTS TO
5062 2C81 19      DAD  D      ; ITEM # LAQ1
5063 2C82 19      DAD  D      ; IN LAQ
5064 2C83 46      MOV  B,M     ;B=JDV
5065 2C84 23      INX  H      ;H,L POINTS TO IDENT
5066 2C85 5E      MOV  E,M     ;D,E
5067 2C86 23      INX  H      ; = PACKAGE
5068 2C87 56      MOV  D,M     ; IDENTIFICATION
5069 2C88 C9      RET          ; POSITIVE RETURN
5070
5071
5072 ;CREATE UDB1 FOR NEXT UPDATE PHOTO. COPY IDENT, JDV.
5073 ; JUPDA, JINDU FROM OLD TCB TO NEW TCB AND INCREMENT JUPDA.
5074 ; CALLED BY IDBPC OR UDBPC.
5075 ;
5076 ;INPUT: B,C POINTS TO BYTE 13 IN OLD TCB.
5077 ;OUTPUT: NEW TCB, H,L POINTS TO BYTE 7 IN NEW TCB.
5078
5079 ; A=JINDU, D,E= OLD JUPDA.
5080 ;ALL REGISTERS EFFECTED.
5081
5080 2C89 CDCB0F    GUDB1:  CALL  GETFS   ;GET FREE BLOCK FOR UDB1
5081 2C8C 110D00    LXI  D,13     ;POINT H.L TO BYTE 13
5082 2C8F 19      DAD  D      ; IN TCB
5083 2C90 0A      LDAX B      ;COPY
5084 2C91 77      MOV  M,A     ; IDENT
5085 2C92 0B      DCX  B      ; FROM
5086 2C93 2B      DCX  H      ; OLD TCB
5087 2C94 0A      LDAX B      ; TO
5088 2C95 77      MOV  M,A     ; NEW TCB
5089 2C96 0B      DCX  B      ;MOVE POINTERS TO
5090 2C97 2B      DCX  H      ; BYTES 11
5091 2C98 0A      LDAX B      ;COPY
5092 2C99 77      MOV  M,A     ; JDV
5093 2C9A 0B      DCX  B      ;MOVE POINTERS TO
5094 2C9B 2B      DCX  H      ; BYTES 10
5095 2C9C 0A      LDAX B      ;A= OLD JUPDA
5096 2C9D 5F      MOV  E,A     ;E= OLD JUPDA FOR UDBPC
5097 2C9E 3C      INR  A      ;A= NEW JUPDA
5098 2C9F 77      MOV  M,A     ;SET JUPDA IN NEW TCB
5099 2CA0 0B      DCX  B      ;MOVE POINTERS TO
5100 2CA1 2B      DCX  H      ; BYTES 9
5101 2CA2 0A      LDAX B      ;COPY
5102 2CA3 77      MOV  M,A     ; JINDU
5103 2CA4 2B      DCX  H      ;POINT H.L TO BYTE 8 IN NEW TCB
5104 2CA5 019424    LXI  B,UDB1  ;EVENT
5105 2CA8 70      MOV  M,B     ; IN TCB
5106 2CA9 2B      DCX  H      ; =
5107 2CAA 71      MOV  M,C     ; UDB1
5108 2CAB C9      RET
5109
5110 ;SET UP COUNT IN PPDB FOR UNBLOCKED PP # C+1 OR A PENDING
5111 ; UNCODED PACKAGE AT PP. FOR STRAIGHT MERGE ONLY.
5112 ; CALLED BY MCLPC OR TECMG.
5113 ;
5114 ;INPUT: C= INDUCT #, PWTNG & COUNT IN PPDB.
5115 ;OUTPUT: COUNT(JINDU) IN PPDB, POINTED TO BY H.L.
5116 ;ALL REGISTERS EXCEPT C EFFECTED.
5117 ;
5118 2CAC 2A4342    PEND:  LHLD  PPDB   ;POINT H.L TO BYTE 1
5119 2CAF 23      INX  H      ; IN PPDB
5120 2CB0 3A033D    LDA  NINDU   ; IF IT IS
5121 2CB3 3D      DCR  A      ; A SINGLE INDUCT
5122 2CB4 C2BE2C    JNZ  PENDS   ; SYSTEM
5123
5124 ;SPECIAL HANDLING FOR SINGLE INDUCT SYSTEMS
5125 ;
5126 2CB7 7E      MOV  A,M     ;PICK UP PWTNG FLAG IN PPDB
5127 2CB8 23      INX  H      ;POINT TO COUNT IN PPDB
5128 2CB9 FE02    CPI  2      ;IF PWTNG FLAG = 2 THEN
5129 2CBB C0      RNZ          ; INCREMENT
5130 2CBC 34      INR  M      ; COUNT
5131 2CBD C9      RET
5132
5133 ;ADJUST COUNT FOR MULTIPLE STRAIGHT INDUCT SYSTEMS
5134 ;
5135 2CBE 110100    PENDS: LXI  D,0001H ;INITIALIZE D = MAX COUNT
5136 ; & E = COUNT INCREMENT

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5137 2CC1 43      MOV   B,E      ;B= INDUCT $, = 1 INITIALLY
5138
5139 2CC2 78      PNWCN: MOV   A,B      ;IF INDUCT
5140 2CC3 B9      CMP   C        ; NOT $
5141 2CC4 CAD72C   JZ    PNXCID   ; JINDU
5142 2CC7 7E      MOV   A,M      ;A= PWTNG IN PPDB
5143 2CC8 FE02    CPI   2        ;IF NOT=
5144 2CCA CACF2C   JZ    PMXCEN   ; 2
5145 2CCD 1E00    MVI   E,0      ;INC COUNT ONLY IF EVERY PWTNG = 2
5146 2CCF 7A      PNXCEN: MOV  A,D      ;A= MAX COUNT
5147 2CD0 23     INX   H        ;POINT H.L TO BYTE 2 IN PPDB
5148 2CD1 BE      CMP   M        ;IF COUNT
5149 2CD2 D2D62C  JNC  PNXCID   ; > MAX COUNT THEN
5150 2CD5 56     MOV   D,M      ;MAX COUNT = COUNT
5151 2CD6 2B     PNXCID: DCX  H      ;POINT H.L TO BYTE 1 IN PPDB
5152 2CD7 D5     PNXCID: PUSH D      ;MOVE H.L
5153 2CD8 112000 LXI  D,32      ; TO
5154 2CDB 19     DAD  D        ; NEXT
5155 2CDC D1     POP  D        ; INDUCT
5156 2CDD 04     INR  B        ;NEXT INDUCT $
5157 2CDE 3A033D LDA  NINDU     ;DO FOR
5158 2CE1 B8     CMP  B        ; ALL
5159 2CE2 D2C22C JNC  PNWCN     ; INDUCTS
5160
5161 2CE5 7A      MOV   A,D      ;B
5162 2CE6 83      ADD  E        ; = MAX COUNT + COUNT INC
5163 2CE7 47     MOV  B,A      ; = NEW COUNT
5164 2CE8 79     MOV  A,C      ;A= INDUCT $
5165 2CE9 CD9927 CALL CPPPP     ;POINT H.L TO
5166 2CEC 23     INX  H        ; BYTE 2
5167 2CED 23     INX  H        ; IN PPDB
5168 2CEE 70     MOV  M,B      ;UPDATE COUNT(JINDU) IN PPDB
5169 2CEF C9     RET
5170
5171 ;PUT JDV & IDENT INTO LAQ IN PPDB FOR INDUCT $ C+1.
5172 ; CALLED BY SNPC.
5173
5174 ;INPUT: C= INDUCT $, JDV,IDENT.
5175 ;OUTPUT: PPDBP POINTS TO BYTE 1 IN PPDB, FLAG Z SET AND
5176 ; A=0 ON LAQ FULL.
5177 ;A
5178 2CF0 79      PLAQ:  MOV  A,C      ;A= INDUCT $
5179 2CF1 CD9927 CALL  CPPPP     ;POINT H.L TO BYTE 1
5180 2CF4 23     INX  H        ; IN PPDB(JINDU)
5181 2CF5 22EB42 SHLD PPDBP     ;OUTPUT POINTER
5182 2CF8 110500 LXI  D,5       ;POINT H.L TO BYTE 6
5183 2CFB 19     DAD  D        ; IN PPDB
5184 2CFC 3A373D LDA  QLIMI     ;A= LAQ SIZE LIMIT DESIRED
5185 2CFF 96     SUB  M        ;IF LAQ NOT
5186 2D00 C8     RZ          ; FULL
5187
5188 ;LAQ NOT YET FULL: INSERT JDV & IDENT
5189
5190 2D01 7E      MOV  A,M      ;PICK UP CURRENT LAQ SIZE
5191 2D02 34     INR  M        ;INC LAQSZ IN PPDB
5192 2D03 23     INX  H        ;POINT H.L TO BYTE 7 IN PPDB
5193 2D04 86     ADD  M        ;COMPUTE LAQ NEW ITEM $
5194 2D05 E607   ANI  07H      ;MASK TO 3 BITS
5195 2D07 5F     MOV  E,A      ;D,E
5196 2D08 1600   MVI  D,0      ; = LAQ NEW ITEM $
5197 2D0A 23     INX  H        ;POINT H.L TO LAQ ITEM 0
5198 2D0B 19     DAD  D        ;POINT H.L TO
5199 2D0C 19     DAD  D        ; NEW ITEM
5200 2D0D 19     DAD  D        ; LOCATION
5201 2D0E EB     XCHG        ;POINT D,E TO NEW LOCATION
5202 2D0F 219642 LXI  H,JDV     ;POINT H.L TO JDV,IDENT
5203 2D12 0603   MVI  B,3      ;B= LOOP COUNT
5204 2D14 CDCF10 CALL MOVEB     ;COPY JDV,IDENT INTO LAQ
5205 2D17 04     INR  B        ;POSITIVE
5206 2D18 C9     RET          ; RETURN
5207
5208 2D19 D5     RSCHD: PUSH D      ;SAVE DELAY COUNT
5209 2D1A E5     PUSH  H      ;SAVE POINTER
5210 2D1B 5E     MOV  E,M      ;PICK UP
5211 2D1C 23     INX  H        ; TCB
5212 2D1D 56     MOV  D,M      ; POINTER
5213 2D1E 7B     MOV  A,E      ;IF NONZERO
5214 2D1F B2     ORA  D        ; THEN EVENT
5215 2D20 CA2D2D JZ   RSSEE     ; PENDING
5216
5217 ;JUNK ANY EVENT ALREADY PENDING
5218
5219 2D23 210700 LXI  H,7       ;POINT TO BYTE 7
5220 2D26 19     DAD  D        ; IN TCB
5221 2D27 118025 LXI  D,JUNK    ;JUNK
5222 2D2A 73     MOV  M,E      ; THE
5223 2D2B 23     INX  H        ; PENDING
5224 2D2C 72     MOV  M,D      ; EVENT
5225

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449
: SCHEDULE NEW EVENT
5226
5227
5228 2D2D CDCB0F RSSEE: CALL GETFS : GET NEW BLOCK FOR TCB
5229 2D30 EB XCHG : POINT DE TO TCB
5230 2D31 E1 POP H : POINT TO TCB
5231 2D32 E5 PUSH H : POINTER
5232 2D33 73 MOV M.E : SET UP
5233 2D34 23 INX H : TCB
5234 2D35 72 MOV M.D : POINTER
5235 2D36 210B00 LXI H.11 : POINT TO BYTE 11
5236 2D39 19 DAD D : IN TCB
5237 2D3A D1 POP D : POIN DE TO TCB POINTER
5238 2D3B 72 MOV M.D : SET BACKWARD
5239 2D3C 2B DCX H : POINTER
5240 2D3D 73 MOV M.E : IN TCB
5241 2D3E 2B DCX H : POINT TO BYTE 9 IN TCB
5242 2D3F 3A9442 LDA JINDU : SAVE INDUCT #
5243 2D42 77 MOV M.A : IN TCB
5244 2D43 2B DCX H : POINT TO BYTE 8 IN TCB
5245 2D44 70 MOV M.B : SET EVENT
5246 2D45 2B DCX H : ENTRY IN
5247 2D46 71 MOV M.C : TCB
5248 2D47 EB XCHG : COMPUTE
5249 2D48 C1 POP B : DE
5250 2D49 2AB642 LHLD CTIME : = EVENT
5251 2D4C 09 DAD B : DUE
5252 2D4D EB XCHG : TIME
5253 2D4E 2B DCX H : POINT TO BYTE 6 IN TCB
5254 2D4F 72 MOV M.D : SET EVENT
5255 2D50 2B DCX H : DUE TIME
5256 2D51 73 MOV M.E : IN TCB
5257 2D52 11FBFF LXI D.0FFFBH : POINT TO
5258 2D55 19 DAD D : TCB
5259 2D56 CD4429 CALL ENQU2 : SCHEDULE EVENT
5260 2D59 C9 RET
5261
5262 2D5A E5 RDELA: PUSH H : SAVE CALLER'S H.L. A
5263 2D5B F5 PUSH PSW : IN STACK
5264 2D5C C5 PUSH B : GET TREE
5265 2D5D D5 PUSH D : PARAMETERS
5266 2D5E AF XRA A : FOR
5267 2D5F CDBA0F CALL GTREP : REAL
5268 2D62 D1 POP D : TIME
5269 2D63 C1 POP B : CLOCK
5270 2D64 C3692D JMP RDLPI
5271
5272
5273 : DELAY D.E PPI PULSES.
5274
5275 : INPUT: TREE PARAMETERS, D,E= # PULSES TO DELAY.
5276 : OUTPUT: TCBP POINTS TO DELAYED EVENT.
5277 : REGISTERS A, B, C, H, L RESTORED, STACK EFFECTED AFTER DELAY.
5278 : TCBP, TREE DATA EFFECTED.
5279
5280 2D67 E5 RDLPP: PUSH H : SAVE CALLER'S H.L. A
5281 2D68 F5 PUSH PSW : IN STACK
5282 2D69 2AB642 RDLPI: LHLD CTIME : SAVE DELAYED EVENT
5283 2D6C 19 DAD D : TIME IN
5284 2D6D E5 PUSH H : STACK
5285 2D6E CDCB0F CALL GETFS : GET FREE BLOCK
5286 2D71 22B442 SHLD TCBP : SAVE POINTER
5287 2D74 110500 LXI D.5 : POINT TO BYTE 5
5288 2D77 19 DAD D : IN TCB
5289 2D78 D1 POP D : SET
5290 2D79 73 MOV M.E : DELAYED
5291 2D7A 23 INX H : EVENT
5292 2D7B 72 MOV M.D : TIME
5293 2D7C 23 INX H : POINT TO BYTE 7 IN TCB
5294 2D7D 119B2D LXI D.R2DEL : SET
5295 2D80 73 MOV M.E : EVENT
5296 2D81 23 INX H : TO BE
5297 2D82 72 MOV M.D : DELA2
5298 2D83 23 INX H : POINT TO BYTE 9
5299 2D84 F1 POP PSW : SAVE CALLER'S A
5300 2D85 77 MOV M.A : IN TCB
5301 2D86 23 INX H : POINT TO BYTE 10
5302 2D87 71 MOV M.C : SAVE CALLER'S
5303 2D88 23 INX H : B.C
5304 2D89 70 MOV M.B : IN TCB
5305 2D8A 23 INX H : POINT TO BYTE 12
5306 2D8B D1 POP D : SAVE CALLER'S
5307 2D8C 73 MOV M.E : H.L
5308 2D8D 23 INX H : IN
5309 2D8E 72 MOV M.D : TCB
5310 2D8F 23 INX H : POINT TO BYTE 14

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5311 2D90 D1      POP      D          ;SAVE
5312 2D91 73      MOV      M,E        ; RETURN
5313 2D92 23      INX      H          ; ADDRESS
5314 2D93 72      MOV      M,D        ; IN TCB
5315 2D94 2AB442  LHLD    TCBP       ;SCHEDULE
5316 2D97 CD4429  CALL    ENQUZ      ; EVENT DELA2
5317 2D9A C9      RET              ;RETURN TO CALLER OF CALLER
5318
5319 ;RESUMPTION OF DELAYED EVENT.
5320 ; PART OF MODULE DELAY.
5321
5322 ;INPUT: TCBP, D,E POINTS TO BYTE 8 OF TCB, TREE PARAMETERS.
5323
5324 2D9B 210700  R2DEL: LXI      H,7      ;POINT TO BYTE 15
5325 2D9E 19      DAD      D          ; IN TCB
5326 2D9F 56      MOV      D,M        ;PUT CALLER'S
5327 2DA0 2B      DCX      H          ; RETURN ADDRESS
5328 2DA1 5E      MOV      E,M        ; IN
5329 2DA2 D5      PUSH    D          ; STACK
5330 2DA3 2B      DCX      H          ;POINT TO BYTE 13
5331 2DA4 56      MOV      D,M        ;SAVE CALLER'S
5332 2DA5 2B      DCX      H          ; H.L
5333 2DA6 5E      MOV      E,M        ; IN
5334 2DA7 D5      PUSH    D          ; STACK
5335 2DA8 2B      DCX      H          ;POINT TO BYTE 11
5336 2DA9 46      MOV      B,M        ;RESTORE
5337 2DAA 2B      DCX      H          ; CALLER'S
5338 2DAB 4E      MOV      C,M        ; B.C
5339 2DAC 2B      DCX      H          ;POINT TO BYTE 10
5340 2DAD 7E      MOV      A,M        ;SAVE CALLER'S A
5341 2DAE F5      PUSH    PSW        ; IN STACK
5342 2DAF 2AB442  LHLD    TCBP       ;RETURN TCB TO
5343 2DB2 CD4F11  CALL    PUTFS      ; FREE STORAGE
5344 2DB5 F1      POP     PSW        ;RESTORE CALLER'S A
5345 2DB6 E1      POP     H          ;RESTORE CALLER'S H.L
5346 2DB7 C9      RET              ;RETURN TO CALLER OF DELAY
5347
5348
5349
5350 ;*****
5351 ; OPTIONAL TRANSLATION AND PROGRAMMING MODULES
5352 ;*****
5353
5354 2DB8          ORG     2F00H      ;TRANSLATION MODULE
5355
5356
5357 ; INITIALIZE TRANSLATION DATABASE AND TABLE.
5358
5359 ;INPUT: INTERRUPTS DISABLED.
5360 ;ALL REGISTERS EFFECTED.
5361
5362 2F00 2A4D42  XLINI:  LHLD    XLDB      ;POINT TO XLDB FOR TABLE # 1
5363 2F03 E5      PUSH    H          ;SAVE POINTER AT TOP OF STACK
5364 2F04 21C03D  LXI      H,XLTBS    ;POINT TO TABLE ADDRESSES
5365 2F07 3A0B3D  LDA      NXLTB      ;SET COUNTER B
5366 2F0A 47      MOV      B,A        ; = # TRANSLATE TABLES
5367 2F0B 5E      XXLDB: MOV      E,M        ;PICK UP
5368 2F0C 23      INX      H          ; TABLE
5369 2F0D 56      MOV      D,M        ; ADDRESS
5370 2F0E 23      INX      H          ;POINT TO NEXT TABLE AD
5371 2F0F E3      XTHL     ;POINT TO BYTE 1 IN XLDB.
5372 2F10 23      INX      H          ; LEAVING TABLE-FULL FLAG TBFU = 0
5373 2F11 73      MOV      M,E        ;SET POINTER
5374 2F12 23      INX      H          ; XLTB
5375 2F13 72      MOV      M,D        ; IN XLDB
5376 2F14 3A0C3D  LDA      TBSIZ      ;COMPUTE D.E = XLTB2

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The embodiments of the invention in which an exclusive property or privilege is claimed are defined as follows.

1. A system for monitoring and controlling the flow of articles in a warehousing system including a sorting conveyer with a plurality of discharge locations in which stored articles are to be batch picked and sorted for shipment comprising:

means for storing data representing each article to be picked, said data including a number assigned to each article which is unique to each article and sort destination information for each article;

indicia mean associated with each article to be picked including said unique number;

scanning means for scanning said indicia means on an article transported within the warehousing system; computer circuit means coupled to said data storing means and to said scanning means for correlating position information of each article with the article identification number;

display means coupled to said computer circuit means for providing display information to an operator of the status of each article to be picked and sorted; a conveyer control circuit for selectively controlling the discharge of articles; and

means coupling said conveyer control circuit to said computer circuit means for controlling the diverting of articles in response to the correlation of their unique identification number with associated sort destination information stored in said data storing means.

2. The system as defined in claim 1 wherein data for each batch of articles to be picked is stored on a recording medium which can be transported from a remote central control area, and wherein said system includes data transfer means coupled to said computer circuit means for reading and transferring data from said recording medium to said data storing means.

3. The system as defined in claim 2 wherein said display means includes an operator terminal coupled to said computer circuit means and including a keyboard and CRT display for selectively interrogating the computer circuit means for status information.

4. The system as defined in claim 3 wherein said display means further includes printer means coupled to said computer circuit means for providing a printout of the status of articles to be picked and sorted.

5. The system as defined in claim 4 wherein said display means further includes additional printer means for printing alarm messages relating to predetermined events occurring during the operation of said system.

6. A system for monitoring the status of articles being processed in a warehouse system including article conveying means comprising:

a plurality of label reading means positioned at selected locations within the warehouse system for reading a unique identification number contained on a label applied to each article moving within the system and uniquely identifying each such article; data storage means including each identification number of articles to be processed and including processing information associated with each such article;

computer circuit means coupled to said label reading means and to said data storage means for monitoring the position of articles travelling within the system and for generating signals representative of the position and therefore the processing status of such articles;

display means coupled to said computer circuit means for displaying the status of articles or groups of articles in the system; and

control means coupled to said computer circuit means for controlling the operation of the warehouse conveying means to control the movement of articles or groups of articles in the system.

7. The system as defined in claim 6 wherein said data storage means includes a data recording medium including each identification number stored thereon together with associated processing information.

8. The system as defined in claim 7 wherein said data storage means further includes data transfer means for reading and writing data to and from said recording medium respectively, said data transfer means coupled to said computer circuit means.

9. The system as defined in claim 8 and further including printer means coupled to said computer circuit means for providing a printout of the status of articles to be processed.

10. The system as defined in claim 9 and further including an operator keyboard and wherein said computer circuit means is programmed to provide reports selectively printed by said printing means when requested by commands from said operator keyboard.

11. A system for monitoring the status of articles being processed comprising:

a plurality of label reading means for reading a unique identification number contained on a label applied to each article moving within the system and uniquely identifying each such article;

data storage means including a data recording medium for storing each identification number together with associated processing information for articles to be processed;

computer circuit means coupled to said label reading means and to said data storage means for monitoring the position of articles travelling within the system and for generating signals representative of the position and therefore the processing status of such articles wherein said data storage means further includes data transfer means for reading and writing data to and from said recording medium respectively, said data transfer means coupled to said computer circuit means;

display means coupled to said computer circuit means for displaying the status of articles or groups of articles in the system said display means including printer means coupled to said computer circuit means for providing a printout of the status of articles to be processed;

an operator keyboard coupled to said computer circuit means and wherein said computer circuit means is programmed to provide reports selectively printed by said printing means when requested by commands from said operator keyboard;

a sorting conveyer with a plurality of discharge chutes;

a conveyer control circuit for selectively controlling the discharge of articles onto said discharge chutes; and

means coupling said conveyer control circuit to said computer circuit means for controlling the diverting of articles in response to the correlation of their unique identification number with associated sort destination information stored in said data storage means.

12. A system for monitoring and controlling the batch order picking and sorting of articles in a ware-

house comprising:

- a plurality of label reading means for reading a unique identification number contained on a label applied to each article moving within the system and uniquely identifying each such article;
- bulk data storage means including each identification number of articles to be sorted and including sort destination information for each such article;
- computer circuit means coupled to said label reading means and to said bulk data storage means for monitoring the physical position of articles travelling within the system and for generating signals representative of the position and therefore the sorting status of such articles;
- display means coupled to said computer circuit means for displaying the status of articles or groups of articles in the system;
- a conveyer control circuit for selectively controlling the discharge of articles at discharge locations in the warehouse; and
- means coupling said conveyer control circuit to said computer circuit means for controlling the diverting of articles in response to the correlation of their unique identification number with associated sort

destination information stored in said data storage means.

13. The system as defined in claim 12 wherein said bulk data storage means receives data from a data recording medium including each identification number stored thereon together with associated sorting information.

14. The system as defined in claim 13 wherein said data storage means includes data transfer means for reading and writing data to and from said recording medium respectively, said data transfer means coupled to said computer circuit means.

15. The system as defined in claim 14 wherein said recording medium is a diskette.

16. The system as defined in claim 15 wherein said display means includes printer means coupled to said computer circuit means for providing a printout of the status of articles to be sorted.

17. The system as defined in claim 16 and further including an operator keyboard and wherein said computer circuit means is programmed to provide reports selectively printed by said printing means when requested by commands from said operator keyboard.

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[54] **INTEGRATED INTERACTIVE RESTAURANT COMMUNICATION METHOD FOR FOOD AND ENTERTAINMENT PROCESSING**

[76] **Inventor:** Lawrence G. Kurland, 26 Farmington La., Melville, N.Y. 11747

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[22] **Filed:** Mar. 14, 1983

[51] **Int. Cl.⁴** G06F 3/04; G06F 15/24; G06F 15/44

[52] **U.S. Cl.** 364/401; 364/410; 364/900

[58] **Field of Search** 364/400-401, 364/404-405, 410-412, 200 MS File, 900 MS File; 235/7 R, 383

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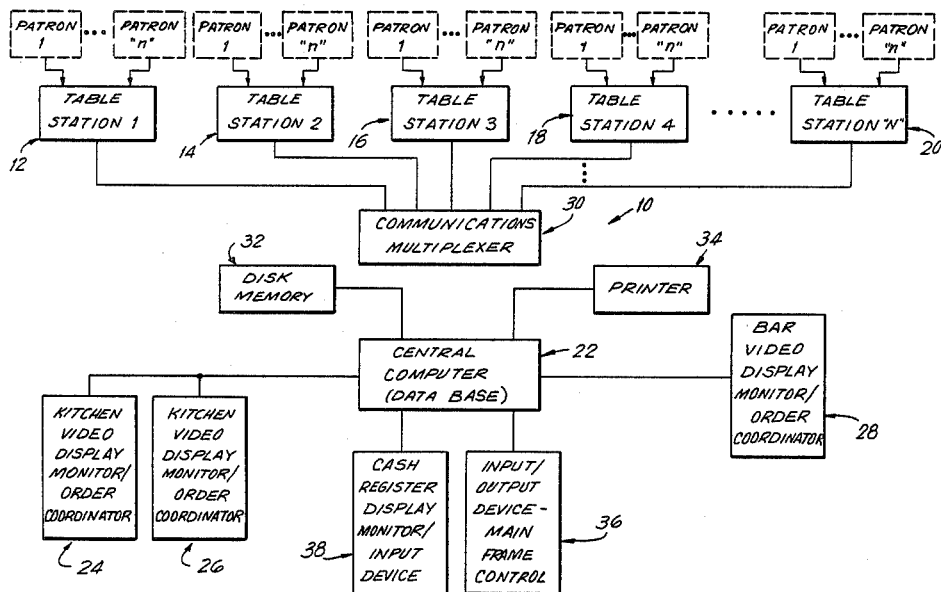
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Primary Examiner—Jerry Smith
Assistant Examiner—Gary V. Harkcom
Attorney, Agent, or Firm—Stiefel, Gross, Kurland & Pavane

[57] **ABSTRACT**

An interactive restaurant communication system (10) provides integrated food and entertainment processing which enables restaurant patrons to accomplish both food selection and select and receive entertainment on a common video monitor (56) at their table. The patrons can obtain menus for individual food selection on the video monitor (56) at their tables and individually enter their orders into a table station "intelligent" terminal (12, 14, 16, 18, 20) at their tables. In addition they can select from and interactively play a variety of remotely retrievable interactive entertainment activities using the video monitor (56) while waiting for the food to arrive, and if desired, where applicable, have the food and entertainment charges automatically added to a composite bill which may be printed at the table station terminal (12, 14, 16, 18, 20) or at a remote central location. The food and entertainment functions of the terminal (12, 14, 16, 18, 20) are down-line loaded from a central data base (22, 32) in response to terminal requests therefor.

11 Claims, 7 Drawing Figures



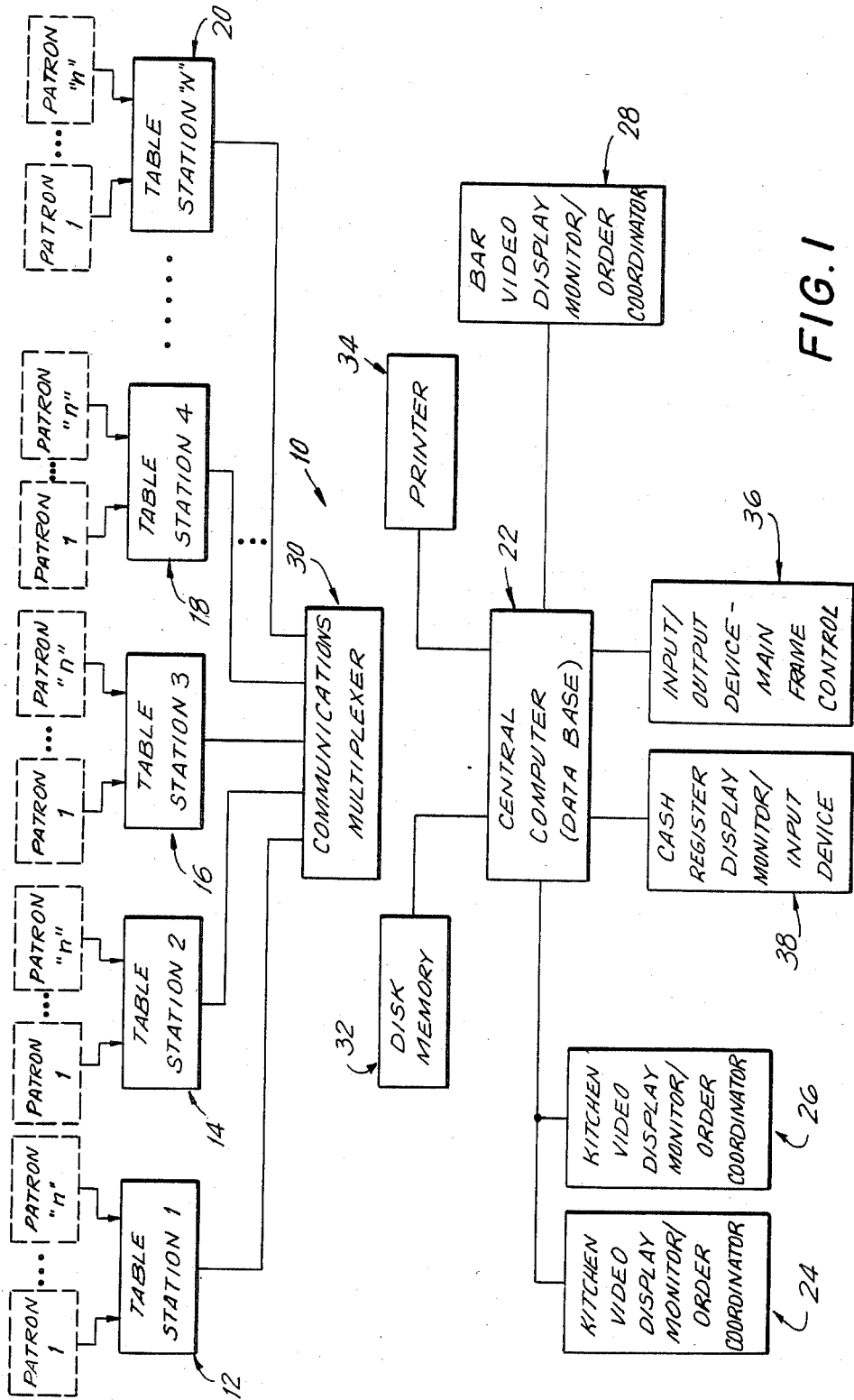


FIG. 1

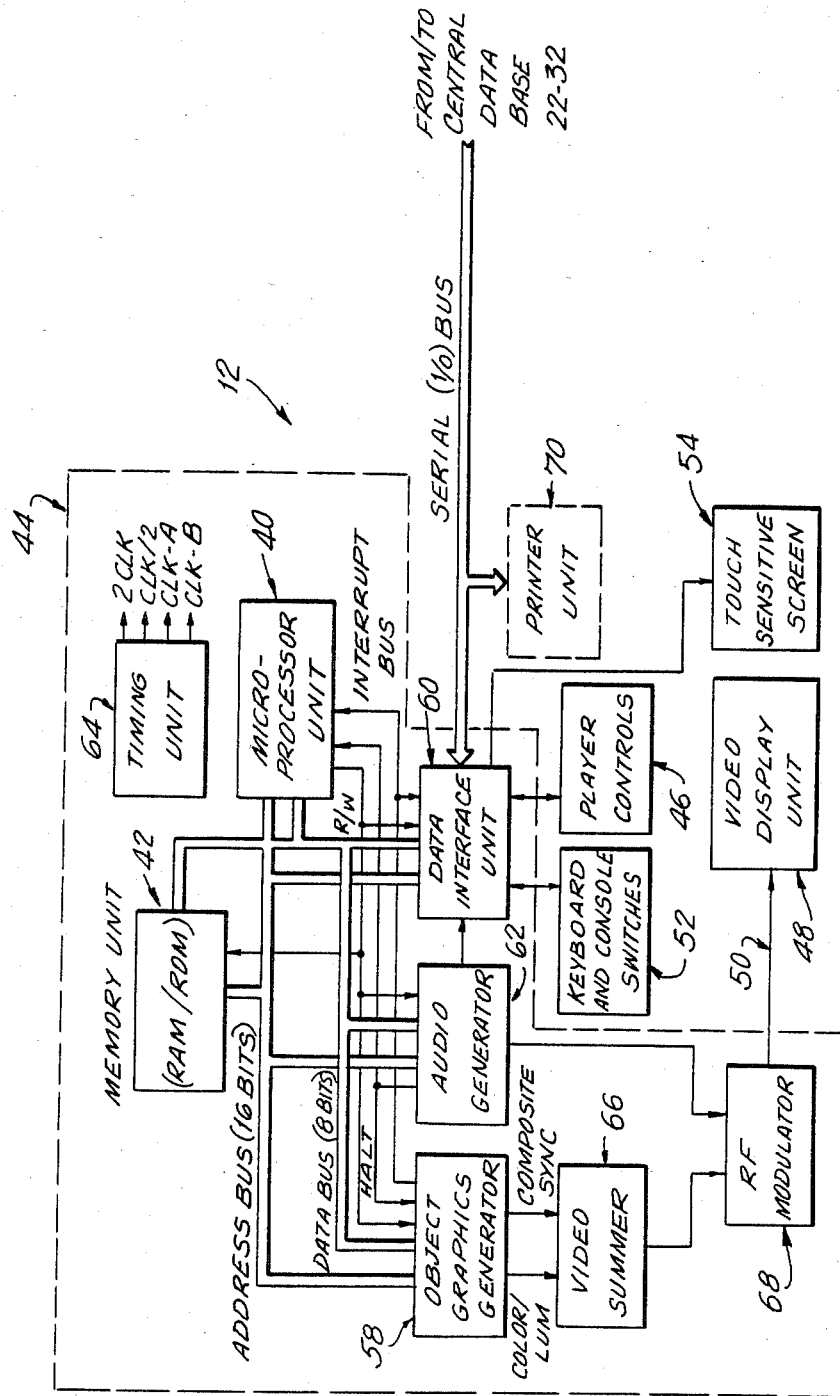


TABLE STATION 1

FIG. 2

FIG. 3

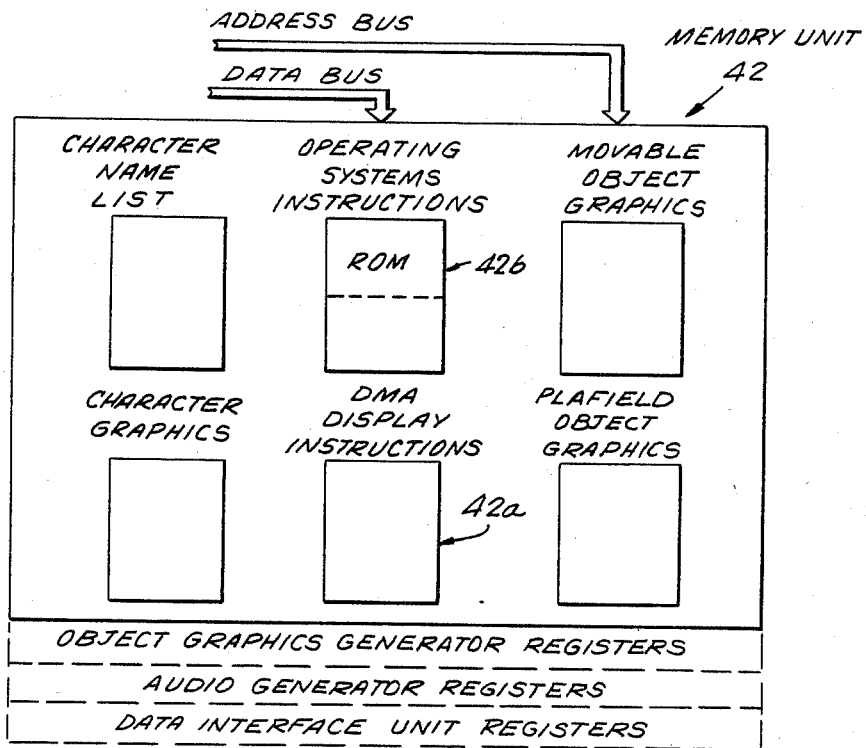


FIG. 4

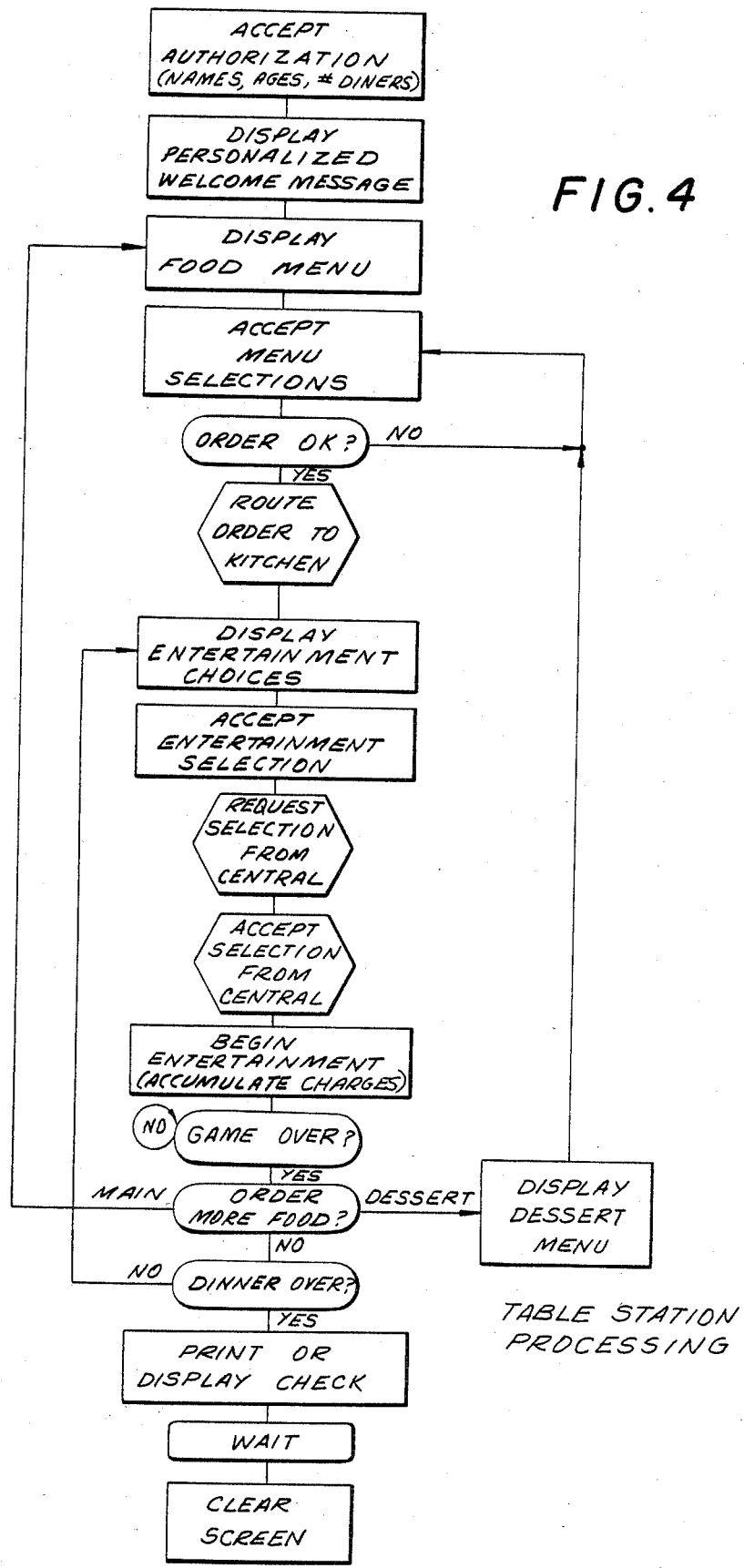
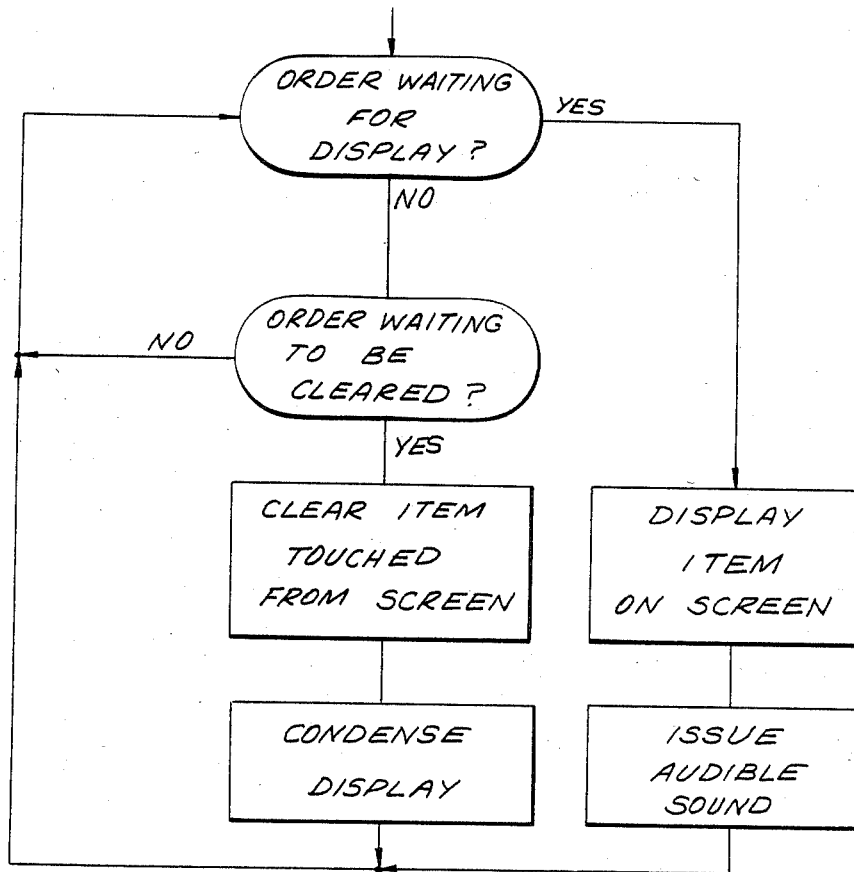


FIG. 5



SAMPLE DISPLAY FORMAT :

TIME ON	TIME LOGGED	TABLE	ITEM
7	6:47	21-1	BURGER PLATTER M/R
7	6:47	21-2	CHEFS SALAD RUSS
7	6:47	21-3	ONION RINGS
12	6:42	14-1	BEEF STEW
12	6:42	14-2	VEAL PARMIGIAN SPAG

KITCHEN/BAR PROCESSING

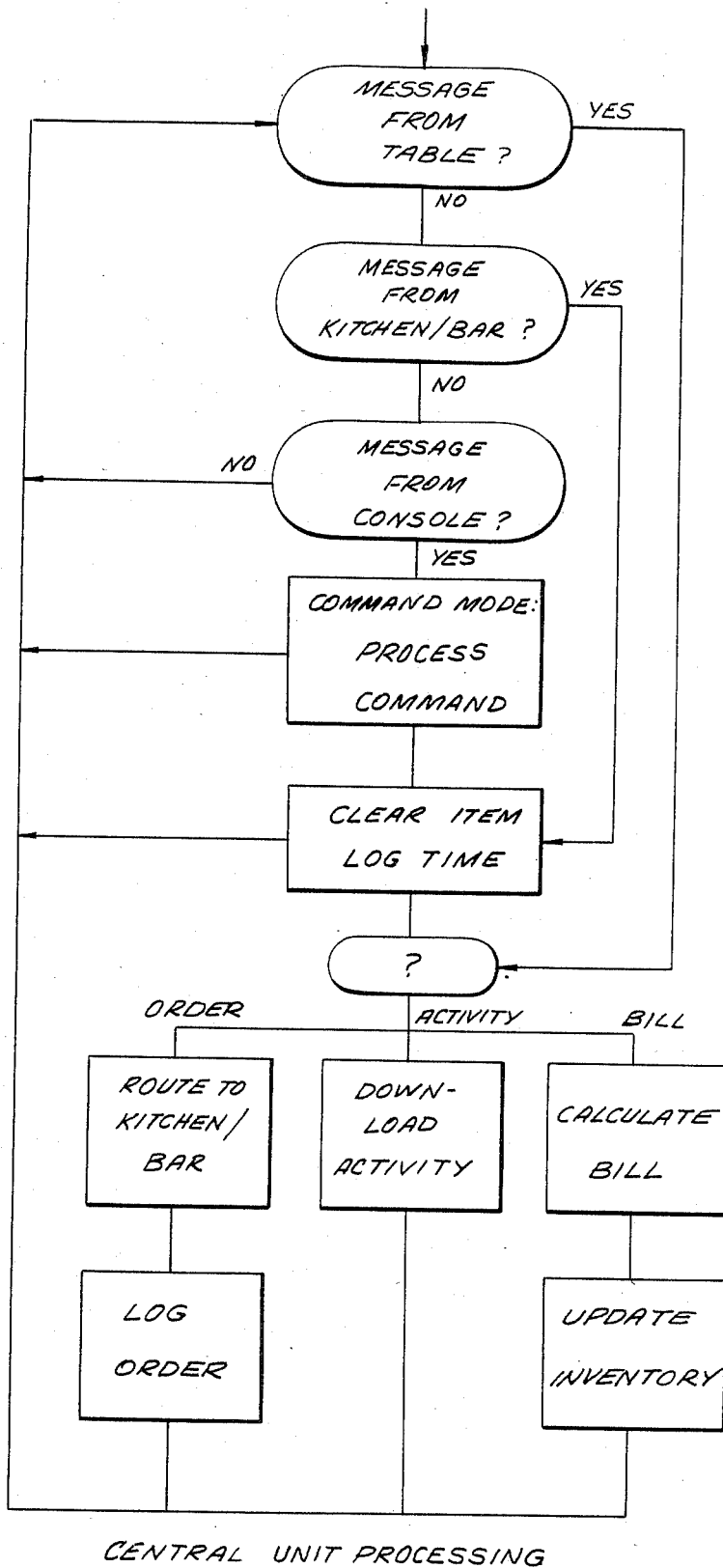


FIG. 6

FIG. 7

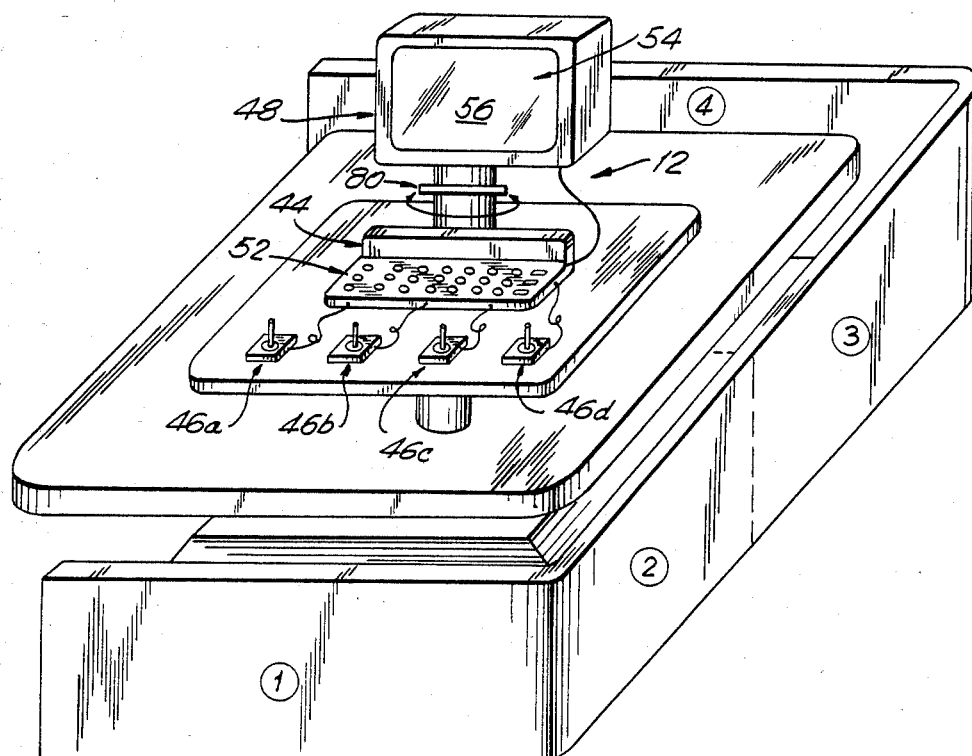


TABLE STATION 1

INTEGRATED INTERACTIVE RESTAURANT COMMUNICATION METHOD FOR FOOD AND ENTERTAINMENT PROCESSING

CROSS-REFERENCE TO RELATED APPLICATIONS

This application is related to the commonly owned copending U.S. patent application entitled "Integrated Interactive Restaurant Communication System for Food and Entertainment Processing", filed Mar. 14, 1983, bearing U.S. Ser. No. 474,984, and naming Lawrence G. Kurland, the named inventor herein, and Eli Gilbert as joint inventors, the contents of which is specifically incorporated by reference herein in its entirety.

TECHNICAL FIELD

The present invention relates to interactive communication systems and particularly to interactive restaurant communication systems.

BACKGROUND ART

Interactive communication systems, such as those employing video display terminals, are well known in the art, such as disclosed, by way of example, in U.S. Pat. Nos. 4,064,490; 4,054,911; 4,296,476; 4,143,360; 4,191,956; 4,122,519; 3,903,402; 4,001,807; 4,001,785; 4,075,686; 4,084,229; 4,251,291; 3,746,780; 3,668,307; 3,836,888; 3,752,908; 3,691,295; 3,968,327; 4,008,369; 4,028,733; 3,757,225; 3,814,841; 4,117,605; 4,264,925 and 4,164,024. Today, this field is a very active one, particularly with the advent and wide spread use of microprocessors in many business applications, as well as in personal computers, such as disclosed in U.S. Pat. No. 4,296,476. These prior art interactive systems have expanded dramatically in the field of cable television, such as disclosed in many of the aforementioned exemplary patents, both in two-way communication systems, such as disclosed in U.S. Pat. Nos. 3,691,295 and 3,668,307 by way of example, and in one-way cyclical transmission systems such as disclosed in U.S. Pat. Nos. 4,064,490 and 4,054,911 by way of example, as well as in real-time interactive systems such as disclosed in U.S. Pat. No. 4,264,925 by way of example. Moreover, point-of-sale terminals have become almost a total replacement for the cash register in today's businesses. The restaurant business, however, has been very slow to modernize in the area of service to the public except for modern fast-food operations, still relying on separate waitress service for order taking and separate entertainment features such as music boxes or video games, and the computer revolution has really not caught up. Rather, increased efficiency in restaurant service has primarily been focused on better manual systems, such as the one developed at the Tiffin Inn in Denver, Colo. where a food coordinator person in the kitchen would interface between the waitresses, who never left the dining room, and the cook using busboys to transport the orders which were assembled on carts. This system, however, although successful, was still a manual system and did not integrate food and entertainment functions. Thus, although cash register type of point-of-sale systems and automatic order entry systems have made in-roads into the restaurant business, as has distributed interactive video game technology, applicant is not aware of any prior art systems which have successfully integrated interactive entertainment and food functions so that a plurality of independent table station terminals, via

down-line loading and two-way communication with a central data base, can accomplish, among other things, transmission of orders to the kitchen, independently selectable down-loading of entertainment modules to the table stations for interactive play at the terminals; automatic computation of composite bills for both food and entertainment, and accumulation of detailed information for restaurant management.

These disadvantages of the prior art are overcome by the system of the present invention.

DISCLOSURE OF THE INVENTION

The present invention relates to an interactive integrated restaurant information communication system for enabling both individualized food and entertainment interactive information communication, such as two-way communication over a common transmission media, between a central remote data base and a plurality of different multipurpose table station terminals located at various table stations throughout the restaurant for use by the restaurant patrons for both food selection, based on down-line loaded food menu modules selectively retrieved from the central data base, and entertainment selection and interactive play, such as video games, also based on down-line loaded entertainment modules selectively retrieved from the central data base. The charges, where applicable, for both the entertainment and food selections can be compositely automatically billed to the table station, with each patron having a unique identification code for billing and/or service purposes, and the bills printed either centrally or at each table station.

The central data base comprises a central main computer which essentially performs the information routing functions, and remote retrievable storage for storing the various food and entertainment programs or sets of control instructions which are retrieved by the various table station terminals in response to selections made by the restaurant patrons, as well as handling coordination or processing and display of food orders in conjunction with kitchen and bar monitors, accumulation of restaurant management information and billing as well as other functions, if desired. Each of the table station terminals comprises a microprocessor and local storage which is down-line loaded with selected sets of control instructions from the central data base, under control of a master control program, in response to patron selection, and a local video display which is utilized to display data for food and entertainment selection as well as to interactively play the game or entertainment selected, with the microprocessor processing incoming data to enable food and entertainments selections to be transmitted to the central data base and to enable the retrieved entertainment to be interactively played at the terminal in response to the retrieved locally stored selected set of control instructions. The food orders, under control of the central computer, are collected from the various table station terminals, and displayed on central kitchen and bar monitors, with the orders being cleared from the monitor screen, such as by using touch-sensitive screens, as they are filled for each table station. If desired, the orders can be assembled by table station, each having its unique patron identification code, and transported to the pertinent table station. Thus, each table station terminal can independently serve to provide both food selection and entertainment functions, interactively with a central data base, in the

integrated restaurant communication system of the present invention, whereby the overall efficiency of the restaurant will be enhanced and better controlled.

BRIEF DESCRIPTION OF DRAWINGS

FIG. 1 is an overall functional block diagram of the presently preferred embodiment of the integrated interactive restaurant communication system of the present invention for providing integrated food and entertainment processing in accordance with the present invention;

FIG. 2 is an overall schematic block diagram of a typical table station terminal for use in the system of FIG. 1;

FIG. 3 is an illustrative diagram of a typical organization of the memory unit portion of the terminal of FIG. 2;

FIG. 4 is a condition responsive logic flow diagram of a typical table station terminal processing control program for the terminal of FIG. 2;

FIG. 5 is a condition responsive logic flow diagram of a typical kitchen or bar monitor processing control program for a typical kitchen or bar monitor in the system of FIG. 1;

FIG. 6 is a condition responsive logic flow diagram of a typical processing control program for the central computer of the system of FIG. 1; and

FIG. 7 is a diagrammatic perspective illustration of a typical table station in the system of FIG. 1.

BEST MODE FOR CARRYING OUT THE INVENTION

Referring now to the drawings in detail, and initially to FIG. 1 thereof, an overall functional block diagram of the presently preferred embodiment of the integrated interactive restaurant communication system, generally referred to by the reference numeral 10, of the present invention is shown. As will be described in greater detail hereinafter, the interactive restaurant communication system 10 provides integrated food and entertainment processing which enables restaurant patrons to obtain menus for individual food selection on a video monitor at their tables, individually enter their orders into a table station "intelligent" terminal at their tables, select from and interactively play a variety of remotely retrievable interactive entertainment activities using the video monitor while waiting for the food to arrive, and if desired, where applicable, have the food and entertainment charges automatically added to a composite bill which may be printed at the table station terminal or at a remote central location. Moreover, as will be further described herein, the system 10 of the present invention is preferably integrated into the overall restaurant operation so that the food orders placed at the various table station terminals, with five such table station terminals 12, 14, 16, 18 and 20 being shown by way of example in FIG. 1, can be collected and routed by a central computer 22 to the kitchen and bar for display on video terminals or monitors 24, 26 and 28, respectively, for the cook and bartender to enable the orders to be promptly filled, which orders, as illustrated in FIG. 5, can also be displayed by time logged into the system 10 to enable the oldest orders to be filled first. In addition, restaurant efficiency can also be monitored by comparing the time the order is logged with the time the order is filled.

As shown and preferred in FIG. 1, each of the table station terminals 12, 14, 16, 18, 20, which will be de-

scribed in greater detail hereinafter with reference to FIGS. 2-4 and 7, is preferably connected via a conventional two-way communications link, such as conventional RS-232 lines, to the central computer or processor 22 through a conventional multiplexed communication port 30. The quantity of table station terminals 12, 14, 16, 18, 20, which may be serviced or controlled by the central computer 22 is essentially limited primarily by the capacity of the central computer 22 and associated memory 32 which is selected, the desired functions to be performed and the desired access time, all of which are considerations within the ordinary skill of the art based on the system 10 description provided herein. The central processor 22 and associated memory 32, such as a conventional disk memory, comprise the central or remote data base 22-32 which is remotely accessed by the various table station terminals 12, 14, 16, 18, 20 in accordance with the present invention. By way of example, the central processor 22, which is selected to be compatible with the table station terminals 12, 14, 16, 18, 20 employed, may be one having the capabilities of a PDP 11/23 processor, available from Digital Equipment Corporation, having, by way of example, 128K of memory, with the associated memory 32 being, by way of example a data storage medium of the type such as available from Data Systems Design as its DSD 880 combination flexible and Winchester disk drive. In addition, by way of example, the communications multiplexers 30 may be of the type such as the DLV 11. As further shown and preferred in FIG. 1, a printer 34, such as a conventional dot matrix printer, such as an Anadex 9625A, is connected to the central processor 22 for printing bills, if desired, as well as restaurant management information and any other reports desired, such as market survey data. In addition, a conventional master console terminal 36, such as an ADM-3A, is preferably connected to the central processor 22 to enable, for example, review of restaurant operations and changes to be made in the data base 22-32. If desired, a conventional type of cash register point-of-sale terminal 38 may also be connected to the central processor 22.

With respect to the kitchen terminals 24, 26 and bar terminal 28 employed with the central processor 22, these terminals 24, 26, 28 are preferably conventional and employ touch-sensitive screens, such as the Elo-graphics E270, placed over the video monitor screen, to enable keyboard free data entry, such as, by way of example, to clear an item from the screen after the order has been filled or to indicate order status. The associated processing to be performed by these conventional data terminal 24, 26, 28, which are also selected to be compatible with the central processor 22 employed, if desired, may preferably be limited to requested order display and clearing of filled orders in response to data entry, such as via the aforementioned associated touch sensitive screen, with central computer 22 logging the time that the order is cleared from the screen if desired. Such an exemplary arrangement is illustrated in FIG. 5 which is a self-explanatory condition responsive logic flow diagram of a typical kitchen or bar monitor or terminal processing control program for a typical kitchen or bar terminal 24, 26, or 28, respectively, in the system 10 of the present invention, which may be conventionally programmed by one or ordinary skill in the art based on the description herein. It should be noted that although only two kitchen terminals 24, 26 and one bar terminal 28 are shown by way of example in the system of FIG. 1, any desired number of such terminals

may be employed dependent upon the needs of the restaurant and the selected capacity of the system 10.

Referring now to FIGS. 2-4 and 7, a typical preferred table station terminal, such as terminal 12 located at table station 1 in the above example, shall now be described in greater detail. As was previously referred to, the table station terminals 12, 14, 16, 18, 20 are selected to be compatible with the central processor 22 selected. In this regard, preferably the table station terminals are functionally similar to the type marketed, by way of example, by Atari, Inc. as its model 800, or such as the type described in U.S. Pat. No. 4,296,476, the contents of which is specifically incorporated by reference herein in its entirety, conventionally modified to provide a data communications interface or capability with the remote data base 22-32. In this regard, as shown and preferred in FIG. 2, which is essentially the type of terminal described in U.S. Pat. No. 4,296,476, which is a terminal employing a microprocessor 40 based data processor having a programmable graphics generator, the peripheral memory devices, such as a local disk drive unit and cassette unit, have been replaced by the down-line loading capability of the terminal 12 from the remote data base 22-32 into the local random access memory or RAM portion 42a of the terminal's memory unit 42 under control of a master control program or executive or supervisory program, such as illustrated by way of example in FIG. 4, stored in the local read only memory or ROM portion 42b of the terminal memory unit 42. The capacity of the ROM and RAM portions of terminal memory unit 42 is preferably conventionally selected so as to permit the desired functions to be performed by the table station terminal 12.

As shown and preferred in FIGS. 2 and 7, the table station terminal 12 includes a console 44, game control apparatus 46, such as the four conventional video game joystick controllers 46a-46d illustrated in FIG. 7, and a video display unit 48, which preferably is a conventional television monitor with console 44 providing a suitable radio frequency signal corresponding to a television raster scan signal to the video display unit 48 via line 50. The table station terminal 12 preferably has two basic modes of operation, the food selection mode and the entertainment mode in which the terminal 12 can function, by way of example, as an interactive video game unit. In the entertainment mode, as will be described in greater detail hereinafter, the table station terminal can be programmed, via down-line loading from the central data base 22-32, to provide various interactive entertainment activities such as interactive video games or interactive educational materials, such as involving text, diagrams and pictures displayed on video display unit 48, as well as audio. An interactive dialogue can be conducted between the restaurant patrons seated at the table station terminal 12, such as at table station 1 in the above example, using a conventional keyboard 52 and/or, if desired, a conventional touch sensitive screen 54, such as the previously mentioned Elographics E270, which could, if desired, preferably be placed over the normal video screen 56, to provide keyboard-free data entry. In either mode, the food selection mode or entertainment mode, the table station terminal is utilized by the restaurant patrons to store or retrieve information from the remotely located central data base 22-32. The video display unit 48 provides the restaurant patrons located at the table station 1 with graphics information, such as an alphanumeric

display and/or pictorial graphics, such as for games, that is conventionally formulated and transmitted to the video display unit 48 by the electronics contained in console 44 via the communicating line 50.

Apart from the master control program or executive or supervisory program for conventionally supervising the overall operation of data manipulation in the table station terminal 12, which is preferably permanently stored in ROM in memory unit 42, the desired operating programs or modules for providing the food selection and entertainment functions of the multipurpose table station terminal 12 are down-line loaded from the central data base 22-32 for storage in the RAM section 42a of memory unit 42 in response to a request therefor which has been input by the restaurant patron via the keyboard 52, or touch sensitive screen 54 if desired. Regardless of the mode in which the table station terminal 12 operates, the operation of the associated internal circuitry illustrated in FIG. 2 remains essentially the same. Thus, the operating program or module, whether a food menu module used for food selection or order entry, or an entertainment module used for entertainment selection and interactive game or other entertainment activity, can display portions of the requisite information or data on video display unit 48. Preferably, the table station terminal 12 includes a conventional programmable object graphics generator 58 which can be called upon to transfer graphics information from memory unit 42 to the video display unit 48. In this regard, the conventional microprocessor unit 40, under direction of the down-line loaded operating program or module, transfers a list of display instructions into the RAM section 42a (FIG. 3) of the memory unit 42, making the display instructions available to the microprocessor 40. The microprocessor 40 can modify portions of the display instructions, such as the addresses of the various instructions, thereby directing the object graphics generator 58 to those sections of the memory unit 42 containing the graphics information to be displayed, such as alphanumeric characters, lines, heading marks and the like. Alternatively, such as in the entertainment mode, the down-line loaded operating program, such as a down-line loaded entertainment module for an interactive video game, may require information to be displayed in graph-like form, such as a playfield display in the form of a cartesian or other coordinate system displayed to the restaurant patron(s) at the table station 1 via video display unit 48. Further, the down-line loaded operating program may call for a moveable cursor to be displayed in which instance the operating program would contain a block of graphics information containing the picture data for the vertical column that the movable object generator 58 constructs on the video display unit 48, including the picture data for the cursor. The microprocessor would then conventionally write into a movable object DMA counter the address of the location within the memory unit 42 of the block of graphics information containing the cursor picture data. The microprocessor 40 would also preferably write a data word to a DMA control register whose contents are used to conventionally inform a DMA control unit that movable object graphics will be displayed. Accordingly, as described in detail in U.S. Pat. No. 4,296,476 specifically incorporated by reference herein, the DMA control unit provides the moveable object DMA counter with signals that cause the DMA counter to sequentially address the memory locations of memory unit 42 containing the graphics information for

the cursor. The data used by the microprocessor 40 in the entertainment mode and/or with respect to cursor movement, in either mode, is preferably provided by the joysticks 46a-46d or operation of the keyboard 52, or the touch-sensitive screen 56, by the restaurant patrons seated at the table station 1. Typically, the joysticks 46a-46d, assuming by way of example one for each of up to four restaurant patrons to be seated at the table station 1, provide user generated position information, and possibly other game information, which is communicated to the microprocessor 40 via a conventional data interface unit 60. The conventional joysticks 46a-46d may typically be of the type disclosed in U.S. Pat. No. 4,091,234, the teachings of which are specifically incorporated by reference herein in its entirety. In addition to the above, as further shown and preferred in FIG. 2, the terminal 12 also includes a conventional audio generator 62 for generating audio signals, a conventional timing unit 64 for generating the various timing signals required such as those illustrated, by way of example, for use in the system of U.S. Pat. No. 4,296,476, a conventional video summer circuit 66 for conventionally summing the color/luminance and composite sync signal output of the object graphics generator 58, and a conventional RF modulator 68 which provides the displayable video signal to the video display unit 48 via line 50. Furthermore, an optional conventional printer unit 70 may also be provided at the table station for local printing of bills, or other data if desired, with the printer 70 being connected to the microprocessor 40 via the data interface unit 60 and a serial (I/O) bus through which the terminal 12 is connected to the central data base 22-32. It should be noted that after the desired food menu module or operating program or entertainment module or operating program is requested by and down-line loaded to the terminal 12, the microprocessor 40 can then process incoming data provided via the keyboard 52, player controls 46 or touch sensitive screen 56 in accordance therewith so that food orders may be selected and transmitted back to the central data base 22-32 for billing and filling and/or interactive games may be played. In this regard, FIG. 3 provides an illustrative diagram of a typical organization of the memory unit 42, such as described in U.S. Pat. No. 4,296,478, and FIG. 4 provides a self-explanatory condition responsive logic flow diagram of a typical processing control program for the table station terminal 12, which may be conventionally programmed by one of ordinary skill in the art based on the description herein.

As shown and preferred in FIG. 7, by way of example, preferably each table station has a unique position identification for each restaurant patron to facilitate accurate order filling and billing. Thus, assuming four restaurant patrons per table station by way of example, the various restaurant patrons at table station 1 illustrated in FIG. 7 would identify themselves to the table station terminal 12 as 1-1, 1-2, 1-3 and 1-4. In addition, by way of example, the video display unit 48 could be mounted on a universal swivel 80 so as to be rotatable for viewing by all of the patrons at the table station 1, with the keyboard 52 and joysticks 46a-46d being connected to the terminal 12 via conventional electronic umbilical cords to facilitate use by each patron at the table station 1. It should be noted that the table station configuration shown in FIG. 7 is merely illustrative and many other such configurations will readily occur to one of ordinary skill in the art.

With respect to the operation of the central data base 22-32 in connection with the processing of food menu and entertainment requests, including the down-line loading of food menu modules or operating programs and entertainment modules or operating programs to the requesting table station terminals 12, 14, 16, 18, 20, the logging and routing of orders to the kitchen and bar terminals 24, 26, 28, the calculation of composite bills, the updating of inventory and the overall processing of messages, a condition responsive logic flow diagram of a typical processing control program for the central computer 22 for accomplishing the above is shown in FIG. 6. The program itself can conventionally readily be written by one of ordinary skill in the art based on the description herein.

Summarizing the overall operation of the preferred system 10 of the present invention, either after or before the diners or restaurant patrons are seated at their respective table station, the table station terminal would preferably display a welcome message and a list of choices for the patrons. The patrons would then select a choice by touching a location on the touch sensitive screen or by pressing keys on the keyboard. The initial choices may be type of food, type of meal, etc. Thereafter, by stepping through various displays, the patrons, using their unique position identification code, could order their meals which would be transmitted to the central data base for billing and filing. The patrons would then be allowed to select entertainment such as games, educational programs, computer generated art displays, etc., and possibly television channels if desired by using a television receiver as the video display unit. Associated charges for both food and entertainment, where applicable, would automatically be added to a composite bill. The central computer downloads the selected activity to the requesting table station terminal, with each terminal being capable of independently requesting and receiving its own operating programs from the central data base. The patrons would then be free to interact with the downloaded operating program, whether it were a food menu module or an entertainment module. If desired, an initialization program could be run daily, or at any time, to set price changes, menu changes, entertainment changes, current time and date, enter deliveries, enter daily specials, etc.

By utilizing the system of the present invention, restaurant efficiency and flexibility are enhanced and better control can be achieved.

It is to be understood that the above described embodiments of the invention are merely illustrative of the principles thereof and that numerous modifications and embodiments of the invention may be derived within the spirit and scope thereof.

What is claimed is:

1. A method of interactively integrating food and entertainment patron selection and interaction for a plurality of table stations in a restaurant communication system comprising the steps of remotely retrievably storing and processing information comprising a plurality of displayable food menu modules and entertainment modules at a central data base, providing a plurality of down-line loadable multipurpose table station terminals at a plurality of said table stations, each of said table station terminals comprising microprocessor means, local storage means, selection and control means and local common video type display means capable of providing a video display of patron selected food menu and entertainment modules, each of said retrievably

stored food and entertainment modules comprising a different set of control instructions corresponding to a different patron selectable food menu or entertainment function, said microprocessor means being operable in accordance with a locally stored one of said selected sets of control instructions;

individually patron selecting a displayable remotely stored food menu module via said table station terminal selection means;

down-line loading said patron selected corresponding set of control instructions to the selecting table station terminal for local storage thereof;

locally displaying food menu selection data corresponding to said patron selected food menu module on said common video type display means in response to said down-line loaded patron selected set of control instructions, said video displayable food menu module data comprising menu selection data corresponding to a choice of food items to be ordered; individually patron selecting at least one of said video displayed food items for remotely patron ordering said selected food items; transmitting said at least one patron ordered food selection back to said central data base for central processing of said order; individually patron selecting a displayable remotely stored entertainment module via said food ordering table station terminal for video display on said common video type display means at said food ordering table station terminal; down-line loading said patron selected corresponding set of control instructions to said food ordering table station terminal for local storage thereof; locally video displaying entertainment data corresponding to said patron selected entertainment module on said food ordering table station terminal common video type display means in response to said down-line loaded selected set of control instructions to said food ordering table station terminal for enabling interaction with said video displayed entertainment data while said food order is being processed; whereby each table station terminal may independently provide both interactive entertainment and food selection on a common video type display terminal in cooperation with a central data base.

2. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 1 further comprising the step of compositely billing charges for a given food ordering table station based on said food and entertainment selections of said down-line loaded information.

3. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 2 wherein said composite billing step comprises the step of centrally printing the composite bill for each of said table stations.

4. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 2 wherein said composite billing step com-

prises the step of locally printing the composite bill for said given table station.

5. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 1 wherein said remote storing step comprises the step of remotely retrievably storing interactive video games as at least a portion of said entertainment modules.

6. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 1 wherein said remote storing step comprises the step of remotely retrievably storing interactive video games as at least a portion of said entertainment modules, said local video display step further comprising the step of locally video displaying a remotely retrieved interactive video game on said food ordering table station terminal common video type display means as said selected entertainment data.

7. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 1 further comprising the step of uniquely identifying each restaurant patron at said food ordering table station and transmitting said unique patron identification to said central data base with at least said food selection information for enabling correlated food distribution based on said central processing thereof.

8. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 7 further comprising the step of compositely billing charges for a given food ordering table station based on said food and entertainment selections of said down-line loaded information.

9. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 1 wherein said central processing step comprises the steps of collecting each food selection from each of said plurality of table station terminals where a selection is made and centrally processing said collected food selections for distribution to said selecting plurality of table stations.

10. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 9 further comprising the step of uniquely identifying each restaurant patron at said food ordering table station and transmitting said unique patron identification to said central data base with at least said food selection information for enabling correlated food distribution based on said central processing thereof.

11. A method of interactively integrating food and entertainment selection for a plurality of table stations in a restaurant communication system in accordance with claim 10 further comprising the step of compositely billing charges for a given food ordering table station based on said food and entertainment selections of said down-line loaded information.

* * * * *

- [54] HIERARCHICAL KNOWLEDGE SYSTEM
- [75] Inventors: James S. Bennett; Jay S. Lark, both of Palo Alto, Calif.
- [73] Assignee: Teknowledge, Inc., Palo Alto, Calif.
- [21] Appl. No.: 628,817
- [22] Filed: Jul. 9, 1984
- [51] Int. Cl.⁴ G06F 15/24
- [52] U.S. Cl. 364/403; 364/468; 364/478; 235/385; 29/703
- [58] Field of Search 364/400-401, 364/403, 468-469, 478, 513, 518; 209/1-2, 546, 552; 235/385; 29/33 K, 33 R, 400 R, 400 M, 428-431, 469, 564, 568, 700-703, 711

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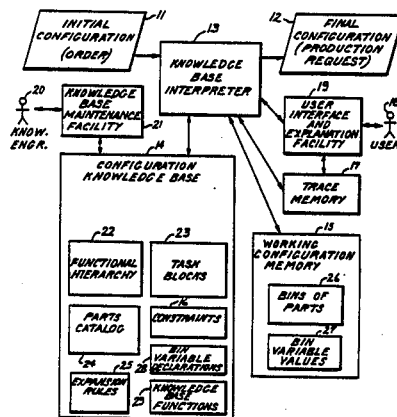
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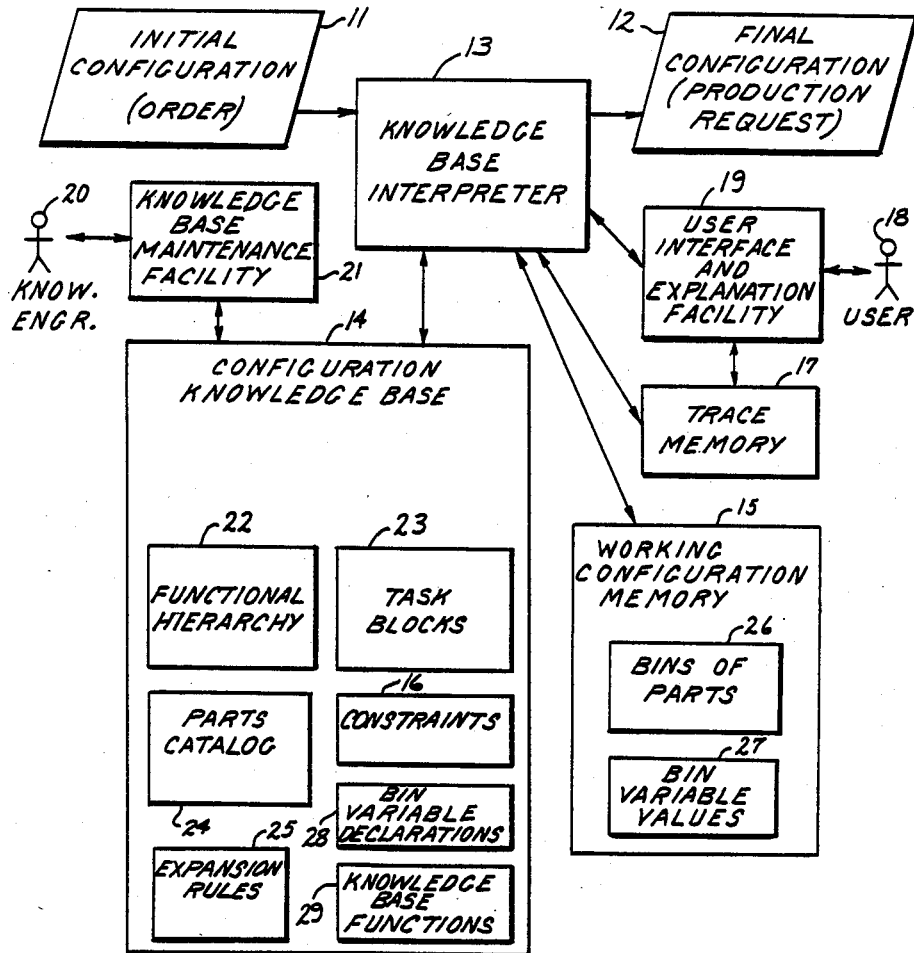
Primary Examiner—Gary V. Harkcom
 Attorney, Agent, or Firm—Leydig, Voit & Mayer

[57] ABSTRACT

A knowledge system has a hierarchical knowledge base comprising a functional decomposition of a set of elements into subsets over a plurality of hierarchical levels, a plurality of predefined functions or conditions of the elements within the subsets of a plurality of the hierarchical levels, and a predefined set of operations to perform on a user-defined set of elements responsive to the functional knowledge base. Preferably, the knowledge base is defined declaratively by assigning parent sets to offspring subsets to define the hierarchy, by indicating the conditions of the subsets which satisfy the predefined functions and by writing task blocks in an imperative language defining the sequence of operations to perform on the user-defined set of elements. Preferably the operations include matching, configuring and expanding the user-defined set of elements into the defined subsets of individual elements and evaluating the predefined functions, and the operations are executed recursively. In a specific embodiment the elements are available components for a system or item of manufacture, and the subsets of elements are sub-assemblies or functionally related components. The predefined functions define condition-action constraints to insure that the sub-assemblies have compatible components. Such a knowledge system has general applicability, is easy to maintain and incrementally modify, has transparent representation of the functional decomposition and the configuration operations, and provides explanation for an assessment of the configuration.

70 Claims, 13 Drawing Figures





10 ↗

Fig. 1.

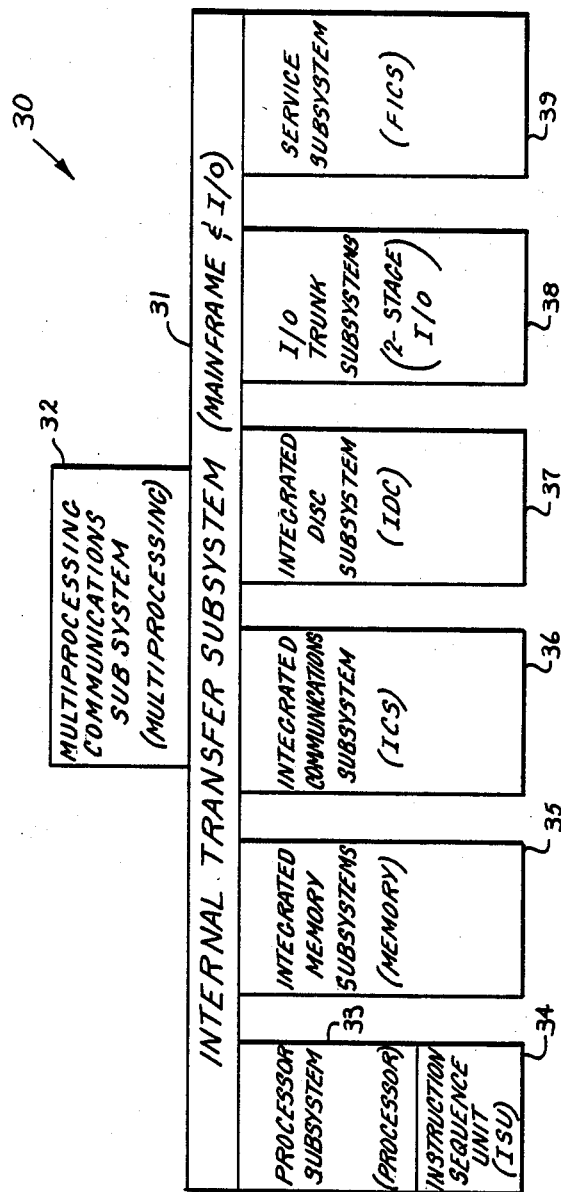


FIG. 2.

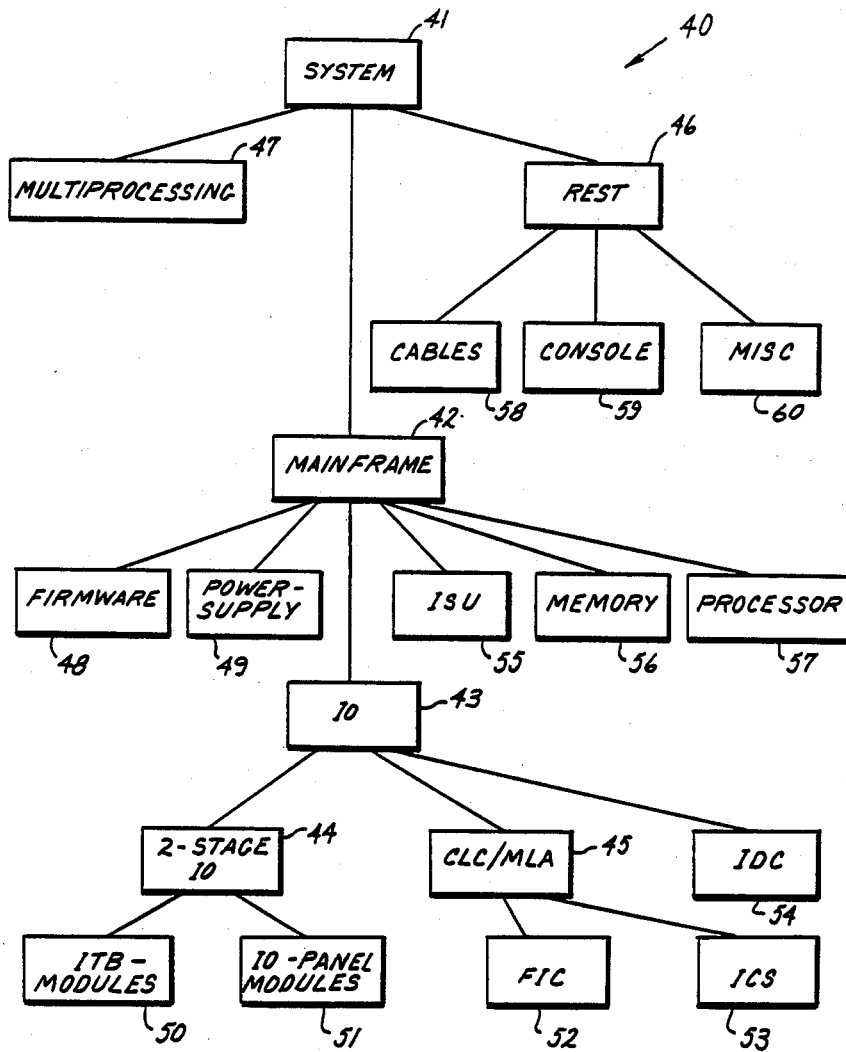


Fig. 3.

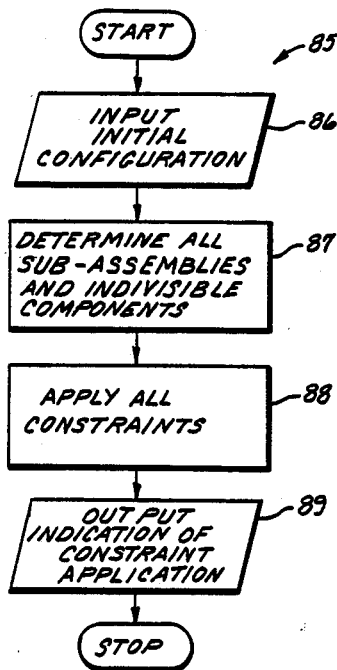
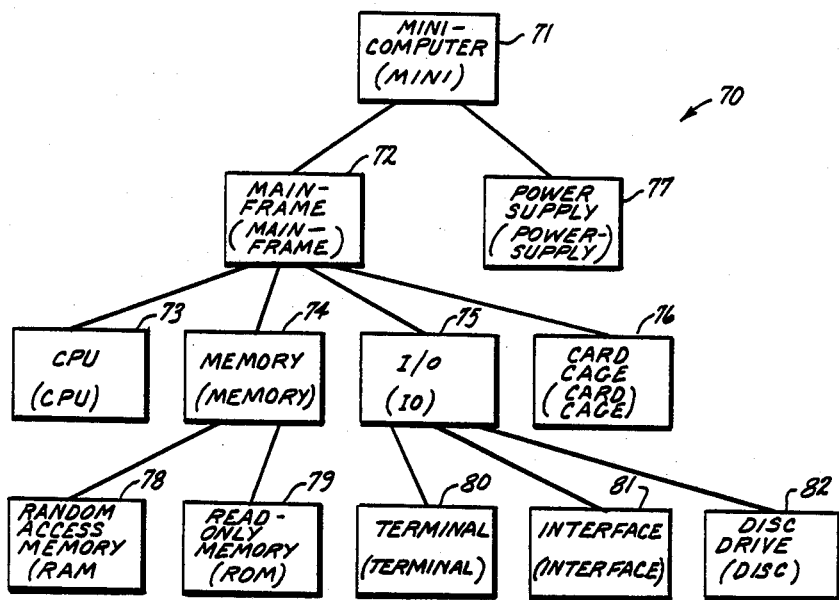


FIG. 4.

FIG. 5.

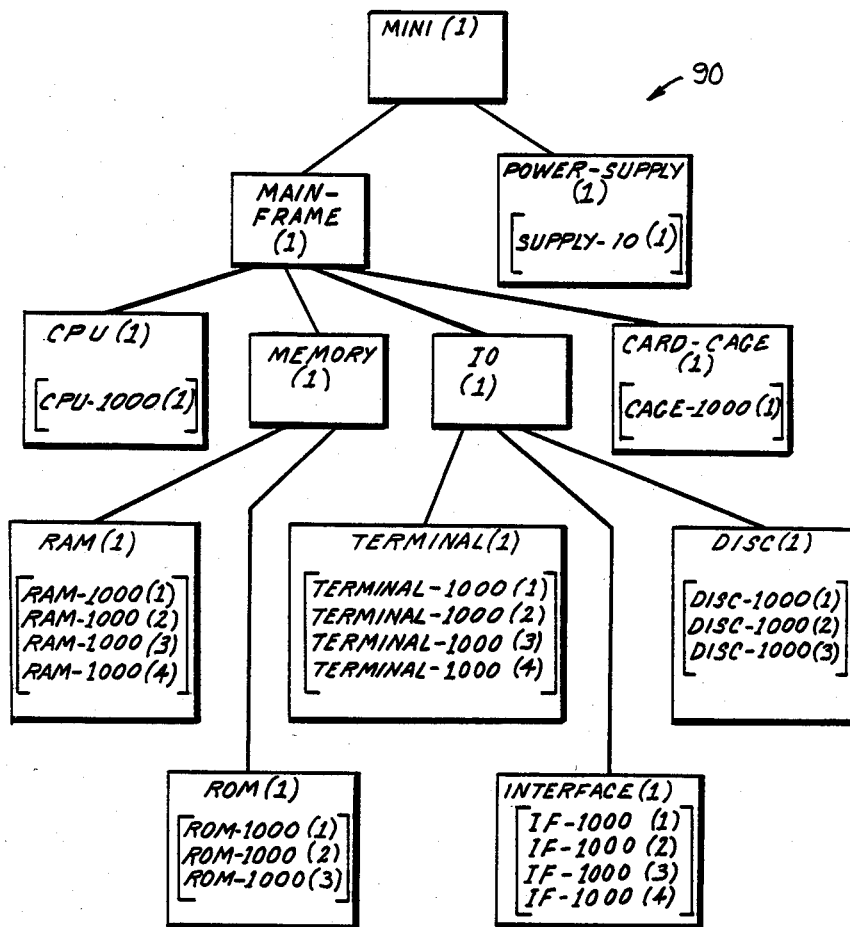


FIG. 6.

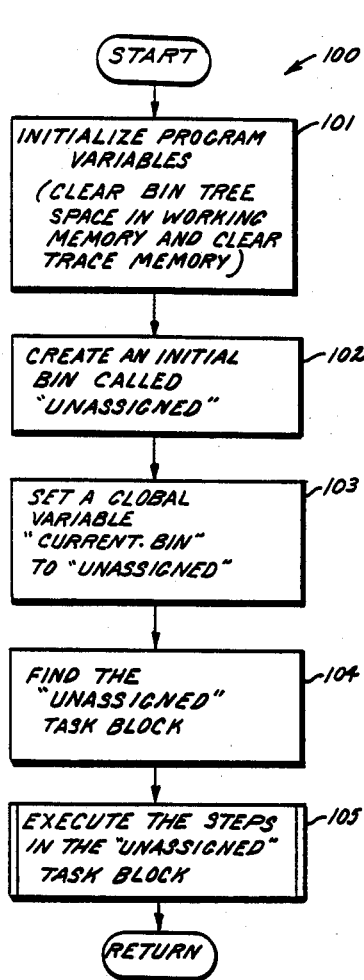


FIG. 7.

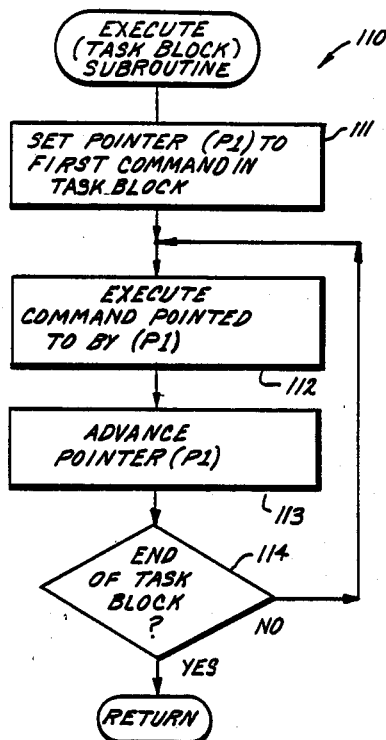
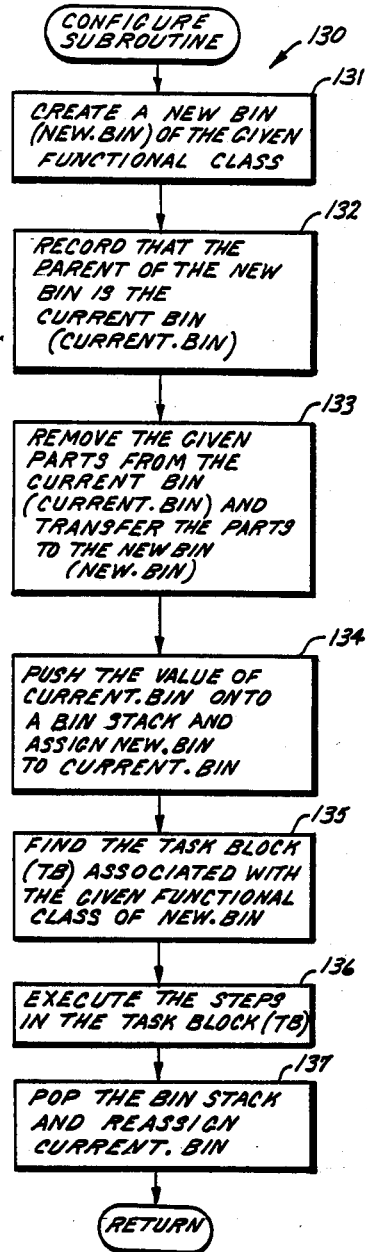
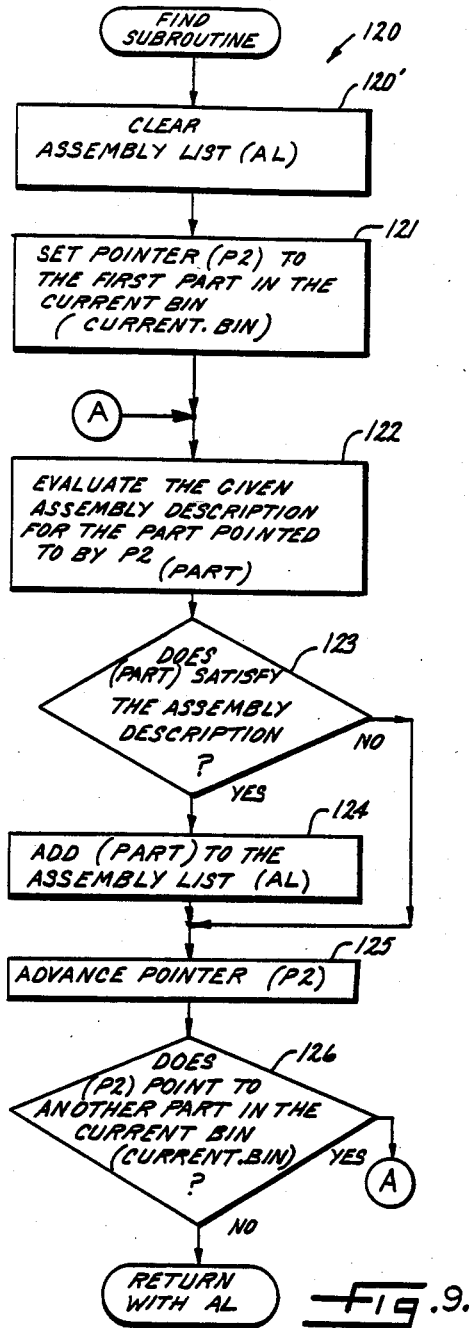


FIG. 8.



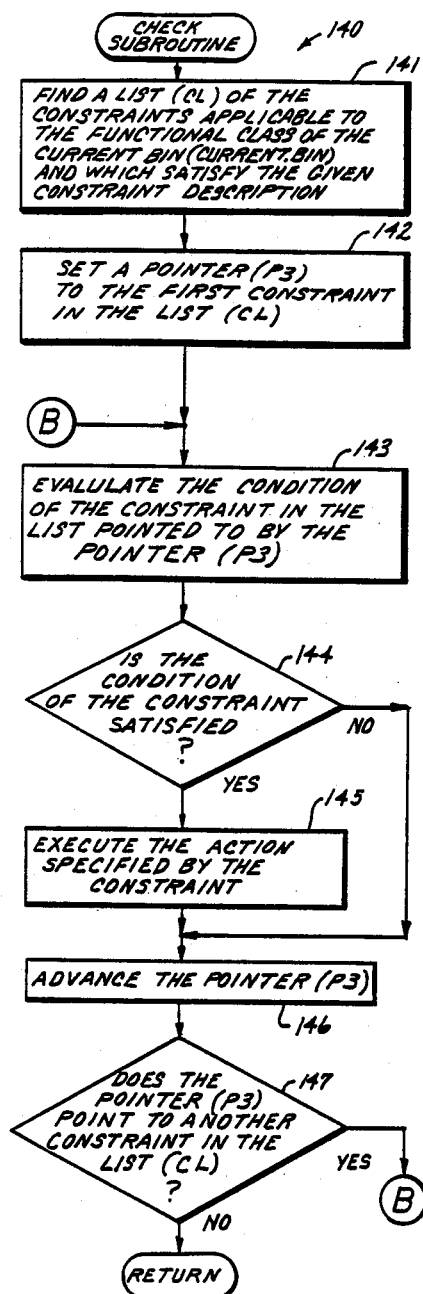


Fig. 11.

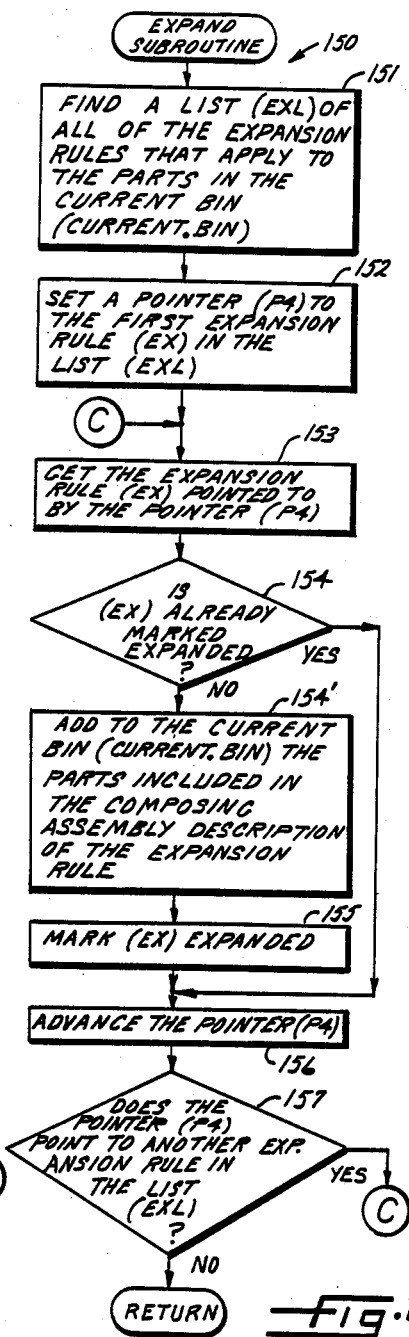


Fig. 12.

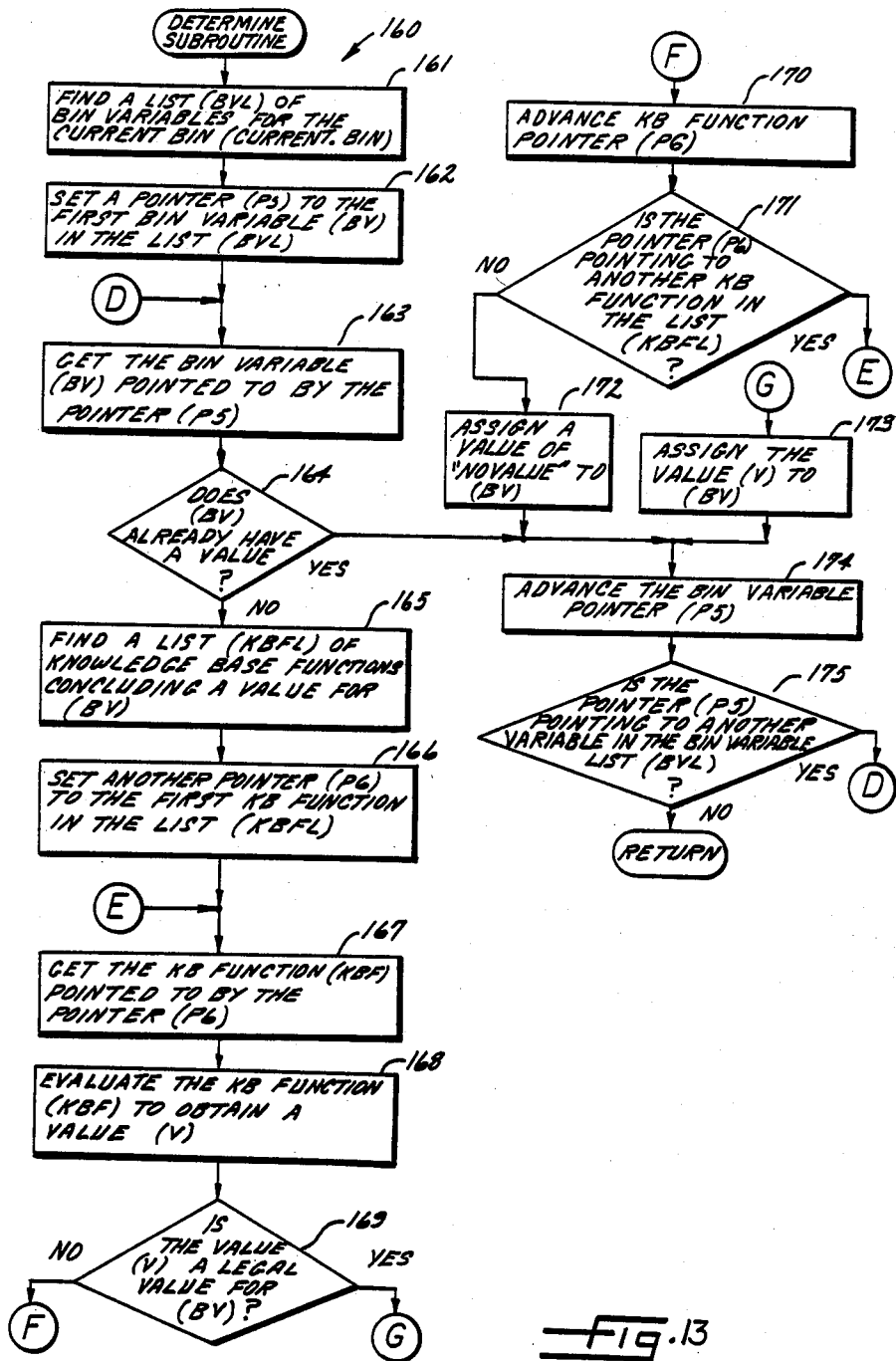


FIG. 13

HIERARCHICAL KNOWLEDGE SYSTEM

FIELD OF THE INVENTION

The present invention relates generally to inventory control and the processing of orders for flexibly assembled systems or items of manufacture, and more specifically to computer systems to aid in the checking of orders for products or systems to be manufactured or assembled.

BACKGROUND OF THE INVENTION

The traditional method of mass production realizes economy of scale due to continuous application of capital and division of labor. In recent years, however, a method of "flexible assembly" has become popular for relatively expensive and complex items such as computers. For these kinds of products, the additional cost of assembling a unique system for a particular customer is offset by the omission of parts that are not desired by the customer. The flexibly assembled product becomes a better match to a specific customer's need than a standard product designed for a particular group of customers.

The typical method of handling an order for a flexibly assembled product is to split the order into secondary orders or "production requests" for groups of standard parts or "kits" that are functionally related or comprise sub-assemblies. Since the production requests are standardized, the flexibly assembled item is sold in terms of features or "packages" which can be made up of the kits. The process of converting the desired packages or features into corresponding production requests is a process of expansion from a list of packages to a list of kits. As this expansion process is carried out, the order is also validated or checked for conformance with marketing and engineering constraints on permissible configurations. In many cases, violations of the constraints are not evident when the order is submitted and are revealed only when the order is partially or totally expanded.

Computer systems have been used previously for expanding and validating orders. The expansion and validation, for example, is part of an inventory control system. Using traditional computer programs, the computer system performs a syntactic check on the list of package numbers comprising an order. The computer system may also use product definition files for separate versions of a generic product as the basis for expanding packages into kits, and may check whether the individual package numbers are legal options within a particular hierarchy under at least some conditions. But global engineering constraints such as power and interconnection are not performed by these computer systems using traditional programming techniques. Although it would be possible to program constraints into the computer program for a specific set of product hierarchies, the computer program would not be easy to update and maintain since the computer program would have to be rewritten for each change in the product hierarchy. Moreover, the constraints are relatively complex and change frequently. It is, therefore, not economical to include global constraint checking in these conventional computer systems.

In order to provide a maintainable computer system for running expansion and validation of an order, knowledge-based systems have been considered wherein the knowledge of what the constraints are and

when the constraints are to be applied is encoded declaratively in a knowledge base. A separate knowledge base interpreter interprets the knowledge base to apply the constraints. Consequently, the knowledge base interpreter need not be modified when the knowledge in the knowledge base is updated or changed. Knowledge systems have been used in general for problems that require diagnosis, recommendation, selection or classification. These problems have traditionally been performed by human experts and are not easily implemented using conventional programming techniques.

Presently there are highly developed commercial tools which may be used to build knowledge systems. The well-known commercial tools (such as KS300 manufactured by Teknowledge Inc., 525 University Avenue, Palo Alto, Calif. 94301) are patterned after a tool called EMYCIN described in *The Emycin Manual* by Van Melle et al., Standford University Report No. STAN-CS-81-885, Standford, Calif. 94305 (October, 1981).

EMYCIN is specifically designed as a domain-independent system for constructing rule-based consultant expert system programs. Domain knowledge is represented in EMYCIN systems primarily as condition-action production rules which are applied according to a goal-directed backward-chaining control procedure. Rules and consultation data are permitted to have associated measures of certainty, and incomplete data entry is allowed. The EMYCIN system includes an explanation facility displaying the line of reasoning followed by the consultation program, and answers questions from the client about the content of its knowledge base. To aid the system designer in producing a knowledge base for a specific domain, EMYCIN provides a terse and stylized language for writing rules; extensive checks to catch common user errors, such as misspellings; and methods for handling all necessary indexing chores.

In addition to production rules, the knowledge base for an EMYCIN system includes a hierarchical structure called a "context tree." The elemental representation of an object or idea is defined as a context-parameter-value triple. The context refers generally to an instance of a particular context type, the parameter refers to an attribute of the context instance and the value refers to the particular value of the parameter for the particular context instance. The context tree is defined by parent and offspring declarations for the context types.

The instantiation of contexts is similar to the invocation of a subroutine for each context, the subroutine in effect being defined by various declarations in the context definition. A consultation is started by instantiating a root context and the branches from this root context define major steps in the consultation during which the offspring contexts of the root node are instantiated. Thus, the context definitions are used to structure the data or evidence required to advise a user about the root context. Besides consultation control, the context tree may be used to organize the distinguished components of some object, or for representing distinguished events or situations that happen to an object.

A rule-based knowledge system has been developed for determining a computer system's configuration from a skeletal specification. A system called XCON has been used by the Digital Equipment Corporation when salesmen take orders and when processing clerks check

and flesh out incoming orders for VAX-11/780 computer systems. Given a customer's order, XCON determines what, if any, modifications have to be made to the order for reasons of system functionality and produces a number of diagrams showing how the various components on the order are to be associated.

A version of XCON called R1 is described in John McDermott, *R1: A Rule-Based Configurer Of Computer Systems*, Department of Computer Science, Carnegie-Mellon University, (April, 1980). It is said that R1 has sufficient knowledge of the configuration domain and of the peculiarities of the various configuration constraints that at each step in the configuration process, R1 simply recognizes what to do. Consequently, it is said that little search is required in order for R1 to configure a computer system.

Although the XCON rule-based system represents a major improvement over traditional computer techniques in the field of computer configuration, it is rather difficult to change the constraint rules in XCON. In particular, to change a constraint rule, one needs to know what occurs before and after the constraint rule is applied. In general, the constraint rules are interlaced and spread, and they are not readily accessible for maintenance since in general they are not conceptually hierarchical.

SUMMARY OF THE INVENTION

The primary object of the invention is to provide a knowledge system for generalized representation and processing of hierarchical assemblies that is easily maintainable and extensible.

Another object of the invention is to provide an intelligible knowledge base representation for hierarchical assemblies and their functionality.

Still another object of the invention is to provide an intelligible knowledge base representation for configuration strategies and actions.

Still another object of the invention is to provide a knowledge system for configuration checking which clearly separates the definition of configuration constraints from configuration checking strategies and actions.

And yet another object of the invention is to provide an improved control procedure for configuration checking.

And still another object of the invention is to provide an explanation facility especially suited for configuration analysis and adaptation.

Briefly, in accordance with the broadest aspect of the invention, a knowledge system for generalized representation and processing of hierarchical assemblies has a hierarchical knowledge base comprising a decomposition of a set of elements into subsets over a plurality of hierarchical levels, a plurality of respective predefined functions or conditions of the elements within the subsets at a plurality of the hierarchical levels, and a predefined set of operations to perform on a user-defined set of elements responsive to the knowledge base. For ease of maintenance and extensibility, the knowledge base is defined declaratively by assigning parent sets to offspring subsets to define the hierarchy, by indicating the conditions of the subsets which satisfy the predefined functions, and by writing task blocks in an imperative language defining the sequence of operations to perform on the user-defined set of elements. Preferably, the operations include operations for matching, configuring and expanding the user-defined set of elements into the

defined subsets of individual elements and for evaluating the predefined functions, and the operations are executed recursively.

In accordance with a preferred embodiment of the invention, a knowledge-based configuration system has separate portions encoding a configuration checking strategy, a description of hierarchical functions of the product to be configured, a catalog of the parts and components which implement those functions, assembly constraints that check whether a given set of components will successfully implement their respective hierarchical functions. The configuration system applies the assembly constraints during configuration checking and if necessary, warns the user or modifies the given set of components to insure compliance with the assembly constraints.

The control knowledge is encoded in an imperative language defining specific configuration control actions to be executed during interpretation of the control knowledge according to a built-in control procedure. Since the configuration control knowledge is explicit and results in modification of the given configuration only in a precisely defined fashion, the configuration system can always explain its conclusions and reasoning, and is intelligible and modifiable.

To provide transparent representation of the control knowledge as well as factual knowledge, the knowledge base is preferably organized into distinct portions including the assembling constraints, task blocks encoding the control strategy knowledge, functional hierarchy classes which become instantiated into bins, bin variables which take on values describing the bins and their contents, the catalog of parts, rules for expanding sub-assemblies into parts, and user-defined functions for computing the values of bin variables. The assembly constraints may be defined in terms of specified bin variables and certain built-in functions responsive to the number of specified parts in specified bins.

A particular task block is provided to be executed at the start of a configuration. The knowledge engineer must provide task blocks which are invoked after instantiation of specified functional hierarchy classes and when the values of the bin variables are to be determined, and may provide explicitly invoked task blocks.

The functional hierarchies are preferably represented by entries defining specific functions that are performed by the configured product, and each entry can be composed of more specific functional entries. Associated with each functional entry is a configuration strategy for checking that the parts in a given configuration will implement the desired function. The system will warn the user when the configuration is incorrect, and can optionally modify the configuration by adding, deleting, and modifying parts in the respective bin for the functional hierarchy. The configuration strategy also describes the order that parts that achieve certain functions should be checked, and describes when to expand sub-assemblies into the same bin of their respective functions.

An explanation facility is provided to allow the user to validate the acceptability of the configuration, to identify unacceptable portions of the configuration, and to identify modifications made to the configuration to satisfy the constraints. A tracing and translation facility is provided to expose the specific steps and the knowledge used in the configuration checking process so that the validity of the configuration strategy is apparent.

BRIEF DESCRIPTION OF THE DRAWINGS

Other objects and advantages of the invention will become apparent upon reading the following detailed description and upon reference to the drawings in which:

FIG. 1 is a schematic diagram of a knowledge system incorporating various aspects of the present invention;

FIG. 2 is a block diagram of an exemplary M1234 computer system which is configured by the knowledge system in FIG. 1 using the knowledge base shown in Appendices II (A)-(E) to the present specification;

FIG. 3 is a tree diagram of the functional hierarchy for the M1234 computer system shown in FIG. 2;

FIG. 4 is a tree diagram of a functional hierarchy for a standard mini computer;

FIG. 5 is a simplified flowchart for the configuration process performed by the knowledge system in FIG. 1;

FIG. 6 is a diagram of a bin tree corresponding to the functional hierarchy of FIG. 4 for the mini computer, including the final configuration of parts in the bins for the working example of Appendix III (A);

FIG. 7 is a flowchart for the executive program in the built-in control procedure of the knowledge system shown in FIG. 1;

FIG. 8 is a flowchart of an EXECUTE subroutine for executing task blocks of control procedure steps;

FIG. 9 is a flowchart of a FIND subroutine for obtaining a list of parts included in a specified bin which match or satisfy a specified assembly description;

FIG. 10 is a flowchart of a CONFIGURE subroutine for creating a new bin of a specified functional class, transferring a specified list of parts from a specified current bin to the new bin, and temporarily passing execution to the task block associated with the functional class of the new bin;

FIG. 11 is a flowchart of a CHECK subroutine for applying a specified subset of the constraints applicable to the functional class of the current bin;

FIG. 12 is a flowchart of an EXPAND subroutine for expanding the parts in the current bin into their component parts by applying predefined expansion rules associated with the current bin; and

FIG. 13 is a flowchart of a DETERMINE subroutine for determining values for bin variables associated with the current bin by applying knowledge base functions associated with the current bin.

While the invention is susceptible to various modifications and alternative forms, a specific embodiment thereof has been shown by way of example in the drawings and will herein be described in detail. It should be understood, however, that it is not intended to limit the invention to the particular form disclosed, but on the contrary, the intention is to cover all modifications, equivalents, and alternatives falling within the spirit and scope of the invention as defined by the appended claims.

DESCRIPTION OF THE PREFERRED EMBODIMENTS

Turning now to FIG. 1, there is shown a block diagram generally designated 10 of a knowledge system for processing an initial configuration of elements or an order 11 to arrive at a final configuration or production request 12. The system 10 is recognized as a knowledge system since it has a domain-independent knowledge base interpreter 13 executing a built-in control procedure for interpreting a domain-dependent knowledge

base 14. The knowledge base 14 encodes a generic configuration in a highly structured and transparent format. In general terms, the knowledge base interpreter 13 matches the elements of the initial configuration in the order 11 to the generic configuration in the knowledge base 14 to structure or configure the elements according to the generic configuration. The elements structured in terms of the generic configuration are stored in a working configuration memory 15.

Once the elements or parts of the order 11 are structured according to the generic configuration, the knowledge base interpreter 13 may apply functions or constraints 16 stored in the knowledge base 14 to the working configuration in the memory 15. The constraints 16, for example, indicate particular aspects of the working configuration 15 that are not apparent from the initial configuration or order 11. The constraints 16 may also indicate desirable changes to make to the working configuration 15. The knowledge base interpreter 13 may execute these indicated changes in order to change the working configuration. At the end of processing by the knowledge base interpreter 13, these executed changes are reflected in the final configuration 12.

The knowledge system 10 has a trace memory 17 for keeping a precise record of the steps performed during the matching of the elements of the initial configuration 11 to the generic configuration in the knowledge base 14 and during the application of the constraints 16 to the working configuration 15. To enable a user or order analyst 18 to inspect relevant portions of the trace 17, a user interface and explanation facility 19 is provided. The user 18 may, for example, request a complete typescript of the contents of the trace memory 17 in an easily understandable form. An example typescript is included in Appendix I at the end of the present specification.

The knowledge base 14 is arranged so that it is easily modified and maintained to reflect desired changes in the generic configuration. The modification and maintenance is performed by a knowledge engineer 20 and the knowledge base 14 is accessed through an editor 21 called a knowledge base maintenance facility. The knowledge engineer 20 initially creates the knowledge base 14 collecting and encoding knowledge about the generic configuration according to a precise format and syntax so that the knowledge is transparent to the knowledge engineer and is machine readable by the knowledge base interpreter 13. For the processing of an order of parts 11 to generate a production request 12 for a system or product, the knowledge engineer 20 collects engineering and marketing data on the generic configuration of the product and its constraints and encodes the information according to the syntax recognized by the knowledge base interpreter 13.

The specific format of the knowledge base 14 is chosen so that the knowledge base can be easily updated to reflect the addition of new products and changes in current products. The generic configuration is represented as a functional hierarchy 22 which describes the product in various degrees of particularity. When a product is modified, only the particular portions of the functional hierarchy 22 need be changed. Moreover, since the hierarchy is in terms of the function of performance of parts of the generic configuration for the product, the changes required in the functional hierarchy 22 are minimal.

For each function in the functional hierarchy 22, a separate task block 23 is provided to encode corresponding steps in the control procedure followed by the knowledge base interpreter 13 to match the parts or elements of the initial configuration 11 to the functional hierarchy 22. A change in the functional hierarchy usually requires only a localized change in the respective task block 23. Also, the constraints 16 correspond to particular functions in the functional hierarchy 22 so that the constraints 16 are also easily modified. Consequently, the constraints associated with a particular function in the functional hierarchy are applied in a manner specified by the respective task block 23.

It should be noted that even though the functional hierarchy 22 might not change during the life of a particular product, the parts or elements making up the product may change during the life of the product. These changes are reflected in a parts catalog 24. Moreover, to simplify the process of taking an order from a customer, an order typically includes packages of parts implementing certain features. To enable the knowledge base interpreter 13 to find the particular parts included in the packages in the order 11, the knowledge base 14 includes a set of expansion rules 25.

In addition to providing a framework for clearly intelligible and maintainable representation of knowledge about a product, the functional hierarchy 22 provides the structure or framework for the control procedure executed by the knowledge base interpreter 13 to match the parts in the order 11 to the generic configuration in the knowledge base 14 and to apply the constraints 16 to validate the order 11 and generate a production request 12. The working configuration memory 15 is organized into respective bins of parts 26 corresponding to the major functional components of the product defined by the functional hierarchy 22. Each bin of parts 26 corresponds to a structural or functional assembly in the product and the matching process is performed by inputting the initial configuration or list of parts 11 into an initial bin of parts 26 and sequentially matching the parts in parent bins to assembly descriptions for offspring bins and transferring the matching parts to the offspring bins. The task blocks 23 specify the particular assembly descriptions and the specific sequence of matching and transferring the parts. Before matching and transferring the parts from certain bins, however, the expansion rules 25 must sometimes be applied to "break open" certain packages of parts in the order 11 since different parts from the same packages sometimes match the assembly descriptions of different offspring bins. At the end of the transfer process, the bins of parts 26 contain respective parts from the initial configuration 11. Hence, the parts in the initial configuration 11 have been structured or configured according to the generic product configuration or functional hierarchy 22 and thus the constraints 16 may be applied to their respective bins of parts 26.

In order to simplify the process of applying the constraints, the constraints 16 may include bin variables defined for particular respective parent bins and which may be referenced in the parent bins and respective offspring bins of the parent bins. The knowledge base 14 includes separate declarations 28 of the bin variables. The respective task blocks 23 include steps for determining the values for the bin variables 27 and the values are stored in the working configuration memory 15. These steps in the task blocks 23 may include function calls to knowledge base functions 29 in the knowledge

base 14. The knowledge base functions 29 typically count parts and perform numerical computations to determine values for bin variables. In addition to conditioning constraints, bin variables may be used to condition the sequence of execution of the steps in the task blocks in order to perform conditional expansion, matching and transfer of parts and assemblies.

Turning now to FIGS. 2 and 3, the process of generating a functional hierarchy for a particular product is illustrated. The knowledge engineer obtains a description of the product from product engineers and technical manuals. The description should include an engineering diagram of the product in the form of a functional block diagram. A functional block diagram generally designated 30 for an exemplary M1234 computer system is shown in FIG. 2. The block diagram 30 shows functional subsystems arranged around an internal transfer subsystem 31 which includes, for example, the mainframe and primary input/output (I/O) bus in the M1234 computer. The internal transfer subsystem 31 exchanges data among the various other subsystems of the computer. These other subsystems include a multiprocessing communications subsystem 32, a processor subsystem 33 including the central processing unit of the computer, an instruction sequence unit 34 for controlling the sequence of instructions executed by the central processor, integrated memory subsystems 35 including the random access memory for the computer, an integrated communication subsystem 36 for providing a number of data links between the computer and external devices, an integrated disk subsystem 37 having a number of disk drives, I/O trunk subsystems 38 interfacing a number of computer peripherals, and a service subsystem 39 for interfacing printers, flexible disk drives, card readers, and computer terminals.

From the block diagram in FIG. 2 and other technical information about the M1234 computer, the knowledge engineer 20 generates a functional hierarchy description of the M1234 computer as illustrated in FIG. 3. In FIG. 3 the hierarchy is arranged in the form of a tree diagram generally designated 40, having an initial or root node 41, intermediate or branch nodes 42-46, and several terminal or leaf nodes 47-60. Each node in the functional hierarchy 40 represents a single functional component. The initial node 41 and the intermediate nodes 42-46 represent functional components that are composed of a number of more specific functional components. Specifically, the functional hierarchy 40 is defined by including declarations of parent functional components in the definitions of the intermediate 42-46 and terminal 47-60 functional components. The functional component MAINFRAME 42, for example, has the parent functional component SYSTEM 41. The functional hierarchy 22 of the knowledge base 14 (FIG. 1) corresponding to the tree diagram 40 in FIG. 3 is listed in Appendix II (A). According to the particular knowledge base syntax of Appendix II (A), it is said that the declared offspring functional component "COMPOSES" its parent functional component.

It should be noted that the consultation typescript of Appendix I was generated by the interpretation of a knowledge base in Appendices II (A)-(E). The task blocks are in Appendix II (B); the parts catalog is in Appendix II (C); the constraints, bin variable declarations, and knowledge base functions are in Appendix II (D); and the expansion rules are in Appendix II (E).

In general the knowledge base is in the form of a list of declarations of task blocks which are invoked and

executed when the system is run, declarations of knowledge base functions which are applied or evaluated during execution of the task blocks to set values of the bin variables, and declarations of knowledge base objects which become instantiated to generate one or more object instances in working memory during execution of the task blocks. Declarations of task blocks and knowledge base functions are preceded by the identifiers DEFTASKBLOCK and DEFKBFUN, respectively, followed by a specified name for the task block or knowledge base function. Declarations of objects are preceded by the identifier DEFCLASS, followed by a specified name for the object and, in parentheses, the class or type of the instantiated object. Object types include BIN (an instance of component or node in the functional hierarchy), PRODUCT.ID (an instance of a particular part), BVI (a bin variable instance), CTI (an instance of applying a constraint or product expansion), and QUEUE. Respective QUEUE instances keep track of the current step of execution in each task block so that the execution of one task block may be interrupted to execute another task block.

The declarations or definitions of object types in the knowledge base include values for instance variables or IVARs associated with instances of the object types. For each offspring BIN, for example, the IVAR "COMPOSES" is set to the name of the BIN's parent BIN. Each BIN and PRODUCT.ID has an IVAR "DESCR" including a description of the respective function or part in English so that the user interface and translation facility 19 may generate comprehensible explanations and traces of the operation of the system 10 (FIG. 1).

Each part in the parts catalog 24 (FIG. 1) implements one of the functions described in the functional hierarchy and each part is annotated to indicate the appropriate functional hierarchy element. In the parts catalog of Appendix IIC, each PRODUCT.ID or individual part has an IVAR "IMPLEMENTS.FN" including the name of the functional component which the part implements. These declarations of the functions implemented by the parts are critical to the operation of the system, since they are used to group and manipulate parts of different functional types during the matching operations.

As noted above, each component of the functional hierarchy has associated with it a particular task block 23 specifying when and how the system should focus on the ordered parts that implement the respective functional component. Also, associated with each functional component in the hierarchy is a set of constraints and bin variables 16 which are evaluated or applied during execution of the task block associated with the functional component.

It is apparent that the functional hierarchy serves both to organize the configuration process and to guide the knowledge engineer in designing and maintaining the knowledge base. By understanding the structure of the functional hierarchy and the effects of the configuration strategy used to process an order, the knowledge engineer can understand how the system will configure an order and how the application of constraints, product definitions, and expansion will affect the processing of that order.

The knowledge system 10 of FIG. 1 has its greatest utility for configuring complex systems such as the M1234 computer as described in FIGS. 2 and 3 and the typescript and knowledge base of Appendices I and II

(A)-(E). For the sake of illustration, however, the internal operation of the system will be described in connection with a more easily understood minicomputer system corresponding to the functional hierarchy generally designated 70 in FIG. 4 and corresponding to the consultation typescripts and knowledge base in Appendices III (A)-(B) and IV (A)-(D).

In the functional hierarchy 70 of FIG. 4, each node such as the initial node 71 is labeled by its functional component and has a class name enclosed in parentheses corresponding to the classes declared in the knowledge base functional hierarchy of Appendix III. The minicomputer 71 is presumed to have two major functional components. The main-frame 72 of the minicomputer encompasses all of the active components of the computer such as the central processing unit 73, a computer memory 74, and input/output devices 75. These active devices 73-75 are physically tied together by a card cage 76. The card cage 76 includes mechanical support for various circuit boards performing the active functions of the computer and also provides a network of electrical connections for the transfer of data among the circuit boards. The second functional component of the minicomputer 71 is a power supply 77 which feeds power to the circuit boards through power supply lines running through the card cage 76.

The minicomputer 71 also has a number of sub-assemblies performing active functions. Random access memory (RAM) 78 records data generated during operation of the minicomputer. Read-only memory (ROM) 79 contains fixed data used, for example, to initially run the minicomputer when the computer is first turned on, and also to run fixed utility programs for input and output and also to interpret, compile, or load a computer program written in a standard programming language.

The input/output function 75 has three separate functional assemblies. The minicomputer may include a number of terminals 80 for interacting with a human user. A terminal includes, for example, a cathode ray tube for displaying data to the operator, and a keyboard for receiving data from the operator. The terminals 80 and any other input/output devices are connected or interfaced to the computer via interface circuits 81 on circuit boards received by the card cage 76. The input/output function 75 may also include disk drives 82 for recording and storing relatively large amounts of data.

As a physical object, the minicomputer 71 is comprised of discrete physical components or parts. For a particular model of minicomputer, the minicomputer should be assembled from only a specified set of parts. These specified parts are listed in the parts catalog of Appendix IV (C). These parts include a CAGE-1000 for implementing the card cage 76, a CPU-1000 central processing unit circuit board for implementing the CPU function, a DISC-1000 disk drive, an IF-1000 interface circuit board, a RAM-1000 random access memory circuit board, a ROM-1000 read-only memory circuit board, a TERMINAL-1000 computer terminal, and three different power supplies, including a five ampere power supply SUPPLY-05, a ten ampere power supply SUPPLY-10, and a fifteen ampere power supply SUPPLY-15. The parts catalog in Appendix IV (C) also assigns a model number MINI-1000 to the minicomputer 71.

An order for a MINI-1000 minicomputer must at least include a set of parts, but the set may be empty. In general, an order for a product consists of three types of information. The first type is called "order information"

consists of such information as the customer name and an order number. It is primarily used to identify the order and does not effect the configuration process. The second type called "order lines" is a list of parts ordered and the quantity of each part desired. In general the order lines include the product model number and additional or optional features requested by the customer. The list of parts ordered and the quantity of each part desired is checked by the knowledge system 10 to ensure compatibility. The third type called "field service parts" is a list of additional part numbers and are merely appended or said to be shipped "on top" of the order. The field service parts are typically replacement or spare parts for the customer. The field service parts are not intended to be built into the system and are not checked either for compatibility among themselves or with the rest of the order.

For the consultation typescript in Appendix III (A), the order consists merely of order lines. These order lines specify that the product to be configured is one MINI-1000 minicomputer comprising four RAM-1000 random access memory circuit boards, three ROM-1000 read only memory circuit boards, four TERMINAL-1000 computer terminals, and three DISC-1000 disk drives.

Once the order is received by the knowledge system 10, the parts in the order are checked for consistency by applying the constraints in the knowledge base (Appendix IV (D)). In general, constraints specify engineering or marketing conditions which must be satisfied for the order lines. The constraints also specify the particular action that should be taken in response to whether the conditions are satisfied when the constraints are applied. For the MINI-1000 minicomputer, the order must include at least one RAM-1000 random access memory board, at least one ROM-1000 read only memory board, at least one TERMINAL-1000 computer terminal, and at least one IF-1000 interface board. The action to take when any of these parts is found to be missing is to add one of the respective missing parts, since these parts are essential to the operation of the minicomputer. The minicomputer also has a number of essential components that must be added if missing, but which can only appear with a quantity of one. The minicomputer must have one and only one power supply (SUPPLY-05, SUPPLY-10, or SUPPLY-15), one and only one CPU-1000 central processing unit card, and one and only one CAGE-1000 card cage.

The MINI-1000 minicomputer also has certain functional constraints conditioned on combinations of components. Due to the limited range of addresses addressed by the central processing unit 73, the memory 74 has a limited size. This memory constraint in turn limits the total number of memory circuit boards to no more than seven. If the number of RAM-1000 cards plus the number of ROM-1000 cards exceeds seven, the user must be told that he has ordered too many memory cards. This kind of constraint is called a warning constraint in contrast to the previous modification constraints which require the working configuration of parts to be changed.

The input/output assembly 75 has a few rather complex constraints. No more than four terminals are permitted for the MINI-1000 computer due to marketing objectives. The customer must, for example, buy a more expensive MINI-2000 computer in order to have a system with more than four terminals. Also, a single interface card can either support four terminals, or one ter-

terminal and one disk drive per card since the DISC-1000 disk drive requires three times as many data lines than the TERMINAL-1000 computer terminal. The order is modified by adding interface cards and deleting terminals so that these input/output conditions are satisfied.

The mainframe 72 has a hardware or interconnection constraint since the CAGE-1000 card cage has room for no more than sixteen cards or circuit boards. The user is warned if the total number of CPU-1000, RAM-1000, ROM-1000, and IF-1000 cards exceeds sixteen.

The minicomputer 71 has a global power constraint due to the fact that the customer may order a power supply that is smaller than that required under the worst case conditions. The power required is a function not only of the number of cards, but of the particular cards that are used. The CPU-1000 card requires two amperes of current, the RAM-1000 card requires one ampere of current, the ROM-1000 card requires 0.5 amperes of current, and the IF-1000 interface card requires 0.2 amperes of current. The knowledge base in Appendix IV (D) includes a special knowledge base function called "COMPUTE.POWER.NEEDED" for calculating the total amount of power required. Based on the required power, a five ampere power supply SUPPLY-05, a ten ampere power supply SUPPLY-10, or a fifteen ampere SUPPLY-15 should be included in the order. The knowledge base in Appendix IV (D) includes a set of modification constraints to ensure that the smallest power supply is included consistent with the total power requirement.

Shown in FIG. 5 is a flowchart generally designated 85 describing the basic steps in the order checking process. In the first step 86, the order or initial configuration 11 is received by the knowledge system 10 (FIG. 1). Then in step 87 the individual components are determined from the order so that the constraints may be applied in step 88. Once the constraints are applied, in step 89 the indication of the application of the constraints is transmitted to the user. Application of warning constraints may result in warnings to the user, and application of the modification constraints may lead to a production request 12 that differs from the original order 11.

For the MINI-1000 minicomputer, the computer's functional hierarchy, parts catalog, and constraints are relatively simple so that the order checking process in FIG. 5 could easily be implemented using a conventional programming language such as the BASIC language. Since an order for a MINI-1000 computer is merely a list of the indivisible components of the computer, the constraints are easily applied. More importantly, the functional class which each part implements is readily apparent from the type of part. For the M1234 computer described in the knowledge base of Appendices II (A)-(E), however, it would be very difficult to use a conventional programming language to implement an order checking computer system and the computer system would be very difficult to maintain.

By using the techniques already described, the components and functionality of a complex product such as a computer can be transparently and explicitly represented in the knowledge base. In accordance with another aspect of the present invention, the definition of configuration constraints is clearly separated from the configuration strategies and actions. In particular, a built-in control procedure is provided which permits the constraints to be defined for particular functional components and separately determines which parts are

applicable to the respective constraints. At the most basic level, this separation requires a process which builds an explicit representation of the assemblies of the product from the parts in the order. Shown in FIG. 6, for example, is an explicit representation generally designated 90 for the MINI-1000 minicomputer and the example order given above. The explicit representation 90 is a tree of bins corresponding to the functional hierarchy of FIG. 4. In FIG. 6 the name of each bin corresponds to the respective functional component or bin class in the hierarchy of FIG. 4. Moreover, the particular instance of the class is identified by a numeral enclosed in parentheses. Although the examples in the appendices only use a single instance of each functional class, the knowledge engineer may write task blocks which create a specified number of bins for a specified functional class.

The bins are used to hold a set of product instances or parts. The product instances are obtained from the order lines, by application of modification constraints or by EXPANSION operations upon parts initially placed in the bins. The MINI-1000 computer, however, does not have any expandable parts in its parts catalog. Typically for marketing reasons, an order may include packages of parts that are distinguished from orders of the individual parts included in the package. For the MINI-1000 computer, for example, it could be economically desirable to define a basic package including the minimum components. Thus, the basic package would include one CAGE-1000, one SUPPLY-05, one CPU-1000, one RAM-1000, one ROM-1000, one TERMINAL-1000, and one IF-1000. The basic package would be sold at a price less than the total price of the components of the package if the components were ordered separately. Of course, the intention of the manufacturer is not to sell the package at a lower cost, but to sell the package at market cost and to increase the price of any additional components purchased on top of the basic package.

To apply the constraints, the packages must be expanded into their component parts since the constraints typically apply to individual parts and not packages. The constraints are preferably independent of the product expansions so that the packages may be redefined merely by changing the expansion rules (25 FIG. 1).

Due to the complexity introduced by packages and modification constraints, it is important that the components of the order are properly associated with corresponding constraints. Preferably, this correspondence is obtained by matching the components in the order to predefined assembly descriptions and then applying the constraints to the components which match the respective assembly descriptions. Also, the preferred method of performing the matching is by transferring or sorting the entire order through the bin tree, starting at the initial bin until the elemental parts end up in the terminal bins. For the MINI-1000 order processed in Appendix III (A), the result of the sorting is illustrated in FIG. 6.

The order in which parts are matched to assembly descriptions, sorted and checked by the constraints is all controlled by the steps in the task blocks. The execution of the steps in the task blocks is started by an executive program 100 shown in FIG. 7. In the first step 101, certain program variables are cleared. By clearing these program variables, space in the working memory is cleared so that a new bin tree may be built. Also, the trace memory is cleared in order to record a new trace.

In step 102 an initial bin is created and is given the name "UNASSIGNED". The name "UNASSIGNED" is given denoting that the particular product being configured is not yet known since the order has not yet been received by the knowledge system.

It should be noted that the knowledge base interpreter 13 is always working on a particular bin at any given time. The name of this particular bin is stored in a global variable called "CURRENT.BIN". In step 103 this global variable "CURRENT.BIN" is set to the value "UNASSIGNED" to record that the configuration process is starting in the "UNASSIGNED" bin. In step 104 "UNASSIGNED" task block is found, and it is executed in step 105 so that the rest of the configuration process is controlled solely by task blocks.

The task blocks are executed by calling an EXECUTE subroutine generally designated 110 in FIG. 8. In general terms, the EXECUTE subroutine sequentially executes all of the commands in a specified task block. In the first step 111, a pointer (P1) is set to the position of the first command in the task block. In step 112 the command pointed to by the pointer (P1) is executed. In step 113 the pointer (P1) is advanced to that it either points to another command or is at the end of the task block. In step 114, the end of the task block is detected by testing whether the pointer is out of the range of the specified task block. If it is not, execution jumps to step 12 to execute the next command. Otherwise, the EXECUTE subroutine is finished and execution returns to the calling program.

Preferably, the task blocks are written in a well-known imperative programming language such as LISP described in the text P. Winston & V. Horn, *LISP*, Addison-Wesley Publishing Co., Reading, Mass. (1981). The "UNASSIGNED" task block is as follows:

```
(DEFTASKBLOCK UNASSIGNED.TB UNASSIGNED ( )
(INPUT.ORDER)
(DETERMINE (TODAYS.DATE ORDERED.PARTS
ORDERED.PRODUCT.LINE))
(IF CHECKING.ORDER THEN
(CONFIGURE (FETCH 'ORDERED.PRODUCT.LINE)
(FIND 'ALL 'CCS))
(CHECK 'EMPTY.BIN))
(DETERMINE (FINAL.PARTS))
(OUTPUT.ORDER)
)
```

Translated into English, the UNASSIGNED task block above first inputs the order of parts. Then it determines the current date, the ordered parts or order lines of the order, and the product line corresponding to the ordered parts. It is presumed that the knowledge system is currently checking the order. Thus, the task block for the ordered product line is fetched from the knowledge base. This task block is, for example, the MINICOMPUTER task block in Appendix IVB. Then a FIND operation is performed to obtain a list of all the parts in the order lines, and the CONFIGURE operation is performed to create the initial bin in the bin tree, to transfer all the parts in the order lines into the initial bin, and to temporarily transfer execution to the MINI-COMPUTER task block. When execution returns, the parts are configured in the bins as shown in the bin tree 90 of FIG. 6. A final set of EMPTY.BIN constraints are applied to ensure that all of the order lines in the order have been properly configured. If so, the FINAL-PARTS are determined and a production request is generated by the OUTPUT.ORDER statement.

The UNASSIGNED task block as well as the task blocks in the knowledge base include a mixture of LISP code written by the knowledge engineer, built-in functions supplied by the LISP compiler, and special functions or subroutines comprising the built-in control procedure of the knowledge base interpreter 13 (FIG. 1). The primary built-in subroutines include the EXECUTE subroutine already described in conjunction with FIG. 8, a FIND subroutine shown in FIG. 9, a CONFIGURE subroutine shown in FIG. 10, a CHECK subroutine shown in FIG. 11, an EXPAND subroutine shown in FIG. 12, and a DETERMINE subroutine shown in FIG. 13.

The basic matching operation is performed by the FIND subroutine of the FIG. 9. The FIND subroutine generally designated 120 obtains a list of parts included in a specified bin which match or satisfy a specified assembly description. An assembly description is a predicate or Boolean function defining the desired set of parts in terms of a conjunction of part names, part types, or other functions which describe the desired parts. A collection of parts satisfies an assembly description when the parts fit that description. In other words, the assembly description serves as a template to select the desired parts.

In the first step 120' of the FIND subroutine, an assembly list (AL) is cleared in order that it may receive the names of the desired parts once the desired parts are found. In step 121 a pointer (P2) is set to the first part in the specified bin. In particular, the FIND subroutine 120 in FIG. 9 is used to find parts in the current bin to transfer them to offspring bins of the current bin. Hence, in step 121 the specified bin is the value of the global variable "CURRENT.BIN". In step 122, the specified assembly description is evaluated for the part in the current bin pointed to by the pointer (P2). The assembly description, in other words, is a Boolean function of the part pointed to by the pointer (P2). In step 123, the Boolean value of the assembly description indicates whether the part satisfies the assembly description. If the part satisfies the assembly description, then the part is one of the desired parts so that in step 124 the part is added to the assembly list (AL). If the part does not satisfy the assembly description, it is not added to the assembly list (AL). Then in step 125 the pointer (P2) is advanced so that it may point to another part in the current bin. The pointer (P2) might not, however, point to another part in the current bin since all of the parts in the current bin might have been tested. In step 126 the pointer (P2) is tested to determine whether it is out of the range of the current bin. If it is not out of range, then the pointer (P2) does point to another part in the current bin and execution jumps back to step 122. Otherwise, all of the parts in the current bin have been checked and the assembly list (AL) will contain the names of all of the parts that satisfy the specified assembly description. Hence, execution returns to the calling task block.

In order to perform the sorting operation described above, a CONFIGURE operation is performed to create a new bin of a specified functional class, to transfer the matched parts from the current bin to the new bin, and to temporarily pass execution to the task block associated with the functional class of the new bin. A flowchart for the CONFIGURE subroutine generally designated 130 is shown in FIG. 10. In step 131 a new bin of the specified functional class is created by allocating space in the working configuration memory 15 (FIG. 1) for the new bin. Space is allocated, for exam-

ple, for instances of all the bin variables declared for the bin and the bin variable instances are marked undetermined to indicate that they do not yet have values. Next, in step 132, it is asserted that the parent of the new bin is the current bin specified by the global variable "CURRENT.BIN". In step 133, the specified parts are removed from the current bin and transferred to the new bin. In step 134 the value of the global variable CURRENT.BIN is pushed onto a bin stack and the global variable is assigned the name of the new bin. This is done in anticipation of interrupting the processing for the current bin in order to process the new bin. Specifically, in step 135, the task block associated with the given functional class of the new bin is found. In step 138 the steps in this task block are executed. Upon completion of executing the steps in this task block, in step 137 the bin stack is popped and the popped value is reassigned to the global variable CURRENT.BIN to continue execution of the task block for the current bin. Execution then returns to the task block which called the CONFIGURE subroutine 130.

The fact that a stack was used in the CONFIGURE subroutine 130 suggests that the CONFIGURE subroutine is advantageously used in a recursive fashion. In general, each task block associated with an intermediate node in the functional class hierarchy will have CONFIGURE operations in its associated task block for creating offspring bins preceded by a FIND operation to specify the parts to transfer to the offspring bins. This ensures that the parts initially placed into the root bin or added by product expansions become sorted through the bin tree to the terminal or leaf bins. Once the parts have filtered down to the terminal leaf bins, the most basic and elementary constraints may be applied to determine, for example, if specified parts or a specified number of parts are included in the terminal bins. To apply these constraints, a CHECK operation is provided to apply a specified subset of the constraints applicable to the functional class of the current bin.

Shown in FIG. 11 is a flowchart of the check subroutine 140. In the first step 141 the knowledge base is searched to obtain a list (CL) of the constraints applicable to the functional class of the current bin and which satisfy the given constraint description. Then, in step 142 a pointer (P3) is set to the first constraint in the constraint list (CL). Next, in step 143 the condition portion of the constraint in the list pointed to by the pointer (P3) is evaluated. The condition of the constraint is typically a Boolean function of the conditions of the various parts in the bin. If the condition of the constraint is satisfied, as determined by testing whether the Boolean value of the conditions is true, then in step 145, the action specified by the constraint is executed. Otherwise, the action specified by the constraint is not executed. In step 146 the pointer (P3) is advanced in order to point to the next constraint in the constraint list (CL). The pointer (P3) must be checked, however, in step 147 to determine whether all of the constraints in the list have been applied. If the pointer (P3) is still within the range of the constraint list (CL), then execution jumps to step 143 to apply the next constraint. Otherwise, the CHECK subroutine 140 is finished and execution returns to the current task block.

It should be noted that during certain steps in the sorting or configuring process, certain kits or groups of parts must be expanded into their component parts. An EXPAND operation is provided for expanding the parts in the current bin into their component parts by

applying predefined expansion rules associated with the current bin. An EXPAND subroutine generally designated 150 is shown in FIG. 12. In the first step 151, a list (EXL) of all of the expansion rules that apply to the parts in the current bin are obtained from the knowledge base. Next, in step 152 a pointer (P4) is set to the first expansion rule (EX) in the expansion rule list (EXL). Then in step 153 the particular expansion rule (EX) pointed to by the pointer (P4) is obtained. The expansion rule (EX) is checked in step 154 to determine whether it has already been marked expanded. It should be noted that a particular expansion rule may have already been applied since a particular expansion rule may apply to more than one bin. If the expansion rule has not already been applied, then in step 154' the expansion is applied by adding to the current bin the parts included in the composing assembly description of the expansion rule. The composing assembly description is a list of the parts included in the respective kit for the respective expansion rule. Once the expansion rule has been applied, it is marked expanded in step 155. Then, in step 156 the pointer (P4) is advanced in order to point to the next expansion rule in the list. However, in step 157 the value of the pointer (P4) must be checked to determine whether it does point to another expansion rule in the list, or whether all of the expansion rules in the list have been applied. If the point (P4) points to another expansion rule in the list, then execution jumps to step 153. Otherwise, execution returns to the current task block.

During the execution of task blocks and the application of constraints, it is in many cases desirable to condition the execution of a particular task block step upon the conditions or attributes of the parts included in the bin for the task block. As noted above, bin variables are provided for this purpose and the bin variables may also be included in the condition portions of constraints. A DETERMINE operation is provided for specifying when to determine values for the bin variable instances associated with the current bin by applying the knowledge base functions associated with the current bin. Shown in FIG. 13 is a flowchart of a DETERMINE subroutine generally designated 160. In the first step 161 a list (BVL) of bin variables for the current bin are obtained from the knowledge base. Next, in step 162 a pointer (P5) is set to the first bin variable in the bin variable list (BVL). Then in step 163, the particular bin variable (BV) pointed to by the pointer (P5) is obtained. Before determining a value for the bin variable instance (BV), in step 164 the bin variable instance is inspected to determine whether it already has a value. If not, then in step 165 a list (KBFL) of knowledge base functions concluding a value for the bin variable (BV) is obtained from the knowledge base. Then in step 166 another pointer (P6) is set to the first knowledge base function in the function list (KBFL). In step 167 the particular knowledge base function (KBF) pointed to by the pointer (P6) is obtained. This knowledge base function (KBF) is evaluated in step 168 to obtain a respective value (V). This respective value (V) might not, however, be a legal value for the particular bin variable (BV). In step 169 the declared legal values for the particular bin variable (BV) are obtained from the knowledge base and compared to the particular value (V) to determine whether the value is a legal value for the bin variable (BV). If not, then in step 170 the knowledge base function pointer (P6) is advanced. In step 171 the pointer is compared to the end of the range for the

function list (KBFL) to determine whether the pointer (P6) is pointing to another knowledge base function. If so, then execution may jump back to step 167 in an attempt to apply this other knowledge base function. Otherwise, there are no more knowledge base functions to try so that in step 172, a value of "NO VALUE" is assigned to the bin variable instance (BV).

If in step 169 the value (V) was a legal value for the particular bin variable (BV), then in step 173 the particular value (V) is assigned to the bin variable instance (BV).

Now that a value for the particular bin variable instance (BV) has been determined, the bin variable pointer may be advanced in step 174 in an attempt to point to the next bin variable in the bin variable list (BVL). In step 175 the pointer (P5) is first compared to the end of the bin variables list to determine whether it is pointing to another bin variable. If so, execution jumps back to step 163 to determine the value for the next bin variable instance. Otherwise, all of the bin variables in the bin variable list (BVL) have values and execution may return to the current task block.

The primary operations of EXECUTE, FIND, CONFIGURE, CHECK, EXPAND, and DETERMINE are used in the task blocks to define a specific configuration strategy suited for the particular product. In general, however, it can be said that the task blocks for the interior bins have similar characteristics and also the task blocks for the leaf or terminal bins have similar characteristics. A generalized task block for an interior bin is shown below:

```

(DEFTASKBLOCK INTERIOR.TB INTERIOR ( )
(CONFIGURE 'OFFSPRING1 (FIND 'ALL 'CCS
(IMPLEMENTS OFFSPRING1)))
(CONFIGURE 'OFFSPRING2 (FIND 'ALL 'CCS
('IMPLEMENTS OFFSPRING2)))
.
.
(CHECK 'REQUIREMENTS)
(EXPAND* '(KITS))
)

```

Translated into English, execution of the interior bin task block first finds all of the components in the interior bin which implements the first offspring function. Then a configure operation creates a bin corresponding to the first offspring function and the parts that are found are transferred to this offspring bin. Execution is then temporarily transferred to the offspring bin. Upon returning from the task block of the first offspring, the components implementing a second offspring function are found and another configuration operation is performed to create a second offspring bin. The parts that are found are transferred to the second offspring bin and execution is passed to the task block for the second offspring function. Upon returning, the parts for any other offspring functions are found and offspring bins are configured. Once all of the offspring bins have been configured, the requirements or constraints for the interior bin are applied. Finally, any kits that may have been added by the application of modification constraints are expanded. The EXPAND* operation is similar to the EXPAND operation except that it specifies a type of part to be expanded (such as any kit) so that the particular kits to be expanded need not be specified.

A generalized task block for leaf bins is shown below:

```
(DEFTASKBLOCK LEAF.TB LEAF()
(CHECK 'CONTENTS)
(EXPAND* '(KITS))
(CHECK 'REQUIREMENTS)
)
```

First, modification constraints are applied in order to ensure that the leaf bin will not be empty. Then, specified kits in the leaf bin are expanded into their component elemental parts. Finally, warning and modification constraints are applied to check that the requirements for the leaf task block have been satisfied.

Once the knowledge engineer encodes a specific configuration strategy into the knowledge base, a test configuration can be run using a "TRACE" option so that the user interface and explanation facility (19 FIG. 1) generates a comprehensible translation of the configuration strategy. For the MINI-1000 minicomputer, for example, the trace in the typescript of Appendix III (A) explains in detail the specific configuration strategy that is followed. This explanation includes the specific task block statements that are executed and any action taken upon application of the constraints. The configuration strategy for the MINI-1000 minicomputer order in Appendix II (A) first looks at the set of ordered parts to determine that the product line is a MINI-1000 minicomputer. These parts are dumped into the root bin MINI(1) in FIG. 6. Then the parts implementing the MAIN-FRAME function are found and the MAIN-FRAME(1) bin is created and these parts are transferred to the MAIN-FRAME(1) bin. In recursive fashion, the parts in the MAIN-FRAME(1) bin are found which implement the-CPU function. A CPU(1) bin is created and these parts are transferred to the CPU(1) bin. The requirements for the CPU(1) bin are checked and it is found that a required CPU board is missing. The knowledge system 10 asks the user for permission to add one CPU-1000 circuit board, thereby completing the CPU(1) bin. Execution returns to the MAIN-FRAME task block. The components implementing the MEMORY function are found and a new MEMORY(1) bin is created to receive these parts. In recursive fashion, the parts in the MEMORY(1) bin implementing the RAM function are found and a new RAM(1) is created to receive these parts. The contents of the RAM(1) bin are checked to ensure that at least one RAM-1000 circuit board is included. Execution then returns to the MEMORY task block which finds the parts implementing the ROM function and creates a new ROM(1) bin to receive these parts. The contents of the ROM(1) bin are checked to ensure that at least one ROM-1000 circuit board is included. Then, execution returns to the MEMORY task block which checks the overall requirements for the MEMORY function. In particular, a MEMORY.CT1 constraint is applied which counts the total number of RAM and ROM boards to ensure that the total number does not exceed seven. When this constraint is applied, it is found that the total number does not exceed seven. Execution then returns to the MAIN-FRAME task block. Next, the MAIN-FRAME task block finds all of the parts in the MAIN-FRAME(1) bin which implement the I/O function, creates a new I/O(1) bin and transfers these parts to the new bin. Execution passes to the I/O task block which finds all of the components in the I/O(1) bin implementing the

TERMINAL task block, the requirements or constraints for the TERMINAL function are applied. It is found that there is at least one terminal in the TERMINAL(1) bin. It is also found that the number of terminals does not exceed four. Execution then returns to the I/O task block.

The I/O task block next configures the DISC function before it configures the INTERFACE function since the interface constraints need to know how many disk drives have been ordered. The DISC function is configured by first finding the components in the I/O(1) bin which implement the DISC function, and creating a new DISC(1) bin to receive these components. Execution is passed to the DISC task block which first checks the requirements for the DISC function. No constraints are found associated with the DISC function. Execution returns to the I/O task block. Since the INTERFACE task block needs to know how many DISC units are included in the order, a bin variable called "IF-CARDS.NEEDED" is defined for the I/O function. A DETERMINE operation is performed to determine the minimum number of interface cards required for the number of terminals and disk drives in the order. After it is determined that there are four terminals and three disk drives so that four interface cards are required, the components in the I/O(1) bin implementing the INTERFACE function are found and a new INTERFACE(1) bin is created to receive these components. Execution then passes to the INTERFACE task block.

The requirements for the INTERFACE function are checked and it is found that no interface cards were ordered. A modification constraint IF.REQ.CT1 is applied which tells the user that the minimum number of interface cards have not been ordered, and ask permission to add four IF-1000 interface cards. These cards are added after permission is received from the user. Execution then returns to the I/O task block which in turn passes execution back to the MAIN-FRAME task block. In the MAIN-FRAME task block, all of the components in the MAIN-FRAME(1) bin are found which implement the CRAD-CAGE function. A new CARD-CAGE(1) is created to receive these components. Execution is passed to the CARD-CAGE task block which checks the CARD-CAGE requirements. The constraint AT-MOST-ONE.CT ensures that no more than one card cage was ordered. The constraint NO-PARTS.CT checks to ensure that at least one card cage was ordered and finds that a card cage was not ordered, and asks permission from the user to add a card cage. After receiving permission from the user, one CAGE-1000 is added to the order. Execution then returns to the MAIN-FRAME task block.

The MAIN-FRAME task block checks the requirements for the MAIN-FRAME function. The MAIN-FRAME.CT1 constraint is applied which ensures that the total number of ROM cards, RAM cards, and IF cards plus one CPU card is no greater than sixteen, so that no action need be taken. Execution then returns to the MINI task block.

In the MINI task block, all of the components in the MINI(1) bin are found that implement the POWER-SUPPLY function, and a new POWER-SUPPLY(1) is created to receive these components. Execution is transferred to the POWER-SUPPLY task block. The initial contents of the POWER-SUPPLY(1) bin are first checked to determine the number of power supplies included in the order. It is found that no power supplies

were ordered. Next, the requirements for the POWER-SUPPLY function are checked and a set of modification constraints are applied in order to determine the particular power supply to be added to the order. Thus, several power supply modification constraints are applied which use a bin variable called "POWER.NEEDED".

It should be noted that the task blocks for the minicomputer do not explicitly have a DETERMINE operation to determine the value of the "POWER.NEEDED" bin variable. A DETERMINE operation is implicitly performed when a constraint is applied that references a bin variable. Upon determining the value of the bin variable "POWER.NEEDED" it is found that the required power to operate the minicomputer is 8.3 amperes. The application of the power supply modification constraints determines that a 10 ampere power supply is needed and one SUPPLY-10 is automatically added to the order. Execution then returns to the MINI task block which in turn passes execution back to the unassigned task block which is listed above in the specification. The final contents of the UNASSIGNED bin are checked to ensure that all of the original components have been considered during the configuration of the minicomputer. It is possible, for example, that some superfluous parts were included in the order, and at this point in the UNASSIGNED task block the superfluous components would be left in the UNASSIGNED bin. It is found, however, that the UNASSIGNED bin is empty, so that no superfluous components were ordered. Next, the list of final parts are determined for the production request. Finally, the final parts are transmitted to the user.

Inspection of the constraints, bin variable declarations and knowledge base functions for the M1234 computer and the MINI-1000 minicomputer knowledge bases in Appendices II (D) and IV (D) reveal that there are many built-in functions available to the knowledge engineer which simplify the knowledge engineer's task of creating the task blocks, constraints, and knowledge base functions. Moreover, the knowledge base is organized so that the knowledge engineer has ready access to the conditions of the parts upon which the constraints and knowledge base functions are responsive and can manipulate the working configuration in any desired fashion.

An operation (ASSIGN parts bin) moves the parts specified by a list of PRODUCT.IDs from the current bin into the specified bin.

If the operation (BVSET binvar newvalue noerror) assigns (newvalue) as the value of (binvar) after checking the (newvalue) correspondence to the format indicated by the LEGAL.VALUE field of the (binvar). The value (newvalue) is typically computed by a KBFUNCTION defined by the knowledge engineer. If (noerror) is true, the checking that the value (binvar) matches the format prescribed in its LEGAL.VALUES field is not performed.

The function (FETCH binvar noerror) returns the value of (binvar) if (binvar) already has a value. If (binvar) does not have a value, FETCH determines that value using the KBFUNCTIONS specified on its HOW.TO.DETERMINE field. If (noerror) is false, the FETCH function checks to make sure that the value of the (binvar) matches the format prescribed in its LEGAL.VALUE field.

The function (MEMBER element list) tests whether (element) is a member of (list). The function MEMBER

returns true (T) if the element is a member of the list, and false (NIL) if not.

The function (FIND number bins description) searches the CONTENTS of (bins) for a certain (number) of parts (PRODUCT.IDs) that satisfy the (description). The FIND function returns a list of PRODUCT.IDs. The (number) of parts defined may be an integer, the word for an integer between ZERO and TEN, or the word ALL; any words must be quoted. The set of (bins) to search are specified by the following set of key words:

CCS	The current bin;
SCS	The immediate parent bin;
CSCS	The current and immediate parent bins;
ACS	The current and all ancestor bins; or
ICS	The current and all descendant bins.

The (description) is a Boolean expression made up of one or more logical clauses and connectors AND and OR. The clause (IMPLEMENT <FUNCTION.CLASS>) succeeds if the part has an IMPLEMENTS.FN that COMPOSES the <FUNCTION.CLASS> specified. The clause (ISA <PART>) succeeds if the part matches this <PART> name. The clause (ISA.SPEC <FUNCTION.CLASS>) succeeds if the part is a PROD.GEN of the <FUNCTION.CLASS> specified.

Many of the built-in functions recognized by the knowledge base interpreter 13 of the knowledge system 10 (FIG. 1) have particular utility in conjunction with constraints. Constraints always apply to a specified set of parts. The parts are specified by a list of BIN-TYPES, which are the FUNCTION.CLASS entries in the functional hierarchy that the constraint applies to. The constraint also includes a STATEMENT of conditions to test against the parts in those bins. The STATEMENTS often test relationships between values of different bin variables, which must be declared in the BIN.VARS.REQUIRED field of the constraint declaration. Also, a list of CHECK.KEYWORDS may be specified to control when certain constraints are applied. These keywords are used by the CHECK operations in the task block steps to select relevant subsets of constraints to apply at different times in a task block.

The knowledge system provides a number of important specializations of the CONSTRAINT knowledge base type, which are further described below. All product type and specific constraints are defined in terms of these specializations. The system also considers product expansions as a particular kind of CONSTRAINT since they add parts to a bin if certain conditions, such as whether a particular product was ordered, are met. In general, the types of constraints include constraints on specific parts, warning constraints, and modification constraints.

There are five types of constraints on specific parts. The type EXPANSION.CT expands a composite part into its component parts if certain conditions are met. The type ALWAYS.EXPAND.CT unconditionally expands a composite part into its component parts. The type FULLY.EXPAND.CT expands a composite part into its component parts, and then expands all of those parts that are expandable. The type EFFECTIVELY.CT allows the knowledge engineer to specify a date after which the specified part is available for the specified product line. The type PHASE-OUT.CT al-

lows the knowledge engineer to specify a date after which the specified part is no longer available for the specified product line.

There are two types of warning constraints. The type `WARN-ON-SUCCESS.CT` warns the user if the conditions in the `STATEMENT` for the constraint succeed. The type `WARN-ON-FAILURE.CT` warns the user if the conditions in the `STATEMENT` fail.

There are two types of modification constraints. The type `MODIFY-ON-SUCCESS.CT` outputs a warning message and modifies the contents of the current bin if the conditions in the `STATEMENT` succeed. The type `MODIFY-ON-FAILURE.CT` outputs a warning message and modifies the contents of the current bin if the conditions in the `STATEMENT` fail.

A constraint `STATEMENT` expresses a set of conditions that the order must meet to be valid. The `STATEMENT` of a constraint is a Boolean expression composed of "clauses", separated by the usual logical connectives `AND` and `OR`. Individual clauses (or the entire `STATEMENT` expression) may be negated by using the form `(NOT <EXPRESSION>)`.

In general the `STATEMENT` clauses test various relationships between expressions of bin variables. The expressions may include the typical numerical relationships `<`, `>`, `=`, `>=`, `<=`, and `<>` (not equal to).

Clauses can also evaluate arbitrary predicates written by the knowledge engineer as `KBFUNCTIONS`. The system provides several predefined functions such as `COUNT.PARTS`, `ORDERED?`, and `SYSTEM` that enhance the transparency of the constraint statements. Moreover, the system can evaluate clauses expressed in a limited, English-like format. These expressions test whether a specified quantity `IS` (or `IS NOT`) `ZERO`, `NULL`, `ATOMIC`, `NUMERIC`, `REAL`, or a `STRING`.

Constraint `STATEMENTS` typically reference the values of bin variables. The bin variables may be referenced directly by the `FETCH` function in a `STATEMENT`. Alternatively, the system can directly interpret references to bin variables when the bin variables are specifically declared in the `BIN.VARS.REQUIRED` field of the constraint declaration.

The function (`COUNT.PARTS` part search) returns the number of (parts) in the current bin. If (search) is true (T), the function `COUNT.PARTS` will search the bin tree below and including the current bin returning the total number of (parts) in these bins.

The function (`ORDERED?` parts search) returns the value true (T) if there is at least one instance of any of the (parts) in the current bin. If the value of (search) is true (T), `ORDERED?` will search the bin tree below (and including) the current bin, returning true (T) if it finds instances of (parts) in these bins. The function (`SYSTEM` system-name1 system-name2 . . .) tests whether the system type or product model of the order is one of the (system-names) indicated in the argument list. Each (system-name) should be a `FUNCTION-CLASS` having a `PRODUCT.GEN` declaration including `SYSTEM`.

The function (`COUNT` list) returns the number of elements in the (list).

For constraints that modify the contents of the current bin, the system provides a special form for specifying a list of parts to add to, delete from, or replace the contents of a bin. This special form is called a `PRODUCT.QTY` or "product quantity". `PRODUCT.QTYs` are specified in the `TO` field of the different `EXPANSION.CTs`, and they are computed by the `ADD-`

`PARTS`, `DEL-PARTS`, and `SET-PARTS` fields of the `MODIFY-ON-SUCCESS.CTs` and `MODIFY-ON-FAILURE.CTs`. Special functions described further below can be used to compute these `PRODUCT.QTY` description for these constraints.

A `PRODUCT.QTY` consists of a list of lists where each sublist specifies a quantity and a part type.

The function (`FIND*` number bins description) searches a set of (bins) for a certain (number) of parts that satisfy the (description). `FIND*` returns a `PRODUCT.QTY` that describes the quantities of the parts found.

The function (`CHOOSE` one parts) is given a specified `PRODUCT.QTY` description of several different parts, and asks the user to select one part from (parts) and returns a `PRODUCT.QTY` describing a single `<SELECTED-PART>`.

As noted above, an expansion rule is a kind of constraint that applies to specific parts. Expansion rules are specified by `EXPANSION.CT` objects in the knowledge base. The `FROM` field of an `EXPANSION.CT` specifies the composite part to be expanded; the `TO` field specifies the list of component parts to add to the bin; the `WHEN` field specifies the conditions under which this expansion should occur; and the `MESSAGE` field specifies the messages that should be printed when the expansion succeeds or fails.

Modification constraints are defined in the knowledge base either as `MODIFY-ON-SUCCESS.CT` objects or a `MODIFY-ON-FAILURE.CT` objects. A warning message is given to the user and the contents of the current bin are modified if the conditions in the `STATEMENT` field succeed or fail, respectively. The knowledge engineer may specify a `MESSAGE` that will warn the user that a modification must be made. The modifications are specified by providing values for one or more of the fields `SET-PARTS`, `ADD-PARTS`, and `DEL-PARTS`. If a list of parts is specified for the field `SET-PARTS`, the current bin is emptied and the list of parts (in the quantities specified) is placed in the bin. Then, if a list of parts is specified in the `ADD-PARTS` field, new parts are added to the bin in the quantities specified. If a list of parts is provided in the `DEL-PARTS` field, existing parts are deleted from the bin in the quantities specified.

In the knowledge base, the `BIN.VAR` declarations must include a `LEGAL.VALUE` field indicating the legal value for the bin variable (e.g., integer, string, etc.), a `HOW.TO.DETERMINE` field including a list of `KBFUNCTIONS` that can be used to determine the value of the bin variable, and a `BIN.TYPES` field including the bin types or function classes with which the variable is associated. It should be noted that values of `BIN.VAR` instances can be referenced by constraints applied by task blocks that are lower in the bin tree than where the `BIN.VAR` instance is created. Thus, the `BIN.TYPES` should associate the bin variable with bins high enough in the bin tree so that the bin variables can be referenced by all of the appropriate constraints.

The system provides two built-in `KBFUNCTIONS` that can be used to ask the user a question to determine the value of specified `BIN.VARS`. The (`ASK.QUESTION`) function uses the `DESCR` field of the `BIN.VAR` to form the question "What is `<DESCR>?`" The `ASK.QUESTION` function uses the optional help field of the `BIN.VAR` declaration to provide help to the user when the user types a question mark (?) in response to the question. Also, the response to the question is

checked against the LEGAL.VALUES expression of the BIN.VAR declaration. The second function (ASK.BOOLEAN.QUESTION) uses the DESCR field of the BIN.VAR declaration to form the question "Is it the case that <DESCR>?" The ASK.BOOLEAN.QUESTION function accepts only responses of YES or NO and prints, "Answer YES or NO please" if any other response is given.

It should be noted that the knowledge base itself may be divided into sections so that the sections may be identified and referenced by the task blocks and constraints. The name for each portion of the knowledge base is preferably composed of a knowledge base prefix and a unique suffix identifying the particular portion or file of the knowledge base. In the preferred embodiment of Appendix II (A)-(E) the suffix FH identifies the file containing the functional hierarchy and the product line task blocks. The suffix NCT identifies the file containing the product line constraints, the bin variables, and the knowledge base functions. The suffix BX contains product line descriptions. The suffix AK identifies the file containing product line packages. The suffix FT identifies the file containing product line features. The suffix KT identifies the product line kits. The suffix KV identifies the file containing the product line kit versions. The suffix XT identifies a file containing miscellaneous parts. The suffix EX identifies the file containing product expansions. A knowledge base file with the suffix MODS is used to record all modifications or changes made to the knowledge base since the last modification by the knowledge engineer using the knowledge base maintenance facility 21 (FIG. 1).

The knowledge base maintenance facility 21 preferably has a number of commands used by the knowledge engineer to create, maintain, and augment the knowledge base 14. These commands include, for example, eight commands for knowledge base data and control functions, three commands for knowledge base editing functions, and four commands for knowledge base file maintenance.

The eight knowledge base data and control commands include SELECT, PP, CHANGES, CONTENTS, CONTROL, IGNORE.CHANGES, REMEMBER.CHANGES, and LIST.KBSETS.

The command SELECT selects a knowledge base to be modified and loads the knowledge base into the host computer system so that it may be further manipulated by the knowledge base maintenance facility 21.

The command PP displays the status of the current knowledge base, including whether changes are being recorded, whether the changes will be automatically saved, how many modifications have been made since the last time the knowledge base was saved, (i.e., written back from the working memory of the knowledge base maintenance facility 21 to the more permanent memory from which the knowledge base was loaded), and how many objects are currently in the current knowledge base.

The CHANGES command displays the modifications made to the knowledge base since the most recent migration, (see MIGRATE below) and indicates which objects have been added, changed, or deleted from the knowledge base.

The command CONTENTS displays the names of all of the objects in the current knowledge base.

The command CONTROL controls whether the system automatically saves modifications after a certain

number of additions, edits, and deletions have been made to the current knowledge base.

The command IGNORE.CHANGES tells the knowledge base maintenance facility 21 to ignore any subsequent additions, modifications, or deletions.

The command REMEMBER.CHANGES turns off the effect of a previous IGNORE.CHANGES command so that subsequent changes will be recorded.

A command LIST.KBSETS displays the contents of specified main knowledge base sets after asking the user to specify the names of the sets.

The three knowledge base editing commands include ADD, EDIT, and DELETE. The command ADD adds new objects to the current knowledge base and prompts the knowledge engineer for the knowledge base type and the name of the new knowledge base object. The command EDIT modifies existing objects in the current knowledge base and prompts for the name of the knowledge base object. The command DELETE removes objects from the current knowledge base and prompts for the name of an existing knowledge base object.

The four knowledge base file maintenance commands include SAVE, MIGRATE, BUILD, and CREATE. The command SAVE saves the copy of the current knowledge base changes in a modification file. The command MIGRATE distributes the current knowledge base modification into the main knowledge base files and writes out a new, empty copy of the modifications file. The command BUILD saves a copy of the current executable system on a file and prompts the knowledge engineer for the name of the file. The command CREATE creates a new knowledge base for a new product line, and automatically prompts for a knowledge base prefix to identify the new product line.

The preferred embodiment of the user interface and explanation facility 19 has several commands to enable the user to obtain an explanation of the operation of the system 10. The user can select options for either obtaining an explanation as a configuration is being run, or after the completion of a configuration. As a configuration is being run, the user can select one or more of six options. The command SHOW-STEPS instructs the explanation facility 19 to print the task block steps as they are executed. The command SHOW-DETERMINATION causes the explanation facility to print the values of bin variables as they are determined. The command SHOW-EXPANSIONS tells the explanation facility to print a description of the expansions applied to parts in the order. The command SHOW-PARTS causes the explanation facility to indicate when parts are added to or deleted from a bin's contents. The command SHOW-FULL-TRACE or ALL causes the explanation facility to indicate all of the events described above.

After an order has been configured and the complete trace has been recorded in the trace memory 17, the user may enter commands to access the trace memory 17. The command SHOW-TRACE causes the explanation facility to replay the trace of the most recent order. The command SHOW-VALUE causes the explanation facility to display the values of the bin variables used in the most recent order. The command SHOW-PART-HISTORY displays the parts history of the most recent order. The command FORGET causes the explanation facility to erase the trace memory 17 so that there is no longer a history of the prior consultation.

The knowledge system 10 has been described above primarily for checking and modifying the order lines in an order for a flexibly assembled product. It should be noted, however, that the modification constraints in part redesign the ordered product. The system 10 is equally useful as a design tool. The design procedure is illustrated by the typescript in Appendix III (B) wherein an order for a MINI-1000 minicomputer merely includes the model number of the desired product. The set of modification constraints in the knowledge base for the MINI-1000 product line (Appendices IV (A)-(D)) designs a minimal system consistent with the modification constraints including one CPU-1000 card, one RAM-1000 card, one ROM-1000 card, one TERMINAL-1000 computer terminal, one IF-1000 card, one CAGE-1000 card cage, and one SUPPLY-05 five ampere power supply.

To use the knowledge system 10 as shown in FIG. 1 as an expert design tool, the initial configuration 11 should include not only an initial set of parts but an initial set of values for bin variables to specify the desired parameters of the final configuration. The initial values for bin variables, for example, may specify the desired conditions and relations between the parts being assembled, and may also specify what modification constraints are to be applied as the product is being built in the working configuration memory 15.

It should also be noted that the knowledge engineer has great flexibility in defining the control procedure for the knowledge base interpreter 13. The knowledge engineer, for example, may write specialized knowledge base functions and task blocks to perform specialized configuration operations. The knowledge engineer may also modify the built-in control procedure of the knowledge base interpreter 13 to modify the control procedure in any desired fashion.

In view of the above, a knowledge system has been described for generalized representation and processing of hierarchical assemblies. The configuration strategies and actions are defined in task blocks which are completely separate from the functional hierarchy which describes the generic configuration of the product and the parts catalog which defines the individual elements available for configuration. Moreover, the constraints are completely separate from the task blocks, functional hierarchy, and parts catalog so that the definition of configuration constraints are clearly separated from configuration checking strategies and actions. The knowledge base interpreter has a built-in control procedure and built-in functions which enable the knowledge engineer to design and implement readily a desired

configuration checking strategy. Since the imperative language of the task blocks clearly defines the configuration checking strategies and actions, a trace of a configuration operation is easily stored in a trace memory including both the actual steps executed and the resulting actions such as warnings or modifications to the configuration. Thus, an explanation facility can be provided which generates an intelligible and comprehensible explanation of the configuration based on the record in the trace memory.

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APPENDIX I

CONSULTATION TYPESCRIPT FOR THE PROCESSING OF AN ORDER FOR AN M1234 COMPUTER

The following is an example M1234 consultation with OCEAN.

Welcome to OCEAN

OCEAN executive> RUN

Processing new order...

Enter name of order file: T

Reading order information...

Order information> ORDER NUMBER: 6330-82-0261

Order information> BRANCH: 1815

Order information> METHOD: SURFACE

Order information>

Reading parts to check...

Order line> M1234

=> 1 M1234

The M1234 product line

Order line> 1234-3003-0690

=> 1 1234-3003-0690

Class model for M1234 system

Order line> 1234-F010

=> 1 1234-F010

Firmware

Order line> 1234-F290

=> 1 1234-F290

Memory 2M, 2 way

Order line> 1234-F361

=> 1 1234-F361

Common trunk

Order line> 1234-F691

=> 1 1234-F691

BSLA

Order line> 1234-F680

=> 1 1234-F680

MFG.PLANT BYTE MULTIPLEX TRUNK

Order line>

Reading parts to be shipped F/S...

F/S order line>

Enter order processing options: ?

Please select from among CHECK, DON'T-CHECK, EDIT-ORDER, QUIETLY, SHOW-COMPILATION, SHOW-CONSTRAINTS, SHOW-DETERMINATION, SHOW-EXPANSIONS, SHOW-FULL-TRACE, SHOW-PARTS, SHOW-STEPS, or other

Enter order processing options: SHOW-FULL-TRACE

Adding the following parts: 1 M1234, 1 1234-3003-0690, 1 1234-F010, 1 1234-F290, 1 1234-F361, 1 1234-F691 and 1 1234-F680

Evaluating: (DETERMINE '(TODAYS.DATE ORDERED.PARTS ORDERED.PRODUCT.LINE)) Today's date of original components is set to: 15-Mar-84 The ordered set of parts of original components is set to: M1234, 1234-3003-0690, 1234-F010, 1234-F290, 1234-F361, 1234-F691 and 1234-F680 The ordered product line of original components is set to: M1234 product line components

Evaluating: (IF CHECKING.ORDER

THEN **COMMENT**

(CONFIGURE (FETCH

'ORDERED.PRODUCT.LINE)

(FIND 'ALL 'CCS))

(CHECK 'EMPTY.BIN))

Configuring M1234 product line components

Initial parts are: 1 M1234, 1 1234-3003-0690, 1 1234-F010, 1 1234-F290, 1 1234-F361, 1 1234-F691 and 1 1234-F680

Evaluating: (BVSET 'SYSTEM.TYPE

(COMPUTE.M1234.SYSTEM.TYPE))

The M1234 system type of original components is set to: M1234 system components

Evaluating: (CONFIGURE (FETCH 'SYSTEM.TYPE)
(FIND 'ALL 'CCS))

Configuring M1234 system components

Initial parts are: 1 M1234, 1 1234-3003-0690, 1
1234-F010, 1 1234-F290, 1 1234-F361, 1 1234-F691
and 1 1234-F680

Evaluating: (EXPAND* '(BG BU.M12))

Checking constraint: 1234-3003-0690.ECT.1

(a ALWAYS.EXPAND.CT)

Expanding 1234-3003-0690 to 1 1234-F262, 1 1234-
F113, 1 1234-K915, 1 1234-K450, 1 1234-K917, 1
1234-K416 and 1 1234-K941 in M1234 system
components.

Adding the following parts: 1 1234-F262, 1 1234-
F113, 1 1234-K915, 1 1234-K450, 1 1234-K917, 1
1234-K416 and 1 1234-K941

Evaluating: (DETERMINE '(CLASS.MODEL.EXPANSION
BASIC.TO.PRICED

MP.ORDER))

The class model parts of M1234 system components is
set to: 1 1234-F262, 1 1234-F113, 1 1234-K915, 1
1234-K450, 1 1234-K917, 1 1234-K416 and 1 1234-K941

The basic-to-priced parts of M1234 system
components is set to: 1 1234-K416, 1 1234-K920, 2
1234-K924, 1 1234-K940 and 1 1234-K941

This is a multi-processing order of M1234 system
components is set to: False

Evaluating: [CONFIGURE 'MULTIPROCESSING
(FIND 'ALL 'CCS '(IMPLEMENTS
MULTIPROCESSING]

Configuring multi-processing components

No initial parts

Evaluating: (CHECK 'CONTENTS)

Checking CONTENTS constraints in multi-
processing components

No CONTENTS constraints were found.

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: (DETERMINE '(MP.PIBS))

PIBS required for MP kits of M1234 system
components is set to: 0

Evaluating: [CONFIGURE 'MAINFRAME (FIND 'ALL 'CCS
'(IMPLEMENTS
MAINFRAME]

Configuring mainframe components

Initial parts are: 1 1234-F010, 1 1234-F290, 1
1234-F361, 1 1234-F691, 1 1234-F680, 1 1234-F262, 1
1234-F113, 1 1234-K915, 1 1234-K450, 1 1234-K917
and 1 1234-K941

Evaluating: [CONFIGURE 'PROCESSOR (FIND 'ALL 'CCS
'(IMPLEMENTS
PROCESSOR)]

Configuring processor components
Initial parts are: 1 1234-K941

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: [CONFIGURE 'MEMORY (FIND 'ALL 'CCS
'(IMPLEMENTS
MEMORY)]

Configuring main memory components
Initial parts are: 1 1234-F290 and 1 1234-F262

Evaluating: (CHECK 'CONTENTS)
Checking CONTENTS constraints in main memory
components

The following constraints apply: AT-MOST-
ONE-MEMORY-FT.CT

Checking constraint: AT-MOST-ONE-MEMORY-FT.CT
(a MODIFY-ON-SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
AND
[COUNT (FIND 'ALL 'CCS
'(ISA FT]
> 1)

succeeded with value T when applied to main
memory components

Removing the following parts: 1 1234-F262

Evaluating: (EXPAND* '(AK FT KT))
Checking constraint: 1234-F290.ECT.1 (a
ALWAYS.EXPAND.CT)

Expanding 1234-F290 to 4 1234-K920, 2 1234-
K921, 8 1234-K924 and 2 1234-K940 in main
memory components.

Adding the following parts: 4 1234-K920, 2
1234-K921, 8 1234-K924 and 2 1234-K940

Evaluating: (DETERMINE '(MAIN-MEMORY.SUPPLIED
MEMORY.INTERLEAVING
MEMORY.PIBS))

The main memory supplied with this order (in
kb) of M1234 system components is set to: 2048

The memory interleaving for this order of
M1234 system components is set to: 2-WAY
PIBS required for main memory w/o extra
plenums of M1234 system components is set to:
16

Evaluating: (CHECK 'INTERLEAVING)
Checking INTERLEAVING constraints in main
memory components

The following constraints apply: 2-WAY-
INTERLEAVING

Checking constraint: 2-WAY-INTERLEAVING (a
INTERLEAVING.CT)

The statement: ((SYSTEM M1234.SYSTEM)
AND
(NOT (MEMORY.INTERLEAVING =
'2-WAY])

failed with value NIL when applied to main
memory components

Evaluating: (CHECK 'MEMORY-CAPACITY)

Checking MEMORY-CAPACITY constraints in main
memory components

No MEMORY-CAPACITY constraints were found.

Evaluating: (CHECK 'MEMORY-INCREMENT)

Checking MEMORY-INCREMENT constraints in main
memory components

The following constraints apply: 1024-
INCREMENT

Checking constraint: 1024-INCREMENT (a
MEMORY-INCREMENT.CT)

The statement: ((SYSTEM M1234.SYSTEM)
AND
(NOT (INCREMENT MAIN-
MEMORY.SUPPLIED 1024)))

failed with value NIL when applied to
main memory components

Evaluating: [CONFIGURE 'ISU (FIND 'ALL 'CCS
'(IMPLEMENTS ISU)]

Configuring ISU memory components

Initial parts are: 1 1234-F113, 1 1234-K915
and 1 1234-K917

Evaluating: (CHECK 'CONTENTS)

Checking CONTENTS constraints in ISU memory
components

The following constraints apply: AT-LEAST-
ONE-FT.CT AT-MOST-ONE-FT.CT

Checking constraint: AT-LEAST-ONE-FT.CT (a
WARN-ON-SUCCESS.CT)

The statement: ([COUNT (FIND 'ALL 'CCS '(ISA
FT]
IS ZERO)

failed with value NIL when applied to ISU
memory components

Checking constraint: AT-MOST-ONE-FT.CT (a
MODIFY-ON-SUCCESS.CT)

The statement: ([COUNT (FIND 'ALL 'CCS
'(ISA FT]
> 1)

failed with value NIL when applied to ISU
memory components

Evaluating: (EXPAND* '(AK FT KT))
 Checking constraint: 1234-F113.ECT.1 (a
 ALWAYS.EXPAND.CT)

Expanding 1234-F113 to 2 1234-K940 and 3
 1234-K913 in ISU memory components.
 Adding the following parts: 2 1234-K940 and 3
 1234-K913

Evaluating: (DETERMINE '(ISU-
 MEMORY.SUPPLIED))
 The isu memory supplied with this order (in
 kb) of M1234
 system components is set to: 24

Evaluating: (CHECK 'REQUIREMENTS)
 Checking REQUIREMENTS constraints in ISU
 memory components
 The following constraints apply: MP.ISU.CT
 Checking constraint: MP.ISU.CT (a WARN-ON-
 SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
 AND MP.ORDER AND (ISU-
 MEMORY.SUPPLIED < 40))
 failed with value NIL when applied to ISU
 memory components

Evaluating: [CONFIGURE 'FIRMWARE (FIND 'ALL 'CCS
 '(IMPLEMENTS
 FIRMWARE]

Configuring firmware components
 Initial parts are: 1 1234-F010

Evaluating: (CHECK 'REQUIREMENTS)
 Checking REQUIREMENTS constraints in firmware
 components
 The following constraints apply:
 I9050.FIRMWARE.CT1
 MP.FIRMWARE.CT FIRMWARE.CT
 Checking constraint: I9050.FIRMWARE.CT1 (a
 MODIFY-ON-SUCCESS.CT)

The statement: [(SYSTEM I9050-1.SYSTEM
 I9050-2.SYSTEM I9050-3.SYSTEM)
 AND
 (NOT (ORDERED? '(1234-F018
 1234-K018])
 failed with value NIL when applied to
 firmware components
 Checking constraint: MP.FIRMWARE.CT (a
 MODIFY-ON-SUCCESS.CT)

The statement: [(SYSTEM M1234:SYSTEM)
AND MP.ORDER AND (NOT
(ORDERED? '5640-F400)]
failed with value NIL when applied to
firmware components
Checking constraint: FIRMWARE.CT (a MODIFY-
ON-SUCCESS.CT)

The statement: [(SYSTEM M1234.SYSTEM)
AND
(NOT MP.ORDER)
AND
(NOT (ORDERED? '(1234-F010
1234-K010 1234-F011
1234-K011)]
failed with value ~~NIL~~ when applied to
firmware components

Evaluating: (CHECK 'CONTENTS)
Checking CONTENTS constraints in firmware
components
The following constraints apply: AT-LEAST-
ONE-FT.CT
AT-MOST-ONE-FT.CT
Checking constraint: AT-LEAST-ONE-FT.CT (a
WARN-ON-SUCCESS.CT)

The statement: ([COUNT (FIND 'ALL 'CCS
'(ISA FT]
IS ZERO)
failed with value NIL when applied to
firmware components
Checking constraint: AT-MOST-ONE-FT.CT (a
MODIFY-ON-SUCCESS.CT)

The statement: ([COUNT (FIND 'ALL 'CCS
'(ISA FT]
> 1)
failed with value NIL when applied to
firmware components

Evaluating: (EXPAND* '(AK FT KT))
Expanding 1234-F010 to 1234-K010 in firmware
components.
Adding the following parts: 1 1234-K010

Evaluating: [CONFIGURE 'IO (FIND 'ALL 'CCS
'(IMPLEMENTS IO)]

Configuring IO components
Initial parts are: 1 1234-F361, 1 1234-F691
and 1 1234-F680

Evaluating: [CONFIGURE 'IDC (FIND 'ALL 'CCS
'(IMPLEMENTS IDC)]

Configuring IDC components
No initial parts

Evaluating: (CHECK 'CONTENTS)
 Checking CONTENTS constraints in IDC
 components

The following constraints apply: NO-
 FTS.CT

Checking constraint: NO-FTS.CT (a MODIFY-
 ON-SUCCESS.CT)

The statement: ([COUNT (FIND 'ALL 'CCS
 '(ISA FT]

IS NOT ZERO)

failed with value NIL when applied to
 IDC components

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: [CONFIGURE 'CLC/MLA
 (FIND 'ALL 'CCS
 '(IMPLEMENTS
 CLC/MLA]

Configuring CLC/MLA components
 No initial parts

Evaluating: [CONFIGURE 'FICS
 (FIND 'ALL 'CCS
 '(IMPLEMENTS
 FICS]

Configuring fast ICS components
 No initial parts

Evaluating: (CHECK 'CONTENTS)
 Checking CONTENTS constraints in fast
 ICS components

No CONTENTS constraints were found.

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: [CONFIGURE 'ICS (FIND 'ALL
 'CCS
 '(IMPLEMENTS
 ICS]

Configuring ICS components
 No initial parts

Evaluating: (CHECK 'CONTENTS)
 Checking CONTENTS constraints in ICS
 components

No CONTENTS constraints were found.

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: (DETERMINE '(#CLC.MLA.LINES
 ICS.PIBS))

The number of clc's of M1234 system
 components is set to: 0

The number of mla's of M1234 system
 components is set to: 0
 The number of clc/mla lines for this order
 of M1234 system
 components is set to: 0
 PIBS required for CLC/MLA of M1234 system
 components is set to: 0

Evaluating: (CHECK 'IO-CAPACITY)
 Checking IO-CAPACITY constraints in
 CLC/MLA components

The following constraints apply:

#CLC.MLA-LINES

Checking constraint: #CLC.MLA-LINES (a
 WARN-ON-SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
 AND
 (#CLC.MLA.LINES > 20))
 failed with value NIL when applied to
 CLC/MLA components

Evaluating: [CONFIGURE '2STAGE-IO
 (FIND 'ALL 'CCS
 '(IMPLEMENTS
 2STAGE-IO]

Configuring two-stage IO components
 Initial parts are: 1 1234-F361, 1 1234-F691
 and 1 1234-F680

Evaluating: [CONFIGURE 'ITB-MODULES
 (FIND 'ALL 'CCS
 '(IMPLEMENTS
 ITB-MODULES]

Configuring ITB components
 Initial parts are: 1 1234-F361

Evaluating: (CHECK 'CONTENTS)
 Checking CONTENTS constraints in ITB
 components
 No CONTENTS constraints were found.

Evaluating: (EXPAND* '(AK FT KT))
 Checking constraint: 1234-F361.ECT1 (a
 ALWAYS.EXPAND.CT)

Expanding 1234-F361 to 2 1234-K931 in
 ITB components.

Adding the following parts: 2 1234-K931

Evaluating: (DETERMINE '(#TRUNK.LINES
 IO.PIBS))

The number of standard io lines for
 this order of M1234
 system components is set to: 2
 PIBS required for IOLC and common trunk
 of M1234 system
 components is set to: 4

Evaluating: [CONFIGURE 'IO-PANEL-MODULES
(FIND 'ALL 'CCS
(IMPLEMENTS
IO-PANEL-MODULES)]

Configuring IO panel components
Initial parts are: 1 1234-F691 and 1
1234-F680

Evaluating: (CHECK 'CONTENTS)
Checking CONTENTS constraints in IO
panel components
No CONTENTS constraints were found.

Evaluating: (EXPAND* '(AK FT KT))
Checking constraint: 1234-F691.ECT1 (a
ALWAYS.EXPAND.CT)
Expanding 1234-F691 to 2 1234-K935, 2
1234-K936 and 2 1234-K937 in IO
panel components.
Adding the following parts: 2 1234-
K935, 2 1234-K936 and 2 1234-K937
Checking constraint: 1234-F680.ECT.1 (a
ALWAYS.EXPAND.CT)
Expanding 1234-F680 to 1 1234-K930 in
IO panel components.
Adding the following parts: 1 1234-K930

Evaluating: (DETERMINE '(#CT.LINES
#IOLC.LINES))

The number of common trunk interface
lines of M1234
system components is set to: 1
The number of iolc interface lines of
M1234 system
components is set to: 2

Evaluating: (CHECK 'IO-CAPACITY)
Checking IO-CAPACITY constraints in two-
stage IO components
The following constraints apply: CT.IO
CT.IOLC.IO CT-LINES IOLC-LINES
Checking constraint: CT.IO (a WARN-ON-
FAILURE.CT)

The statement: (#CT.LINES <=
TRUNK.LINES)
succeeded with value T when applied to
two-stage IO components
Checking constraint: CT.IOLC.IO (a WARN-
ON-FAILURE.CT)

The statement: (#TRUNK.LINES =
#IOLC.LINES + #CT.LINES)

failed with value NIL when applied to
two-stage IO components

*** WARNING - The number of IOLC and common trunk I/F lines
does not match the number of HS, MS, and LS IO lines.

Checking constraint: CT-LINES (a IO-
LINES.CT)

The statement: ((SYSTEM M1234.SYSTEM)
AND
(NOT (BETWEEN #CT.LINES
0 6)))

failed with value NIL when applied to
two-stage IO components

Checking constraint: IOLC-LINES (a IO-
LINES.CT)

The statement: ((SYSTEM STSYSTEM)
AND
(NOT (BETWEEN
#IOLC.LINES 0 6)))

failed with value NIL when applied to
two-stage IO components

Evaluating: [CONFIGURE 'PLENUM (FIND 'ALL 'CCS
'(IMPLEMENTS
PLENUM]

Configuring plenum components

No initial parts

Evaluating: (CHECK 'CONTENTS)
Checking CONTENTS constraints in plenum
components

The following constraints apply: AT-MOST-
ONE-FT.CT

Checking constraint: AT-MOST-ONE-FT.CT (a
MODIFY-ON-SUCCESS.CT)

The statement: ([COUNT (FIND 'ALL 'CCS
'(ISA FT]
> 1)

failed with value NIL when applied to
plenum components

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: [CONFIGURE 'POWER-SUPPLY
(FIND 'ALL 'CCS
'(IMPLEMENTS
POWER-SUPPLY]

Configuring power supply components

No initial parts

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in power
supply components

The following constraints apply: MP.POWER-SUPPLY.CT

Checking constraint: MP.POWER-SUPPLY.CT (a MODIFY-ON-SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
AND MP.ORDER AND
(NOT (ORDERED? '(1234-F400
1234-K400]

failed with value NIL when applied to power supply components

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: [CONFIGURE 'REST (FIND 'ALL 'CCS
'(IMPLEMENTS REST]

Configuring remaining components

Initial parts are: 1 1234-K416

Evaluating: [CONFIGURE 'CABLE (FIND 'ALL 'CCS
'(IMPLEMENTS CABLE]

Configuring cable components

No initial parts

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: [CONFIGURE 'CONSOLE (FIND 'ALL 'CCS
'(IMPLEMENTS
CONSOLE]

Configuring console components

Initial parts are: 1 1234-K416

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: [CONFIGURE 'MISC (FIND 'ALL 'CCS
'(IMPLEMENTS MISC]

Configuring miscellaneous components

No initial parts

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: (CHECK 'REQUIREMENTS)

Checking REQUIREMENTS constraints in miscellaneous components

The following constraints apply: CLC/MLA-DISPLAY

Checking constraint: CLC/MLA-DISPLAY (a MODIFY-ON-SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
AND
((COUNT.PARTS '1234-K545 T)
= 1)
AND
(BETWEEN #CLC.MLA.LINES 11
20))

failed with value NIL when applied to miscellaneous components

Evaluating: (EXPAND* '(AK FT KT))

Evaluating: (CHECK 'POWER)

Checking POWER constraints in M1234 system components

The following constraints apply: EXTRA.MINUS5
EXTRA.PLUS5 MINUS5 PLUS5

Checking constraint: EXTRA.MINUS5 (a WARN-ON-SUCCESS.CT)

The power used by this order of M1234 system components is set to: +5V: 33 -2V: 0 -5V: 226

The statement: ((SYSTEM M1234.SYSTEM)
AND
(NOT MP.ORDER)
AND
(ORDERED? '1234-F400 T)
AND
(POWER.USED:MINUS.5 <= 250))

failed with value NIL when applied to M1234 system components

Checking constraint: EXTRA.PLUS5 (a WARN-ON-SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
AND
(ORDERED? '1234-F413 T)
AND
(POWER.USED:PLUS.5 <= 60))

failed with value NIL when applied to M1234 system components

Checking constraint: MINUS5 (a MODIFY-ON-SUCCESS.CT)

The power supplied for this order of M1234 system components is set to: +5V: 0 -2V: 0 -5V: 0

The statement: ((SYSTEM M1234.SYSTEM)
AND
(POWER.USED:MINUS.5 -
POWER.SUPPLIED:MINUS.5)
> 250)

failed with value NIL when applied to M1234 system components

Checking constraint: PLUS5 (a MODIFY-ON-SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
AND
(POWER.USED:PLUS.5 -
POWER.SUPPLIED:PLUS.5)
> 60)

failed with value NIL when applied to M1234 system components

Evaluating: (CHECK 'PIBS)

Checking PIBS constraints in M1234 system components

The following constraints apply: ICS.PIBS.CT
IO.PIBS.CT PIBS.CT1 PIBS.CT2

Checking constraint: ICS.PIBS.CT (a WARN-ON-SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
 AND
 (ICS.PIBS > 8))
 failed with value NIL when applied to M1234
 system components
 Checking constraint: IO.PIBS.CT (a WARN-ON-
 SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
 AND
 (IO.PIBS + ICS.PIBS > 12))
 failed with value NIL when applied to M1234
 system components
 Checking constraint: PIBS.CT1 (a WARN-ON-
 SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
 AND
 (ORDERED? '1234-K926)
 AND
 (TOTAL.PIBS > 38))
 failed with value NIL when applied to M1234
 system components
 Checking constraint: PIBS.CT2 (a MODIFY-ON-
 SUCCESS.CT)

The statement: ((SYSTEM M1234.SYSTEM)
 AND
 (NOT (ORDERED? '1234-K926))
 AND
 (TOTAL.PIBS > 38))
 failed with value NIL when applied to M1234
 system components
 Evaluating: (EXPAND* '(AK FT KT))

Evaluating: (CHECK 'EFFECTIVITY)
 Checking EFFECTIVITY constraints in M1234 system
 components
 No EFFECTIVITY constraints were found.

Evaluating: (CHECK 'PHASE-OUT)
 Checking PHASE-OUT constraints in M1234 system
 components
 No PHASE-OUT constraints were found.
 Checking EMPTY.BIN constraints in M1234 system
 components

The following constraints apply:
 UNASSIGNED.EMPTY.CT
 Checking constraint: UNASSIGNED.EMPTY.CT (a WARN-
 ON-SUCCESS.CT)

The statement: (BIN:CONTENTS IS NOT NULL)
 failed with value NIL when applied to original
 components

Evaluating: (DETERMINE '(FINAL.PARTS))

The final set of parts of original components is set to: M1234, 1234-3003-0690, 1234-K450, 1234-K941, 1234-F290, 1234-K920, 1234-K921, 1234-K924, 1234-K940, 1234-F113, 1234-K915, 1234-K917, 1234-K913, 1234-F010, 1234-K010, 1234-F361, 1234-K931, 1234-F691, 1234-F680, 1234-K935, 1234-K936, 1234-K937, 1234-K930 and 1234-K416

Evaluating: (IF (MEMBER (FETCH 'ORDERED.PRODUCT.LINE)
'(PRODUCT.LINE
I9050.PRODUCT.LINE))
THEN (OUTPUT.ORDER))

What type of output form would you like:

PRODUCTION.WORKSHEET

Order number: 6330-82-0261

Branch: 1815

Method: SURFACE

Remarks:

*** WARNING - The number of IOLC and common trunk I/F lines does not match the number of HS, MS, and LS IO lines.

1	1234-3003-0690	Class model for M1234 system	F/S
Basic To Priced			
1	1234-K416	CONSOLE INDICATOR INTERFACE	F/A
1	1234-K920	MSU EXPANSION MODULE	F/A
2	1234-K924	MSU 256K ARRAY	F/A
1	1234-K940	ISU/MSU RESISTOR	F/A
1	1234-K941	CRYSTAL 84/56 NS	F/A
Customer BTO			
1	1234-K450	VS HARDWARE	F/A
3	1234-K920	MSU EXPANSION MODULE	F/A
2	1234-K921	MSU, INTERFACE	F/A
6	1234-K924	MSU 256K ARRAY	F/A
3	1234-K940	ISU/MSU RESISTOR	F/A
1	1234-K915	READ-ONLY MEMORY AND TERMINATOR	F/A
1	1234-K917	SERVICE PROCESSOR, ISU NO.1	F/A
3	1234-K913	ISU, 8K FAST	F/A
1	1234-K010	Firmware	F/A
2	1234-K931	DIRECT MEMORY ACCESS TRUNK	F/A
2	1234-K935	BIT SERIAL LINK CABLES	F/A
2	1234-K936	CONNECTOR MODULE	F/A
2	1234-K937	CT BSL ADAPTER	F/A
1	1234-K930	COMMON TRUNK DATA INTERFACE	F/A

What type of output form would you like:

OCEAN executive>

APPENDIX II (A)

KNOWLEDGE BASE FOR M1234 COMPUTER
FUNCTIONAL HIERARCHY

```

(* XTEST-FHS definition)
(DEFINSTANCE XTEST-FHS FHS
  SET.OF XTESTKB
  FILED.ON XTESTFH
  ELEMENTS          (2STAGE-IO CABLE CLC/MLA CONSOLE FICS
                     FIRMWARE ICS IDC IO IO-PANEL-MODULES ISU
                     ITB-MODULES MAINFRAME MEMORY MISC MULTI-
                     PROCESSING PLENUM POWER-SUPPLY PROCESSOR
                     PRODUCT.LINE REST SP M1234.SYSTEM))

(DEFCLASS 2STAGE-IO (BIN)
  ()
  (METAClass FUNCTION.CLASS
   DESCR "two-stage IO"
   COMPOSES (IO))
  )

(DEFCLASS CABLE (BIN)
  ()
  (METAClass FUNCTION.CLASS
   DESCR "cable"
   COMPOSES (REST))
  )

(DEFCLASS CLC/MLA (BIN)
  ()
  (METAClass FUNCTION.CLASS
   DESCR "CLC/MLA"
   COMPOSES (IO))
  )

(DEFCLASS CONSOLE (BIN)
  ()
  (METAClass FUNCTION.CLASS
   DESCR "console"
   COMPOSES (REST))
  )

(DEFCLASS FICS (BIN)
  ()
  (METAClass FUNCTION.CLASS
   DESCR "fast ICS"
   COMPOSES (CLC/MLA))
  )

(DEFCLASS FIRMWARE (BIN)
  ()
  (METAClass FUNCTION.CLASS
   DESCR "firmware"
   COMPOSES (MAINFRAME))
  )

```

```
(DEFCLASS ICS (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   DESCR "ICS"
   COMPOSES (CLC/MLA))
)

(DEFCLASS IDC (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   DESCR "IDC"
   COMPOSES (IO))
)

(DEFCLASS IO (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   DESCR "IO"
   COMPOSES (MAINFRAME))
)

(DEFCLASS IO-PANEL-MODULES (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   DESCR "IO panel"
   COMPOSES (2STAGE-IO))
)

(DEFCLASS ISU (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   DESCR "ISU memory"
   COMPOSES (MAINFRAME))
)

(DEFCLASS ITB-MODULES (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   DESCR "ITB"
   COMPOSES (2STAGE-IO))
)

(DEFCLASS MAINFRAME (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   DESCR "mainframe"
   COMPOSES (SYSTEM))
)

(DEFCLASS MEMORY (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   DESCR "main memory"
   COMPOSES (MAINFRAME))
)

(DEFCLASS MISC (BIN)
```

```
(
(METACLASS FUNCTION.CLASS
  DESCR "miscellaneous"
  COMPOSES (REST))
)

(DEFCLASS MULTIPROCESSING (BIN)
  (
  (METACLASS FUNCTION.CLASS
    DESCR "multi-processing"
    COMPOSES (SYSTEM))
  )
)

(DEFCLASS PLENUM (BIN)
  (
  (METACLASS FUNCTION.CLASS
    DESCR "plenum"
    COMPOSES (MAINFRAME))
  )
)

(DEFCLASS POWER-SUPPLY (BIN)
  (
  (METACLASS FUNCTION.CLASS
    DESCR "power supply"
    COMPOSES (MAINFRAME))
  )
)

(DEFCLASS PROCESSOR (BIN)
  (
  (METACLASS FUNCTION.CLASS
    DESCR "processor"
    COMPOSES (MAINFRAME))
  )
)

(DEFCLASS PRODUCT.LINE (BIN)
  (
  (METACLASS FUNCTION.CLASS
    DESCR "M1234 product line"
    PROD.GEN (PRODUCT.LINE))
  )
)

(DEFCLASS REST (BIN)
  (
  (METACLASS FUNCTION.CLASS
    DESCR "remaining"
    COMPOSES (SYSTEM))
  )
)

(DEFCLASS SP (BIN)
  (
  (METACLASS FUNCTION.CLASS
    DESCR "SP"
    COMPOSES (MAINFRAME))
  )
)
```

```
(DEFCCLASS M1234.SYSTEM (BIN)
  ()
  (METACCLASS FUNCTION.CLASS
    PROD.GEN (SYSTEM))
  )
```

```
(* XTEST-TBS definition)
(DEFINSTANCE XTEST-TBS TBS
  SET.OF XTESTKB
  FILED.ON XTESTFH
  ELEMENTS
    (2STAGE-IO.TB CABLE.TB CLC/MLA.TB
     CONSOLE.TB FICS.TB FIRMWARE.TB ICS.TB
     IDC.TB IO-PANEL-MODULES.TB IO.TB ISU.TB
     ITB-MODULES.TB MAINFRAME.TB MEMORY.TB
     MISC.TB MULTIPROCESSING.TB PLENUM.TB
     POWER-SUPPLY.TB PROCESSOR.TB
     PRODUCT.LINE.TB REST.TB SP.TB SYSTEM.TB))
```

APPENDIX II (B)

KNOWLEDGE BASE FOR M1234 COMPUTER
TASK BLOCKS

```
(DEFTASKBLOCK 2STAGE-IO.TB 2STAGE-IO ()
  (* jsb: " 4-Aug-83 23:17")
  (CONFIGURE 'ITB-MODULES (FIND 'ALL 'CCS '(IMPLEMENTS ITB-
MODULES)))
  (CONFIGURE 'IO-PANEL-MODULES (FIND 'ALL 'CCS '(IMPLEMENTS
IO-PANEL-MODULES)))
  (CHECK 'IO-CAPACITY))

(DEFTASKBLOCK CABLE.TB CABLE () (* jsb: " 4-Aug-83 23:13")
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK CLC/MLA.TB CLC/MLA ()
  (* jsb: " 7-AUG-83 20:32")
  (CONFIGURE 'FICS (FIND 'ALL 'CCS '(IMPLEMENTS FICS)))
  (CONFIGURE 'ICS (FIND 'ALL 'CCS '(IMPLEMENTS ICS)))
  (DETERMINE '(#CLC.MLA.LINES ICS.PIBS))
  (CHECK 'IO-CAPACITY))

(DEFTASKBLOCK CONSOLE.TB CONSOLE ()
  (* jsb: " 4-Aug-83 23:13")
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK FICS.TB FICS () (* jsb: " 4-Aug-83 23:16")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK FIRMWARE.TB FIRMWARE ()
  (* jsb: " 7-AUG-83 17:24")
  (CHECK 'REQUIREMENTS)
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT)))
```

```

(DEFTASKBLOCK ICS.TB ICS ()          (* jsb: " 4-Aug-83 23:17")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK IDC.TB IDC ()          (* jsb: " 4-Aug-83 23:14")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK IO-PANEL-MODULES.TB IO-PANEL-MODULES ()
  (* jsb: " 7-AUG-83 20:30")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT))
  (DETERMINE '(#CT.LINES #IOLC.LINES)))
(DEFTASKBLOCK IO.TB IO ()          (* jsb: " 4-Aug-83 23:12")
  (CONFIGURE 'IDC (FIND 'ALL 'CCS '(IMPLEMENTS IDC)))
  (CONFIGURE 'CLC/MLA (FIND 'ALL 'CCS '(IMPLEMENTS
CLC/MLA)))
  (CONFIGURE '2STAGE-IO (FIND 'ALL 'CCS '(IMPLEMENTS 2STAGE-
IO))))

(DEFTASKBLOCK ISU.TB ISU ()          (* jst: "11-AUG-83 08:49")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT))
  (DETERMINE '(ISU-MEMORY.SUPPLIED))
  (CHECK 'REQUIREMENTS))

(DEFTASKBLOCK ITB-MODULES.TB ITB-MODULES ()
  (* jsb: " 7-AUG-83 20:29")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT))
  (DETERMINE '(#TRUNK.LINES IO.PIBS)))

(DEFTASKBLOCK MAINFRAME.TB MAINFRAME ()
  (* jsb: " 4-Aug-83 23:07")
  (CONFIGURE 'PROCESSOR (FIND 'ALL 'CCS '(IMPLEMENTS
PROCESSOR)))
  (CONFIGURE 'MEMORY (FIND 'ALL 'CCS '(IMPLEMENTS MEMORY)))
  (CONFIGURE 'ISU (FIND 'ALL 'CCS '(IMPLEMENTS ISU)))
  (CONFIGURE 'FIRMWARE (FIND 'ALL 'CCS '(IMPLEMENTS
FIRMWARE)))
  (CONFIGURE 'IO (FIND 'ALL 'CCS '(IMPLEMENTS IO)))
  (CONFIGURE 'PLENUM (FIND 'ALL 'CCS '(IMPLEMENTS PLENUM)))
  (CONFIGURE 'POWER-SUPPLY (FIND 'ALL 'CCS '(IMPLEMENTS
POWER-SUPPLY))))

(DEFTASKBLOCK MEMORY.TB MEMORY ()
  (* jsb: " 7-AUG-83 20:28")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT))
  (DETERMINE '(MAIN-MEMORY.SUPPLIED MEMORY.INTERLEAVING
MEMORY.PIBS))
  (CHECK 'INTERLEAVING)
  (CHECK 'MEMORY-CAPACITY)
  (CHECK 'MEMORY-IEMENT))

```



```

(DEFTASKBLOCK MISC.TB MISC () (* jsb: " 9-AUG-83 14:45")
  (EXPAND* '(AK FT KT))
  (CHECK 'REQUIREMENTS)
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK MULTIPROCESSING.TB MULTIPROCESSING ()
  (* jsb: " 7-AUG-83 20:31")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT))
  (DETERMINE '(MP.PIBS)))
(DEFTASKBLOCK PLENUM.TB PLENUM ()
  (* jsb: " 4-Aug-83 23:13")
  (CHECK 'CONTENTS)
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK POWER-SUPPLY.TB POWER-SUPPLY ()
  (* jsb: " 7-AUG-83 17:24")
  (CHECK 'REQUIREMENTS)
  (EXPAND* '(AK_FT KT)))

(DEFTASKBLOCK PROCESSOR.TB PROCESSOR ()
  (* jsb: " 4-Aug-83 23:08")
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK PRODUCT.LINE.TB PRODUCT.LINE ()
  (* jsb: " 3-AUG-83 18:17")
  (BVSET 'SYSTEM.TYPE (COMPUTE.SYSTEM.TYPE))
  (CONFIGURE (FETCH 'SYSTEM.TYPE)
    (FIND 'ALL 'CCS)))

(DEFTASKBLOCK REST.TB REST () (* jsb: " 4-Aug-83 23:14")
  (CONFIGURE 'CABLE (FIND 'ALL 'CCS '(IMPLEMENTS CABLE)))
  (CONFIGURE 'CONSOLE (FIND 'ALL 'CCS '(IMPLEMENTS
CONSOLE)))
  (CONFIGURE 'MISC (FIND 'ALL 'CCS '(IMPLEMENTS MISC))))

(DEFTASKBLOCK SP.TB SP () (* jsb: " 4-Aug-83 23:08")
  (EXPAND* '(AK FT KT)))

(DEFTASKBLOCK SYSTEM.TB SYSTEM ()
  (* jst: "11-AUG-83 09:11")
  (EXPAND* '(BG BU.V85))
  (DETERMINE '(CLASS.MODEL.EXPANSION BASIC.TO.PRICED
MP.ORDER))
  (CONFIGURE 'MULTIPROCESSING (FIND 'ALL 'CCS '(IMPLEMENTS
MULTIPROCESSING)))
  (CONFIGURE 'MAINFRAME (FIND 'ALL 'CCS '(IMPLEMENTS
MAINFRAME)))
  (CONFIGURE 'REST (FIND 'ALL 'CCS '(IMPLEMENTS REST)))
  (CHECK 'POWER)
  (CHECK 'PIBS)
  (EXPAND* '(AK FT KT))
  (CHECK 'EFFECTIVITY)
  (CHECK 'PHASE-OUT))

```

APPENDIX II (C)

KNOWLEDGE BASE FOR M1234 COMPUTER
PARTS CATALOG

```
(* XTEST-BSS definition)
(DEFINSTANCE XTEST-BSS BSS
  SET.OF XTESTKB
  FILED.ON XTESTBX
  ELEMENTS (M1234))

(DEFCLASS M1234 (PRODUCT.ID)
  ()
  (METACLASS BS
    DESCR "The M1234 product line"
    MFG.PLANT 1804
    IMPLEMENTS.FN PRODUCT.LINE
    BASIC.GROUPS NIL)
  )

(* XTEST-BGS definition)
(DEFINSTANCE XTEST-BGS BGS
  SET.OF XTESTKB
  FILED.ON XTESTBX
  ELEMENTS (M1234-II))

(DEFCLASS M1234-II (PRODUCT.ID)
  ()
  (METACLASS BG
    DESCR "M1234-II SYSTEM "
    IMPLEMENTS.FN M1234.SYSTEM
    BASIC.UNITS NIL)
  )

(* XTEST-BUS definition)
(DEFINSTANCE XTEST-BUS BUS
  SET.OF XTESTKB
  FILED.ON XTESTBX
  ELEMENTS (1234-0066-0690 1234-3003-0690))

(DEFCLASS 1234-0066-0690 (PRODUCT.ID)
  ()
  (METACLASS BU
    DESCR "Vanilla model for M1234 system"
    MFG.PLANT 1804
    IMPLEMENTS.FN SYSTEM
    ACCESSORIES NIL
    REQUIRED.IVARS NIL)
  )

(DEFCLASS 1234-3003-0690 (PRODUCT.ID)
  ()
  (METACLASS BU.M12
    DESCR "Class model for M1234 system"
```

```

IMPLEMENTS.FN M1234.SYSTEM
BASIC.TO.PRICED ((1 . 1234-K416)
                 (1 . 1234-K920)
                 (2 . 1234-K924)
                 (1 . 1234-K940)
                 (1 . 1234-K941))
VANILLA.MODEL 1234-0066-0690)

```

```

(* XTEST-AKS definition)
(DEFINANCE XTEST-AKS AKS
 SET.OF XTESTKB
 FILED.ON XTESTAK
 ELEMENTS (1234-P413 1234-P723 1234-P724 1234-P742 1234-
           P743 1234-P954 1234-P955))

```

```

(DEFCLASS 1234-P413 (PRODUCT.ID)
 ()
 (METACCLASS AK
  MFG.PLANT 1804
  DESCR "Additional Power Supply"
  IMPLEMENTS.FN POWER-SUPPLY)
 )

```

```

(DEFCLASS 1234-P723 (PRODUCT.ID)
 ()
 (METACCLASS AK
  DESCR "1024K TO 2048K MEM ADD (CONT TWO 1234-K920, FOUR
                                     1234-K924)"
  MFG.PLANT 1804
  IMPLEMENTS.FN MEMORY)
 )

```

```

(DEFCLASS 1234-P724 (PRODUCT.ID)
 ()
 (METACCLASS AK
  DESCR "2048K TO 3072 MEMORY ADDITION (CONTAINS FOUR 1234-
                                           K924)"
  MFG.PLANT 1804
  IMPLEMENTS.FN MEMORY)
 )

```

```

(DEFCLASS 1234-P742 (PRODUCT.ID)
 ()
 (METACCLASS AK
  DESCR "VHST (INCL. 1234-K930 & K931) "
  MFG.PLANT 1804
  IMPLEMENTS.FN 2STAGE-IO)
 )

```

```

(DEFCLASS 1234-P743 (PRODUCT.ID)
 ()
 (METACCLASS AK
  DESCR "I/O LINK CONTROL "
  MFG.PLANT 1804

```

COMMENT (NOTE CRIT-KITS-0006-NOTE-09)
 IMPLEMENTS.FN 2STAGE-IO)

)

(DEFCLASS 1234-P954 (PRODUCT.ID)

()

(METACCLASS AK

DESCR "FIRST CLC/MLA COMB (CONTAINS 1234-K905, K901 AND
 K903)"

MFG.PLANT 1804

IMPLEMENTS.FN CLC/MLA)

)

(DEFCLASS 1234-P955 (PRODUCT.ID)

()

(METACCLASS AK

DESCR "SECOND CLC/MLA COMB (CONTAINS 1234-K905, K901 AND
 K903)"

MFG.PLANT 1804

IMPLEMENTS.FN CLC/MLA)

)

(* XTEST-FTS definition)

(DEFINSTANCE XTEST-FTS FTS —

SET.OF XTESTKB

FILED.ON XTESTFT

ELEMENTS (1234-F010 1234-F113 1234-F262 1234-F290 1234-
 F360 1234-F361 1234-F362 1234-F400 1234-F413
 1234-F504 1234-F510 1234-F545 1234-F680 1234-
 F690 1234-F691 1234-K010))

(DEFCLASS 1234-F010 (PRODUCT.ID)

()

(METACCLASS FT

DESCR "Firmware"

IMPLEMENTS.FN FIRMWARE

UNIQUE.EXPANSION 1234-K010)

)

(DEFCLASS 1234-F113 (PRODUCT.ID)

()

(METACCLASS FT

IMPLEMENTS.FN ISU

DESCR "ISU 113K"

PIBS (0 0)

POWER (0.0 8.0 33.0))

)

(DEFCLASS 1234-F262 (PRODUCT.ID)

()

(METACCLASS FT

DESCR "Memory 1M, 2 way"

IMPLEMENTS.FN MEMORY

POWER (6.0 0.0 34.0)

PIBS (4 0))

)

```

(DEFCLASS 1234-F290 (PRODUCT.ID)
  ()
  (METACCLASS FT
    DESCR "Memory 2M, 2 way"
    IMPLEMENTS.FN MEMORY
    POWER (44.0 0.0 12.0)
    PIBS (8 0))
  )

(DEFCLASS 1234-F360 (PRODUCT.ID)
  ()
  (METACCLASS FT
    DESCR "MFG.PLANT COMMON TRUNK "
    IMPLEMENTS.FN ITB-MODULES
    POWER (0.0 8.6 13.5)
    PIBS (2 2))
  )

(DEFCLASS 1234-F361 (PRODUCT.ID)
  ()
  (METACCLASS FT
    DESCR "Common trunk"
    IMPLEMENTS.FN ITB-MODULES
    POWER (0.0 17.2 27.0))
  )

(DEFCLASS 1234-F362 (PRODUCT.ID)
  ()
  (METACCLASS FT
    DESCR "Common trunk"
    IMPLEMENTS.FN ITB-MODULES
    POWER (0.0 25.8 40.5))
  )

(DEFCLASS 1234-F400 (PRODUCT.ID)
  ()
  (METACCLASS FT
    IMPLEMENTS.FN POWER-SUPPLY
    MFG.PLANT 1804
    DESCR "-5V power supply")
  )

(DEFCLASS 1234-F413 (PRODUCT.ID)
  ()
  (METACCLASS FT
    DESCR "MFG.PLANT POWER SUPPLY "
    IMPLEMENTS.FN POWER-SUPPLY
    POWER (0 0 -62)
    PIBS (0 0))
  )

(DEFCLASS 1234-F504 (PRODUCT.ID)
  ()
  (METACCLASS FT
    DESCR "MFG.PLANT FIRST CLC/MLA "

```

```
IMPLEMENTS.FN CLC/MLA
POWER (2 1 5.5)
PIBS (2 2))
)
```

```
(DEFCLASS 1234-F510 (PRODUCT.ID)
```

```
( )
(METACCLASS FT
  DESCR "Fast ICS, 2CLC, 2MLA"
  IMPLEMENTS.FN FICS
  POWER (11.0 2.0 4.0)
  PIBS (4 4))
)
```

```
(DEFCLASS 1234-F545 (PRODUCT.ID)
```

```
( )
(METACCLASS FT
  DESCR "MFG.PLANT ICS LIGHT DISPLAY "
  IMPLEMENTS.FN MISC)
)
```

```
(DEFCLASS 1234-F680 (PRODUCT.ID)
```

```
( )
(METACCLASS FT
  DESCR "MFG.PLANT BYTE MULTIPLEX TRUNK "
  IMPLEMENTS.FN IO-PANEL-MODULES
  POWER (7 .4 .5)
  PIBS (0 0))
)
```

```
(DEFCLASS 1234-F690 (PRODUCT.ID)
```

```
( )
(METACCLASS FT
  DESCR "MFG.PLANT BSLA"
  IMPLEMENTS.FN IO-PANEL-MODULES
  POWER (7.8 .4 .5)
  PIBS (0 0))
)
```

```
(DEFCLASS 1234-F691 (PRODUCT.ID)
```

```
( )
(METACCLASS FT
  DESCR "BSLA"
  IMPLEMENTS.FN IO-PANEL-MODULES
  POWER (15.6 .8 1.0))
)
```

```
(DEFCLASS 1234-K010 (PRODUCT.ID)
```

```
( )
(METACCLASS FT
  DESCR "Firmware"
  IMPLEMENTS.FN FIRMWARE)
)
```

```
(* XTEST-KTS definition)
(DEFINSTANCE XTEST-KTS KTS
  SET.OF XTESTKB
  FILED.ON XTESTKT
  ELEMENTS      (1234-K010 1234-K400 1234-K413 1234-K416 1234-
                 K450 1234-K901 1234-K903 1234-K905 1234-K913
                 1234-K915 1234-K917 1234-K920 1234-K921 1234-
                 K924 1234-K930 1234-K931 1234-K935 1234-K936
                 1234-K937 1234-K939 1234-K940 1234-K941))
```

```
(DEFCLASS 1234-K010 (PRODUCT.ID)
```

```
( )
```

```
(METACLASS FT
```

```
  DESCR "Firmware"
```

```
  IMPLEMENTS.FN FIRMWARE)
```

```
)
```

```
(DEFCLASS 1234-K400 (PRODUCT.ID)
```

```
( )
```

```
(METACLASS KT
```

```
  IMPLEMENTS.FN POWER-SUPPLY
```

```
  MFG.PLANT 1804
```

```
  DESCR "-5V power supply")
```

```
)
```

```
(DEFCLASS 1234-K413 (PRODUCT.ID)
```

```
( )
```

```
(METACLASS KT
```

```
  DESCR "+5V ADD. POWER SUPPLY      "
```

```
  MFG.PLANT 1804
```

```
  IMPLEMENTS.FN POWER-SUPPLY)
```

```
)
```

```
(DEFCLASS 1234-K416 (PRODUCT.ID)
```

```
( )
```

```
(METACLASS KT
```

```
  DESCR "CONSOLE INDICATOR INTERFACE  "
```

```
  MFG.PLANT 1804
```

```
  IMPLEMENTS.FN CONSOLE)
```

```
)
```

```
(DEFCLASS 1234-K450 (PRODUCT.ID)
```

```
( )
```

```
(METACLASS KT
```

```
  DESCR "VS HARDWARE      "
```

```
  MFG.PLANT 1804
```

```
  IMPLEMENTS.FN MAINFRAME)
```

```
)
```

```
(DEFCLASS 1234-K901 (PRODUCT.ID)
```

```
( )
```

```
(METACLASS KT
```

```
  DESCR "MULTILINE ADAPTER  "
```

```
  MFG.PLANT 1804
```

```
  IMPLEMENTS.FN CLC/MLA)
```

```
)
```

(DEFCLASS 1234-K903 (PRODUCT.ID)

()

(METACCLASS KT

DESCR "CLC EXPANSION MODULE"

MFG.PLANT 1804

IMPLEMENTS.FN CLC/MLA)

)

(DEFCLASS 1234-K905 (PRODUCT.ID)

()

(METACCLASS KT

DESCR "FAST CLC "

MFG.PLANT 1804

IMPLEMENTS.FN CLC/MLA)

)

(DEFCLASS 1234-K913 (PRODUCT.ID)

()

(METACCLASS KT

DESCR "ISU, 8K FAST "

MFG.PLANT 1804

IMPLEMENTS.FN ISU)

)

(DEFCLASS 1234-K915 (PRODUCT.ID)

()

(METACCLASS KT

DESCR "?? Mystery kit ?? "

MFG.PLANT 1804

IMPLEMENTS.FN ISU)

)

(DEFCLASS 1234-K917 (PRODUCT.ID)

()

(METACCLASS KT

DESCR "SERVICE PROCESSOR, ISU NO.1 "

MFG.PLANT 1804

IMPLEMENTS.FN ISU)

)

(DEFCLASS 1234-K920 (PRODUCT.ID)

()

(METACCLASS KT

DESCR "MSU EXPANSION MODULE"

MFG.PLANT 1804

IMPLEMENTS.FN MEMORY)

)

(DEFCLASS 1234-K921 (PRODUCT.ID)

()

(METACCLASS KT

DESCR "MSU, INTERFACE "

MFG.PLANT 1804

IMPLEMENTS.FN MEMORY)

)


```

(DEFCLASS 1234-K924 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "MSU 256K ARRAY "
    MFG.PLANT 1804
    IMPLEMENTS.FN MEMORY)
  )
)
(DEFCLASS 1234-K930 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "COMMON TRUNK DATA INTERFACE "
    MFG.PLANT 1804
    IMPLEMENTS.FN IO-PANEL-MODULES)
  )
)
(DEFCLASS 1234-K931 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "DIRECT MEMORY ACCESS TRUNK "
    MFG.PLANT 1804
    IMPLEMENTS.FN ITB-MODULES)
  )
)
(DEFCLASS 1234-K935 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "BIT SERIAL LINK CABLES "
    MFG.PLANT 1804
    IMPLEMENTS.FN IO-PANEL-MODULES)
  )
)
(DEFCLASS 1234-K936 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "CONNECTOR MODULE "
    MFG.PLANT 1804
    IMPLEMENTS.FN IO-PANEL-MODULES)
  )
)
(DEFCLASS 1234-K937 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "CT BSL ADAPTER "
    MFG.PLANT 1804
    IMPLEMENTS.FN IO-PANEL-MODULES)
  )
)
(DEFCLASS 1234-K939 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "BSL ENTRY MODULE "
    MFG.PLANT 1804
    IMPLEMENTS.FN IO-PANEL-MODULES)
  )
)

```

```
(DEFCLASS 1234-K940 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "ISU/MSU RESISTOR      "
    MFG.PLANT 1804
    IMPLEMENTS.FN MEMORY)
  )
)
```

```
(DEFCLASS 1234-K941 (PRODUCT.ID)
  (
  (METACCLASS KT
    DESCR "CRYSTAL 84/56 NS      "
    MFG.PLANT 1804
    IMPLEMENTS.FN PROCESSOR)
  )
)
```

```
(* XTEST-KVS definition)
(DEFINSTANCE XTEST-KVS KVS
  SET.OF XTESTKB
  FILED.ON XTESTKV
  ELEMENTS (1234-K400-V001))
```

```
(DEFCLASS 1234-K400-V001 (PRODUCT.ID)
  (
  (METACCLASS KV
    IMPLEMENTS.FN POWER-SUPPLY
    MFG.PLANT 1804
    DESCR "-5V power supply")
  )
)
```

APPENDIX II (D)

KNOWLEDGE BASE FOR M1234 COMPUTER CONSTRAINTS & BIN VARIABLES KNOWLEDGE BASE FUNCTIONS

```
(* XTEST-CTS definition)
(DEFINSTANCE XTEST-CTS CTS
  SET.OF XTESTKB
  FILED.ON XTESTNCT
  ELEMENTS (AT-LEAST-ONE-FT.CT AT-MOST-ONE-FT.CT AT-MOST-
    ONE-MEMORY-FT.CT NO-FTS.CT NO-MP-
    FTS.CT #CLC.MLA-LINES CLC/MLA-
    DISPLAY CT.IO CT.IOLC.IO
    EXTRA.MINUS5 EXTRA.PLUS5 FIRMWARE.CT
    ICS.PIBS.CT IO.PIBS.CT MINUS5
    MP.FIRMWARE.CT MP.ISU.CT MP.POWER-
    SUPPLY.CT PLUS5 M1234.PIBS.CT1
    M1234.PIBS.CT2
  ))
```

```
(DEFCLASS AT-LEAST-ONE-FT.CT (CTI)
  (
  (METACCLASS WARN-ON-SUCCESS.CT
    BIN.TYPES (ISU FIRMWARE)
```

```

STATEMENT ((COUNT (FIND 'ALL 'CCS '(ISA FT)))
           IS ZERO)
MESSAGE ("The customer must order some part that
         implements " (SAND CURRENT.BIN
                       'GET.TRANSLATION)

```

```

"; this order is
incorrect.")

```

```

CHECK.KEYWORDS (CONTENTS)
)

```

```

(DEFCLASS AT-MOST-ONE-FT.CT (CTI)

```

```

()

```

```

(METACLASS MODIFY-ON-SUCCESS.CT

```

```

  BIN.TYPES (CLC/MLA PLENUM ISU FIRMWARE)

```

```

  CHECK.KEYWORDS (CONTENTS)

```

```

  STATEMENT ((COUNT (FIND 'ALL 'CCS '(ISA FT)))
             > 1)

```

```

  DEL-PARTS (ELIMINATE.BTP (FIND* 'ALL 'CCS '(ISA FT)))
)

```

```

(DEFCLASS AT-MOST-ONE-MEMORY-FT.CT (CTI)

```

```

()

```

```

(METACLASS MODIFY-ON-SUCCESS.CT

```

```

  CHECK.KEYWORDS (CONTENTS)

```

```

  BIN.TYPES (MEMORY)

```

```

  STATEMENT ((SYSTEM M1234.SYSTEM)

```

```

    AND

```

```

    (COUNT (FIND 'ALL 'CCS '(ISA FT)))

```

```

    > 1)

```

```

  DEL-PARTS (ELIMINATE.BTP (FIND* 'ALL 'CCS '(ISA FT)))
)

```

```

(DEFCLASS NO-FTS.CT (CTI)

```

```

()

```

```

(METACLASS MODIFY-ON-SUCCESS.CT

```

```

  BIN.TYPES (IDC)

```

```

  CHECK.KEYWORDS (CONTENTS)

```

```

  STATEMENT ((COUNT (FIND 'ALL 'CCS '(ISA FT)))
             IS NOT ZERO)

```

```

  MESSAGE ("The customer is not allowed to order products
           which implement "

```

```

           (SAND CURRENT.BIN 'GET.TRANSLATION)

```

```

           ". This order is not correct.")

```

```

  DEL-PARTS (FIND* 'ALL 'CCS '(ISA FT))
)

```

```

(DEFCLASS NO-MP-FTS.CT (CTI)

```

```

()

```

```

(METACLASS MODIFY-ON-SUCCESS.CT

```

```

  CHECK.KEYWORDS (CONTENTS)

```

```

  BIN.TYPES (MULTIPROCESSING)

```

```

  STATEMENT ((SYSTEM)

```

```

    AND

```

```

    (COUNT (FIND 'ALL 'CCS '(ISA FT)))

```

```

    IS NOT ZERO)

```

```

  DEL-PARTS (FIND* 'ALL 'CCS '(ISA FT))
)

```

```

(DEFCLASS #CLC.MLA-LINES (CTI)
()
(METACLASS WARN-ON-SUCCESS.CT
  BIN.TYPES (CLC/MLA)
  BIN.VARS.REQUIRED (#CLC.MLA.LINES)
  STATEMENT ((SYSTEM M1234.SYSTEM)
    AND
    (#CLC.MLA.LINES > 20))
  MESSAGE ("The order specifies too many CLC/MLA lines.")
  CHECK.KEYWORDS (IO-CAPACITY)
)
(DEFCLASS CLC/MLA-DISPLAY (CTI)
()
(METACLASS MODIFY-ON-SUCCESS.CT
  BIN.TYPES (MISC)
  BIN.VARS.REQUIRED (#CLC.MLA.LINES)
  STATEMENT ((SYSTEM M1234.SYSTEM)
    AND
    ((COUNT.PARTS '1234-K545 T)
    = 1)
    AND
    (BETWEEN #CLC.MLA.LINES 11 20))
  MESSAGE (
    "This order uses more than 10 CLC/MLA lines. However, the
    customer only ordered 1 Remote ICS display (1234-F545).")
  )
  ADD-PARTS (QUOTE ((1 . 1234-F545)))
  CHECK.KEYWORDS (REQUIREMENTS)
)
(DEFCLASS CT.IO (CTI)
()
(METACLASS WARN-ON-FAILURE.CT
  BIN.VARS.REQUIRED (#TRUNK.LINES #CT.LINES)
  CHECK.KEYWORDS (IO-CAPACITY)
  BIN.TYPES (2STAGE-IO)
  STATEMENT (#CT.LINES <= #TRUNK.LINES)
  MESSAGE ("There are not enough common trunk I/F lines for
the HS, MS, and LS lines ordered.")
)
)
(DEFCLASS CT.IOLC.IO (CTI)
()
(METACLASS WARN-ON-FAILURE.CT
  BIN.VARS.REQUIRED (#TRUNK.LINES #CT.LINES #IOLC.LINES)
  CHECK.KEYWORDS (IO-CAPACITY)
  BIN.TYPES (2STAGE-IO)
  STATEMENT (#TRUNK.LINES = #IOLC.LINES + #CT.LINES)
  MESSAGE (
    "The number of IOLC and common trunk I/F lines does not
    match the number of HS, MS, and LS IO lines.")
  )
)

```

```
(DEFCLASS EXTRA.MINUS5 (CTI)
  ()
  (METAClass WARN-ON-SUCCESS.CT
   BIN.TYPES (SYSTEM)
   BIN.VARS.REQUIRED (POWER.USED MP.ORDER)
   STATEMENT ((SYSTEM M1234.SYSTEM)
              AND
              (NOT MP.ORDER)
              AND
              (ORDERED? '1234-F400 T)
              AND
              (POWER.USED:MINUS.5 <= 250))
   MESSAGE ("The customer ordered a -5V power supply which
is not required to run this system.")
   CHECK.KEYWORDS (POWER))
  )
```

```
(DEFCLASS EXTRA.PLUS5 (CTI)
  ()
  (METAClass WARN-ON-SUCCESS.CT
   BIN.TYPES (SYSTEM)
   BIN.VARS.REQUIRED (POWER.USED)
   STATEMENT ((SYSTEM M1234.SYSTEM)
              AND
              (ORDERED? '1234-F413 T)
              AND
              (POWER.USED:PLUS.5 <= 60))
   MESSAGE ("The customer ordered a +5V power supply which
is not required to run this system.")
   CHECK.KEYWORDS (POWER))
  )
```

```
(DEFCLASS FIRMWARE.CT (CTI)
  ()
  (METAClass MODIFY-ON-SUCCESS.CT
   BIN.TYPES (FIRMWARE)
   BIN.VARS.REQUIRED (MP.ORDER)
   CHECK.KEYWORDS (REQUIREMENTS)
   STATEMENT ((SYSTEM M1234.SYSTEM)
              AND
              (NOT MP.ORDER)
              AND
              (NOT (ORDERED? '(1234-F010 1234-K010 1234-F011
1234-K011))))
   MESSAGE (
"NO firmware has been ordered for this non-MP system. Please
check with the customer and determine what operating system
will be used with this system. If it is VS-1, they should
order 1 1234-F011; otherwise if it is RS-1, they should
order 1 1234-F010. Please select one of these features for
this order.")
   SET-PARTS (CHOOSE.ONE '((1. 1234-F010) (1 . 1234-F011))))
  )
```

```
(DEFCLASS ICS.PIBS.CT (CTI)
()
(METACLASS WARN-ON-SUCCESS.CT
BIN.TYPES (SYSTEM)
BIN.VARS.REQUIRED (ICS.PIBS)
STATEMENT ((SYSTEM M1234.SYSTEM)
AND
(ICS.PIBS > 8))
MESSAGE ("The order uses more ICS PIBS than is allowed
for this system.")
CHECK.KEYWORDS (PIBS))
)
```

```
(DEFCLASS IO.PIBS.CT (CTI)
()
(METACLASS WARN-ON-SUCCESS.CT
BIN.TYPES (SYSTEM)
BIN.VARS.REQUIRED (IO.PIBS ICS.PIBS)
STATEMENT ((SYSTEM M1234.SYSTEM)
AND
(IO.PIBS + ICS.PIBS > 12))
MESSAGE ("The order uses more IO and ICS PIBS combined
than is allowed for this system.")
CHECK.KEYWORDS (PIBS))
)
```

```
(DEFCLASS MINUS5 (CTI)
()
(METACLASS MODIFY-ON-SUCCESS.CT
BIN.TYPES (SYSTEM)
BIN.VARS.REQUIRED (POWER.SUPPLIED POWER.USED)
STATEMENT ((SYSTEM M1234.SYSTEM)
AND
(POWER.USED:MINUS.5 - POWER.SUPPLIED:MINUS.5)
> 250)
MESSAGE ("The order uses more -5V power than is
supplied.")
ADD-PARTS (QUOTE ((1 . 1234-F400)))
CHECK.KEYWORDS (POWER))
)
```

```
(DEFCLASS MP.FIRMWARE.CT (CTI)
()
(METACLASS MODIFY-ON-SUCCESS.CT
BIN.TYPES (FIRMWARE)
BIN.VARS.REQUIRED (MP.ORDER)
STATEMENT ((SYSTEM M1234.SYSTEM)
AND MP.ORDER AND (NOT (ORDERED? '5640-F400)))
MESSAGE ("All MP orders require MP firmware.")
SET-PARTS (QUOTE ((1 . 5640-F400)))
CHECK.KEYWORDS (REQUIREMENTS))
)
```

```
(DEFCLASS MP.ISU.CT (CTI)
()
(METACLASS WARN-ON-SUCCESS.CT
```

```

BIN.TYPES (ISU)
CHECK.KEYWORDS (REQUIREMENTS)
STATEMENT ((SYSTEM M1234.SYSTEM )
           AND MP.ORDER AND (ISU-MEMORY.SUPPLIED < 40))
BIN.VARS.REQUIRED (MP.ORDER ISU-MEMORY.SUPPLIED)
MESSAGE (
"This MP order requires at least 40KB of ISU memory.
If this system will be running VS2 V03, please select 1234-
F122 or F132; otherwise, if it will be running VS2 V04,
please select 1234-F123 or F128.")
)

```

```

(DEFCLASS MP.POWER-SUPPLY.CT (CTI)
()
(METACLASS MODIFY-ON-SUCCESS.CT
BIN.TYPES (POWER-SUPPLY)
CHECK.KEYWORDS (REQUIREMENTS)
STATEMENT ((SYSTEM M1234.SYSTEM)
           AND MP.ORDER AND (NOT (ORDERED? '(1234-F400 1234-
           K400))))
BIN.VARS.REQUIRED (MP.ORDER)
ADD-PARTS (QUOTE ((1 . 1234-F400)))
MESSAGE ("This MP order requires an additional +5V power
supply."))
)

```

```

(DEFCLASS PLUS5 (CTI)
()
(METACLASS MODIFY-ON-SUCCESS.CT
BIN.TYPES (SYSTEM)
BIN.VARS.REQUIRED (POWER.SUPPLIED POWER.USED)
STATEMENT ((SYSTEM M1234.SYSTEM)
           AND
           (POWER.USED:PLUS.5 - POWER.SUPPLIED:PLUS.5)
           > 60)
MESSAGE ("The order uses more +5V power than is
supplied.")
ADD-PARTS (QUOTE ((1 . 1234-F413)))
CHECK.KEYWORDS (POWER))
)

```

```

(DEFCLASS M1234.PIBS.CT1 (CTI)
()
(METACLASS WARN-ON-SUCCESS.CT
BIN.TYPES (SYSTEM)
BIN.VARS.REQUIRED (TOTAL.PIBS)
CHECK.KEYWORDS (PIBS)
STATEMENT ((SYSTEM M1234.SYSTEM)
           AND
           (ORDERED? '1234-K926)
           AND
           (TOTAL.PIBS > 38))
MESSAGE (
"The order includes a LH plenum and uses more than 38 PIBS;
this is not a legal order.")
)

```

```

(DEFCLASS M1234.PIBS.CT2 (CTI)
  ()
  (METACLASS MODIFY-ON-SUCCESS.CT
    BIN.TYPES (SYSTEM)
    BIN.VARS.REQUIRED (TOTAL.PIBS)
    CHECK.KEYWORDS (PIBS)
    STATEMENT ((SYSTEM M1234.SYSTEM)
      AND
      (NOT (ORDERED? '1234-K926))
      AND
      (TOTAL.PIBS > 38))
    MESSAGE ("The order uses more than 38 PIBS; must add LH
plenum.")
    ADD-PARTS (QUOTE ((1 . 1234-F441))))
)

(* XTEST-NCTS definition)
(DEFINSTANCE XTEST-NCTS NCTS
  SET.OF XTESTKB
  FILED.ON XTESTNCT
  ELEMENTS (1024-INCREMENT 2-WAY-INTERLEAVING CT-LINES IOLC-
    LINES))

(DEFCLASS 1024-INCREMENT (CTI)
  ()
  (METACLASS MEMORY-INCREMENT.CT
    BIN.TYPES (MEMORY)
    STATEMENT ((SYSTEM M1234.SYSTEM)
      AND
      (NOT (INCREMENT MAIN-MEMORY.SUPPLIED 1024))))
)

(DEFCLASS 2-WAY-INTERLEAVING (CTI)
  ()
  (METACLASS INTERLEAVING.CT
    BIN.TYPES (MEMORY)
    STATEMENT ((SYSTEM M1234.SYSTEM)
      AND
      (NOT (MEMORY.INTERLEAVING = '2-WAY))))
)

(DEFCLASS CT-LINES (CTI)
  ()
  (METACLASS IO-LINES.CT
    BIN.TYPES (2STAGE-IO)
    STATEMENT ((SYSTEM M1234.SYSTEM)
      AND
      (NOT (BETWEEN #CT.LINES 0 6))))
)

(DEFCLASS IOLC-LINES (CTI)
  ()
  (METACLASS IO-LINES.CT
    BIN.TYPES (2STAGE-IO)
    STATEMENT ((SYSTEM M1234.SYSTEM)
      AND
      (NOT (BETWEEN #IOLC.LINES 0 6))))
)

```



```

(* XTEST-BVS definition)
(DEFINSTANCE XTEST-BVS BVS _
  SET.OF XTESTKB
  FILED.ON XTESTNCT
  ELEMENTS (#CLC #CLC.MLA.LINES #CT.LINES #IOLC.LINES #MLA
            #TRUNK.LINES 59.POSITION.BACKPANEL
            BASIC.TO.PRICED CAP.PROCESSOR
            CLASS.MODEL.EXPANSION ICS.PIBS IO.PIBS
            ISU-MEMORY.SUPPLIED MAIN-MEMORY.SUPPLIED
            MEMORY.INTERLEAVING MEMORY.PIBS
            MP.ORDER MP.PIBS OPERATING.SYSTEM
            POWER.SUPPLIED POWER.USED TOTAL.PIBS))

(DEFCLASS #CLC (BVI)
  ()
  (METAClass BIN.VAR
   BIN.TYPES (SYSTEM)
   DESCR "the number of CLC's"
   HOW.TO.DETERMINE (COMPUTE.#CLC)
   LEGAL.VALUES INTEGER)
  )

(DEFCLASS #CLC.MLA.LINES (BVI)
  ()
  (METAClass BIN.VAR
   BIN.TYPES (SYSTEM)
   DESCR "the number of CLC/MLA lines for this order"
   HOW.TO.DETERMINE (COMPUTE.#CLC.MLA.LINES)
   LEGAL.VALUES (AND INTEGER (SATISFIES (LAMBDA (value)
                                           (BETWEEN value 0 20))))))
  )

(DEFCLASS #CT.LINES (BVI)
  ()
  (METAClass BIN.VAR
   BIN.TYPES (SYSTEM)
   DESCR "the number of common trunk interface lines"
   HOW.TO.DETERMINE (COMPUTE.#CT.LINES)
   LEGAL.VALUES (AND INTEGER (SATISFIES (LAMBDA (value)
                                           (BETWEEN value 0 8))))))
  )

(DEFCLASS #IOLC.LINES (BVI)
  ()
  (METAClass BIN.VAR
   BIN.TYPES (SYSTEM)
   DESCR "the number of IOLC interface lines"
   HOW.TO.DETERMINE (COMPUTE.#IOLC.LINES)
   LEGAL.VALUES (AND INTEGER (SATISFIES (LAMBDA (value)
                                           (BETWEEN value 0 8))))))
  )

(DEFCLASS #MLA (BVI)
  ()
  (METAClass BIN.VAR
   BIN.TYPES (SYSTEM)

```

```
DESCR "the number of MLA's"
HOW.TO.DETERMINE (COMPUTE.#MLA)
LEGAL.VALUES INTEGER)
)
```

```
(DEFCLASS #TRUNK.LINES (BVI)
```

```
(
(METAClass BIN.VAR
BIN.TYPES (SYSTEM)
DESCR "the number of standard IO lines for this order"
HOW.TO.DETERMINE (COMPUTE.#TRUNK.LINES)
LEGAL.VALUES (AND INTEGER (SATISFIES (LAMBDA (value)
(BETWEEN value 0 8))))))
)
```

```
(DEFCLASS 59.POSITION.BACKPANEL (BVI)
```

```
(
(METAClass BIN.VAR
BIN.TYPES (SYSTEM)
DESCR "this order will use a 59-position backpanel (T/N >
1762)"
HOW.TO.DETERMINE (COMPUTE.59.POSITION.BACKPANEL)
LEGAL.VALUES BOOLEAN)
)
```

```
(DEFCLASS BASIC.TO.PRICED (BVI)
```

```
(
(METAClass BIN.VAR
BIN.TYPES (SYSTEM)
DESCR "the basic-to-priced parts"
HOW.TO.DETERMINE (COMPUTE.BASIC.TO.PRICED)
LEGAL.VALUES PRODUCT.QTYS)
)
```

```
(DEFCLASS CAP.PROCESSOR (BVI)
```

```
(
(METAClass BIN.VAR
BIN.TYPES (SYSTEM)
DESCR "this is a CAP processor"
HOW.TO.DETERMINE (ASK.BOOLEAN.QUESTION)
LEGAL.VALUES BOOLEAN)
)
```

```
(DEFCLASS CLASS.MODEL.EXPANSION (BVI)
```

```
(
(METAClass BIN.VAR
BIN.TYPES (SYSTEM)
DESCR "the class model parts"
HOW.TO.DETERMINE (COMPUTE.CLASS.MODEL.EXPANSION)
LEGAL.VALUES PRODUCT.QTYS)
)
```

```
(DEFCLASS ICS.PIBS (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "PIBS required for CLC/MLA"
    HOW.TO.DETERMINE (COMPUTE.ICS.PIBS)
    LEGAL.VALUES INTEGER)
  )

(DEFCLASS IO.PIBS (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "PIBS required for IOLC and common trunk"
    HOW.TO.DETERMINE (COMPUTE.IO.PIBS)
    LEGAL.VALUES INTEGER)
  )

(DEFCLASS ISU-MEMORY.SUPPLIED (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "the ISU memory supplied with this order (in KB)"
    HOW.TO.DETERMINE (COMPUTE.ISU-MEMORY.SUPPLIED)
    LEGAL.VALUES INTEGER)
  )

(DEFCLASS MAIN-MEMORY.SUPPLIED (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "the main memory supplied with this order (in KB)"
    HOW.TO.DETERMINE (COMPUTE.MAIN-MEMORY.SUPPLIED)
    LEGAL.VALUES INTEGER)
  )

(DEFCLASS MEMORY.INTERLEAVING (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "the memory interleaving for this order"
    HOW.TO.DETERMINE (COMPUTE.MEMORY.INTERLEAVING)
    LEGAL.VALUES (AND ATOM (MEMBEROF (1-WAY 2-WAY 4-WAY))))
  )

(DEFCLASS MEMORY.PIBS (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "PIBS required for main memory w/o extra plenums"
    HOW.TO.DETERMINE (COMPUTE.MEMORY.PIBS)
    LEGAL.VALUES INTEGER)
  )
```

```

(DEFCLASS MP.ORDER (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "this is a multi-processing order"
    HOW.TO.DETERMINE (COMPUTE.MP.ORDER)
    LEGAL.VALUES BOOLEAN)
  )
(DEFCLASS MP.PIBS (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "PIBS required for MP kits"
    HOW.TO.DETERMINE (COMPUTE.MP.PIBS)
    LEGAL.VALUES INTEGER)
  )
(DEFCLASS OPERATING.SYSTEM (BVI)
  ()
  (METACLASS ENUMERATED.BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "the operating system for this order"
    ENUMERATED.SET (RS-1 VS-1 MP))
  )
(DEFCLASS POWER.SUPPLIED (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "the power supplied for this order"
    HOW.TO.DETERMINE (COMPUTE.POWER.SUPPLIED)
    LEGAL.VALUES POWER.SPEC)
  )
(DEFCLASS POWER.USED (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "the power used by this order"
    HOW.TO.DETERMINE (COMPUTE.POWER.USED)
    LEGAL.VALUES POWER.SPEC)
  )
(DEFCLASS TOTAL.PIBS (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (SYSTEM)
    DESCR "the total PIBS required for this order"
    HOW.TO.DETERMINE (COMPUTE.TOTAL.PIBS)
    LEGAL.VALUES INTEGER)
  )

```



```

(SANDC (FIND 'ONE 'CCS '(ISA BU.M12))
'GET.BASIC.TO.PRICED))

(DEFKBFUN COMPUTE.CLASS.MODEL.EXPANSION () (* edited:
"11-AUG-83 06:02")
(PROG (CTS:(LISTOF CONSTRAINT)
CT:CONSTRAINT)
(CTS :=(SANDC (FIND 'ONE 'CCS '(ISA BU.M12))
'GET.EXPANSION.CTS))
(IF CTS
THEN (CT :=(SAND (CAR CTS)
'GET.TO)))
(RETURN CT)))

(DEFKBFUN COMPUTE.ICS.PIBS () (* edited: "10-AUG-83 20:39")
(PROG (ANS:INTEGER)
(ANS :=(FETCH '#CLC)+(FETCH '#MLA))
(IF (IREMAINDER ANS 2)= 1
THEN (ANS := ANS + 1))
(RETURN ANS)))

(DEFKBFUN COMPUTE.IO.PIBS () (* edited: "11-AUG-83 02:12")
(2*(FETCH '#TRUNK.LINES)))

(DEFKBFUN COMPUTE.ISU-MEMORY.SUPPLIED () (* jst: " 3-AUG-
83 10:57")
(2*(COUNT.PARTS '1234-K912)+ 8*(COUNT.PARTS '1234-K913)+
32*(COUNT.PARTS '1234-K909)))

(DEFKBFUN COMPUTE.MAIN-MEMORY.SUPPLIED () (* jst: "3-AUG-
83 10:57")
(256*(COUNT.PARTS '1234-K924)+ 64*(COUNT.PARTS '1234-
K922)))

(DEFKBFUN COMPUTE.MEMORY.INTERLEAVING () (* edited: "11-
AUG-83 06:19")
(SELECTQ (COUNT.PARTS '1234-K921)
(1 '1-WAY)
(2 '2-WAY)
((4 0)
'4-WAY)
(SHOULDNT)))

(DEFKBFUN COMPUTE.MEMORY.PIBS () (* jst: "26-JUL-83 02:02")
(2*(COUNT.PARTS '1234-K924)))

(DEFKBFUN COMPUTE.MP.ORDER () (* jst: "11-AUG-83 07:39")
(IF (FIND 'ONE 'CCS '(IMPLEMENTS MULTIPROCESSING)) OR
(ORDERED? '5640-F400 T)
THEN T
ELSE NIL))

(DEFKBFUN COMPUTE.MP.PIBS () (* edited: "11-AUG-83 01:11")
((COUNT.PARTS '1234-K917)+(COUNT.PARTS '1212-C003-0091)+
2*(COUNT.PARTS '1401-C058-0090)))

```

```

(DEFKBFUN COMPUTE.POWER.SUPPLIED () (* edited: "10-AUG-83
                                         07:54")
  (PROG (POWER.PART:POWER.SPEC SUM:POWER.SPEC)
    (SUM :=(SEND SUM NEW.SELF))
    (FOR POWER.SUPPLY IN '(1234-F400 1234-F413 1234-
      F396) WHEN (ORDERED? POWER.SUPPLY T)
      DO (SUM +(SAND POWER.SUPPLY 'GET.POWER)))
    (RETURN SUM:INVERTED)))

(DEFKBFUN COMPUTE.POWER.USED () (* jst: "11-AUG-83 07:42")
  (PROG (SUM:POWER.SPEC IO:NUMBER MEM:NUMBER CLC:NUMBER
    MLA:NUMBER FUDGE:NUMBER)
    (SUM :=(SEND SUM NEW.SELF))
    (MEM :=(FETCH 'MAIN-MEMORY.SUPPLIED)/ 256)
    (IO :=(FETCH '#TRUNK.LINES))
    (CLC :=(FETCH '#CLC))
    (MLA :=(FETCH '#MLA))

    (FUDGE :=(SELECTQ (FETCH 'SYSTEM.TYPE)
      (M1234.SYSTEM 180.0)
      (SHOULDNT)))
    (SUM:MINUS.5 := FUDGE + 2.0*MEM + 15.0*IO + 2.0*CLC
      + 0.0*MLA)
    (SUM:PLUS.5 := 6.0 + 1.5*MEM + 7.5*IO + 2.0*CLC +
      4.0*MLA)
    (RETURN SUM)))

(DEFKBFUN COMPUTE.TOTAL.PIBS () (* edited: "11-AUG-83
                                         05:42")
  (PROG (BASIC:NUMBER MEM:NUMBER PERF:BOOLEAN LH:BOOLEAN
    RH:BOOLEAN)
    (MEM :=(COUNT.PARTS '1234-K924 T))
    (BASIC :=(FETCH 'MP.PIBS)+(FETCH 'IO.PIBS)+(FETCH
      'ICS.PIBS))
    (PERF :=(ORDERED? '1234-K941 T))
    (LH :=(ORDERED? '1234-K926 T))
    (RH :=(ORDERED? '1234-K928 T))
    (RETURN (SELECTQ (FETCH 'SYSTEM.TYPE)
      ((M1234.SYSTEM)
        (10 + BASIC +(IF LH
          THEN 0
          ELSE MEM)))
      (SHOULDNT))))))

(DEFKBFUN COMPUTE.SYSTEM.TYPE () (* jsb: " 7-AUG-83 20:04")
  (PROG (SYSTEM:FUNCTION!CLASS)
    (IF (SYSTEM :=(FIND 'ONE 'CCS '(AND (ISA BG)
      (IMPLEMENTS SYSTEM))))
      THEN (SYSTEM :=(SANDC SYSTEM 'GET.IMPLEMENTS.FN))
    ELSEIF (SYSTEM :=(FIND 'ONE 'CCS '(ISA BU.M12)))
      THEN (SYSTEM :=(SANDC SYSTEM 'GET.IMPLEMENTS.FN))
    (IF SYSTEM IS NULL
      THEN (printout T

```

"I can't find either a basic group or a class model.
I will configure this order as a generic M1234 system."

```

      T)
      (SYSTEM := 'SYSTEM))
    (RETURN SYSTEM))

(DEFKBFUN ELIMINATE.BTP (CURRENT.CONTENTS:PRODUCT.QTYS)
                      (* jsb: "26-JAN-84 09:34")
  (PROG (CMPARTS:PRODUCT.QTYS REMOVE.PARTS:PRODUCT.QTYS
        REMAINING.PARTS:PRODUCT.QTYS
        PART:PRODUCT.QTY CMPART:PRODUCT.QTY WORKING.PARTS:
        PRODUCT.QTYS)
    (CMPARTS :=(FETCH 'CLASS.MODEL.EXPANSION T))
    (REMOVE.PARTS :=(FOR PART IN CURRENT.CONTENTS WHEN
      (CMPART :=(SEND CMPARTS MEMBER PART))
    COLLECT (A PRODUCT.QTY WITH NAME = PART:NAME QTY
      =(MAX PART:QTY
        CMPART:QTY))))
    (WORKING.PARTS :=(COPY CURRENT.CONTENTS))
    (REMAINING.PARTS :=(SEND WORKING.PARTS SUBTRACT
      REMOVE.PARTS))
    (PART :=(CAR REMAINING.PARTS))
    (IF (CDR REMAINING.PARTS) IS NULL AND PART:QTY = 1
      THEN (RETURN REMOVE.PARTS)
    ELSE (WORKING.PARTS :=(COPY CURRENT.CONTENTS))
      (RETURN (SEND WORKING.PARTS SUBTRACT (CHOOSE.ONE
        REMAINING.PARTS))))))

(DEFKBFUN INCREMENT (VALUE:INTEGER INCREMENT:INTEGER)
                    (* jsb: " 3-AUG-83 10:57")
  ((IREMAINDER VALUE INCREMENT)
   IS ZERO))

```

APPENDIX II (E)

KNOWLEDGE BASE FOR M1234 COMPUTER
EXPANSION RULES

```

(* XTEST-ECTS definition)
(DEFINSTANCE XTEST-ECTS ECTS
  SET.OF XTESTKB
  FILED.ON XTESTEX
  ELEMENTS (1234-0066-0690.ECT1 1234-3003-0690.ECT.1 1234-
    F113.ECT.1 1234-F262.E CT.1 1234-F290.ECT.1
    1234-F360.ECT.1 1234-F361.ECT1 1234-F362.ECT1
    1234-F400.ECT 1234-F413.ECT.1 1234-F504.ECT.1
    1234-F510.ECT1 1234-F545.ECT1 1234-F680.ECT.1
    1234-F690.ECT.1 1234-F690.ECT.2 1234-F691.ECT1
    1234-K400.ECT1 1234-P413.ECT 1234-P723.ECT.1
    1234-P724.ECT.1 1234-P742.ECT.1 1234-
    P743.ECT.1 1234-P954.ECT.1 1234-P955.ECT.1
    M1234-II.ECT.1))

```



```
(DEFCLASS 1234-0066-0690.ECT1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-0066-0690
TO ((3 . 1234-K913-V001)
(1 . 1234-K920-V002)
(1 . 1234-K450-V001)
(1 . 1234-K915-V001)
(1 . 1234-K917-V001)
(1 . 1234-K921-V003)
(2 . 1234-K924-V001)
(3 . 1234-K940-V001)))
)
```

```
(DEFCLASS 1234-3003-0690.ECT.1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-3003-0690
TO ((1 . 1234-F262)
(1 . 1234-F113)
(1 . 1234-K915)
(1 . 1234-K450)
(1 . 1234-K917)
(1 . 1234-K416)
(1 . 1234-K941)))
)
```

```
(DEFCLASS 1234-F113.ECT.1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F113
TO ((2 . 1234-K940)
(3 . 1234-K913)))
)
```

```
(DEFCLASS 1234-F262.ECT.1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F262
TO ((2 . 1234-K920)
(2 . 1234-K921)
(4 . 1234-K924)
(2 . 1234-K940)))
)
```

```
(DEFCLASS 1234-F290.ECT.1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F290
TO ((4 . 1234-K920)
(2 . 1234-K921)
(8 . 1234-K924)
(2 . 1234-K940)))
)
```

```
(DEFCLASS 1234-F360.ECT.1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F360
TO ((1 . 1234-K931)))
)
```

```
(DEFCLASS 1234-F361.ECT1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F361
TO ((2 . 1234-K931)))
)
```

```
(DEFCLASS 1234-F362.ECT1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F362
TO ((3 . 1234-K931)))
)
```

```
(DEFCLASS 1234-F400.ECT (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
TO ((1 . 1234-K400))
FROM 1234-F400)
)
```

```
(DEFCLASS 1234-F413.ECT.1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F413
TO ((1 . 1234-K413)))
)
```

```
(DEFCLASS 1234-F504.ECT.1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F504
TO ((1 . 1234-K905)
(1 . 1234-K901)
(1 . 1234-K903)))
)
```

```
(DEFCLASS 1234-F510.ECT1 (CTI)
()
(METACLASS ALWAYS.EXPAND.CT
FROM 1234-F510
TO ((2 . 1234-K905)
(2 . 1234-K901)
(2 . 1234-K903)))
)
```

```

(DEFCLASS 1234-F545.ECT1 (CTI)
  ()
  (METACLASS ALWAYS.EXPAND.CT
    FROM 1234-F545
    TO ((1 . 1234-K545)))
)

(DEFCLASS 1234-F680.ECT.1 (CTI)
  ()
  (METACLASS ALWAYS.EXPAND.CT
    FROM 1234-F680
    TO ((1 . 1234-K930)))
)

(DEFCLASS 1234-F690.ECT.1 (CTI)
  ()
  (METACLASS ALWAYS.EXPAND.CT
    FROM 1234-F690
    TO ((1 . 1234-K935)
        (1 . 1234-K936)
        (1 . 1234-K937)))
)

(DEFCLASS 1234-F690.ECT.2 (CTI)
  ()
  (METACLASS EXPANSION.CT
    FROM 1234-F690
    TO ((1 . 1234-K935)
        (1 . 1234-K936)
        (1 . 1234-K939))
    BIN.VARS.REQUIRED (TRACER.NUMBER)
    WHEN (TRACER.NUMBER <= 1769))
)

(DEFCLASS 1234-F691.ECT1 (CTI)
  ()
  (METACLASS ALWAYS.EXPAND.CT
    FROM 1234-F691
    TO ((2 . 1234-K935)
        (2 . 1234-K936)
        (2 . 1234-K937)))
)

(DEFCLASS 1234-K400.ECT1 (CTI)
  ()
  (METACLASS ALWAYS.EXPAND.CT
    TO ((1 . 1234-K400-V001))
    FROM 1234-K400)
)

```

```
(DEFCLASS 1234-P413.ECT (CTI)
()
(METACCLASS ALWAYS.EXPAND.CT
FROM 1234-P413
TO ((1 . 1234-F413)))
)
```

```
(DEFCLASS 1234-P723.ECT.1 (CTI)
()
(METACCLASS ALWAYS.EXPAND.CT
FROM 1234-P723
TO ((1 . 1234-F290)))
)
```

```
(DEFCLASS 1234-P724.ECT.1 (CTI)
()
(METACCLASS ALWAYS.EXPAND.CT
FROM 1234-P724
TO ((1 . 1234-F295)))
)
```

```
(DEFCLASS 1234-P742.ECT.1 (CTI)
()
(METACCLASS ALWAYS.EXPAND.CT
FROM 1234-P742
TO ((1 . 1234-F360)
(1 . 1234-F680)))
)
```

```
(DEFCLASS 1234-P743.ECT.1 (CTI)
()
(METACCLASS ALWAYS.EXPAND.CT
FROM 1234-P743
TO ((1 . 1234-F360)
(1 . 1234-F690)))
)
```

```
(DEFCLASS 1234-P954.ECT.1 (CTI)
()
(METACCLASS ALWAYS.EXPAND.CT
FROM 1234-P954
TO ((1 . 1234-F504)))
)
```

```
(DEFCLASS 1234-P955.ECT.1 (CTI)
()
(METACCLASS ALWAYS.EXPAND.CT
FROM 1234-P955
TO ((1 . 1234-F504)))
)
```

```
(DEFCLASS M1234-II.ECT.1 (CTI)
  ()
  (METACCLASS ALWAYS.EXPAND.CT
   FROMM1234-II
   TO ((1 . 1234-3003-0690)))
)
```

APPENDIX III (A)

MINICOMPUTER ORDER CHECKING TYPESCRIPT

(MINI) COAST executive> RUN

Evaluating: (INPUT.ORDER)

Processing new order...

Enter name of order file: T

Reading order information...

Order information>

Reading parts to check...

Order line> 1 MINI-1000

Order line> 4 RAM-1000

Order line> 3 ROM-1000

Order line> 4 TERMINAL-1000

Order line> 3 DISC-1000

Order line>

Reading parts to be shipped F/S...

F/S order line>

Enter order processing options: SHOW-FULL-TRACE

Adding the following parts: 1 MINI-1000, 4 RAM-1000, 3
ROM-1000, 4 TERMINAL-1000 and 3 DISC-1000

Evaluating: (DETERMINE '(TODAYS.DATE ORDERED.PARTS
ORDERED.PRODUCT.LINE))

Today's date of original components is set to: 20-Jun-84

The ordered set of parts of original components is set
to: MINI-1000, RAM-1000, ROM-1000, TERMINAL-1000 and
DISC-1000

The ordered product line of original components is set
to: minicomputer components

Evaluating: (IF CHECKING.ORDER

THEN **COMMENT**

(CONFIGURE (FETCH

'ORDERED.PRODUCT.LINE)

(FIND 'ALL 'CCS))

(CHECK 'EMPTY.BIN))

Configuring minicomputer components

Initial parts are: 1 MINI-1000, 4 RAM-1000, 3 ROM-1000,
4 TERMINAL-1000 and 3 DISC-1000

Evaluating: [CONFIGURE 'MAIN-FRAME (FIND 'ALL 'CCS
'(IMPLEMENTS MAIN-FRAME)]

Configuring mainframe components

Initial parts are: 4 RAM-1000, 3 ROM-1000, 4
TERMINAL-1000 and 3 DISC-1000

Evaluating: [CONFIGURE 'CPU (FIND 'ALL 'CCS
'(IMPLEMENTS CPU)]

Configuring cpu components

No initial parts

Evaluating: (CHECK 'REQUIREMENTS)

Checking REQUIREMENTS constraints in cpu
components

The following constraints apply: AT-MOST-
ONE.CT NO-PARTS.CT

Checking constraint: AT-MOST-ONE.CT (a MODIFY-
ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
> 1)

failed with value NIL when applied to cpu
components

Checking constraint: NO-PARTS.CT (a MODIFY-ON-
SUCCESS.CT)

The statement: ((COUNT CONTENTS)
IS ZERO)

succeeded with value T when applied to cpu
components

Adding required component to CPU.1

Adding 1 CPU-1000 to cpu components ; OK? YES

Adding the following parts: 1 CPU-1000

Evaluating: [CONFIGURE 'MEMORY (FIND 'ALL 'CCS
'(IMPLEMENTS MEMORY)]

Configuring memory components

Initial parts are: 4 RAM-1000 and 3 ROM-1000

Evaluating: [CONFIGURE 'RAM (FIND 'ALL 'CCS
'(IMPLEMENTS RAM)]

Configuring random access memory components

Initial parts are: 4 RAM-1000

Evaluating: (CHECK 'REQUIREMENTS)

Checking REQUIREMENTS constraints in random
access memory components

The following constraints apply: NO-
PARTS.CT

Checking constraint: NO-PARTS.CT (a MODIFY-ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
IS ZERO)

failed with value NIL when applied to
random access memory components

Evaluating: [CONFIGURE 'ROM (FIND 'ALL 'CCS
'(IMPLEMENTS ROM]

Configuring read-only memory components
Initial parts are: 3 ROM-1000

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in read-
only memory components

The following constraints apply: NO-
PARTS.CT

Checking constraint: NO-PARTS.CT (a MODIFY-
ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
IS ZERO)

failed with value NIL when applied to read-
only memory components

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in memory
components

The following constraints apply: MEMORY.CT1

Checking constraint: MEMORY.CT1 (a WARN-ON-
SUCCESS.CT)

The number of ram boards ordered of minicomputer
components is set to: 4

The number of rom boards ordered of minicomputer
components is set to: 3

The statement: (#RAM + #ROM > 7)
failed with value NIL when applied to memory
components

Evaluating: [CONFIGURE 'IO (FIND 'ALL 'CCS
'(IMPLEMENTS IO]

Configuring I/O components

Initial parts are: 4 TERMINAL-1000 and 3 DISC-
1000

Evaluating: [CONFIGURE 'TERMINAL (FIND 'ALL 'CCS
'(IMPLEMENTS TERMINAL]

Configuring terminal components

Initial parts are: 4 TERMINAL-1000

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in terminal components

The following constraints apply: NO-PARTS.CT MAX.TERMINALS.CT

Checking constraint: NO-PARTS.CT (a MODIFY-ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
IS ZERO)

failed with value NIL when applied to terminal components

Checking constraint: MAX.TERMINALS.CT (a
MODIFY-ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
> 4)

failed with value NIL when applied to terminal components

Evaluating: [CONFIGURE 'DISC (FIND 'ALL 'CCS
'(IMPLEMENTS DISC]

Configuring disk drive components

Initial parts are: 3 DISC-1000

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in disk drive components

No REQUIREMENTS constraints were found.

Evaluating: (DETERMINE '(IF-CARDS.NEEDED))
The number of discs ordered of I/O components is set to: 3

The number of interface cards required for this order of I/O components is set to: 4

Evaluating: [CONFIGURE 'INTERFACE (FIND 'ALL
'CCS '(IMPLEMENTS INTERFACE]

Configuring interface components

No initial parts

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in interface components

The following constraints apply: IF.REQ.CT1

Checking constraint: IF.REQ.CT1 (a MODIFY-ON-SUCCESS.CT)

The statement: (IF-CARDS.NEEDED = CONTENTS)

succeeded with value T when applied to interface components

You have not ordered the minimum number of interface cards.
Replacing contents of interface components with 4 IF-1000
; OK? YES

Removing the following parts:

Adding the following parts: 4 IF-1000

Evaluating: [CONFIGURE 'CARD-CAGE (FIND 'ALL 'CCS '(IMPLEMENTS CARD-CAGE]

Configuring card cage components

No initial parts

Evaluating: (CHECK 'REQUIREMENTS)

Checking REQUIREMENTS constraints in card cage components

The following constraints apply: AT-MOST-ONE.CT NO-PARTS.CT

Checking constraint: AT-MOST-ONE.CT (a MODIFY-ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)

> 1)

failed with value NIL when applied to card cage components

Checking constraint: NO-PARTS.CT (a MODIFY-ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)

IS ZERO)

succeeded with value T when applied to card cage components

Adding required component to CARD-CAGE.1

Adding 1 CAGE-1000 to card cage components ; OK? YES

Adding the following parts: 1 CAGE-1000

Evaluating: (CHECK 'REQUIREMENTS)

Checking REQUIREMENTS constraints in mainframe components

The following constraints apply: MAINFRAME.CT.1

Checking constraint: MAINFRAME.CT.1 (a WARN-ON-SUCCESS.CT)

The number of interface cards of minicomputer components is set to: 3

The statement: (#ROM + #RAM + #IF-CARDS + 1 > 16)

failed with value NIL when applied to mainframe components

Evaluating: [CONFIGURE 'POWER-SUPPLY (FIND 'ALL 'CCS
'(IMPLEMENTS POWER-SUPPLY]

Configuring power supply components
No initial parts

Evaluating: (CHECK 'CONTENTS)
Checking CONTENTS constraints in power supply
components

The following constraints apply: PS.AMBIGUOUS.CT
Checking constraint: PS.AMBIGUOUS.CT (a MODIFY-ON-
SUCCESS.CT)

The statement: ((COUNT CONTENTS)
> 1)

failed with value NIL when applied to power
supply components

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in power supply
components

The following constraints apply: PS.REQ.CT1
PS.REQ.CT2 PS.REQ.CT3
Checking constraint: PS.REQ.CT1 (a MODIFY-ON-
SUCCESS.CT)

The power required to operate this order of
minicomputer components is set to: 8.3

The statement: (POWER.NEEDED <= 5 AND (NOT
(ORDERED? 'SUPPLY-05 T))
failed with value NIL when applied to power
supply components

Checking constraint: PS.REQ.CT2 (a MODIFY-ON-
SUCCESS.CT)

The statement: [POWER.NEEDED > 5 AND POWER.NEEDED
<= 10 AND (NOT (ORDERED? 'SUPPLY-10)]
succeeded with value T when applied to power
supply components

Adding 1 SUPPLY-10 to power supply components ; OK? YES
Adding the following parts: 1 SUPPLY-10
Checking constraint: PS.REQ.CT3 (a MODIFY-ON-
SUCCESS.CT)

The statement: [POWER.NEEDED > 10 AND
POWER.NEEDED <= 15 AND (NOT
(ORDERED? 'SUPPLY-15]

failed with value NIL when applied to power
supply components
Checking EMPTY.BIN constraints in power supply
components

The following constraints apply:
UNASSIGNED.EMPTY.CT
Checking constraint: UNASSIGNED.EMPTY.CT (a WARN-
ON-SUCCESS.CT)

The statement: (BIN:CONTENTS IS NOT NULL)
failed with value NIL when applied to original
components

Evaluating: (DETERMINE '(FINAL.PARTS))

The final set of parts of original components is set to:
MINI-1000, CPU-1000, RAM-1000, ROM-1000, TERMINAL-1000,
DISC-1000, IF-1000, CAGE-1000 and SUPPLY-10

Evaluating: (OUTPUT.ORDER)

What type of output form would you like: FINAL.PARTS

```
1 CPU-1000
4 RAM-1000
3 ROM-1000
4 TERMINAL-1000
3 DISC-1000
4 IF-1000
1 CAGE-1000
1 SUPPLY-10
```

What type of output form would you like:

Do you want to use this order to reconfigure something in
finished goods?

Would you like an explanation of this order?

APPENDIX III (B)

MINICOMPUTER DESIGN TYPESCRIPT

(MINI) COAST executive> RUN

Evaluating: (INPUT.ORDER)

Processing new order...

Enter name of order file: T

Reading order information...

Order information>

Reading parts to check...

Order line> 1 MINI-1000

Order line>

Reading parts to be shipped F/S...

F/S order line>

Enter order processing options: SHOW-FULL-TRACE

Adding the following parts: 1-MINI-1000

Evaluating: (DETERMINE '(TODAYS.DATE ORDERED.PARTS
ORDERED.PRODUCT.LINE))

Today's date of original components is set to: 20-Jun-84

The ordered set of parts of original components is set
to: MINI-1000

The ordered product line of original components is set
to: minicomputer components

```

Evaluating: (IF CHECKING.ORDER
              THEN  **COMMENT**
                  (CONFIGURE (FETCH
                              'ORDERED.PRODUCT.LINE)
                          (FIND 'ALL 'CCS))
                  (CHECK 'EMPTY.BIN))

```

Configuring minicomputer components
Initial parts are: 1 MINI-1000

```

Evaluating: [CONFIGURE 'MAIN-FRAME (FIND 'ALL 'CCS
              '(IMPLEMENTS MAIN-FRAME)]

```

Configuring mainframe components
No initial parts

```

Evaluating: [CONFIGURE 'CPU (FIND 'ALL 'CCS
              '(IMPLEMENTS CPU)]

```

Configuring cpu components
No initial parts

```

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in cpu
components
The following constraints apply: AT-MOST-
ONE.CT NO-PARTS.CT
Checking constraint: AT-MOST-ONE.CT (a MODIFY-
ON-SUCCESS.CT)

```

```

The statement: ((COUNT CONTENTS)
                > 1)
failed with value NIL when applied to cpu
components
Checking constraint: NO-PARTS.CT (a MODIFY-ON-
SUCCESS.CT)

```

```

The statement: ((COUNT CONTENTS)
                IS ZERO)
succeeded with value T when applied to cpu
components

```

Adding required component to CPU.1
Adding 1 CPU-1000 to cpu components ; OK? YES
Adding the following parts: 1 CPU-1000

```

Evaluating: [CONFIGURE 'MEMORY (FIND 'ALL 'CCS
              '(IMPLEMENTS MEMORY)]

```

Configuring memory components
No initial parts

```

Evaluating: [CONFIGURE 'RAM (FIND 'ALL 'CCS
              '(IMPLEMENTS RAM)]

```

Configuring random access memory components
No initial parts

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in random
access memory components
The following constraints apply: NO-PARTS.CT
Checking constraint: NO-PARTS.CT (a MODIFY-
ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
IS ZERO)
succeeded with value T when applied to
random access memory components

Adding required component to RAM.1
Adding 1 RAM-1000 to random access memory components ; OK?
YES

Adding the following parts: 1 RAM-1000

Evaluating: [CONFIGURE 'ROM (FIND 'ALL 'CCS
'(IMPLEMENTS ROM]

Configuring read-only memory components
No initial parts

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in read-
only memory components
The following constraints apply: NO-
PARTS.CT
Checking constraint: NO-PARTS.CT (a MODIFY-
ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
IS ZERO)
succeeded with value T when applied to
read-only memory components

Adding required component to ROM.1
Adding 1 ROM-1000 to read-only memory components ; OK? YES
Adding the following parts: 1 ROM-1000

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in memory
components
The following constraints apply: MEMORY.CT1
Checking constraint: MEMORY.CT1 (a WARN-ON-
SUCCESS.CT)

The number of ram boards ordered of minicomputer
components is set to: 1
The number of rom boards ordered of minicomputer
components is set to: 1

The statement: (#RAM + #ROM > 7)
failed with value NIL when applied to memory
components

Evaluating: [CONFIGURE 'IO (FIND 'ALL 'CCS
'(IMPLEMENTS IO)]

Configuring I/O components
No initial parts

Evaluating: [CONFIGURE 'TERMINAL (FIND 'ALL 'CCS
'(IMPLEMENTS TERMINAL)]

Configuring terminal components
No initial parts

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in terminal
components

The following constraints apply: NO-
PARTS.CT MAX.TERMINALS.CT

Checking constraint: NO-PARTS.CT (a MODIFY-
ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
IS ZERO)

succeeded with value T when applied to
terminal components

Adding required component to TERMINAL.1

Adding 1 TERMINAL-1000 to terminal components ; OK? YES

Adding the following parts: 1 TERMINAL-1000

Checking constraint: MAX.TERMINALS.CT (a
MODIFY-ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS)
> 4)

failed with value NIL when applied to
terminal components

Evaluating: [CONFIGURE 'DISC (FIND 'ALL 'CCS
'(IMPLEMENTS DISC)]

Configuring disk drive components
No initial parts

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in disk
drive components

No REQUIREMENTS constraints were found.

Evaluating: (DETERMINE '(IF-CARDS.NEEDED))

The number of discs ordered of I/O components is
set to: 0

The number of interface cards required for this order of I/O components is set to: 1

Evaluating: [CONFIGURE 'INTERFACE (FIND 'ALL 'CCS '(IMPLEMENTS INTERFACE]

Configuring interface components
No initial parts

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in interface components
The following constraints apply: IF.REQ.CT1
Checking constraint: IF.REQ.CT1 (a MODIFY-ON-SUCCESS.CT)

The statement: (IF-CARDS.NEEDED = CONTENTS)
succeeded with value T when applied to interface components

You have not ordered the minimum number of interface cards.
Replacing contents of interface components with 1 IF-1000
; OK? YES

Removing the following parts:
Adding the following parts: 1 IF-1000

Evaluating: [CONFIGURE 'CARD-CAGE (FIND 'ALL 'CCS '(IMPLEMENTS CARD-CAGE]

Configuring card cage components
No initial parts

Evaluating: (CHECK 'REQUIREMENTS)
Checking REQUIREMENTS constraints in card cage components
The following constraints apply: AT-MOST-ONE.CT NO-PARTS.CT
Checking constraint: AT-MOST-ONE.CT (a MODIFY-ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS) > 1)
failed with value NIL when applied to card cage components
Checking constraint: NO-PARTS.CT (a MODIFY-ON-SUCCESS.CT)

The statement: ((COUNT CONTENTS) IS ZERO)
succeeded with value T when applied to card cage components

Adding required component to CARD-CAGE.1
Adding 1 CAGE-1000 to card cage components ; OK? YES
Adding the following parts: 1 CAGE-1000

Evaluating: (CHECK 'REQUIREMENTS)
 Checking REQUIREMENTS constraints in mainframe
 components

The following constraints apply: MAINFRAME.CT.1
 Checking constraint: MAINFRAME.CT.1 (a WARN-ON-
 SUCCESS.CT)

The number of interface cards of minicomputer
 components is set to: 1

The statement: (#ROM + #RAM + #IF-CARDS + 1 > 16)
 failed with value NIL when applied to mainframe
 components

Evaluating: [CONFIGURE 'POWER-SUPPLY (FIND 'ALL 'CCS
 '(IMPLEMENTS POWER-SUPPLY]

Configuring power supply components
 No initial parts

Evaluating: (CHECK 'CONTENTS)
 Checking CONTENTS constraints in power supply
 components

The following constraints apply: PS.AMBIGUOUS.CT
 Checking constraint: PS.AMBIGUOUS.CT (a MODIFY-ON-
 SUCCESS.CT)

The statement: ((COUNT CONTENTS)
 > 1)

failed with value NIL when applied to power
 supply components

Evaluating: (CHECK 'REQUIREMENTS)
 Checking REQUIREMENTS constraints in power supply
 components

The following constraints apply: PS.REQ.CT1
 PS.REQ.CT2 PS.REQ.CT3

Checking constraint: PS.REQ.CT1 (a MODIFY-ON-
 SUCCESS.CT)

The power required to operate this order of
 minicomputer components is set to: 3.7

The statement: (POWER.NEEDED <= 5 AND (NOT
 (ORDERED? 'SUPPLY-05 T)))
 succeeded with value T when applied to power
 supply components

Adding 1 SUPPLY-05 to power supply components ; OK? YES

Adding the following parts: 1 SUPPLY-05

Checking constraint: PS.REQ.CT2 (a MODIFY-ON-
 SUCCESS.CT)

The statement: [POWER.NEEDED > 5 AND POWER.NEEDED
 <= 10 AND (NOT (ORDERED? 'SUPPLY-10)]

failed with value NIL when applied to power
 supply components

Checking constraint: PS.REQ.CT3 (a MODIFY-ON-
 SUCCESS.CT)

The statement: [POWER.NEEDED > 10 AND
POWER.NEEDED <= 15 AND (NOT (ORDERED?
'SUPPLY-15]

failed with value NIL when applied to power
supply components

Checking EMPTY.BIN constraints in power supply
components

The following constraints apply:
UNASSIGNED.EMPTY.CT

Checking constraint: UNASSIGNED.EMPTY.CT (a WARN-
ON-SUCCESS.CT)

The statement: (BIN:CONTENTS IS NOT NULL)
failed with value NIL when applied to original
components

Evaluating: (DETERMINE '(FINAL.PARTS))

The final set of parts of original components is set to:
MINI-1000, CPU-1000, RAM-1000, ROM-1000, TERMINAL-1000, IF-
1000, CAGE-1000 and SUPPLY-05

Evaluating: (OUTPUT.ORDER)

What type of output form would you like: FINAL.PARTS

```
1 CPU-1000
1 RAM-1000
1 ROM-1000
1 TERMINAL-1000
1 IF-1000
1 CAGE-1000
1 SUPPLY-05
```

What type of output form would you like:

Do you want to use this order to reconfigure something in
finished goods?

Would you like an explanation of this order?

APPENDIX IV (A)

KNOWLEDGE BASE FOR MINICOMPUTER FUNCTIONAL HIERARCHY

```
(* MINI-FHS definition)
(DEFINSTANCE MINI-FHS FHS
  SET.OF MINIKB
  FILED.ON MINIFH
  ELEMENTS (CARD-CAGE CPU DISC INTERFACE IO MAIN-FRAME
  MEMORY MINICOMPUTER POWER-SUPPLY RAM ROM TERMINAL))
(DEFCLASS CARD-CAGE (BIN)
  ()
  (METACLASS FUNCTION.CLASS
  COMPOSES (MAIN-FRAME)
  DESCR "card cage")
  )
```

```
(DEFCLASS CPU (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   COMPOSES (MAIN-FRAME)
   DESCR "cpu")
)

(DEFCLASS DISC (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   COMPOSES (IO)
   DESCR "disk drive")
)

(DEFCLASS INTERFACE (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   COMPOSES (IO)
   DESCR "interface")
)

(DEFCLASS IO (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   COMPOSES (MAIN-FRAME)
   DESCR "I/O")
)

(DEFCLASS MAIN-FRAME (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   COMPOSES (MINICOMPUTER)
   DESCR "mainframe")
)

(DEFCLASS MEMORY (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   COMPOSES (MAIN-FRAME)
   DESCR "memory")
)

(DEFCLASS MINICOMPUTER (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   PROD.GEN (PRODUCT.LINE)
   DESCR "minicomputer")
)

(DEFCLASS POWER-SUPPLY (BIN)
  ()
  (METACLASS FUNCTION.CLASS
   COMPOSES (MINICOMPUTER)
   DESCR "power supply")
)
```

```

(DEFCLASS RAM (BIN)
  (
  (METACLASS FUNCTION.CLASS
    COMPOSES (MEMORY)
    DESCR "random access memory")
  )
)

(DEFCLASS ROM (BIN)
  (
  (METACLASS FUNCTION.CLASS
    COMPOSES (MEMORY)
    DESCR "read-only memory")
  )
)

(DEFCLASS TERMINAL (BIN)
  (
  (METACLASS FUNCTION.CLASS
    COMPOSES (IO)
    DESCR "terminal")
  )
)

```

APPENDIX IV (B)

KNOWLEDGE BASE FOR MINICOMPUTER
TASK BLOCKS

```

(* MINI-TBS definition)
(DEFINSTANCE MINI-TBS TBS
  SET.OF MINIKB
  FILED.ON MINIFH
  ELEMENTS (CARD-CAGE.TB CPU.TB DISC.TB INTERFACE.TB IO.TB
  MAIN-FRAME.TB MEMORY.TB
            MINICOMPUTER.TB POWER-SUPPLY.TB RAM.TB ROM.TB
            TERMINAL.TB))

(DEFTASKBLOCK CARD-CAGE.TB CARD-CAGE () (CHECK
  'REQUIREMENTS)
)

(DEFTASKBLOCK CPU.TB CPU () (CHECK 'REQUIREMENTS)
)

(DEFTASKBLOCK DISC.TB DISC () (CHECK 'REQUIREMENTS)
)

(DEFTASKBLOCK INTERFACE.TB INTERFACE () (CHECK
  'REQUIREMENTS)
)

(DEFTASKBLOCK IO.TB IO ()
)

```

(* The order of the CONFIGURE statements is
significant below.

In particular, INTERFACE should be configured last!)

```

(CONFIGURE 'TERMINAL (FIND 'ALL 'CCS '(IMPLEMENTS
  TERMINAL)))
(CONFIGURE 'DISC (FIND 'ALL 'CCS '(IMPLEMENTS DISC)))
(DETERMINE '(IF-CARDS.NEEDED))
(CONFIGURE 'INTERFACE (FIND 'ALL 'CCS '(IMPLEMENTS
  INTERFACE)))

(DEFTASKBLOCK MAIN-FRAME.TB MAIN-FRAME () (CONFIGURE 'CPU
(FIND 'ALL 'CCS '(IMPLEMENTS CPU)))
  (CONFIGURE 'MEMORY (FIND 'ALL 'CCS '(IMPLEMENTS MEMORY)))
  (CONFIGURE 'IO (FIND 'ALL 'CCS '(IMPLEMENTS IO)))
  (CONFIGURE 'CARD-CAGE (FIND 'ALL 'CCS '(IMPLEMENTS CARD-
    CAGE)))
  (CHECK 'REQUIREMENTS))

(DEFTASKBLOCK MEMORY.TB MEMORY () (CONFIGURE 'RAM (FIND 'ALL
  'CCS '(IMPLEMENTS RAM)))
  (CONFIGURE 'ROM (FIND 'ALL 'CCS '(IMPLEMENTS ROM)))
  (CHECK 'REQUIREMENTS))

(DEFTASKBLOCK MINICOMPUTER.TB MINICOMPUTER () (CONFIGURE
  'MAIN-FRAME
  (FIND 'ALL 'CCS '(IMPLEMENTS
    MAIN-FRAME)))
  (CONFIGURE 'POWER-SUPPLY (FIND 'ALL 'CCS '(IMPLEMENTS
    POWER-SUPPLY))))

(DEFTASKBLOCK POWER-SUPPLY.TB POWER-SUPPLY () (CHECK
  'CONTENTS)
  (CHECK 'REQUIREMENTS))

(DEFTASKBLOCK RAM.TB RAM () (CHECK 'REQUIREMENTS)
)

(DEFTASKBLOCK ROM.TB ROM () (CHECK 'REQUIREMENTS)
)

(DEFTASKBLOCK TERMINAL.TB TERMINAL () (CHECK 'REQUIREMENTS)
)

```

APPENDIX IV (C)

 KNOWLEDGE BASE FOR MINICOMPUTER
 PARTS CATALOG

```

(* MINI-PRODS definition)
(DEFINANCE MINI-PRODS PRODS
  SET.OF MINIKB
  FILED.ON MINIPROD
  ELEMENTS (CAGE-1000 CPU-1000 DISC-1000 IF-1000 MINI-1000
  RAM-1000 ROM-1000 SUPPLY-05 SUPPLY-10 SUPPLY-15 TERMINAL-
  1000))

```

```
(DEFCLASS CAGE-1000 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN CARD-CAGE)
  )
```

```
(DEFCLASS CPU-1000 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN CPU)
  )
```

```
(DEFCLASS DISC-1000 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN DISC)
  )
```

```
(DEFCLASS IF-1000 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN INTERFACE)
  )
```

```
(DEFCLASS MINI-1000 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN MINICOMPUTER)
  )
```

```
(DEFCLASS RAM-1000 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN RAM)
  )
```

```
(DEFCLASS ROM-1000 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN ROM)
  )
```

```
(DEFCLASS SUPPLY-05 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN POWER-SUPPLY)
  )
```

```
(DEFCLASS SUPPLY-10 (PRODUCT.ID)
  ()
  (METACCLASS PID
    IMPLEMENTS.FN POWER-SUPPLY)
  )
```

```

(DEFCLASS SUPPLY-15 (PRODUCT.ID)
  ()
  (METACLASS PID
    IMPLEMENTS.FN POWER-SUPPLY)
  )

(DEFCLASS TERMINAL-1000 (PRODUCT.ID)
  ()
  (METACLASS PID
    IMPLEMENTS.FN TERMINAL)
  )

(* MINI-ECTS definition)
(DEFINSTANCE MINI-ECTS ECTS
  SET.OF MINIKB
  FILED.ON MINIPROD
  ELEMENTS NIL)

```

APPENDIX IV (D)

KNOWLEDGE BASE FOR MINICOMPUTER
CONSTRAINTS & BIN VARIABLES
KNOWLEDGE BASE FUNCTIONS

```

(* MINI-CTS definition)
(DEFINSTANCE MINI-CTS CTS
  SET.OF MINIKB
  FILED.ON MINICT
  ELEMENTS (AT-MOST-ONE.CT IF.REQ.CT1 MAINFRAME.CT.1
    MAX.TERMINALS.CT MEMORY.CT1 NO-PARTS.CT PS.AMBIGUOUS.CT
    PS.REQ.CT1 PS.REQ.CT2 PS.REQ.CT3))

(DEFCLASS AT-MOST-ONE.CT (CTI)
  ()
  (METACLASS MODIFY-ON-SUCCESS.CT
    SET-PARTS (LIST (A PRODUCT.QTY WITH QTY = 1 NAME =
  (GETCLASS (CAR (CURRENT.BIN:CONTENTS))))))
  STATEMENT ((COUNT CONTENTS)
    > 1)
  BIN.TYPES (CPU CARD-CAGE)
  CHECK.KEYWORDS (REQUIREMENTS))
  )

(DEFCLASS IF.REQ.CT1 (CTI)
  ()
  (METACLASS MODIFY-ON-SUCCESS.CT
    SET-PARTS (LIST (A PRODUCT.QTY WITH QTY = (FETCH 'IF-
    CARDS.NEEDED)
    NAME = 'IF-1000))
  MESSAGE ("You have not ordered the minimum number of
  interface cards.")

```

```

BIN.VARS.REQUIRED (IF-CARDS.NEEDED)
STATEMENT (IF-CARDS.NEEDED = CONTENTS)
BIN.TYPES (INTERFACE)
CHECK.KEYWORDS (REQUIREMENTS)
)

(DEFCLASS MAINFRAME.CT.1 (CTI)
()
(METACLASS WARN-ON-SUCCESS.CT
MESSAGE ("You have ordered more than 16 cards for this
         mainframe.")
BIN.VARS.REQUIRED (#ROM #RAM #IF-CARDS)
STATEMENT (#ROM + #RAM + #IF-CARDS + 1 > 16)
BIN.TYPES (MAIN-FRAME)
CHECK.KEYWORDS (REQUIREMENTS))
)

(DEFCLASS MAX.TERMINALS.CT (CTI)
()
(METACLASS MODIFY-ON-SUCCESS.CT
SET-PARTS (LIST (A PRODUCT.QTY WITH QTY = 4 NAME =
'TERMINAL-1000))
MESSAGE ("You have ordered too many terminals; the
         maximum is 4.")
STATEMENT ((COUNT CONTENTS)
          > 4)
BIN.TYPES (TERMINAL)
CHECK.KEYWORDS (REQUIREMENTS))
)

(DEFCLASS MEMORY.CT1 (CTI)
()
(METACLASS WARN-ON-SUCCESS.CT
MESSAGE ("You have ordered too many memory (RAM and ROM)
         cards.")
BIN.VARS.REQUIRED (#RAM #ROM)
STATEMENT (#RAM + #ROM > 7)
BIN.TYPES (MEMORY)
CHECK.KEYWORDS (REQUIREMENTS))
)

(DEFCLASS NO-PARTS.CT (CTI)
()
(METACLASS MODIFY-ON-SUCCESS.CT
ADD-PARTS (LIST (A PRODUCT.QTY WITH QTY = 1 NAME =
                (SELECTQ (GETCLASS CURRENT.BIN)
                (RAM 'RAM-1000)
                (ROM 'ROM-1000)
                (TERMINAL 'TERMINAL-1000)
                (INTERFACE 'IF-1000)
                (CARD-CAGE 'CAGE-1000)
                (CPU 'CPU-1000)
                (SHOULDNT))))
MESSAGE ("Adding required component to " CURRENT.BIN)

```

```

STATEMENT ((COUNT CONTENTS)
           IS ZERO)
BIN.TYPES (RAM ROM TERMINAL CPU CARD-CAGE)
CHECK.KEYWORDS (REQUIREMENTS))
)

(DEFCLASS PS.AMBIGUOUS.CT (CTI)
  ()
  (METACLASS MODIFY-ON-SUCCESS.CT
    SET-PARTS (MAKE.PRODUCT.QTYS CURRENT.BIN:CONTENTS)
    STATEMENT ((COUNT CONTENTS)
               > 1)
    BIN.TYPES (POWER-SUPPLY)
    CHECK.KEYWORDS (CONTENTS))
  )

(DEFCLASS PS.REQ.CT1 (CTI)
  ()
  (METACLASS MODIFY-ON-SUCCESS.CT
    ADD-PARTS (LIST (A PRODUCT.QTY WITH QTY = 1 NAME =
                    'SUPPLY-05))
    BIN.VARS.REQUIRED (POWER.NEEDED)
    STATEMENT (POWER.NEEDED <= 5 AND (NOT (ORDERED? 'SUPPLY-
              05 T)))
    BIN.TYPES (POWER-SUPPLY)
    CHECK.KEYWORDS (REQUIREMENTS))
  )

(DEFCLASS PS.REQ.CT2 (CTI)
  ()
  (METACLASS MODIFY-ON-SUCCESS.CT
    ADD-PARTS (LIST (A PRODUCT.QTY WITH QTY = 1 NAME =
                    'SUPPLY-10))
    BIN.VARS.REQUIRED (POWER.NEEDED)
    STATEMENT (POWER.NEEDED > 5 AND POWER.NEEDED <= 10 AND
              (NOT (ORDERED? 'SUPPLY-10)))
    BIN.TYPES (POWER-SUPPLY)
    CHECK.KEYWORDS (REQUIREMENTS))
  )

(DEFCLASS PS.REQ.CT3 (CTI)
  ()
  (METACLASS MODIFY-ON-SUCCESS.CT
    ADD-PARTS (LIST (A PRODUCT.QTY WITH QTY = 1 NAME =
                    'SUPPLY-15))
    BIN.VARS.REQUIRED (POWER.NEEDED)
    STATEMENT (POWER.NEEDED > 10 AND POWER.NEEDED <= 15 AND
              (NOT (ORDERED? 'SUPPLY-15)))
    BIN.TYPES (POWER-SUPPLY)
    CHECK.KEYWORDS (REQUIREMENTS))
  )

```



```

(* MINI-BVS definition)
(DEFINSTANCE MINI-BVS BVS
  SET.OF MINIKB
  FILED.ON MINICT
  ELEMENTS (#DISCS #IF-CARDS #RAM #ROM #TERMINALS IF-
    CARDS.NEEDED POWER.NEEDED))

(DEFCLASS #DISCS (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (IO)
    HOW.TO.DETERMINE (COMPUTE.#DISC ASK.QUESTION)
    LEGAL.VALUES INTEGER
    DESCR "the number of discs ordered")
  )

(DEFCLASS #IF-CARDS (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (MINICOMPUTER)
    HOW.TO.DETERMINE (COMPUTE.#IF-CARDS ASK.QUESTION)
    LEGAL.VALUES INTEGER
    DESCR "the number of interface cards")
  )

(DEFCLASS #RAM (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (MINICOMPUTER)
    HOW.TO.DETERMINE (COMPUTE.#RAM ASK.QUESTION)
    LEGAL.VALUES INTEGER
    DESCR "the number of RAM boards ordered")
  )

(DEFCLASS #ROM (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (MINICOMPUTER)
    HOW.TO.DETERMINE (COMPUTE.#ROM ASK.QUESTION)
    LEGAL.VALUES INTEGER
    DESCR "the number of ROM boards ordered")
  )

(DEFCLASS #TERMINALS (BVI)
  ()
  (METACLASS BIN.VAR
    BIN.TYPES (IO)
    HOW.TO.DETERMINE (COMPUTE.#TERMINALS ASK.QUESTION)
    LEGAL.VALUES INTEGER
    DESCR "the number of terminals ordered")
  )

```

```

(DEFCLASS IF-CARDS.NEEDED (BVI)
  ()
  (METACCLASS BIN.VAR
    BIN.TYPES (IO)
    HOW.TO.DETERMINE (COMPUTE.IF-CARDS.NEEDED ASK.QUESTION)
    LEGAL.VALUES INTEGER
    DESCR "the number of interface cards required for this
           order")
  )

(DEFCLASS POWER.NEEDED (BVI)
  ()
  (METACCLASS BIN.VAR
    BIN.TYPES (MINICOMPUTER)
    HOW.TO.DETERMINE (COMPUTE.POWER.NEEDED ASK.QUESTION)
    LEGAL.VALUES NUMBER
    DESCR "the power required to operate this order")
  )

(* MINI-BVMS definition)
(DEFINSTANCE MINI-BVMS BVMS
  SET.OF MINIKB
  FILED.ON MINICT
  ELEMENTS (COMPUTE.#DISC COMPUTE.#IF-CARDS COMPUTE.#RAM
    COMPUTE.#ROM COMPUTE.#TERMINALS
    COMPUTE.IF-CARDS.NEEDED COMPUTE.POWER.NEEDED))

(DEFKBFUN COMPUTE.#DISC ()(COUNT.PARTS 'DISC-1000 T)
  )

(DEFKBFUN COMPUTE.#IF-CARDS ()(COUNT.PARTS 'IF-1000 T)
  )

(DEFKBFUN COMPUTE.#RAM ()(COUNT.PARTS 'RAM-1000 T)
  )

(DEFKBFUN COMPUTE.#ROM ()(COUNT.PARTS 'ROM-1000 T)
  )

(DEFKBFUN COMPUTE.#TERMINALS ()(COUNT.PARTS 'TERMINAL-1000
  T)
  )

(DEFKBFUN COMPUTE.IF-CARDS.NEEDED ()

  (* Compute the minimum number of interface cards
    required by this order.
    This calculation assumes that the number of terminals
    and discs are correct and acceptable so
    configure those first.)

  (PROG (#TERMINALS:INTEGER #DISCS:INTEGER)
    (#TERMINALS :=(SAND CURRENT.BIN 'FETCH '#TERMINALS))
    (#DISCS :=(SAND CURRENT.BIN 'FETCH '#DISCS))
    (RETURN (SELECTQ #DISCS

```

```

(0 1)
(1 (IF #TERMINALS = 1
THEN 1
ELSE 2))
(2 (IF #TERMINALS = 4
THEN 3
ELSE 2))
(3 (IF #TERMINALS = 4
THEN 4
ELSE 3))
#DISCS)))

```

```

(DEFKBFUN COMPUTE.POWER.NEEDED ()(PROG (#RAM:INTEGER
#ROM:INTEGER #IF-CARDS:INTEGER)
(#RAM :=(SAND CURRENT.BIN 'FETCH '#RAM))
(#ROM :=(SAND CURRENT.BIN 'FETCH '#ROM))
(#IF-CARDS :=(SAND CURRENT.BIN 'FETCH '#I
F-CARDS))
(RETURN (2 + 1*#RAM + .5*#ROM + .2*#IF-CA
RDS)))
)

```

APPENDIX V GLOSSARY

ASSEMBLY

A specific set of product instances or parts.

ASSEMBLY DESCRIPTION

A Boolean function or predicate defining a set of parts in terms of a set of part names, names of assemblies or part types. A collection of parts satisfies an assembly description when the parts fit that description.

BIN

A set of product instances created during execution of a CONFIGURE operation in a task block. The product instances are obtained from the order lines, by an EXPANSION operation upon parts initially placed in the bin, or by application of a modification constraint.

BIN TREE

A hierarchical structure of bins corresponding to the hierarchy of the functional classes of the bins.

CHECK

An operation which applies a specified set of constraints to the product instances or parts in the current bin in order to check whether there are any configuration problems with the parts in the bin.

CONDITION

A Boolean function of the elements in a defined set.

CONFIGURE

An operation which creates a new bin of a specified functional class and moves a respective subset of parts from the current bin to the new bin, and then invokes execution of the task block associated with the type of the new bin.

CONFIGURATION

The grouping of a set of elements into subsets.

CONSTRAINTS

A statement of a condition that certain parts must satisfy and a respective action to be conditionally executed in response to whether the condition is

satisfied when the constraint is applied.

EXPAND

An operation which selects a specified subset of parts from the current bin and adds any new parts that correspond to the "product expansions" or sub-components of the specified subset of parts.

FIELD SERVICE PARTS

A list of parts, such as replacement and spare parts, that are included in an order but which are not checked for compatibility with themselves or with the rest of the order.

FIND

An operation which matches the parts in the current bin to an assembly description and copies the matching parts to an assembly list.

FUNCTION

(1) The purpose or performance of an assembly.

(2) A numerical or logical value responsive to the domain of the function. A function of the parts in an assembly, for example, may indicate whether a specified number of specified parts are included in the assembly or whether the parts have specified conditions or attributes. In a computer program, functions are applied or evaluated at only specified instants of time.

FUNCTIONAL HIERARCHY

The decomposition of a configuration or product into its major functional components of sub-assemblies.

INITIAL BIN

The root of the bin tree

MODIFICATION CONSTRAINT

A constraint which changes the set of product instances or parts in the current bin, and may also issue a warning message to the user.

OPERATION

A specific action in a control procedure or computer program which, when executed, modifies the sequence of execution of the control procedure or modifies the data or configuration elements processed by the control procedure.

ORDER

The input to the order processing system including order information, order lines, and field service parts.

ORDER INFORMATION

Information which primarily identifies an order, such as customer name and order number, and does not affect the configuration process.

ORDER LINES

A list of parts ordered and the quantity of each part that will be checked.

PRODUCT EXPANSION

For a specified part or product instance, the set of corresponding sub-parts or product instances making up the specified part or product instance.

PRODUCT INSTANCE

An explicit representation of an individual part considered by the order processing system.

TASK BLOCK

A set of imperative language statements specifying operations to perform with respect to associated bins.

UNASSIGNED

A type of bin denoting that the configuration system has not yet determined what product the parts of the bin will implement and therefore which knowledge base should be used to configure the order.

WARNING CONSTRAINT

A constraint which issues a message, informing the user of a problem with the configuration.

What is claimed is:

1. A knowledge system comprising a computer having a memory storing a knowledge base, said knowledge base including

predefined descriptions of assemblies of predefined components for indicating whether a specified component is permitted in a specified assembly, and

for at least some of said descriptions of assemblies, sets of predefined conditions of different properties of the components in the assembly for indicating whether certain predefined properties of a specified set of components are compatible within the assembly,

said memory comprising a control procedure for first obtaining a predetermined initial set of predefined components,

thereafter matching the initial set of components to the predefined descriptions of assemblies to determine respective sets of matching components including the components in the initial set which are also permitted in the respective assemblies,

thereafter applying the sets of conditions to the respective sets of matching components for the respective assemblies, and

indicating whether the respective conditions are satisfied,

said computer including means for executing said control procedure to thereby first obtain said predetermined initial set of predefined components, thereafter match said initial set of components to said predefined descriptions of assemblies, thereafter apply said sets of conditions to the respective sets of matching components for the respective assemblies and generate an indication of whether the conditions are satisfied.

2. The knowledge system as claimed in claim 1, wherein said descriptions of assemblies include the names of at least some of said components in the initial set of components.

3. The knowledge system as claimed in claim 1, wherein said knowledge base further includes a catalog of components including component names and respective names of assemblies comprising said components, and wherein said descriptions of assemblies include the names of some of said assemblies in said catalog of components.

4. The knowledge system as claimed in claim 3, wherein said descriptions of assemblies comprise Boolean functions of names of assemblies.

5. The knowledge system as claimed in claim 1, wherein said control procedure includes a set of control steps for obtaining attributes of the components in said initial set of components, and wherein said descriptions of assemblies include respective Boolean functions of said attributes.

6. The knowledge system as claimed in claim 1, wherein said control procedure includes a set of control steps for obtaining a predetermined initial set of predefined components by inputting a predetermined list of components.

7. The knowledge system as claimed in claim 1, wherein the knowledge base further includes respective actions for said conditions, and wherein said control procedure includes a set of control steps for carrying out said actions for indicating whether the respective conditions are satisfied.

8. The knowledge system as claimed in claim 7, wherein said actions comprise adding and deleting specified components from said initial set of components.

9. The knowledge system as claimed in claim 1, wherein the descriptions of assemblies of predefined components include at least one definition of a relationship between at least two of said descriptions of assemblies, and wherein the control procedure includes a set of control steps for matching the initial set of components to at least one of said two descriptions of assemblies in response to the descriptions of the two assemblies and in response to the defined relationship between the two assemblies.

10. The knowledge system as claimed in claim 9, wherein said relationship is a structural relationship between said two assemblies.

11. The knowledge system as claimed in claim 9, wherein said relationship is a functional relationship between said two assemblies.

12. The knowledge system as claimed in claim 1, wherein some of said assemblies are sub-assemblies of other of said assemblies, and wherein said knowledge base includes a definition of a hierarchy of said descriptions of assemblies defining offspring-parent relationships between said descriptions of sub-assemblies and the respective descriptions of assemblies which comprise said sub-assemblies, and wherein said control procedure for matching the initial set of components to the predefined descriptions of assemblies includes means for first matching the list of components to the parent descriptions of assemblies, and then matching the list of the respecting matching components for the respective parent assemblies to their respective offspring descriptions of sub-assemblies.

13. The knowledge system as claimed in claim 12, wherein some of said parent descriptions of assemblies represent structures comprising their respective sub-assemblies physically connected together.

14. The knowledge system as claimed in claim 12, wherein some of said parent descriptions of assemblies represent functions collectively performed by the respective offspring sub-assemblies.

15. The knowledge system as claimed in claim 1, wherein said knowledge base further comprises a set of expansion rules defining expandable components including some of said components in the initial set of components in terms of respective sub-components of the expandable components and wherein said control procedure further includes a set of control steps for selecting expandable components from the initial set of components and adding the respective sub-components corresponding to the selected expandable components to the initial set of components.

16. The knowledge system as claimed in claim 15, wherein said set of control steps for selecting expandable components includes control steps for searching the initial set of components for specified components and when any of the specified expandable components are found, expanding those expandable components by adding their respective sub-components to the initial set of components.

17. The knowledge system as claimed in claim 1, further comprising a working configuration portion of said memory, and wherein said set of control steps for matching the initial set of components to the defined assemblies includes control steps for recording the respective matching components in respective bin portions of said working configuration portion of said memory.

18. The knowledge system as claimed in claim 17, wherein said knowledge base further includes respective modification actions for making specified changes to the recorded set of matching components in the respective bin portions of said working configuration portion of said memory, and wherein the said control procedure for applying the sets of conditions includes a set of control steps for conditionally executing said changes in response to whether the respective conditions are found to be satisfied for the respective assemblies when the respective conditions are applied.

19. The knowledge system as claimed in claim 18, wherein said conditions include descriptions of assemblies including specific components.

20. The knowledge system as claimed in claim 18, wherein said descriptions of assemblies include Boolean functions of specified attributes of components.

21. The knowledge system as claimed in claim 1, wherein said control procedure includes
 a built-in control procedure independent of the descriptions of assemblies and said predefined conditions, and
 task blocks defining control steps for said matching of said components to specified descriptions of assemblies and applying said conditions to the respective matching components, said task blocks being executed when specified steps in said built-in control procedure are reached.

22. The knowledge system as claimed in claim 21, wherein said task blocks include task blocks associated with respective assemblies and executed when said components are matched with the respective descriptions of assemblies of the task blocks.

23. The knowledge system as claimed in claim 22, wherein said descriptions of assemblies of predefined components include Boolean functions of specified attributes of components.

24. The knowledge system as claimed in claim 22, wherein said control steps defined by said task blocks include control steps for applying specified subsets of said conditions corresponding to the specified descriptions of assemblies.

25. The knowledge system as claimed in claim 21, wherein said task blocks include a specified task block to be executed before the other task blocks are executed.

26. The knowledge system as claimed in claim 21, wherein said conditions include Boolean functions of specified attributes of specified components, and wherein said knowledge base includes a knowledge base function associated with a respective one of said specified attributes, said knowledge base function being implicitly invoked when said condition is applied including said Boolean function of said specified attribute, so that the knowledge base function determines a value for its respective specified attribute.

27. The knowledge system as claimed in claim 21, wherein said built-in control procedure includes a set of control steps for generating a trace of the operations performed by said control procedure including the sequence of control steps performed when said task blocks are executed, and wherein said control procedure also includes a set of control steps for generating a user-understandable explanation from said trace and transmitting said explanation to a user.

28. The knowledge system as claimed in claim 1, wherein said control procedure includes a set of control steps for generating a trace of the results of said matching of said components and said application of said conditions and also includes a set of control steps for generating an explanation understandable to a human user from said trace and transmitting said explanation to a human user.

29. A knowledge system comprising a computer having a memory storing a knowledge base, said knowledge base including

predefined descriptions of assemblies of predefined components for indicating whether a specified component is permitted in a specified assembly, and

corresponding sets of assembly constraints including predefined conditions of different properties of the components in their respective assemblies for indicating whether certain predefined properties of a specified set of components are compatible within the assembly, and also including predefined actions to perform when the constraints are applied and the conditions are satisfied including outputting a warning to a user indicating specified conditions are satisfied,

said memory further comprising a control procedure for

first obtaining a predetermined list of components, thereafter matching the components in said list to the predefined descriptions of assemblies to determine respective sets of matching components containing the components in said list which are permitted in the respective assemblies,

thereafter applying the sets of assembly constraints to the respective sets of matching components from the list for the respective assemblies, and

conditionally performing the respective actions in response to whether their respective conditions are satisfied when the assembly constraints are applied, said computer including means for executing said control procedure to thereby first obtain said predetermined list of components, thereafter match the components in said list to the predefined descriptions of assemblies, thereafter apply said sets of constraints to the respective sets of matching components from said list of components, and con-

ditionally perform said respective actions in response to whether their respective conditions are satisfied when the assembly constraints are applied.

30. The knowledge system as claimed in claim 29 wherein said descriptions of assemblies define a hierarchy of descriptions of assemblies and sub-assemblies, and said control procedure for matching includes a procedure for matching said components in the respective assemblies to the descriptions of specified sub-assemblies to determine the components for which said constraints are applied.

31. The knowledge system as claimed in claim 29, wherein said constraints include conditions responsive to whether specified parts have been matched to the respective descriptions of assemblies.

32. The knowledge system as claimed in claim 29, wherein said constraints include conditions responsive to whether specified parts have specified attributes.

33. The knowledge system as claimed in claim 32, wherein the knowledge base includes definitions of variables for specifying the attributes of components included in the descriptions of respective assemblies.

34. The knowledge system as claimed in claim 33, wherein said knowledge base includes knowledge base functions for at least some of said variables specifying steps for determining values for the respective attributes of said components, and said control procedure includes a set of control steps for implicitly invoking and executing the respective knowledge base functions when applying constraints including conditions referencing the respective variables having the knowledge base functions.

35. The knowledge system as claimed in claim 29, further comprising a working configuration portion of said memory, and wherein said control procedure for matching the components in said list to the descriptions of assemblies includes a set of control steps for recording the respective matching components in respective bin portions of said working configuration portion of said memory.

36. The knowledge system as claimed in claim 35, wherein said control procedure for applying the sets of assembly constraints includes a set of control steps for applying a specified set of constraints to the matching components recorded in a specified bin portion of said working configuration portion of said memory.

37. The knowledge system as claimed in claim 35, wherein said knowledge base includes expansion rules defining sub-components for respective components, and wherein said control procedure includes a set of control steps for searching a specified bin for at least some of said components having sub-components defined by said expansion rules, and recording the respective sub-components in the specified bin.

38. The knowledge system as claimed in claim 35, wherein the actions of the assembly constraints include actions specifying modifications to the set of components recorded in a specified bin.

39. The knowledge system as claimed in claim 35, wherein said working configuration portion of said memory includes a current bin memory location for storing a value specifying a particular bin.

40. The knowledge system as claimed in claim 39, wherein the control procedure includes task blocks including control steps, at least some of the bins having an associated task block for specifying operations to perform in connection with their respective bins, and

the control procedure includes a set of steps for executing the task block associated with a specified bin.

41. The knowledge system as claimed in claim 40, wherein the set of steps for executing the task block associated with a specified bin include the steps of pushing the value of the current bin memory location onto a stack, setting the value of the current bin memory location to specify the specified bin, executing the control steps in the task block for the specified bin, popping the stack and assigning the popped value to the current bin memory location.

42. The knowledge system as claimed in claim 40, wherein the task blocks include imperative language statements specifying distinct steps in said control procedure.

43. The knowledge system as claimed in claim 42, wherein the knowledge base includes expansion rules specifying that at least some of said components are expandable and are comprised of sub-components and wherein the imperative language statements include a separate statement for specifying that for each expandable component in a specified bin, its respective sub-components are to be recorded in the specified bin.

44. The knowledge system as claimed in claim 42, wherein the imperative language statements include a separate statement for specifying that specified assembly constraints are to be applied to the components in a specified bin.

45. The knowledge system as claimed in claim 42, wherein the imperative language statements include a separate statement for executing a specified task block.

46. The knowledge system as claimed in claim 42, wherein the control procedure includes a set of control steps for generating a sequential record of the imperative language statements that are executed, the components added to bins, the constraints applied, and the actions that are executed when the conditions of the applied constraints are satisfied.

47. The knowledge system as claimed in claim 40, wherein said knowledge base includes a plurality of separate portions for processing predetermined lists to configure different respective products, and wherein the control procedure includes control steps for obtaining said predetermined list by an input operation and for selecting the knowledge base portion corresponding to a product number in said predetermined list.

48. The knowledge system as claimed in claim 35, wherein said control procedure for matching the components in said list includes a set of control steps for finding the components in a specified bin which satisfy a specified assembly description.

49. The knowledge system as claimed in claim 48, wherein said control steps for finding the components in a specified bin are responsive to specified assembly descriptions including Boolean functions of part names, names of assemblies and part types.

50. A knowledge system comprising a computer having memory storing

a knowledge base including knowledge about a set of related elements, and

an initial list of certain ones of said elements,

said knowledge about said set of related elements including

a predefined declaration of hierarchical decomposition of said set of related elements into separately defined subsets of said elements, some of said separately defined subsets being composed of other of

said separately defined subsets having fewer elements, and

for said separately defined subsets of said elements, respective predefined functions of the respective elements in the subsets, the domain of each function thereby being its respective subset, such function having a predefined result obtained when the function is applied responsive to the specified elements within its domain,

said memory including a working configuration portion for storing lists of said elements,

said memory storing a predefined matching procedure for matching a specified list of certain ones of said elements to a specified one of said subsets to thereby obtain a list of matching elements which are the elements in said specified list which are also elements of said specified subset, said list of matching elements being stored in said working configuration portion of said memory, and said memory further storing a predefined control procedure for successively applying said matching procedure to match said initial list to said subsets of elements composing said set of related elements, and for each of said subsets of elements composed of said subsets of fewer elements applying said matching procedure to match the previously obtained respective list of matching elements to each of said subsets of fewer elements composing said subset of elements to thereby obtain lists of matching elements representing the configuration of said initial list according to said hierarchical decomposition, and

applying said functions to the respective lists of matching elements within the domains of the respective functions and obtained by matching to the respective subsets, and

said computer including means for executing said control procedure to thereby obtain said matching lists of elements representing a configuration of said initial list according to said hierarchical decomposition, and for applying said functions to the respective matching lists within the domains of the respective functions.

51. The knowledge system as claimed in claim 50, wherein said functions include functions responsive to whether specified elements of the initial list are included within said respective lists of matching elements within the domains of the respective functions.

52. The knowledge system as claimed in claim 50, wherein the knowledge base includes predetermined changes to the composition of said lists of matching elements stored in said working configuration portion of said memory, and wherein said functions include respective Boolean condition functions for indicating said changes, and wherein the control procedure includes control steps for executing the indicated changes to thereby conditionally change said lists of matching elements.

53. The knowledge system as claimed in claim 52, further comprising an explanation facility for generating an ordered list of the conditions for which changes were indicated and the changes which were executed, and transmitting the ordered list to a user.

54. The knowledge system as claimed in claim 50 wherein said subsets correspond to hierarchical functions collectively performed by the combinations of their respective elements, and the definitions of said subsets explicitly stage their respective functions.

55. The knowledge system as claimed in claim 54, wherein said respective functions of the respective elements in the subsets include at least one function defining a variable, and at least one of said task blocks includes a statement specifying when the function is applied to determine a value for its respective variable.

56. The knowledge system as claimed in claim 50, wherein said declaration of hierarchical decomposition defines classes of subsets of said elements, and said control procedure includes configuring control steps for creating specified instances of said classes of subsets and adding corresponding elements of said initial list to the instances of said subsets.

57. The knowledge system as claimed in claim 56, wherein said control procedure includes a set of control steps for recursively applying said configuring control steps.

58. The knowledge system as claimed in claim 57, wherein said set of control steps for recursively applying starts recursive application at the uppermost level of the hierarchical decomposition.

59. The knowledge system as claimed in claim 50, wherein said knowledge base includes a plurality of task blocks specifying in part the operation of said control procedure, and said subsets are associated with particular ones of said task blocks, and said task blocks specify operations to perform with respect to their associated subsets.

60. The knowledge system as claimed in claim 59, wherein said task blocks include imperative language statements including separate statements for creating and storing in said working configuration portion of memory the list of matching elements obtained by matching to a specified one of said subsets, and for applying a specified set of functions to the respective lists of matching elements within the domains of said functions.

61. A knowledge system for checking a production request for a flexibly-assembled product, said request including a list of part names and respective quantities, said knowledge system comprising a computer having a memory storing a knowledge base, said knowledge base including

a declaration of a hierarchy of functional assemblies including parent-offspring relationships between respective assemblies and sub-assemblies,

a parts catalog including a predefined set of parts and identification of the functional assemblies which the respective parts compose, and

a set of conditions applicable to respective assemblies for indicating whether parts configured into the respective assemblies are compatible,

said memory also including

a working configuration portion of said memory subdivided into respective bins for receiving the names of parts configured into the respective assemblies, and

a control procedure for obtaining said request,

configuring the part names in the request into at least one parent assembly by referencing the parts to the parts catalog to determine whether the parts compose the parent assembly and when they compose the parent assembly adding the respective part names to the parent bin,

configuring the part names in the request to respective offspring assemblies by referencing the parts in the respective parent bins to the parts catalog

to determine whether the parts comprise the offspring assemblies and when they comprise the respective offspring assemblies adding the respective part names to the respective offspring bins, and

applying said conditions to the part names in the respective assemblies, and said computer including means for executing said control procedure to obtain said request, configure the part names in said request to respective offspring assemblies, and apply said conditions to the part names in the respective assemblies to thereby check whether the request includes compatible parts for all of the assemblies in the product.

62. The knowledge system as claimed in claim 61 wherein the knowledge base further comprises task blocks for respective ones of said assemblies, said task blocks including imperative language statements for specifying control procedure steps responsive to the parts in the bins for the respective assemblies.

63. The knowledge system as claimed in claim 62 wherein the knowledge base further comprises a separate set of expansion rules specifying the composition of expandable parts in terms of sub-parts, and wherein the task blocks include imperative language statements for adding the sub-parts of expandable parts in the bins to the respective bins.

64. The knowledge system as claimed in claim 61, wherein said knowledge base further comprises a set of actions specifying changes to the set of parts configured into at least some of said assemblies in response to the respective conditions for said assemblies, and wherein said control procedure includes a set of control steps for changing the sets of parts configured into the respective assemblies in the fashion indicated by said actions in response to whether the respective conditions are satisfied when the respective conditions are applied.

65. The knowledge system as claimed in claim 64 wherein the control procedure includes a set of control steps for generating and transmitting to a user a list of said changes made to the set of parts configured into the respective assemblies in the fashion indicated by said actions and an explanation of reasons why the changes were made.

66. The knowledge system as claimed in claim 64 wherein some of said actions specify that certain additional information is to be requested from a user and that certain changes are to be conditionally made in response to information received from the user.

67. A knowledge system for designing a product including assemblies of predetermined parts, said knowledge system comprising a computer having a portion of memory storing a knowledge base and a

portion of memory subdivided into respective bins for receiving selected names of said parts for comprising the respective assemblies,

said knowledge base including a control procedure for

obtaining a set of assembly constraints including conditions applicable to respective assemblies for indicating whether parts in the respective bins have desired properties and respective actions indicating predetermined changes to the sets of parts in the respective bins for obtaining said desired properties in response to whether the respective conditions are satisfied, and

applying said conditions of said assembly constraints to the respective assemblies and conditionally executing the respective changes in response to whether the conditions are satisfied,

wherein said knowledge base further includes a hierarchy defining at least some of said assemblies as offspring sub-assemblies of respective parent assemblies, and wherein said assembly constraints associated with said parent assemblies have conditions referencing the conditions of parts in the respective offspring bins, and

said computer including means for executing said control procedure to thereby obtain said set of assembly constraints, apply said assembly constraints, and conditionally execute said respective changes to thereby obtain a design for said product, said design being indicated by the list of names of the parts in the bins after said conditions are applied and said changes are conditionally executed.

68. The knowledge system as claimed in claim 67 wherein the constraints of said offspring assemblies reference the conditions of parts in the respective parent bins via respective bin variables having respective scopes including the respective parent assembly conditions and the respective offspring assembly conditions.

69. The knowledge system as claimed in claim 67 wherein the knowledge base further comprises task blocks for respective ones of said assemblies, said task blocks including imperative language statements for specifying control procedure steps responsive to the parts in the bins for the respective assemblies.

70. The knowledge system as claimed in claim 67 wherein the control procedure includes a set of control steps for generating and transmitting to a user an ordered list of said changes in the sequence that the changes are executed, and an explanation of why said changes were made.

* * * * *

[54] MATERIAL REQUIREMENTS PLANNING SYSTEM AND PROCEDURES FOR USE IN PROCESS INDUSTRIES

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[51] Int. Cl.⁴ G06F 15/00

[52] U.S. Cl. 364/403; 364/468

[58] Field of Search 364/403, 468

[56] References Cited

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Attorney, Agent, or Firm—Parmelee, Bollinger & Bramblett

[57] ABSTRACT

A system for controlling the flow of semiconductor

products and their components through a production facility including assembly and final testing of a large number of different products with multiple product grades. The system stores information on the demand and inventory of all product grades together with grade distribution data giving the yield of all co-products of a product family from testing the common component of that family. The system includes means for calculating the number of common components to be tested to cover the requirements for all co-products of a family. The distribution data also gives the yield of all by-products of a test for particular products, and the system operates to adjust the inventory status for other products corresponding to such by-products, to avoid excess production of such other products in meeting projected demand. For cases where the data indicate for any product family a surplus of a higher grade and a shortage of a lower grade, the system includes means to decide whether to downgrade a portion of the higher grade to the lower grade; the projected inventory status of each grade is adjusted accordingly.

10 Claims, 7 Drawing Figures

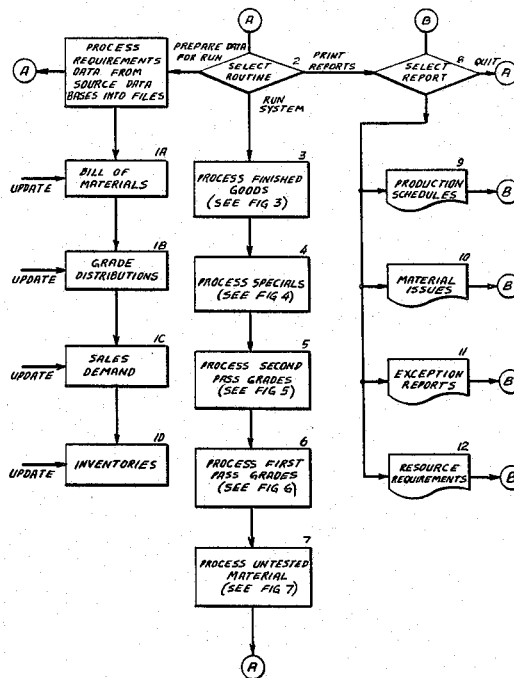


Fig. 1.

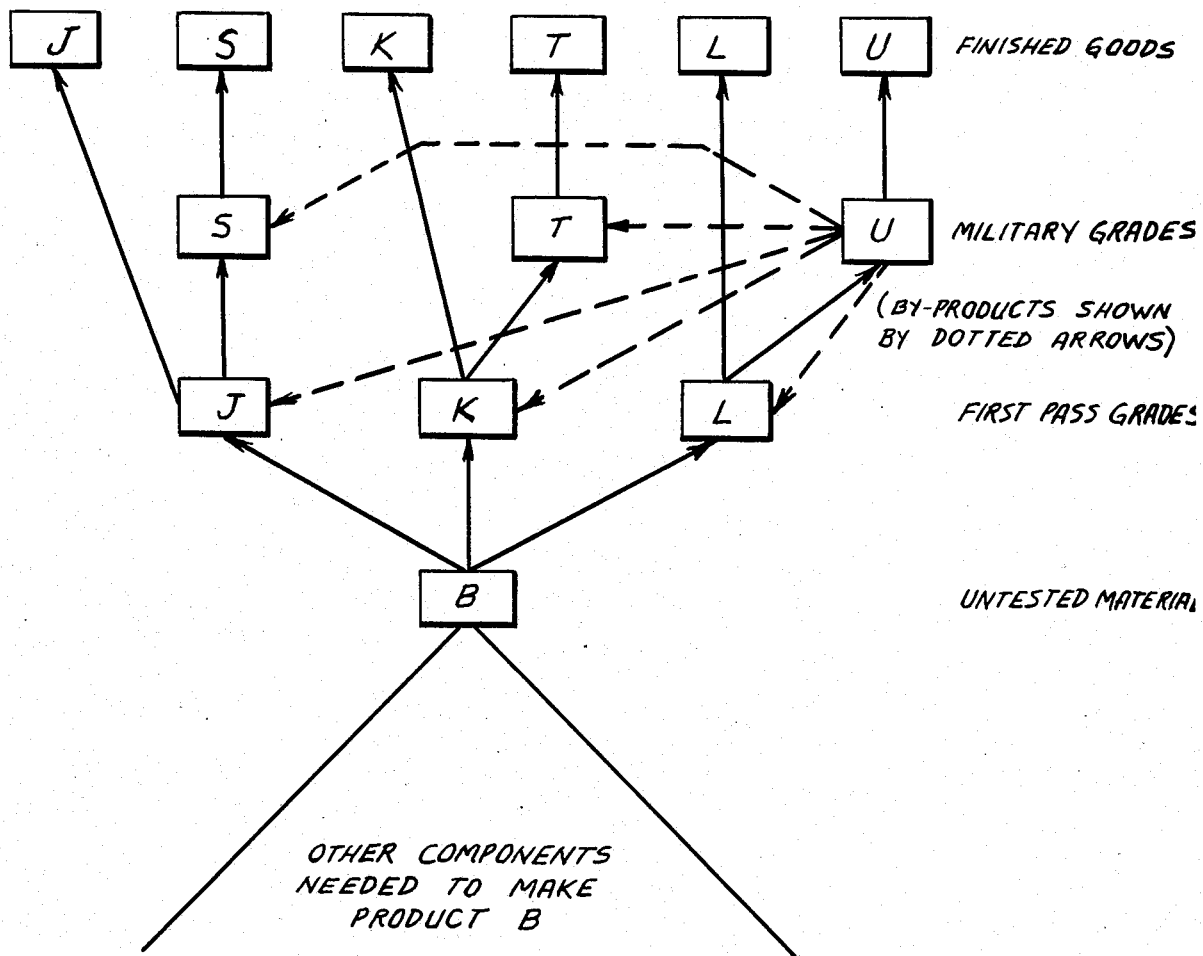
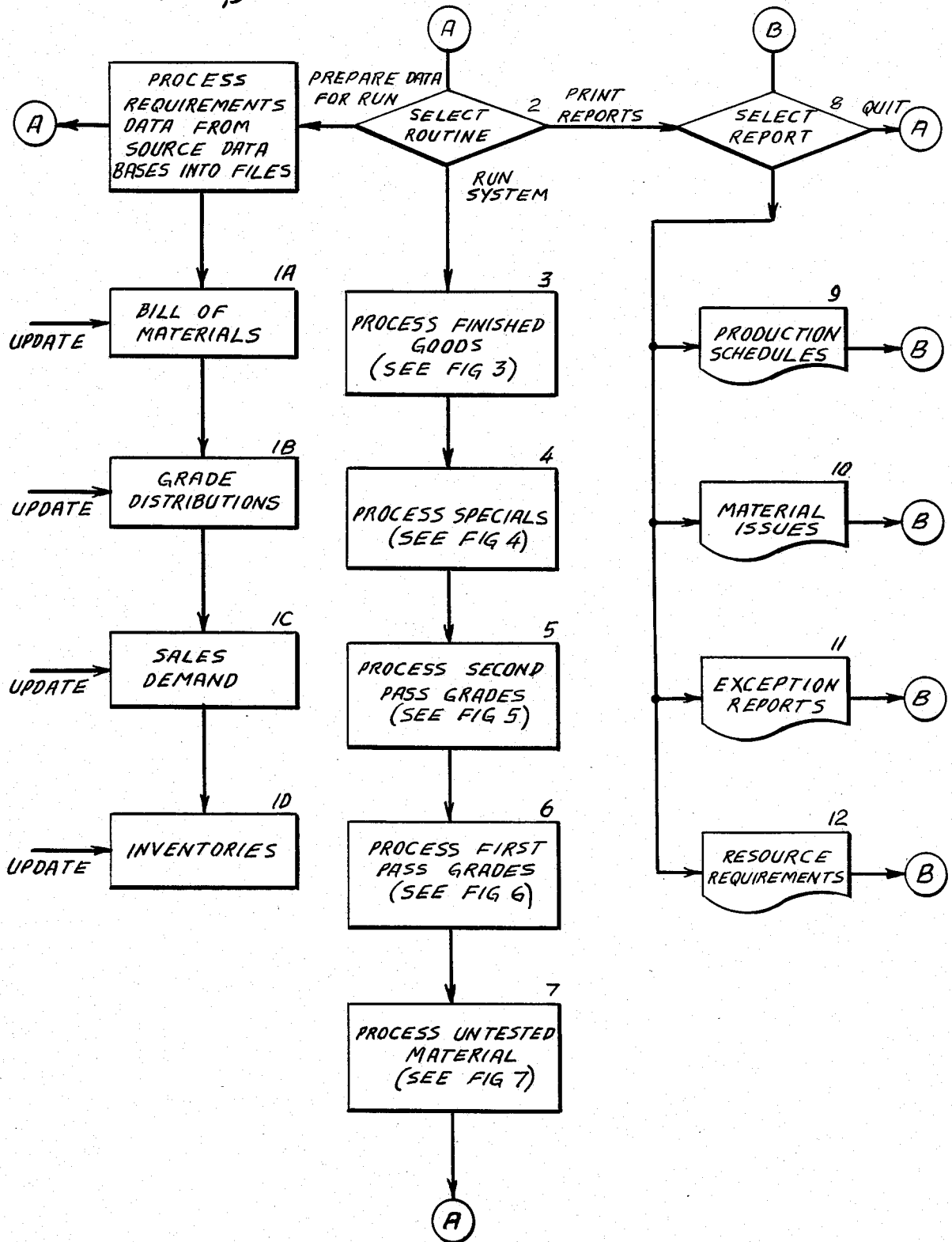
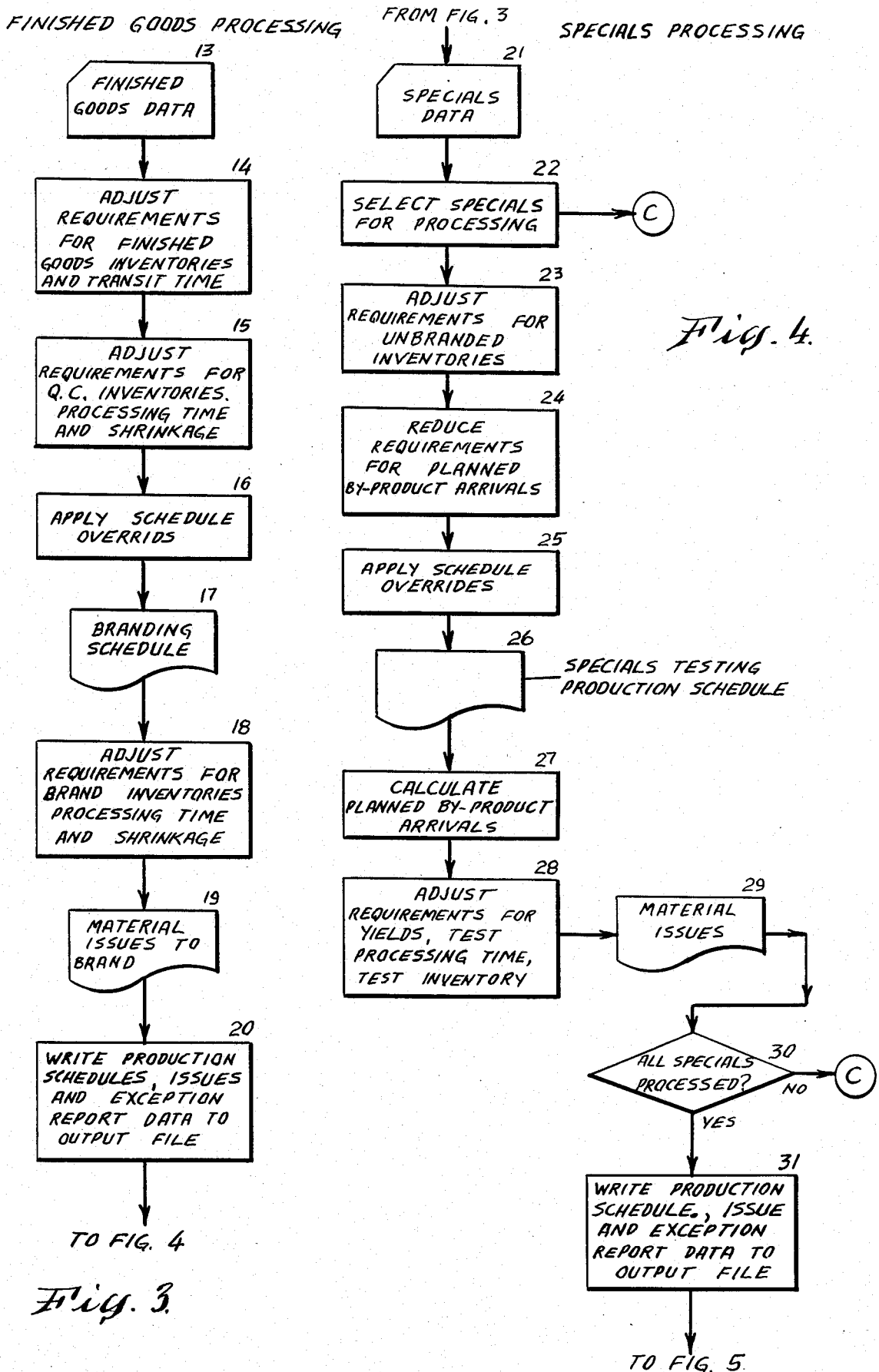
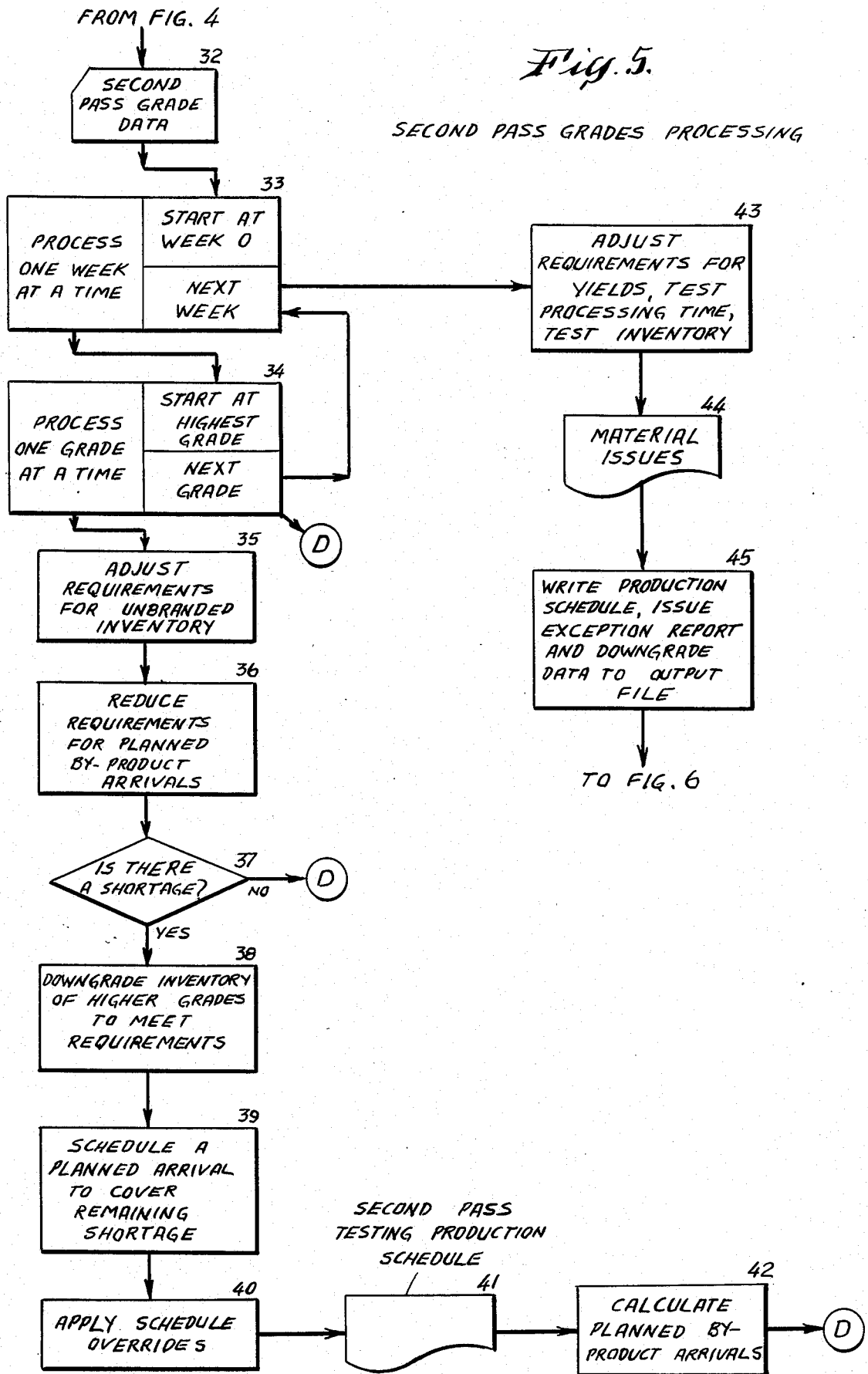
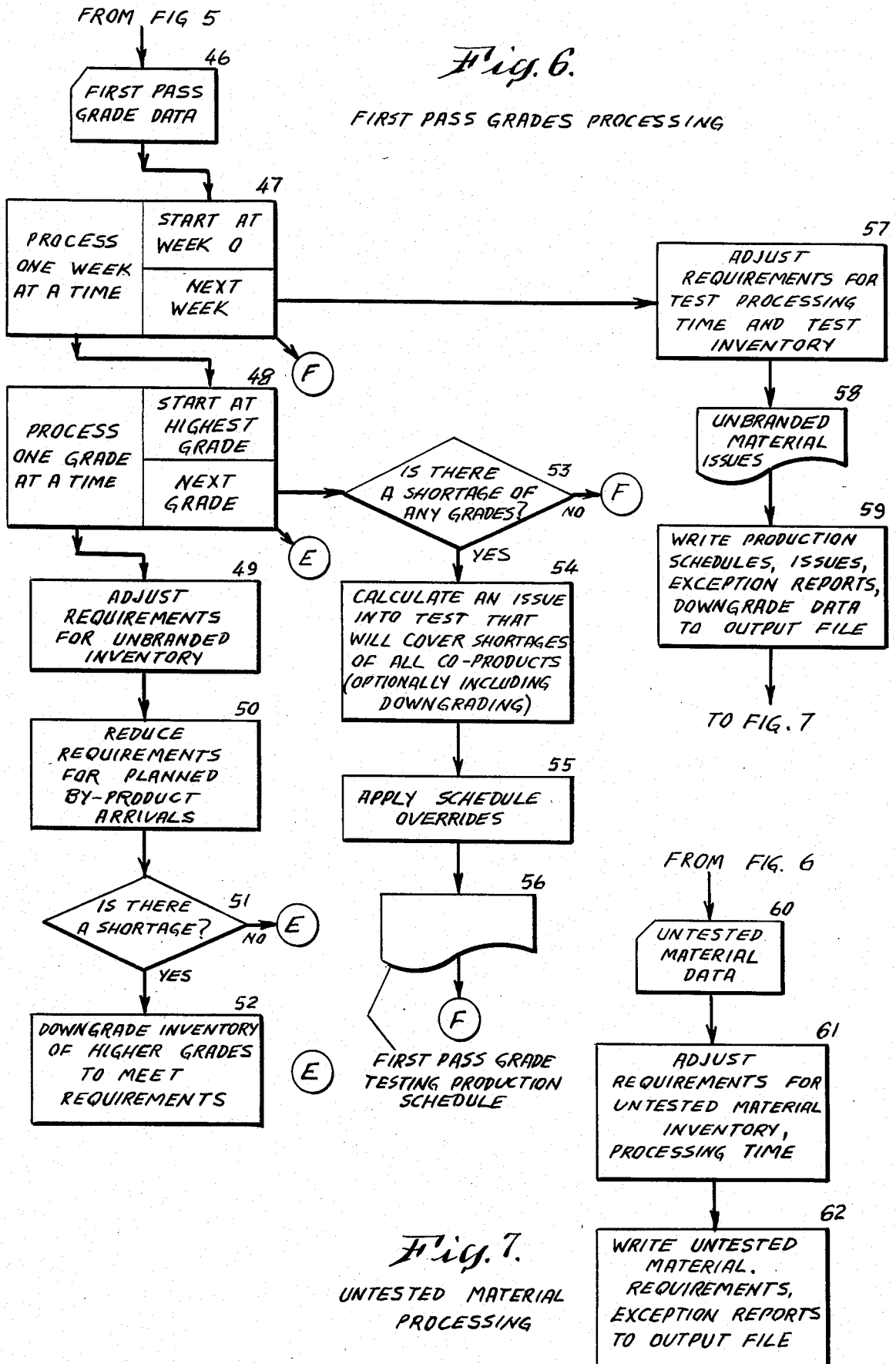


Fig. 2.









MATERIAL REQUIREMENTS PLANNING SYSTEM AND PROCEDURES FOR USE IN PROCESS INDUSTRIES

BACKGROUND OF THE INVENTION

1. Field of the Invention

This invention relates to a computerized system for planning the testing of products as part of a manufacturing process. The invention is illustrated for use in planning the testing of semiconductors in the form of integrated circuit (IC) chips.

Semiconductor chips require testing and grading against specifications before they can be sold. The testing of any lot of product will also sort the functional units into separate grades depending on their electrical characteristics. More rigorous or selective testing can then be performed on these grades to produce products which are guaranteed to perform in the more strenuous environment demanded by military or aerospace applications. In addition, certain customers may require special tests to be performed on the product for their specific use.

Production of semiconductors begins in wafer fabrication during which the electronic circuits are grown on thin wafers of silicon in a series of alternating photographic and deposition processes which create hundreds or thousands of identical circuits on each wafer. This process can take from two to upwards to ten weeks depending on the complexity of the process and the number of "layers" involved.

The effectiveness of the wafer fabrication process is dependent on many environmental factors over which the manufacturer has varying degrees of control but which have a definite effect on the functionality and electrical characteristics of the resulting product. In order to eliminate further expensive processing of non-functional circuits, the wafers are probed at this point—i.e. each circuit is powered up, tested and the failures are marked with ink to be later discarded.

After probing, the individual dice are separated, inspected for visual defects and mounted in the individual packages which allow them to communicate with the real world. These packages come in many shapes and sizes from two-lead metal cans to upwards of sixty-four lead dual in-line pin ceramic or plastic packages. Wires finer than a human hair connect the terminals on the dice to the leads on the packages which are then sealed to protect the chip from environmental and mechanical damage.

After this assembly process is complete, the units are tested and graded according to their electrical performance against the range of specifications. The simplest grades are measured at room temperature and are known as commercial. The more difficult grades are those for military and aerospace applications. These must pass strict government standards after being subjected to extreme environmental stress, including continual operation for a week under high temperatures. In addition, the critical characteristics of the higher grade products must not drift out of specification over the whole operating temperature range.

The testing process can be considered as a series of hurdles which the product must pass. Product which fails can be saved to retest against a different set of hurdles, sold as the highest grade for which it has quali-

fied or conceivably sold as a lower grade if no other demand for the product exists.

The process is made especially complex by the large number of different types and grades of product. For example, in one commercial plant, 80 wafer types are assembled into 160 different package configurations which yield over 2500 different end products after testing. The simplest product has 2 grades while the most complex had 88 different grades of standard catalog items and customer specials.

The task for production control is to identify and schedule the correct sequence of tests and an adequate supply of untested material coming out of assembly to support the demand for all of the items. The goal is to accomplish this with minimum inventories and optimal utilization of scarce and expensive computerized test equipment while maintaining timely and accurate deliveries.

2. Description of the Prior Art

The standard prior art technique for planning the type of production used for semiconductor chips—that is, the technique for specifically determining the requirements for production of the source materials needed in a higher level product—is called Material Requirements Planning (MRP). Such planning is carried out by a computer operating under software control. The computer computations typically are based on a Bill of Materials which defines the source components of any product, the sub-components of such components, and so forth down to the lowest level materials—usually raw materials.

Briefly, Material Requirements Planning performs its planning function by considering the demand for an item and subtracting out any inventories to determine additional production required. This net requirement is then "exploded", that is, multiplied by the appropriate quantities of each component and advanced by the time needed to produce the item. The result is the demand for each component. The process is repeated for each component in turn until the lowest level is reached.

SUMMARY OF THE INVENTION

In the testing process described herein, all untested products are given a first pass test to eliminate non-functional units and produce a preliminary grading. The output from the test becomes input to the military or second pass test. Optionally, the output from either the first pass or second pass test may become input to a test for so-called customer specials. Any product which passes any of the tests can be branded with its identifying grade and sold as a finished product.

The procedure is controlled by a software program operating with a data base defining the proper components for finished goods, customer specials (if used), second pass grades and first pass grades. The procedure resolves three issues in the testing environment which standard MRP does not address. These are: by-product planning, co-product planning, and downgrading, which may be described briefly as follows:

BY-PRODUCT PLANNING

When testing for military grades, some of the product will pass the test. The product which fails is still good product, however, and is called a by-product. Normally this by-product will be another item already in the Bill of Materials and most often the component item itself. Standard MRP does not account for the production of these by-products nor does it consider them as available

inventory to be netted out of additional production requirements.

CO-PRODUCT PLANNTING

A special case occurs for first pass grades where two or more grades are simultaneously produced from the same test, and are referred to as co-products. In effect these grades are by-products of each other. Standard MRP which calculates items one at a time sequentially cannot simultaneously plan the component quantity which must be tested to meet the demand for all of these grades together.

DOWNGRADING

Because grading is a natural phenomenon independent of demand, its distribution does not necessarily match the distribution of demand. This may result in surplus inventories of higher grades and shortages of lower grades. There are circumstances where it makes economic sense to downgrade these higher grades to cover the shortages rather than to produce more of the lower grades. Standard MRP always covers shortages by additional production. In the system described herein, additional processing is employed to determine which products, when and in what quantities should be downgraded.

A general purpose computer system operating under software control as described herein resolves the above three issues through logic procedures inserted into standard MRP software.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a simplified diagrammatic representation of hierarchical levels to which items can be assigned for testing;

FIG. 2 is a flow diagram giving an overview of the entire MRP system as modified to incorporate the present invention;

FIG. 3 is a flow diagram showing the logic of finished goods scheduling represented by functional block 3 of FIG. 2. (Note: the logic is standard prior art MRP and is shown herein only for completeness);

FIG. 4 is a flow diagram showing the logic of customer specials represented by the functional block 4 of FIG. 2;

FIG. 5 is a flow diagram showing the logic of second pass test planning represented by functional block 5 of FIG. 2;

FIG. 6 is a flow diagram showing the logic of first pass test planning represented by functional block 6 of FIG. 2; and

FIG. 7 is a flow diagram showing the logic for planning untested material represented by the functional block 7 of FIG. 2.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

The present invention is based on the concept that in some manufacturing industries there are relationships between the products being processed through testing which make it important and advantageous to plan the testing of certain product families together. To that end, the system to be described categorizes products into different levels so that each level can be processed using the distinctive procedures which apply to that level. A multi-level product structure used in the system to be described is as follows:

Level 0—Finished Goods

Level 1—Customer Specials

Level 2—Second Pass Tested Standard Items

Level 3—First Pass Tested Standard Items

Level 4—Untested Product

Defining these in reverse order (i.e. in the direction in which product actually moves), Level 4 contains untested material as it comes from assembly, and Level 3 contains the grade distribution resulting from the first ambient test of newly assembled product. Typically in the semiconductor industry there are three grades—J, K and L, with J the lowest grade and L the highest. These grades may be branded and shipped as their level 0 branded standard commercial counterparts or issued for higher grade testing.

Level 2 contains the military versions of the commercial grades. The product is stressed, burned in and retested for its drift over a temperature range. An L so tested which retains its electrical characteristics might be called a U. A K might become T and a J and S. The L which fails as a U might still be an L or perhaps have degraded to a T or S, K or J. The exact distribution would be spelled out on a "lot traveller" accompanying the group of items to be tested and specifying characteristics derived from the Bill of Materials stored in the data base of the computer which controls the flow of product testing.

Level 1 contains customer specials which are produced by additional selective testing of Level 2 or 3 standard products for some customer specific parameters. This level may also contain specials which are made from other specials and thus can be divided into several sublevels. Level 1 can be considered optional, and the processing in Level 1 will not be discussed in detail herein.

Level 0 contains finished goods which are the final branded versions of the products manufactured at one of the lower levels. No additional testing is required to move a product to this level. Due to branding variations, however, several finished products can be made from the same unbranded component.

While this concept of levelling serves to group together products which interact in the planning process and which are planned using the same logic, it also retains the original concept of low level standard MRP code, namely that products are always at a higher level than their components. Level 0 items are always made from Level 1 2 or 3 items. Level 1 items are always made from other Level 1 items or from Level 2 or 3 items. Level 2 is always made from Level 3 which is always made from Level 4.

FIG. 1 presents a simplified exemplary illustration of the type of product-level hierarchy described above (omitting any showing of the "specials" level). It will be seen that the initial group of untested items referred to as B are assembled from other components (not identified in the diagram but listed in the Bill of Materials). A planned first pass test will be expected to produce products graded as J, K, L (identifying prescribed electrical characteristics), in pre-established ratios in accordance with experience data stored in the data base. Rejects will be discarded.

If it is desired to produce a number of items of military grade U, a group of grade L products will be tested. That is, military grade U requires the superior electrical characteristics of commercial grade L, but must be capable of such performance under military environmental conditions. Upon such testing, some of grade L will pass, to be designated grade U. Others will

pass the military, requirements but with lower electrical characteristics, and be designated grades S or T (as indicated by the dotted lines). Still others will fail the military requirements, but still pass the electrical tests to be accepted back as grade L or as grades J or K. The final graded products can be branded and transferred to "finished goods".

In standardizing the product flow through test in such a level-oriented fashion, it is desirable to close out all test lots to a stock room after each level of testing is completed. This improves inventory control and lot traceability, and makes it possible to capture the actual grade distribution of all intermediate tests as inventory transactions which later become the basis for up-dating the grade distribution data base.

The grade distribution data base (which is not used in standard MRP planning) contains the planned yield distribution to all grades resulting from testing a component. If it is desired to plan a test for item SC from component JC, the yield must be known. For example, if 100 units of SC are needed and the yield is 50%, the demand for component JC will be 200. In addition, however, the present system takes into account the yield and identity of the by-products of this test. In this case there might be a 45% yield to JC (components which failed the SC test) while the other 5% are scrap, destroyed in the testing process.

A typical grade distribution data base entry might include the following:

	PRODUCT	COMPONENT	BY-PRODUCTS		
Item	TC	KC	SC	KC	JC
Yield	.40		.10	.30	.10

The grade distribution data base defines for all items in Levels 1, 2 and 3 the yield to that item from its component. Also defined is the identity of all by-products (which may be on the same or lower levels), and the yield to each. The system updates these grade distributions periodically through batch analysis of the historical stockroom grading inventory transactions.

The present system also incorporates a function referred to as downgrading, which will be described in detail hereinbelow. In order to plan downgrading, it is necessary to know the hierarchy of grades, that is, which standard grades within Levels 2 and 3 can be substituted for lower grades with no additional testing. This information also is stored in the grade distribution data base.

Referring now to FIG. 2, it will be seen that the overall system flow diagram breaks into three main routines. The diamond 2 represents the select routine decision, to prepare the data for run, to run the system, or to select a report. The letters in circles represent return points for the system loops.

Block 1 indicates the portion of the system which prepares the raw data for processing including the Bill of Materials (Block 1A), grade distribution data (Block 1B), sales demand (Block 1C) and inventories (Block 1D). This information is stored in data base files and used by the operating part of the system when it is being run to prepare the plan for assembly and testing.

In a standard Bill of Materials used in prior art MRP systems, products are assigned to levels based on their relationship to other products. Finished goods are assigned to the first level (usually called Level 0); their components are assigned to Level 1 etc., down to the

lowest level. The general rule is that an item's level number is equal to the highest level number of any item it is used in plus one. Standard MRP processes items by level starting at 0 thus assuring that for any item all demand for it as a component of other items will be present at the time that item is processed.

In the Bill of Materials used in the present system, items are not assigned to relative levels but rather to absolute levels based on their testing classification, namely: Finished Goods, Customer Specials, Second Pass, First Pass and Untested Product. This classification does not violate the standard rule that an item comes after any item of which it is a component, but it does assign to the same level items which are to be planned together because of by-product or co-product relationships. This assignment is also advantageous because it defines exactly the proper sequence for planning the products and makes it possible to match each level or set of products to the appropriate modified logic. Block 1A represents this Bill of Material data base information. A simplified example of the basic type of information contained in the Bill of Materials is represented by the diagram of FIG. 1.

Block 1B represents a series of tables stored in the data base giving the percentage yield to by-product or co-product items which result from testing a component. (A by-product is defined as: a usable unbranded grade that has failed the test for a desired item. A co-product is defined as: one of several related usable unbranded grades classified by the same test.)

Block 1C represents projected sales demand such as that based on marketing sales estimates. Block 1D represents the inventory data. The information of Blocks 1C and 1D would be found in one form or other in any MRP system.

Blocks 9 through 12 represent output reports generated by the system. Such reports would be generated in any standard MRP system.

Blocks 3 through 7 represent the logical processing procedures which are at the heart of the present system. The processing steps for each of these blocks is described in expanded detail in FIGS. 3-7, and will be explained in more detail hereinafter. It may be noted, however, that each block 3-7 represents a single level as defined above. The processing at each of these blocks applies the appropriate logic to the items at its level and all items in a level are finished before the next level is begun. The end result is a projected assembly and test plan for a relatively long time frame typically subdivided into weeks, and providing projections for all products indicating the number of each to be assembled and to be input into the various test levels.

FIG. 3 shows the data processing logic for the finished goods level (Level 0) represented by Block 3 of FIG. 2. This logic is standard MRP logic and is included only for completeness. NO testing, yielding or grading is involved in the brand/final-QC process. Independent demand is introduced at Block 13 as a combination of backlog, forecast and safety stock for all standard and special finished goods. This demand, net of finished goods inventories, becomes dependent demand for those items' unbranded components and all items are processed before proceeding to Level 1.

The resulting requirements are not part of the production schedule for test, but rather a transfer plan which drives the test schedule, serving as the master schedule. This transfer plan is exploded through the test planning logic sequences to produce a testing schedule

for Levels 1, 2 and 3 and a Level 4 requirement for assembly. Test capacity constraints for the test schedule levels and material constraints in assembly can then be identified if necessary.

Blocks 14, 15 and 18 execute normal inventory netting, lead time adjustments and shrinkage calculations. Block 16 allows the schedules to modify the system generated schedules using a standard industry technique known as a firm planned order. Blocks 17 and 19 create reports which can be printed as normal MRP outputs and Block 20 saves the exploded demand for components which will be part of the input at the lower levels.

FIG. 4 shows the data processing logic for the customer specials level. Block 21 assembles the demand, yield, inventory, and Bill of Materials data for all of the items which have been assigned to this level. Although the processing for this level is optional, depending upon the nature of the industry, it does include process steps which are used at lower levels, and therefore certain parts of the customer specials procedure will be described herein.

Block 22 selects the special products to be processed first. This is necessary because some specials may have other specials as their component or as their by-product. In such a case it is necessary to process the higher level item first and then the component or by-product item. Block 22 first selects all of the specials which are neither components nor by-products of other items for processing at this time. After these items have been completed and eliminated from the list, the next execution of Block 22 will select the next level of components for processing. This process is repeated until all of the specials have been processed.

For those items selected in Block 22 the processing continues in Block 23. Here the inventory is subtracted from the demand as in normal MRP.

Block 24 introduces part of the new logical processing which deals with by-product planning, specifically that of netting out by-products as an inventory supply. This is necessary at the specials level because the by-products of some specials are other specials. Block 23 subtracted inventory on hand from the first time period's demand. Block 24 subtracts out by-product supplies from the demand in the period when they will be available. (Note: The by-product supply data comes from other portions of the system, as will be described later.) For example, if there is demand of 500 units each for two periods and an on-hand inventory of 50 units and by-product supplies of 100 units each period, the net demand for each period would be calculated as follows:

	Period 1	Period 2
Demand	500	500
Inventory	50	
By-product	100	100
Net demand	350	400

This computation has the effect of covering demand with an alternate supply of material (i.e. the by-product from prior planning calculations) before calling for more production. Standard MRP does not do this.

Blocks 25 and 26 concern firm planned orders and standard MRP reports, which are conventional, and therefore will not be described.

Block 27 introduces the remaining part of the new by-product planning logic—the calculation of the supply of by-products resulting from the tests which are

being planned. (This result may be used in Block 24 in a later processing sequence for another item at this level.)

This by-product computation requires two steps. Step one computes the quantity of the component which must be tested to cover the demand. This is done by dividing the net demand by the yield. For example, 100 units of product A with a 50% yield requires the testing of 200 units of component B.

Step two computes the by-products. For example, if the grade distribution data base for product A above shows two by-products B (yield 20%) and C (yield 25%), then block 27 generates a supply of 40 units of B and 50 units of C for the same period as the projected demand of 100 units of A. Note that B can be both a component and a by-product of the same test, but because the component demand is advanced to an earlier period by the length of the testing process the demand for 200 units of B as a component will precede the supply of B as a by-product.

Blocks 28 and 29 continue the normal MRP logic. Block 30 checks to see whether or not there are additional levels of specials to be processed. If so, the program branches back to Block 22 as indicated by the letter C in a circle. If not, it proceeds to Block 31 which saves the exploded demand and any by-product supplies for the second or first pass grades, then passes control to the second level on FIG. 4. If no specials processing is included in the system, the finished goods processing of FIG. 2 will be followed immediately by that of FIG. 4.

FIG. 4 specifies the processing of second pass grades. Here, unlike the specials level, all of the second pass grades of the same product family are processed together, one week at a time. This is necessary in order to correctly account for by-product supplies to a lower grade from a higher grade and to allow for downgrading. This simultaneous processing of multiple products is a significant departure from standard MRP and highly important in the functioning of the new system.

Block 32 functions in the usual way to assemble the data for the products at this level. This data comes from the data files of Block 1 (FIG. 2) and comprises inventories, projected demand, grade distribution and the Bill of Materials. Blocks 33 and 34 control the sequence of processing one week at a time and, within each week, one grade at a time, proceeding from the highest to the lowest. Block 35 nets out the inventory using standard MRP logic. Block 36 represents the new logical processing used to net out by-product arrivals which may either come from the specials level or from higher second-pass grades. The calculations are exactly the same as described above with respect to Block 24 and therefore will not be repeated here. Block 37 branches to the next lower grade if there is no requirement for the current one.

Block 38 introduces the logical processing for the downgrading procedure discussed previously. Because grading is a natural phenomenon independent of demand its distribution does not necessarily match the distribution of demand. This may result in surplus inventories of higher grades and shortages of lower grades. There are circumstances where it makes economic sense to downgrade these higher grades to cover the shortages rather than produce more of the lower grades. The processing procedures in Block 38 decide whether or not to downgrade any surplus inventory of higher grades in the same family to cover the require-

ment of the current grade. This downgrading is a very significant departure from standard MRP.

Any downgrading decision involves three questions, asked from the point of view of the receiving item, i.e. the lower grade:

1. Is there a shortage—if not, no downgrading is necessary.
2. Is there a higher grade with a surplus—if not, no downgrading is possible.
3. Should the higher grade be downgraded—this decision involves a comparison of the relative demands and yields of the two grades.

For example, in the case of three grades L, K and J where L is the highest grade and J the lowest, the procedure first determines whether there is a shortage of K and if so, whether or not to downgrade L. Then a shortage of J is determined and, if so, a decision is made whether or not to downgrade first K and then L.

1. A shortage exists if the demand for a grade in the current period is larger than the total of the beginning inventory and the by-product supply.

2. A surplus exists if the demand for the grade in the current period is smaller than the total of the beginning inventory and the by-product supply.

3. The formula for determining whether to downgrade has three parts:

(a) For the higher (donor) grade compute:

$$\sum_{I=W}^N (D_I - B_I) \times (N - I + 1) / Y$$

where

I is a counter beginning with the current week W

N is the total number of periods (weeks) in the plan

D_I is the demand for the grade in the period

B_I is the by-product supply in the period

Y is the yield to the grade (N - I + 1) is a factor which gives greater weight to near term demand than future demand.

It can be seen that the outcome from this formula will be increased by large demand, near term demand, small by-product supply, and low yield. The outcome will be decreased by small demand, far future demand, large by-product supply and high yield.

(b) Compute the same formula for the lower (receiver) grade.

(c) Compare the results. If the result for the donor is lower than for the receiver, the higher grade should be downgraded.

For example:		3	4	5	6	7	8 (N)
Period (I)							
L	Beginning Inventory	200					
	Demand	50	150	100	300	500	100
	By-Product Supply	100	100	100	100	100	100
K	Beginning Inventory	0					
	Demand	600	500	800	300	900	700
	By-Product Supply	50	100	50	40	200	300

-continued

For example:		3	4	5	6	7	8 (N)
Period (I)							
(N - I + 1)		6	5	4	3	2	1
1.	K has a shortage of 550 in week 3 (600 - 50)						
2.	L has a surplus of 250 in week 3 (200 + 100 - 50)						
3.	(a) Donor formula = (6 × -50) + (5 × 50) + (4 × 0) + (3 × 200) + (2 × 400) + (1 × 0) / 3 4500						
	(b) Receiver formula = (6 × -550) + (5 × 400) + (4 × 750) + (3 × 260) + (2 × 700) + (1 × 300) / 4 27200						
	(c) Since (b) is larger than (a) the surplus of 250 L should be downgraded to K leaving 300 more K to make from new testing.						

Once the decision to downgrade has been made, then the quantity to downgrade is the surplus of the higher grade or the net requirement of the lower grade, whichever is less. Downgrading should only take place when the demand for the higher grade is low or far enough in the future (relative to the demand for the lower grade), and the yield (the ability to produce more of the higher grade later) is reasonably high.

The net effect of downgrading is to minimize the cost of carrying the excess inventory and the cost of current testing.

Block 39 calculates the planned production to cover the balance of the current grade (i.e. after the downgrading transfer) using normal MRP logic. Block 40 applies firm planned orders and Block 41 prepares the schedule report, as discussed previously.

Block 42 calculates by-product supplies to lower second-pass grades or to first-pass grades, using the same calculations as described above with respect to Block 27. Blocks 43, 44 and 45 perform the same standard MRP functions as Blocks 28, 29 and 31, referred to above. All of the components are on the first-pass level.

After all grades have been completed for one week, the processing proceeds to the next week. After all families on the second-pass level have been completed, the processing proceeds to the first pass level of FIG. 6.

The first-pass processing is very similar to the second-pass processing. The sequence proceeds from week to week, and from high grade to low grade, governed by Blocks 47 and 48. This sequence provides the basis for the co-product planning at the first-pass level. Blocks 49-52 represent logical processing steps identical to those of Blocks 35-39 at the second-pass level.

Once the net requirements for all grades in the family have been calculated, including the projection of inventory downgrades as described above, co-product processing is carried out to determine the proper quantity of the component common to all grades which will yield enough of each grade to cover all of their requirements. That is, it is necessary to compute the issue quantity of the common component which will cover the yield to the most limiting of the grades in the set. Most simply this means dividing the net demand for each grade by its yield and issuing the largest of the three resulting quantities.

Set forth below is an example of such co-product planning using three grades (L, K and J) where L is the highest grade.

Item	Demand	Yield	Test Qty.	Output
L	100	.20	500	200 L
K	300	.30	1000*	300 K
J	450	.50	900	500 J

The demand for the three grades are 100, 300 and 450 respectively, and their yields are 20, 30 and 50%. The processing divides the net demand for each grade by its yield and issues an order for the largest of the three resulting quantities. In this case, it would calculate an issue of 1000 which will just cover the requirement for K, and result in an extra 100 L's and an extra 50 J's.

It may be advantageous also to employ downgrading in co-product planning. By using surplus yield of higher grades to cover lower grades, all the demand can be satisfied while testing fewer parts, using less capacity and generating less inventory.

Set forth below is a table giving an example where downgrading is applied fully to the same three products with the same initial demand and yield as listed above.

Item	Demand	Yield	Test Qty.	Output	Downgrading
L	100	.20	500	170 L	-70
K	400	.50	800	255 K	+45
J	850	1.00	850*	425 J	+25

In this case, however, the computation takes into account cumulative demand and cumulative yield for each successive grade. For example, the demand for K is cumulative of L and K ($100+300=400$) with yield of L and K combined ($0.20+0.30=0.50$) before dividing to determine the limiting test quantity. While 500 will cover the L alone and 800 will cover L and K combined, 850 are needed to cover all three grades. This is the amount to be tested—less than the 1000 needed in the no-downgrade option. This yields an extra 70 L's, and 45 of them will be used to cover the K deficit (300 were needed) and 25 to cover the J deficit (450 were needed).

Alternatively, the downgrading algorithm (see above) can be used to decide which grades should be downgraded to which. These downgrade combinations can then be planned together by adding their demands and yields and computing the issue quantity which covers the combination. After this downgrading adjustment has been made, the planned issue quantity is the largest needed to support any of the combinations or uncombined grades in the set. It will be found that use of the downgrading algorithm for co-product planning will produce a result intermediate the no-downgrading case (1000 units) and the full-downgrading case (850) units, which may be a more satisfactory decision from a practical point of view.

The resulting yield to each grade and the appropriate downgrade transactions then are computed, followed by the ending balances which will roll into the following period.

Block 55 applies firm planned orders as described above. Blocks 56-59 are the same as Blocks 41, 43 44 and 45 at the second-pass level.

After all first-pass families have been processed for all weeks, there results a test schedule for Level 3 items and dependent demand for Level 4, i.e. the components of untested material. This net requirement is passed on to the standard MRP logic at the lower levels of the product structure, as indicated in FIG. 7.

To summarize, processing procedures of the disclosed system are carried out at five levels of the product structure as set forth below:

0. FINISHED BRANDED DEVICES

This module first calculates a schedule for receipts into finished goods by netting out existing on-hand inventories against demand (backlog and forecast) and inventory targets. Against this requirement it applies inventories in factory bonded stock and final QC and offsets by lead time to compute the schedule for branding and transfer to final QC. This schedule is saved and printed and can be modified with "firm planned orders" which affect the schedule at lower levels and can be saved to apply to subsequent runs of the system.

These requirements net of in process inventories at brand are captured as issues from unbranded stock of the item's component and will be used as demand for that component at the appropriate level (1, 2 or 3).

Exception reports highlighting current and projected inventory imbalances are generated at this and all subsequent levels.

1. PRODUCTS SPECIALLY SELECTED FROM OTHER PRODUCTS

In this level, items and their components can both exist on the same level of the product structure. Items are processed in successive batches so that the requirements for an item's component are generated before that component is processed.

The processing consists of the normal calculations, i.e. demand (for the Level 0 or other Level 1 item) net of existing unbranded inventory becomes the scheduled test completion for the item. That requirement, divided by the yield, offset by a lead time, and net of existing work in process inventory is the demand for the item's component which is another Level 1 or a Level 2 or 3 item.

This level (a) considers by-product arrivals of the item from other Level 1 tests before computing additional tests for the item and, (b) calculates by-products to other products (on level 1, 2 or 3) which will result from the test for this item.

2. SECOND PASS TESTED STANDARD GRADES

This level employs the by-product computations to consider arrivals from Level 1 and generates by-product calculations for other Level 2 and 3 items.

This level also employs downgrading to apply excess inventories of higher grades to demand for lower grades of the same product family before scheduling new testing of the lower grade. This level therefore is processed one week at a time and within each week from the highest to the lowest grades.

3. FIRST PASS STANDARD GRADES

This level includes consideration of by-product arrivals from Levels 1 and 2 tests and inventory downgrading logic introduced in Level 2.

In addition, it calculates the correct quantity of the untested component which will yield the proper quantities of each of the several first pass grades to meet their demands. In so doing it may determine when and in what amounts higher grades yielded from the tests should be downgraded to meet the demand for lower grades thus reducing the total requirement for testing.

4. UNTESTED MATERIAL

This level uses standard MRP logic to determine the required receipts of untested material to meet the Level 3 requirements.

Although a preferred embodiment of the invention has been disclosed herein in detail, it is to be understood that this is for the purpose of illustrating the invention, and should not be construed as necessarily limiting the invention since those of skill in this art can readily make various changes and modifications thereto without departing from the scope of the invention as reflected in the claims hereof.

What is claimed is:

1. In combination in a system for controlling the flow of a plurality of products and components thereof through the assembly and testing procedures of a manufacturing operation comprising plural sequential process steps, including the development of an assembly and test schedule for the successive periods of a forward planning time frame, and wherein the products are arranged in families comprising different grades meeting respective performance specifications corresponding to different quality ratings, said system including data file means for storing information representing the inventory status of each product grade and the expected demand for each product grade during each of said periods, said system being of the type which is arranged to make calculations with respect to such stored information so as to determine the net requirements for each product grade during each of said successive time periods;

that improvement to such system wherein said data file means stores grade distribution information giving the co-product yield for each grade of each product family from testing the common component of the product family; and

said system further comprising means to calculate, based on said grade distribution information, the number of such common components to be tested for each of said periods in order to meet the net requirements projected for the co-product grades for those periods.

2. A system as in claim 1, wherein said data file means further includes stored by-product information giving for at least certain of the products the identities and percentages of items which will be produced as by-products when testing for any particular one of said certain products selected;

said system further including means to calculate and to store for a projected test of such particular product the number of each such by-products to be produced; and

means responsive to such stored by-product number data for updating the effective inventory of a corresponding item when the net requirements for that item are being developed subsequently.

3. A system as in claim 1, including means operable in making said calculations with respect to said stored information, for determining the net requirements for at least two grades of the same product family, to downgrade a surplus portion of the higher of said grades to the lower grade so as to minimize the need for additional assembly or testing to produce the lower grade.

4. A system as in claim 3, wherein said downgrading decision is made by taking into account the relative demands and yields of the two grades.

5. A system as in claim 1, including means operable, in calculating said number of common components, to decide whether to downgrade a portion of a higher grade to cover the requirements for a lower grade of the product family.

6. A system as in claim 5, wherein the downgrading decision is based on the relative demands and yields of the two grades.

7. In combination in a system for controlling the flow of a plurality of products and components thereof through testing procedures of a manufacturing operation comprising plural sequential process steps, said system including means for developing an assembly and test schedule setting forth the planning goals for each of the products for the successive periods of a forward planning time frame, each of the products being required to meet respective performance specifications to be designated as such product, said system including data file means for storing information representing the inventory status of each product and the expected demand for each product during each of said periods, said system being of the type which is arranged to make calculations with respect to such stored information so as to determine the net requirements for each product;

that improvement to such system wherein said data file means stores information giving for at least certain of the products the identities and yields of items which are expected to be produced as by-products when testing for any particular one of said certain products, said items corresponding to other products being manufactured in said manufacturing operation;

means to calculate, based on said by-product information, the number of such by-product items which will result from testing for such particular product; and

means for adjusting the effective inventory data for said other products in accordance with said by-product number data.

8. In combination in a system for controlling the flow of a plurality of products and components thereof through the assembly and testing procedures of a manufacturing operation comprising plural sequential process steps, including the development of an assembly and test schedule for the successive periods of a forward planning time frame, and wherein the products are organized in families comprising different grades meeting respective performance specifications corresponding to different quality ratings, said system including data file means for storing information representing the inventory status of each product grade and the expected demand for each product grade during each of said periods, said system being of the type which is arranged to make calculations with respect to such stored information so as to determine the net requirements for each product grade during each of said successive time periods;

that improvement to such system comprising means operable in making said calculations with respect to said stored information, for determining the net future requirements for at least two grades of the same product family, to downgrade a surplus portion of the higher grade to the lower grade and to adjust the effective inventory of both grades.

9. A system as in claim 8, including means responsive to the relative demands and yields of the two grades for making the decision as to whether to downgrade.

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10. A system as in claim 9, wherein said responsive means comprising means to calculate, for the two products, the following:

$$\sum_{I=W}^N (D_I - B_I) \times (N - I + 1)/Y$$

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where

I is a counter beginning with a current week W

N is a total number of periods (weeks) in a plan

D_I is a demand for the grade in the period

B_I is a by-product supply in the period

Y is a yield to the grade

(N-I+1) is a factor which gives greater weight to near term demand than future demand.

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- [54] **PROJECT CONSTRUCTION WITH DEPICTION MEANS AND METHODS**
- [76] Inventor: **Stuart Ockman**, 210 Copples La., Wallingford, Pa. 19086
- [21] Appl. No.: **641,160**
- [22] Filed: **Aug. 16, 1984**

Related U.S. Application Data

- [63] Continuation-in-part of Ser. No. 559,706, Dec. 9, 1983, abandoned.
- [51] Int. Cl.⁴ **G06F 15/40; G06F 15/60**
- [52] U.S. Cl. **364/518; 364/512; 364/468; 434/108**
- [58] **Field of Search** **364/518, 521, 512, 505, 364/400, 401, 419, 468; 334/108; 239/89 R, 69**

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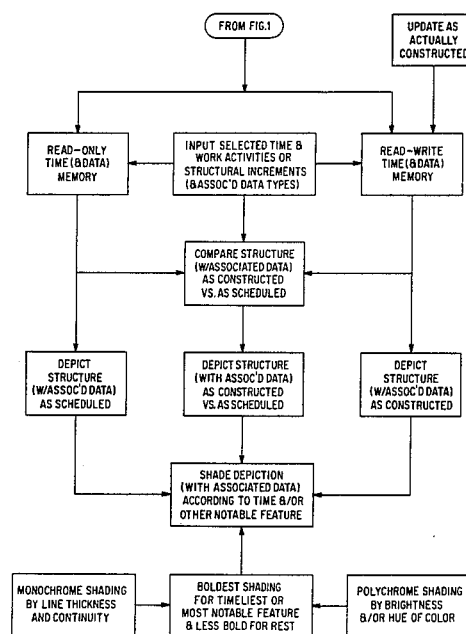
It Can't Fasten, Cut, Hammer, Drill, Dig, Bend, Level or Lift, IBM Sales Information Advertisement.

Primary Examiner—Errol A. Krass
Assistant Examiner—Kevin J. Teska
Attorney, Agent, or Firm—Charles A. McClure

[57] **ABSTRACT**

Construction of projects such as buildings, bridges, dams, industrial plants, means of transport, or the like with the aid of means and methods for incremental depiction of such buildings, etc. Complex structural features of the construction project are divided into discrete increments and characterized by ordered work activities essential to their construction, then are depicted as in construction drawings to assist in supervision of the project itself. Such characterization and depiction enable urgencies to be emphasized, as by shading structural features due to be started or finished by some given time in the construction schedule differently from those not due until later—or to have been completed earlier. Likewise, during actual construction, deviations from the project schedule can be distinguished readily, as by depicting structural features due (or overdue) in boldest shading and by diminishing the boldness of shading the more time is available within which to complete subsequent features. The boldness of monochromatic shading may be by line thickness and/or continuity, whereas for polychromatic representation it may be by hue and/or brightness of color. Storage and retrieval of structural information, work activities, and other data are conveniently machine-assisted, as by an appropriately programmed general-purpose computer or by an otherwise similar computer especially dedicated to such purpose.

34 Claims, 10 Drawing Figures



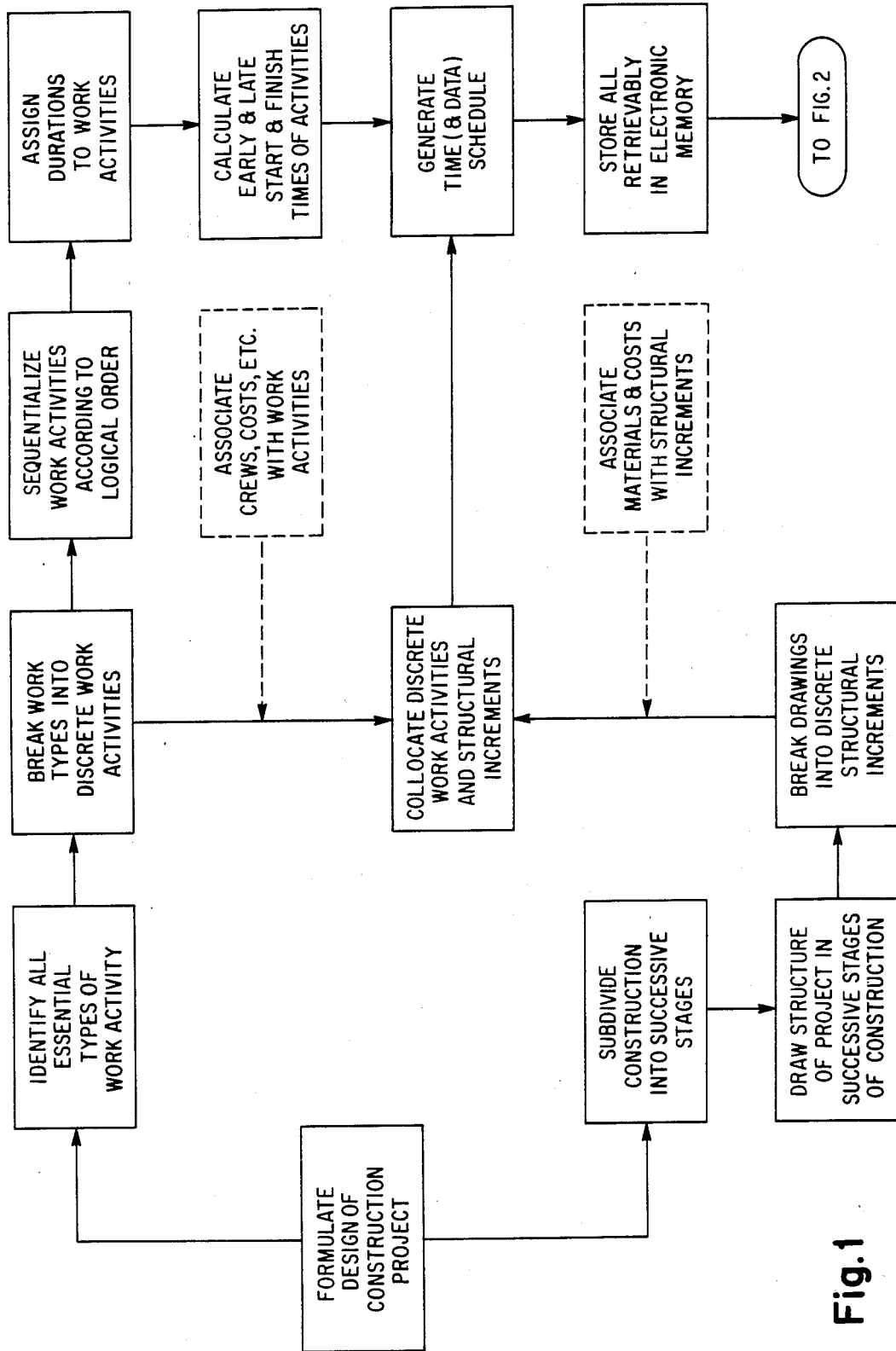
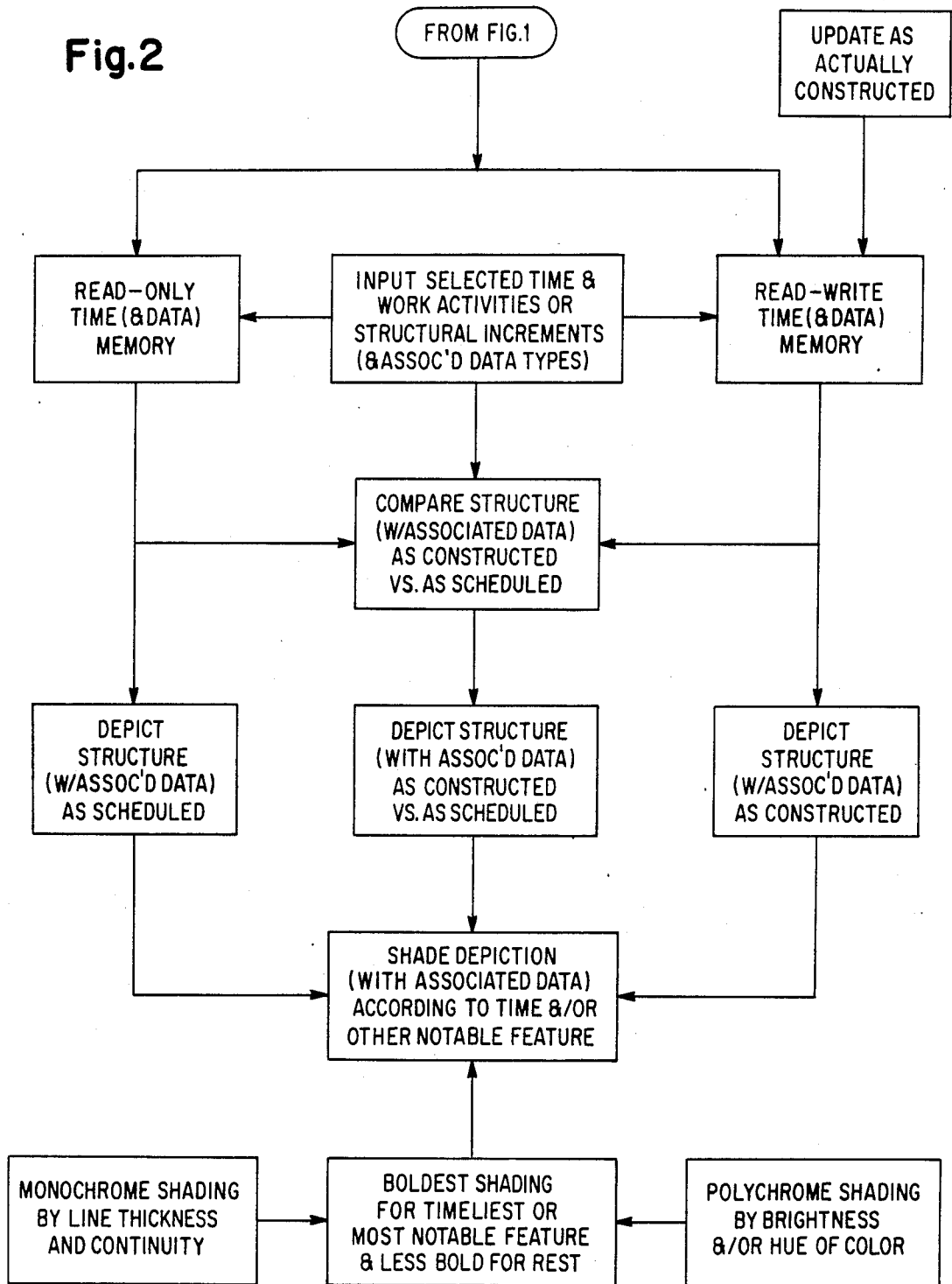


Fig. 1

Fig.2



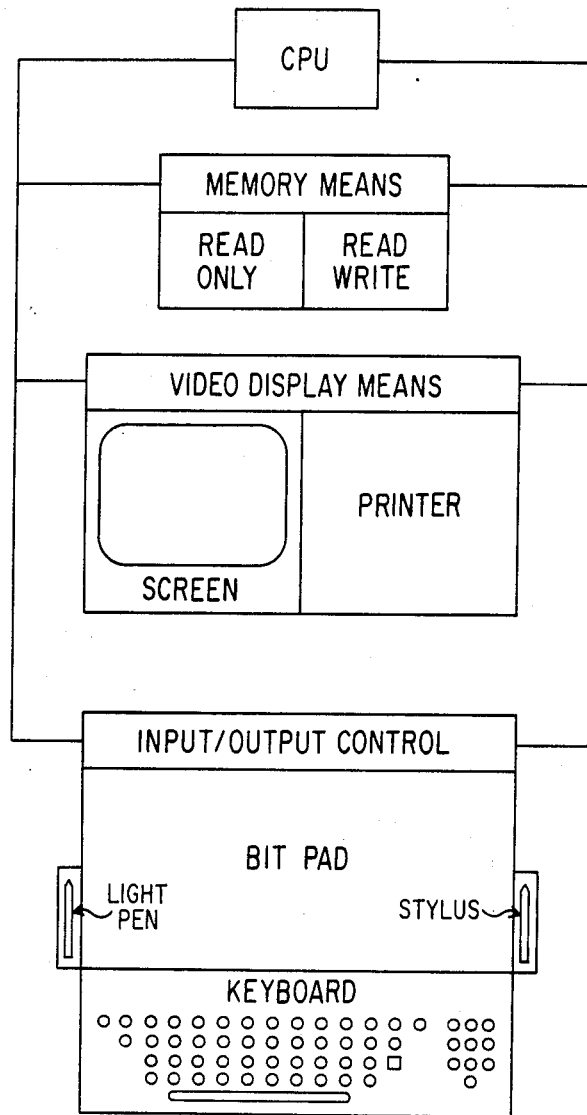


Fig. 3

FIG. 3A

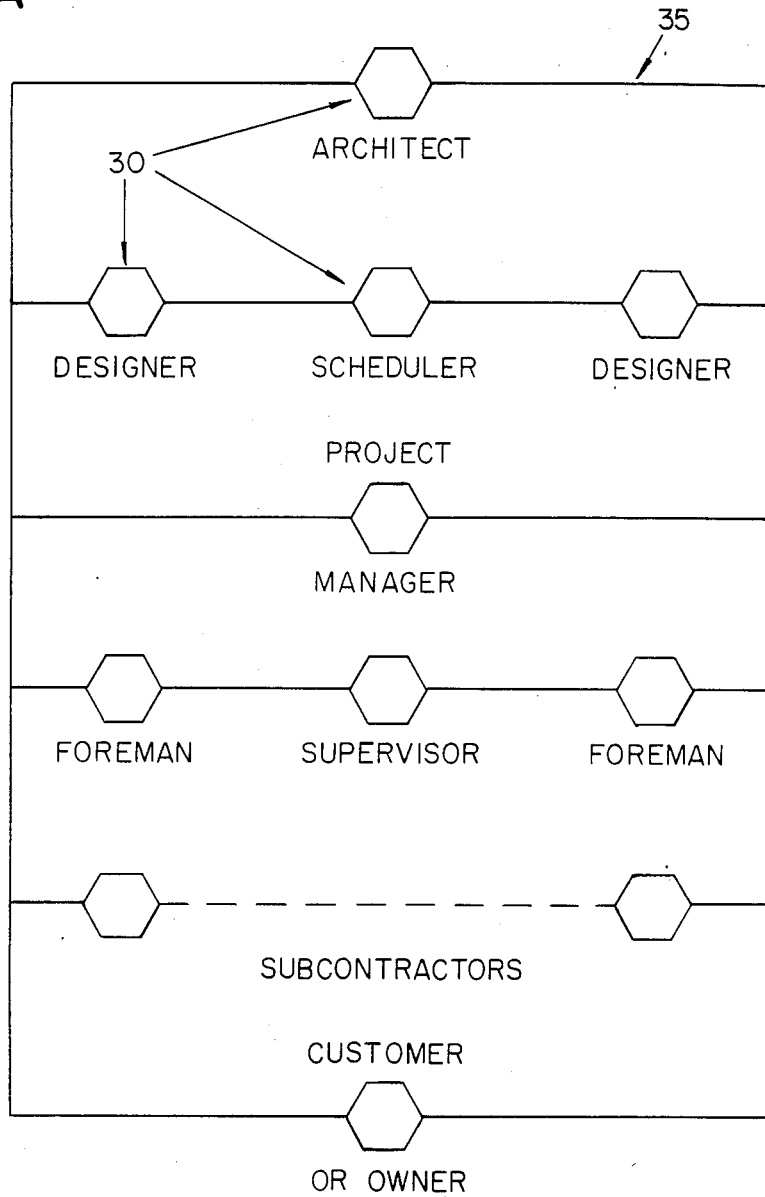
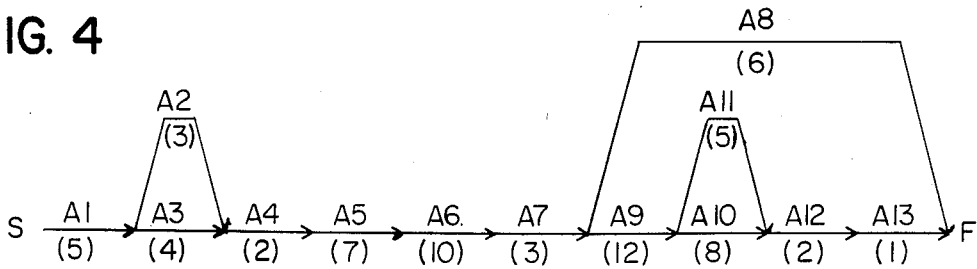


FIG. 4



LEGEND

E = EARLY
 L = LATE
 S = START
 F = FINISH

EF / ES LS \ LF

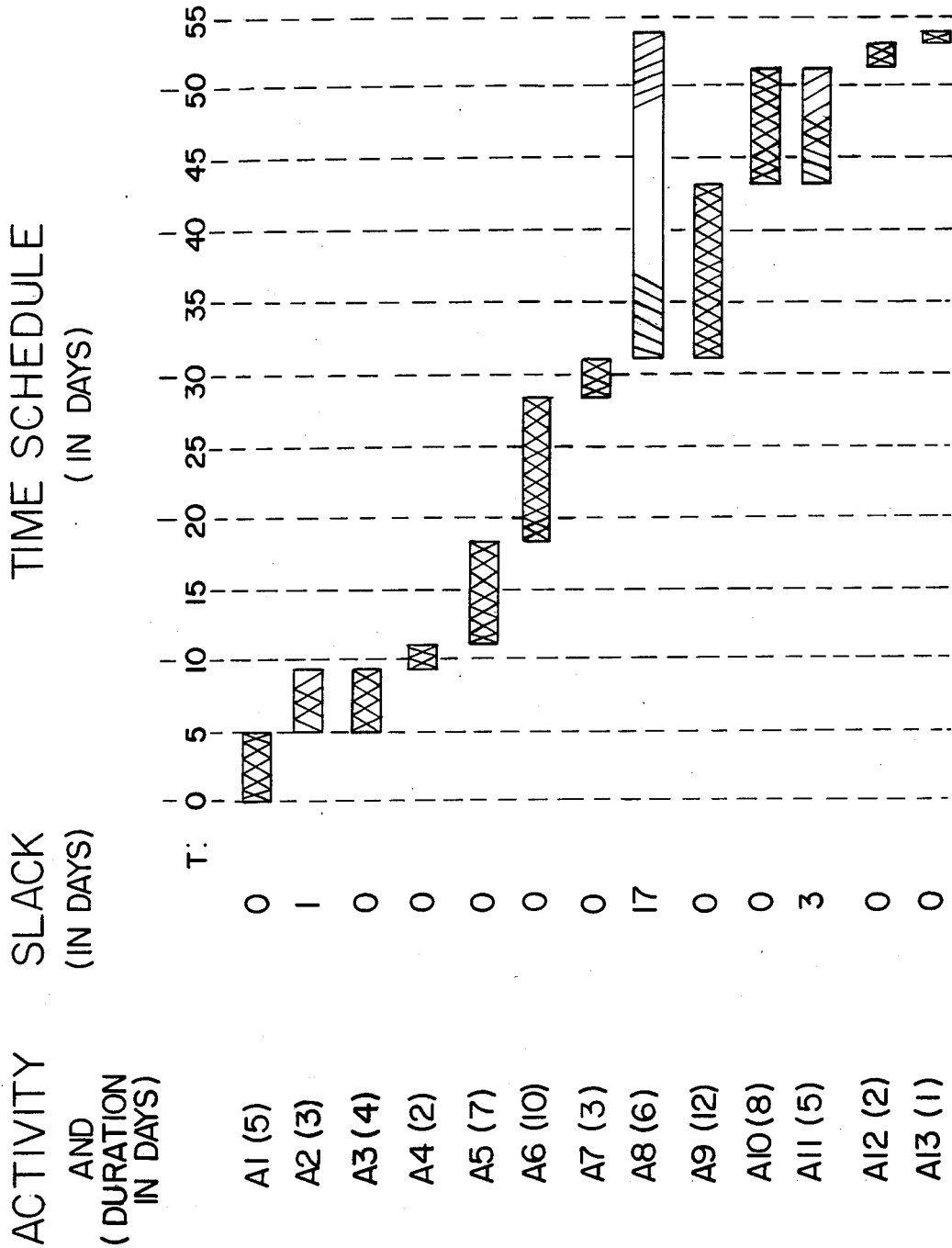


FIG. 5

FIG. 6

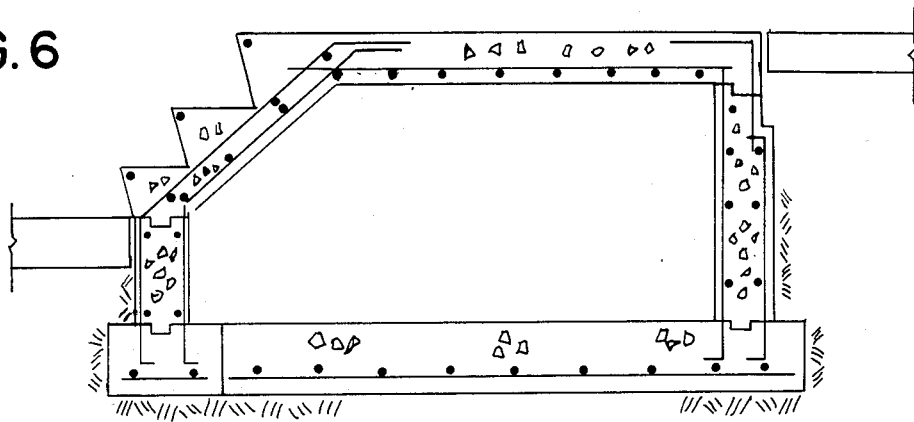


FIG. 7

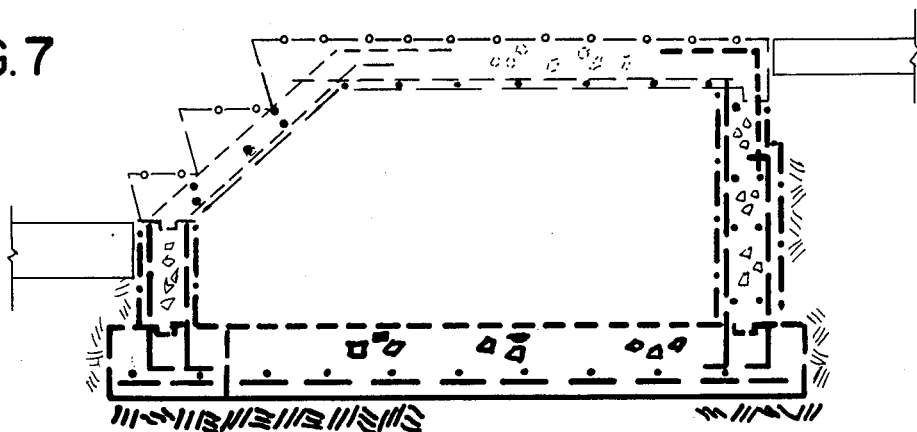
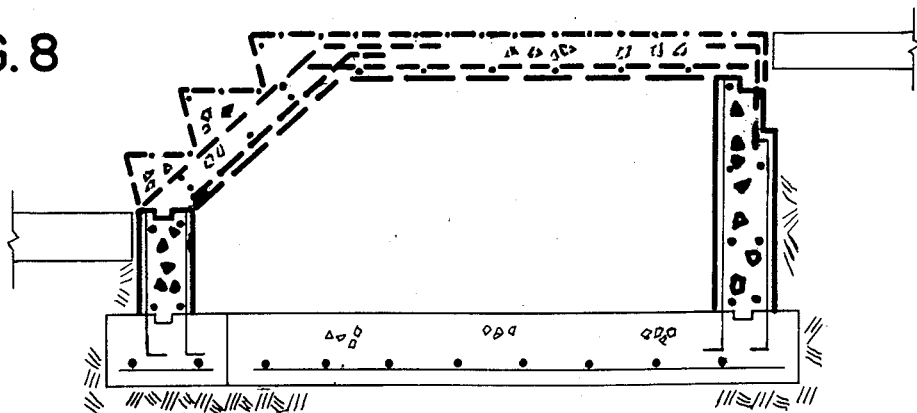


FIG. 8



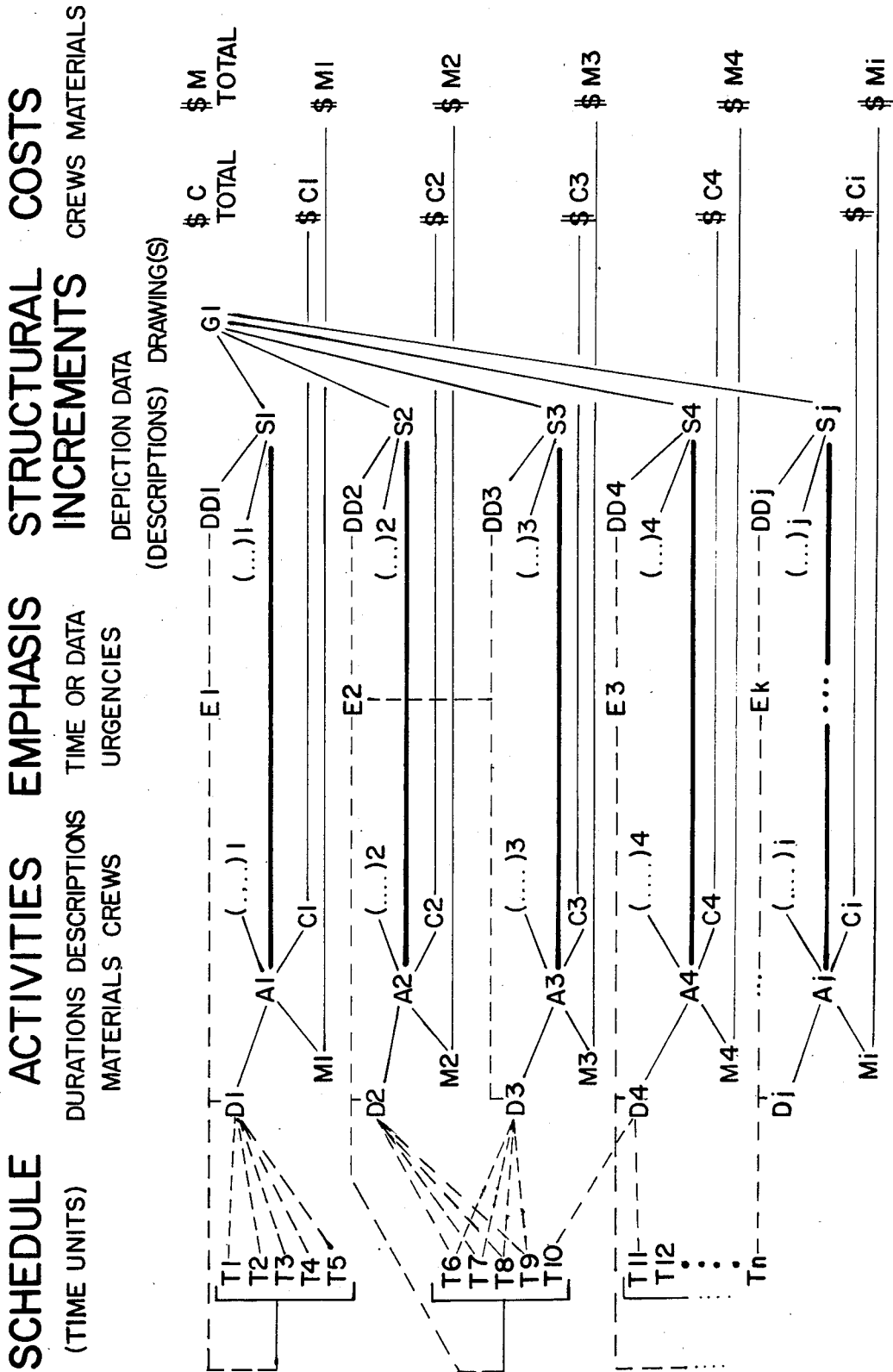


FIG. 9

PROJECT CONSTRUCTION WITH DEPICTION MEANS AND METHODS

This is a continuation-in-part of my similarly entitled copending application, Ser. No. 559,706 filed Dec. 9, 1983 and subsequently abandoned.

This invention relates to novel visual depiction of structural and optional indication of related features of construction projects in their planned and actual construction, and is applicable to construction of buildings, bridges, dams, industrial plants, means of transport, and the like.

Complex construction projects comprise too many components and involve too many structural and other interrelationships to be kept in mind without visual aids; such aids include drawings of the intended results, graphical and tabular schedules of the order in which components have to become available and be put into place relative to one another, and budgets, for example. However, it is difficult to assemble the pertinent information in an intelligible manner, especially for such diverse audience types as artisans, architects, contractors, engineers, financiers, and others. Also, in international construction projects, diversity of language is often a barrier to common understanding of what has to be done. Ordinarily, whenever a project requires a large set of drawings, it (and they) are divided into manageable subsets, often with the aid of hierarchical levels of detail in the various drawings, but such subdivided drawings require additional notations to render their interrelationships fully intelligible to individual viewers.

A primary object of the present invention is to provide novel visual depiction in representation of structural and related information about construction projects, especially complex ones.

Another object of this invention is to render intelligible at a glance the scheduled status of selected increments of structure at any given stage in a construction project.

A further object of the invention is to focus attention, in visual pictorial representations of projects undergoing construction, upon actual status of structural increments versus scheduled status.

Yet another object of this invention is to depict structural features of a construction project, whether scheduled only or actually being constructed, in such manner as to convey information about related non-structural features.

A still further object of the invention is to provide in the management of a construction project—whether scheduled only or in construction—a data base comprising indicia of discrete increments of structure, collocated indicia of sequentialized work activities, a superimposed time schedule for the project, and related data.

Other objects of the present invention, together with methods and means for attaining the various objects, will be apparent from the following description and the accompanying diagrams, which are presented by way of example rather than limitation.

FIG. 1 is a block diagram of steps taken in practicing the present invention;

FIG. 2 is a block diagram of further steps taken in practicing this invention; and

FIG. 3 is a largely schematic diagram of apparatus useful in practicing the invention; and

FIG. 3A is a schematic diagram of a network of workstations, each comprising apparatus such as is shown in the preceding view.

FIG. 4 is a project-scheduling network diagram of activities for a portion of a building construction project; and

FIG. 5 is a construction project schedule corresponding to the scheduled activities network diagram of the preceding view.

FIG. 6 is a substantially conventional construction drawing of the same scheduled portion of a building construction project;

FIG. 7 is a similar drawing of the same building portion, but shaded in accordance with the scheduled start of construction; and

FIG. 8 is a similar drawing of the same building portion, but shaded according to the project status well into the schedule, with part of the scheduled construction overdue for completion.

FIG. 9 is a schematic diagram of interrelationships between and among data elements useful according to the present invention.

In general, the objects of this invention are accomplished, in building and similar construction projects, by depicting selected structural features characterized by times of the work activities essential to their construction, plus (optionally) additional related data. Such characterization highlights urgencies, as by depicting structural features due to be started or to be completed by any given time in the construction schedule, more boldly than structural features not then due. At any given time during actual construction, structural features due (or overdue) for completion can be emphasized in similar manner, with little or no resort to verbal representation, in a mnemonically helpful way that is advantageous in reducing time, effort and cost or taking other remedial measures.

In summary detail, one is enabled to attain those objects by a sequence of steps including identifying, in advance, discrete increments of structure from start to finish of a given project, depicting the structural increments as in construction drawings, identifying discrete work activities essential to constructing the respective structural increments, assigning durations to the given activities, sequentializing the work activities among themselves, calculating start and finish times for the respective activities so sequentialized, preparing a time (and data) schedule from start to finish of the project, and storing the foregoing for retrieval.

Thus, the invention provides a data base management for and uses construction projects, with a time schedule as a key by which to relate indicia of sequentialized work activities and discrete increments of resulting structure, as well as other similarly related data. The resulting relational data base enables data associated with either the work activities or the structural increments to be related to the other by way of the time schedule—whether expressed in days, weeks, or other units of time—and enables the structural increments to be depicted with differential emphasis dependent upon urgencies of time or other important data item, such as cost.

FIG. 1 illustrates these and related steps in more detail. The usual preliminary task is to formulate the design of a given construction project, as in terms of specific design criteria. The intended construction is subdivided into successive stages, and corresponding structural drawings are prepared, subdivided into dis-

crete increments of structure. Likewise, kinds of work activity essential to the construction are identified and are broken down into discrete work activities. Optionally, as suggested within broken lines, related data may be associated with work activities (e.g., budgeted labor costs, crew sizes, and availability) or with structural increments (e.g., budgeted materials costs, quantities, and availability). Discrete structural increments and discrete work activities are collocated; that is, cross-referred to each other, in a one-to-one relationship.

Too, as noted in FIG. 1, durations are assigned to respective work activities, which are sequentialized among themselves in their logical order (i.e., which of them must precede which, and both early and late start and finish times are calculated for those activities. The resulting time data and the collocated discrete work activities and structural increments, together with other related data, are then used to generate a time (and data) schedule for the entire construction project from start to finish, whereupon the schedule and related information are stored for subsequent retrieval, as in electronic memory means.

It will be understood that, in generality, the foregoing steps can be carried out by hand, with pencil and paper, but that large construction projects are so complex that it is convenient to have the assistance of inanimate apparatus, such as a digital computer, capable of speeding up the necessary logical and mathematical calculations and, optionally, the actual drafting process, so as to complete them within a reasonable time. Storage in electronic memory means follows naturally upon such machine-assisted steps.

FIG. 2, in similar manner to FIG. 1, shows subsequent steps that, although also conceivably performable by hand, actually are reasonably feasible only with similar machine assistance. Indeed, FIG. 2 begins by showing the storing of information from FIG. 1 twice: once in a read-only memory, which represents the scheduled construction, and additionally in a read-write memory that can be updated from time to time in accordance with actual construction during the project. One or more intermediate types of memory (not shown), that may be overwritten only after due precautions are met may be interposed to represent target schedules bearing authorized changes from the original schedule, as will be readily apparent. However, in the interest of clarity no such embodiment is described further here or illustrated in the drawings, nor are means shown for storing various operating programs, or for enabling essential calculation to be performed, as is conventional.

As shown in FIG. 2, the read-write memory is updated during the project with whatever departures from the schedule actually occur in the process of construction. The operator may retrieve upon request from either memory a structural depiction, as of any selected time within the schedule, characterized by time urgency or other data so stored. Also, the two memories can be compared in any such regards to disclose any discrepancies between the two, as where actual construction is ahead of schedule (or under budget) or—usually more important—is behind schedule (or over budget), for example. Furthermore, either of the memories can be compared with one or more intermediate memories to identify any differences between original or current plans and interim target schedules—and to depict such differences in terms of the corresponding structure.

As indicated further in FIG. 2, according to this invention the structural increments of interest (or related to other items of interest) are depicted so as to emphasize their status, such as urgency, timeliness, or expense, as well as (if desired) annotated with related numerical data or verbal comments. Such emphasis is accomplished by what is called "shading" and more particularly as quasi-mnemonic "boldness" of shading to signify comparable (but unlike) characteristics in monochromatic and polychromatic video representations.

Because of black's optical dullness it is especially suitable for both monochromatic and polychromatic depiction of structural features already completed or otherwise not requiring emphasis. As shown subsequently, boldness and monochrome shading is conveniently represented by line thickness and continuity; thus, continuous or solid lines are bolder than broken ones, and thick lines are bolder than normally thin ones. In polychrome, shading is represented by hue and/or brightness of color. Structure overdue for completion takes the boldest representation, such as a double-thick line in monochrome, or the boldest color (e.g., red) in polychrome. Those structures with later due dates drop off incrementally to less bold colors, as in spectral order (e.g., orange, yellow green, forest green, sky blue, navy blue, violet—or dark brown); for which suitable monochrome analogs are long-dash, short-dash, and dash-dot double-thick lines, followed by long-dash, short-dash, and dash-dot normal or single-thickness lines, respectively. For convenience, these relationships are summarized in the table below, although it will be recognized that others may be substituted with like effect.

TABLE I

PRIORITIES	PRIORITIES AND SHADINGS	
	POLYCHROME	MONOCHROME
Greatest	Red	Continuous double-thick
Next greatest	Orange	Long-dash double-thick
Not so great	Yellow Green	Short-dash double-thick
Intermediate	Forest Green	Dash-dot double-thick
Less	Sky Blue	Long-dash single thickness
Even less	Navy Blue	Short-dash single thickness
Least	Dark Brown	Dash-dot single thickness

Such graduated boldness of shading of the structural increments themselves conveys at a glance language-independent information that otherwise could be conveyed if at all—only by relatively less effective methods. Of course, more or fewer gradations could be used, as could other shading methods in color or in monochrome. Users differ in preference for monochrome or polychrome emphasis of lines, but where area shading is desired the usual choice is color.

A useful supplementary representational aid is an intensity difference sufficient to be readily distinguishable from existing shadings. Whether continuous or intermittent, such differences are useful for emphasizing areas, lines, or points of particular interest. For example, features to be removed may be illustrated as above, but at reduced (or increased) intensity, all or part of the time.

A blinking or flashing of the video display, such as between normal intensities, is useful for noting the presence of slack or float time in an activity for constructing a given structural increment. Blinking or flashing is also analogously useful to show temporary features (e.g., forms) that are both placed and removed during the course of the project. The showing of such temporary

features preferably will alternate between the respective shadings characterizing their priorities of placement and removal, with the latter preferably diminished in intensity. Another use of blinking or flashing of the display, such as between normal and higher intensities, is to draw attention to structural features characterized by some data element (e.g., availability of work crews or materials, budgeted or actual costs or variances between them, or progress payments) not usually found in drawings. Verbal or graphical addenda may be included to similar effect, optionally with the aid of arrows, flags, or other pointers.

FIG. 3 shows computer apparatus 30 comprising components to enable this invention to be practiced with a high degree of machine assistance. Despite the largely schematic diagram here, persons ordinarily skilled in the art of digital computers will readily understand what is meant, so as to become enabled to practice this invention with relative ease. Shown centrally is the video display means, which conveniently comprises a video screen and optionally a printer capable of displaying both alphanumerics and graphics. A plotter or photographic device may replace or supplement the printer as a static or non-transient display device, whereas video screens or equivalent devices can provide a dynamic or transient display, such as blinking or alternating over time between more than one representation, which is sometimes desirable. Connected to the video means is electronic memory means including both read-only and read-write memories, as previously described. Also connected thereto are a central processing unit (CPU) and input/output control means, as is conventional.

The input/output means conveniently includes not only the customary keyboard but also an electronic drafting surface or "bit pad" to transmit into the memory drawings made thereon, as with the indicated stylus. Alternatively, a light pen may be used similarly in conjunction with the screen. Not shown is conventional scanning means that may be used with pen-and-ink construction drawings to convert them into readily stored digital signals. The foregoing apparatus may be of general-purpose type or may be dedicated especially to practicing this invention.

FIG. 3A shows schematically a number of computer workstations 30 connected together into communications network 35, which itself is useful according to this invention. One or more workstations 30 can readily be made available to major functional groups involved in any construction project (e.g., proposal, design, scheduling, procurement, construction, and startup); or trades (e.g., concrete, carpentry, electrical, plastering, and plumbing); and at levels of management from foremen through intermediate supervisory personnel to the project manager; and optionally even to higher line managers and possibly to pertinent staff personnel (e.g. financial, legal, or writers). This diagram fragmentarily illustrates representative availability of such apparatus in actual construction operations, specifically to the architect, designers, scheduler, contractor's managers, subcontractors, and customer. Of course, authority to revise the design or the schedule and update work performed and structure built should be limited to authorized persons. Great care should be taken for security of the data base.

It will be understood, of course, that such networking is well within the current state of the art and yet may take advantage of future improvements, inasmuch as the

proprietary nature of this invention lies more in what is done than in the specific means by which it is accomplished. The same is true of the computer means. Individual stations 30 preferably have self-contained computer units or "nodes" with their own memory banks and so interconnected that what is stored in the memory of each is available to every other unit, and that unavailability of any unit does not interrupt the network but only renders the unit and its memory unavailable until again on line. Also desirable are high resolution of display so as to show the structural increments in considerable detail, and rapid response because of the quantity of detail to be processed. A computer that meets these needs is readily available from Apollo Computer Inc. (of Chelmsford, Mass.) under the name DOMAIN.

FIG. 4 shows a project scheduling network diagram resulting from sequentialization of work activities for constructing a minor portion of a building as shown in subsequent diagrams. About a dozen activities (designated by arrows) suffice to go from early to late stages, such as excavation, forming, placing, and stripping of foundation, stairs, and landing. The respective activities are designated by An above the arrow, where n is a numerical indicator of a given work activity; and by (d) below the arrow, where d is the number of days from start to finish of such activity. Thus, the first activity arrow is marked A1/(5); the 1 indicating that it is the first, and the 5 indicating that it is to require five days.

Thus, FIG. 4 shows the critical path as a horizontal line from left to right, made up of arrows to which are juxtaposed activity indicia A1, A3, A4, A5, A6, A7, A9, A10, A12, and A13; whereas activities A2, A8, and A11 all require less time to complete than do one or more parallel activities and so have slack time available to them, as appears more definitely in the next view. Activities may be identified in more complex manner (e.g., with indicia of trade type or other information), but these simple indicia will suffice here.

Tabulated below are the work activities and their indicia.

TABLE 2

EXAMPLE OF WORK ACTIVITIES	
INDICIA	DESCRIPTIONS
A1	Excavate for foundation slab
A2	Place reinforcing for foundation slab
A3	Place forms for foundation slab
A4	Place foundation slab concrete
A5	Place reinforcing for walls
A6	Place forms for walls
A7	Place wall concrete
A8	Strip forms already placed, backfill
A9	Place forms for stairs and landing
A10	Place reinforcing for stairs
A11	Place reinforcing for landing
A12	Place stairs and landing concrete
A13	Strip remaining forms

FIG. 5 shows a resulting construction schedule, with a time scale (in days) at the top, and at the left a list of the various activities by indicia (An, where n is the number assigned to the given activity) plus duration (in days). Activity early start and late finish dates (days), respectively, are indicated by the left and right ends of the horizontal bars extending to the right of the respective activity indicia. Inside the bars are both forward and backward slant (or slash) characters, whose left-most ends indicate the respective early and late start times and whose right ends indicate the respective early and late finish times, as stated in the Legend included in

the upper right corner of the drawing. A bar completely cross-hatched by the oppositely slanted characters indicates the duration of an activity without slack, whereas a bar that is not completely cross-hatched indicates presence of slack. Thus, activity A2 has one day of float or slack, activity A8 has seventeen such days, and A11 has three of them. Such a bar chart is an accepted scheduling aid, but its graphical character is only indirectly (verbally) related to the intended resulting structure.

FIG. 6 shows in sectional elevation the corresponding portion of the scheduled construction project, comprising a plurality of structural increments uncharacterized as to the actual or scheduled order in which they were (or were to have been) completed, together with forms, reinforcing bars, etc. to be used in the construction. Included are a foundation, walls, stairs, and a landing—all rather fragmentarily shown—but adequate for an illustrative example. In more complex drawings, features to be shown might be grouped by type of work activity, location on a drawing, or by subcontractor, for example. As will be apparent, ordering or ranking of pictorial structural increments may be by starting times or completion times (or both can be shown alternately, or on two separate display devices simultaneously). For convenience, the examples here favor completion times as a matter of choice. Shown in normal continuous lines at the middle left and upper right are adjacent structural features not part of the present schedule.

FIG. 7 depicts the same structural portion (of FIG. 6) shaded in monochrome according to the scheduled activity completion times (of FIG. 5). Scheduled first as most urgent is the excavation for the foundation, shown accordingly in continuous double-thick lines (which in polychrome would be red, but of merely single thickness). Scheduled next are reinforcing bars—horizontal for the foundation and vertical for the walls—which appear in long-dash thick lines (orange if in color). Next, in short-dash double-thick lines (cf. yellow green) is the concrete foundation itself. Scheduled next are wall forms, shown in thick dot-dash lines (cf. forest green), which—because they are to be removed subsequently—should appear intermittently in diminished intensity, preferably coded to their removal deadline if in an electronic video display (but not here). Structural increments resulting from the next few activities are shown in long-dash single-thickness lines (cf. sky blue): the foundation walls (formed and placed, stripped and backfilled). However, because of the slack in the stripping activity, it may well appear alternately as an item of least urgency (dot-dash, cf. dark brown). Forms for both the stairs and the landing, and the reinforcement for the stairs, scheduled next, are shown in normally thick (thin) short-dash lines (cf. navy blue). The concrete stairs and landing, as well as final stripping, being of lesser priority, are also dot-dashed (cf. dark brown) normal single-thickness lines.

FIG. 8 shows in similar fashion the same structural portion, in its actual stage of construction on a given day (T29, the start of the second half), but shaded to emphasize construction lagging the schedule. Reference to FIG. 5 indicates that the first six of the activities were to have been completed by T28. A1–A5 are done, so their resultant structural increments are shown in continuous black single-width lines (similar if in color), confirming that the project is on schedule to that extent. However, FIG. 8 shows by its use of solid double-thick lines (cf. red) that A6 (placing forms for the walls) is behind schedule and, thus, very likely to delay subse-

quent activities, especially those having no scheduled slack (A7, placing wall concrete; A9, forms for stairs and landing; and A10, stairs and landing concrete). Such visual emphasis encourages prompt remedial measures, such as shortening the duration of one or more of the activities, as by enlarging work crews or paying them for overtime—doubtless at increased cost, which can be similarly highlighted in a separate depiction as an aid to decision-making.

In FIG. 8 slack-rich form-stripping step A8 is shaded like A13 (the one-day final form-stripping step) because of their identical completion dates, shown in FIG. 5. Activity A11 (placing landing reinforcing), which also has some slack, is shaded like the similar step of placing reinforcing for the landing because they also have identical required completion dates. According to this invention, as already noted, such slack could be shown readily in a video Display—not here—by flashing between the respective shadings for the activity's scheduled early start and permissible late finish. Two-stage flashing to indicate slack preferably differs from that for showing placing and removal of a structural feature by being of equal intensity in both stages of its showing (compared with a diminished intensity in one of the stages for a removal feature).

FIG. 9 shows schematically main data elements under headings: SCHEDULE, ACTIVITIES, EMPHASIS, STRUCTURAL INCREMENTS, and COSTS. The elements are shown in the form of alphabetical indicia, which in general form are T_n for schedule times, A_i for work activities, E_k for shading emphasis, S_j for structural increments, $\$C_i$ for crew costs, and $\$M_i$ for materials costs. In specific occurrences numbers are substituted for the appended lower-case letters, and in the drawing such numbers are more or less sequential. Mentioning of the data elements frequently by name (rather than by indicia) in further discussion of this diagram will be understood as compatible (rather than in conflict) with the principle that the indicia stand for their respective elements in this data base and are often meant when the context fails to indicate otherwise.

Interrelationships among the indicated data elements in FIG. 9 are shown by lines (of various degrees of boldness) joining them. Primary linkage between a given discrete work activity, on the one hand, and the increment(s) of structure produced by that activity, on the other hand, are represented in bold horizontal lines. Each activity has linked to it, as shown by ordinary lines, its duration D_i , its crew C_i and materials M_i requirements, and its description (. . .) $_i$ —all expressed here in general terms. Each structural increment has similarly linked to it: its graphical depiction data DD_j and usually also verbal description (. . .) $_j$. Correspondingly numbered descriptions of activities and structural increments may be merged into composite descriptions, if desired. A many-to-one relationship of structural increments to their source drawings is common, and here they all are linked to only one drawing (G1).

In FIG. 9 the duration of each work activity is linked by dashed lines at the left to the time schedule (shown fragmentarily as a succession of time units T1 to T12 . . . T $_n$). One such linking line appears for each time unit required by the activity duration. Early and late times—where not identical—would unduly complicate this diagram and, thus, are omitted here for clarity. Successive time units are bracketed into five-unit intervals for convenience of illustration and as an example of

a common range (a 5-day week). Other user-defined ranges may be substituted as the user may prefer for review of the project or schedule, and the degrees of emphasis, e.g., shorter ranges in the near term and longer ones further away. Here the listed activities (A1 through A4) and their durations are those of the same example treated in FIG. 5 and corresponding text.

The middle heading in FIG. 9 is EMPHASIS, a concept most often based upon time urgency but sometimes on some other connected data element, such as cost. Here emphasis indicia (for degrees of time urgency or data variance) E1, E2, E3, . . . Ek scale downward, linked to both the duration indicia and the time schedule by broken lines, to indicate that the linkage is subject to changes as time passes. Successive time units (e.g., days) may be deemed to define a fixed reference system, whereas emphasis can be expected to slide fairly regularly past time and relative to the various activities as well, though perhaps not so regularly in the instance of budgeted costs.

Also commonly in many-to-one relations are types (and numbers) of workers and materials, but it is convenient to cross-reference each work activity, in one-to-one manner, to its own crew and materials requirements (shown in FIG. 9 as Ci and Mi). It is apparent that the time schedule is linked to crews, materials, and cost data via the activity indicia (and durations). Crews and materials are linked directly to their costs—shown generally as \$Ci and \$Mi (at the lower right)—below their individual values and (at the upper right) their totals \$C and \$M. Of course, averages, dispersions, variances, other statistical and accounting measures, and the like can be calculated for various data categories.

In such a data base representation (as FIG. 9) any given time intersects, somewhat indirectly, at least one activity (usually many more) and structural increment(s), whereupon the array of time intersections with activities and structures constitutes a calendar of the project and may be reconstituted more directly as such—as suggested in dashed lines on the drawing. Moreover, as previously noted, the scheduled and actual construction times for the various structural features provide calendars for them in like manner, and comparison of the structures as scheduled and as actually built readily highlights departures of construction from the original schedule—which may have to be revised from time to time.

Placement of suitable relational links (shown in part in FIG. 9) in digital electronic computer memories (as via memory addresses) is well known in the art and is within the skill of software designers. Also within programming skill are the steps of converting drawn elements of structure to depiction in another medium such as a video display, shading such display with reference to a schedule or related time characteristics, comparing scheduled and actual items and determining differences between them, and annotating a display of structural features graphically with arrows or other symbols and optionally with textural information. The same is true of devising ways of varying such displays by diminished or enhanced intensity, blinking or flashing between normal and varied intensity (including on and off) and other visual modifications that will come to mind.

The drawings to be stored electronically need not be prepared differently from what draftsmen normally provide for construction projects; indeed, some customary annotation of relationships can be eliminated because they will be apparent in the video display. This is

not to say that improvements in draftsmanship will not come in handy in preparing drawings for use according to this invention.

Examples of improvements include performing the drafting with implements adapted to store a drawing as soon as it is made (or while it is being made), and utilizing prepared graphical or pictorial symbols for components that can be called up from storage rather than being drawn from scratch every time they are required, also keying the views so that they can be combined into more comprehensive illustrations—all as are becoming well known in the art of computer-aided drafting. These and other measures reduce an already large memory requirement and permit advantageous reduction in paper storage, use, and disposal.

It should be understood that “structural” is used generally in this specification to denote whatever makes up a constructed work (building, bridge, etc.) rather than being interpreted in a more limited structural engineering sense. Thus, electrical, plumbing, and finishing features of construction, for example, may well be treated as structural increments in the practice of the invention, as may collateral activities essential to a construction project, such as hiring personnel, buying or leasing equipment, expediting delivery, and receiving materials on site.

As is well known in the scheduling art, the “early start” date of any given work activity is the earliest date by which all essential prior activities can be (or have been) completed, and the “early finish” date is obtained by adding the minimum duration of that activity to its early start date. The “late finish” date of a given work activity is the latest date by which that activity can be completed without delaying completion of the project, and the “late start” date is obtained by subtracting the minimum duration of that activity from its late finish date. Time units other than days may be more suitable for some projects. The assigned duration for any work activity, if not a fixed number of time units, may be expressed in a range with a stated minimum and a stated maximum, with or without a stated intermediate, such as their mean duration, a defined “most likely” duration, etc. Software for handling these variants will come readily to mind for persons skilled in the art of computer programming.

Accordingly, programs as such do not constitute part of the invention claimed here, whereas marking for display and displaying timely structural features (with or without related data) are subject methods or steps of this invention. It is also apparent that a report generator (program) may be devised so as to enable a human operator to obtain one or more of a wide variety of temporary displays or permanent printouts at the touch of a key or (with voice-recognition equipment) just by an oral request. Indeed, computer programs or software devised or adapted especially for use in effecting the present objectives may be particularly useful—and may be patentable (or not).

Advantages and benefits of this invention have been mentioned above; other advantages may become self-evident to readers of this specification; and ultimately the greatest advantages and benefits can be expected to accrue to those who undertake to practice it. Prominent among the advantages and benefits are increased facility in project scheduling and control, decrease in required personnel and communications, and consequent reduction in costs all around. Some modifications in this invention have been noted above. Others may be made, as

by adding, deleting, combining, or subdividing parts or steps, while retaining advantages and benefits of the invention, which itself is defined in the following claims.

I claim:

1. In means to accomplish a building construction project provided with a schedule of discrete work activities to construct increments of structure of the project and provided with drawings of such structural increments,

means depicting such structural increments as of a given time in the schedule, and

means shading the depicted structural increments with degrees of boldness determined by time characteristics of the depicted structural increments.

2. Construction project means according to claim 1, including display means for depicting selected structural increments with differential boldness.

3. Construction project means according to claim 2, wherein the display means is monochromatic, and boldness of shading comprises lateral thickness in depicted lines.

4. Construction project means system according to claim 2, wherein the display means is polychromatic, and boldness of shading comprises hue of color in depicted lines or areas.

5. Construction project means according to claim 1, wherein the shading means is effective to shade most boldly the increments of structure due for completion by the given time in the schedule, to shade less boldly the increments of structure due for completion within a selected increment of time after the given time, and to shade even less boldly the increments of structure not due for completion until thereafter.

6. Construction project means according to claim 1, wherein the given time is a time during the actual project construction, wherein the shading means is effective to shade most boldly the increments of structure due for completion by the given time, to shade less boldly the increments of structure due for completion within a selected increment of time after the given time, and to shade with an unlike degree of boldness the increments of structure already completed.

7. Construction project means according to claim 1, wherein the given time is a time during the actual project construction, including means for determining the time relationship between the originally scheduled and currently projected completion times for the respective structural increments, and wherein the shading means is effective to shade more boldly the increments of structure whose currently projected completion times are later than as scheduled, and to shade less boldly the increments of structure whose currently projected completion times are as scheduled.

8. Construction project means according to claim 1, wherein the given time is a time during the actual project construction, and wherein the shading means is effective to shade more boldly the increments of structure whose currently projected costs are over the budget, and to shade less boldly the increments of structure whose currently projected costs are within the budget.

9. Construction project means according to claim 1, wherein the given time is a time after the actual construction of the project, and the shading means is effective to shade more boldly the increments of structure whose actual costs exceeded the budget, and to shade less boldly the increments of structure whose actual costs were within the budget.

10. Structural graphics representation system for construction projects, comprising operatively interconnected means for depicting discrete increments of structure for a given project,

memory means for retrievably storing data about activities essential to construction of the discrete structural increments and data representing durations of the respective work activities and their sequencing,

processing means to calculate therefrom early and late start and finish times for the respective activities and to calculate a resulting time schedule for the period from start to finish of the construction project,

input/output means for use in inputting data and instructions to the memory means and to the processing means and in selectively retrieving data therefrom for display including structural increment data to depict selected structural increments thereon,

and means for shading structural increments differently in accordance with their scheduled times of completion.

11. Structural graphics system according to claim 10, wherein instructions and data are retrievably stored in the memory means, for effecting diverse shading of these structural increments scheduled to have been completed before a selected time, those increments scheduled for completion then or within a selected time thereafter, and those increments not scheduled for completion until later.

12. Structural graphics system according to claim 11, wherein the stored instructions for time-dependent diverse shading prescribe the boldest diverse such shading for depiction of those structural increments due for completion at—or within the selected time after—the given time, and less bold such shading for increments not due until later.

13. Structural graphics system according to claim 12, wherein stored instructions for monochromatic shading prescribe boldness in terms of longitudinal continuity and lateral thickness of line.

14. Structural graphics system according to claim 12, wherein stored instructions for polychromatic shading prescribe boldness in terms of colors ordered in selected hues and brightnesses.

15. Structural graphics system according to claim 11, wherein stored instructions prescribe diminished intensity of depiction for a structural feature to be removed during the project.

16. Structural graphics system according to claim 11, wherein stored instructions prescribe intermittent switching of depiction, for a structural feature to be placed during the project and then removed later during the project, between one such diverse shading for the feature's placement time and another such diverse shading for its removal time.

17. In means for construction of a building, bridge, dam, industrial plant, or means of transport,

whereby a schedule of discrete work activities is prepared for constructing corresponding increments of structure of the project,

and whereby depictions of such resulting structural increments in construction drawings are also prepared,

the improvement comprising

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means for depicting the corresponding structure as in such construction drawings, as of a given time within the schedule,

including means for shading the respective structural increments with degrees of boldness corresponding to respective time relationships relative to the given time.

18. Construction project means according to claim 17, including means for displaying such depiction as a non-transient image.

19. Construction project means system according to claim 9, wherein a structural increment is scheduled for removal, and including means in the shading means for showing it shaded in diminished intensity.

20. Construction project means according to claim 18, including as display means a printer.

21. Construction project means according to claim 17, including means for displaying such depiction as a dynamic transient image.

22. Construction project means system according to claim 12, wherein a given structural increment is scheduled to be placed before being removed, and including means in the shading means for showing it alternately shaded according to its placement time and its removal time.

23. Construction project means according to claim 21, including as display means a video screen.

24. Means for use in construction of a building, bridge, dam, industrial plant, or means of transport, wherein work activities are performed to build increments of structure according to drawings thereof and a time schedule therefore, such means comprising

means for storing the schedule and the structural drawings for retrieval,

means for retrieving such structural drawings corresponding to the status of the project as of any given time in the schedule, and

means for shading the respective structural increments as of such given time with degrees of boldness corresponding to selected construction characteristics.

25. Construction project means according to claim 24, including in the storage means a read-write memory containing data defining the actual construction.

26. Construction project means according to claim 25, including means for comparing the read-write memory of the actual construction with the read-only memory of the scheduled construction in terms of the selected construction characteristics.

27. Construction project means according to claim 24, including in the storage means a read-only memory containing data defining the scheduled construction.

28. In means for construction of a building, bridge, dam, industrial plant, or means of transport, wherein certain work activities are performed to build discrete

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increments of structure according to a given time schedule, the improvement comprising

comparing scheduled and actual structure of such increments at a given time in the schedule during actual construction,

depicting most boldly incomplete structural increments scheduled for completion within a selected incremental period of time from the given time, and

altering the construction rate to bring the completion of incomplete structural increments into closer agreement with their scheduled construction times.

29. Project construction process according to claim 28, including enlarging work crews to increase the actual rate of construction.

30. In means for construction of a building, bridge, dam, industrial plant, or means of transport, wherein certain costs are budgeted and work activities are performed to build discrete increments of structure according to a given time schedule, the improvement comprising

comparing scheduled and actual structure of such increments at a given time in the schedule during actual construction,

depicting most boldly incomplete structural increments scheduled for completion within a selected incremental period of time from the given time but running over budgeted costs, and

altering the costs of work activities for incomplete structural increments to bring the costs of completion into closer agreement with their budgeted construction costs.

31. Project construction process according to claim 30, including reducing work crews to decrease their contribution to construction costs.

32. In accomplishing a building construction project wherein discrete work activities are performed to construct increments of structure of the project according to drawings of such structural increments and according to a construction schedule, the steps of

depicting as of a given time the existing extent of the structural increments under construction, shaded with more or less boldness to indicate more or less discrepancy between their actual extent and scheduled extent when behind schedule, and

takin appropriate steps to reduce such degree of discrepancy of the most boldly shaded structural increments.

33. Building construction method according to claim 32, wherein such appropriate steps include reducing the overall time to construct the most boldly shaded structural increments.

34. Building construction method according to claim 33, wherein such appropriate steps include increasing the actual man-hours beyond the man-hours scheduled for the work activities to construct the most boldly shaded structural increments.

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